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## Assignment 11

Due date: Thursday, 5 December 2024, 23:59

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### Exercise 11.4, Combining Proof Systems (★)

(8 Points)

Let

$$\Sigma = (\mathcal{S}, \mathcal{P}, \tau, \phi)$$

be a complete and sound proof system.

a) Define  $\mathcal{P}'$  and  $\phi'$  so that

$$\Sigma' = (\mathcal{S} \times \mathcal{S} \times \mathcal{S}, \mathcal{P}', \tau', \phi')$$

is a complete and sound proof system (and prove it!), where

$$\tau'((s_1, s_2, s_3)) = 1 \iff \text{at least 2 among } \tau(s_1), \tau(s_2), \tau(s_3) \text{ are equal to 1}$$

b) Let

$$\bar{\Sigma} = (\mathcal{S}^2, \bar{\mathcal{P}}, \bar{\tau}, \bar{\phi})$$

be a complete and sound proof system with

$$\begin{aligned} \bar{\tau}((s_1, s_2)) = 1 &\iff \text{exactly 1 of the statements is true in } \Sigma, \\ &\text{that is, } \tau(s_1) = 1 \text{ or } \tau(s_2) = 1, \text{ but not both.} \end{aligned} \tag{1}$$

Define  $\mathcal{P}^*$  and  $\phi^*$  so that  $\Sigma^* = (\mathcal{S}, \mathcal{P}^*, \tau^*, \phi^*)$  is a complete and sound proof system (and prove it!), where

$$\tau^*(s) = 1 \iff \tau(s) = 0$$

a)

Remark  $\mathcal{S}' = \mathcal{S} \times \mathcal{S} \times \mathcal{S} = \{(s_1, s_2, s_3) \mid \forall i \in \{1, 2, 3\} s_i \in \mathcal{S}\}$ , where  $|\mathcal{S}'| = |\mathcal{S}|^3$ .

Let

$$\begin{aligned} \mathcal{P}' &= \mathcal{P} \times \mathcal{P} \times \mathcal{P} = \{(p_1, p_2, p_3) \mid \forall i \in \{1, 2, 3\} p_i \in \mathcal{P}\} \\ \phi' : \mathcal{S}' \times \mathcal{P}' &\rightarrow \{0, 1\}, \quad s' = (s_1, s_2, s_3) \in \mathcal{S}', p' = (p_1, p_2, p_3) \in \mathcal{P}' \\ (s', p') &\mapsto \begin{cases} 1 & \text{if } \forall i \in \{1, 2, 3\} (\phi(s_i, p_i) = 1 \wedge \text{at least 2 of } \tau(s_i) = 1) \\ 0 & \text{otherwise} \end{cases} \end{aligned}$$

**Completeness:**

Assume  $\tau'(s') = 1$ . Then, at least 2 of  $\tau(s_1), \tau(s_2), \tau(s_3)$  must equal 1 by definition of  $\tau'$  in **a**).

By completeness of proof system  $\Sigma$ , every true statement has a proof, i.e.  $\forall i \in \{1, 2, 3\}, \exists p_i \in \mathcal{P}, \phi(s_i, p_i) = 1$ .

Hence, by grouping any 3 arbitrary statements in 3-tuples there must exist at least a corresponding arbitrary grouping of proofs in 3-tuples proving (at least 2) such statements, i.e.  $\forall s' \in \mathcal{S}', s' = (s_1, s_2, s_3), \exists p' \in \mathcal{P}', p' = (p_1, p_2, p_3) \wedge \phi'(s', p') = 1$ .

Thus,  $\Sigma'$  is complete.

**Soundness:** Assume  $\phi'(s', p') = 1$ . Then,  $\forall i \in \{1, 2, 3\} \phi(s_i, p_i) = 1$  and at least 2 of  $\tau(s_1), \tau(s_2), \tau(s_3)$  equal 1 by definition of  $\phi'$ .

By soundness of proof system  $\Sigma$ , no false statement has a valid proof proving it true, i.e.  $\forall i \in \{1, 2, 3\}, \forall p_i \in \mathcal{P}, \forall s_i \in \mathcal{S}, \phi(s_i, p_i) = 1 \implies \tau(s_i) = 1$ .

Hence, no proof for false statement exists, specifically satisfying  $\tau'$  definition of having at least 2 of  $\tau(s_i) = 1$ .

Thus,  $\Sigma'$  is sound.

**b)**

**Completeness: Soundness:**