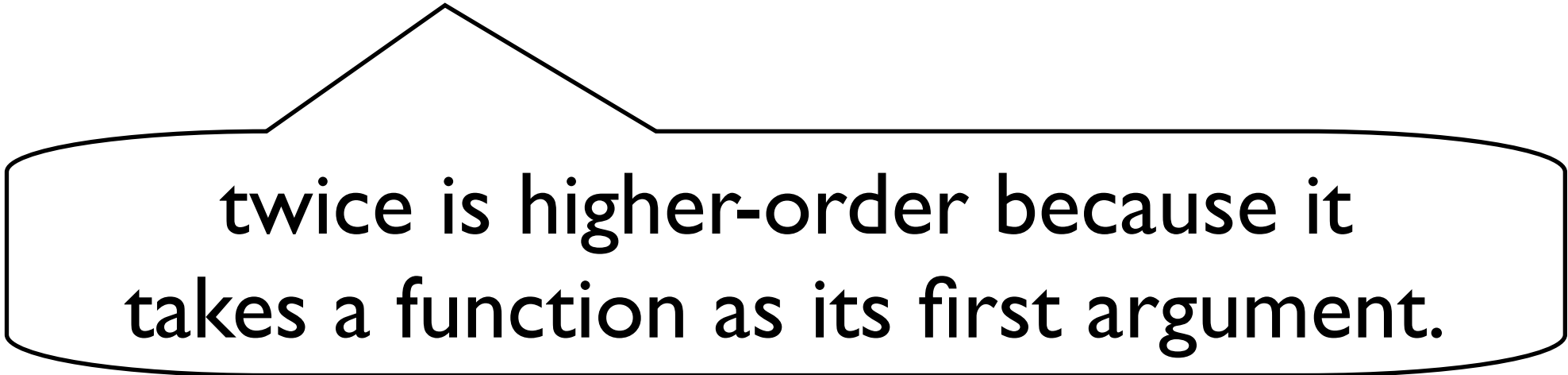


Higher Order Functions

A function is called higher-order if it takes a function as an argument or returns a function as a result.

```
twice      :: (a → a) → a → a  
twice f x = f (f x)
```



twice is higher-order because it takes a function as its first argument.

- Common programming idioms, such as applying a function twice, can be encapsulated as general purpose higher-order functions;
- Special purpose languages can be defined using higher-order functions, such as for list processing, interaction, or parsing;
- Algebraic properties of higher-order functions can be used to reason about programs.

The higher-order library function called map applies a function to every element of a list.

```
map :: (a → b) → [a] → [b]
```

For example:

```
> map (+1) [1,3,5,7]  
[2,4,6,8]
```

The map function can be defined in a particularly simple manner using a list comprehension:

```
map f xs = [f x | x ← xs]
```

Alternatively, using recursion:

```
map f []      = []  
map f (x:xs) = f x : map f xs
```

```
map :: (a -> b) -> [a] -> [b]
map f xs = [ f x | x <- xs]
```

```
map (+1) [1,3,5,7]
```

```
=
```

```
[(+1) x | x <- [1,3,5,7]]
```

```
=
```

```
[(+1) 1] ++ [(+1) 3] ++ [(+1) 5] ++ [(+1) 7]
```

```
=
```

```
[2] ++ [4] ++ [6] ++ [8]
```

```
=
```

```
[2,4,6,8]
```

$\text{map} :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]$
 $\text{map } f [] = []$
 $\text{map } f (x:xs) = f x : \text{map } f xs$

$\text{map } (+1) [1, 3, 5, 7]$
 $=$
 $\text{map } (+1) (1 : (3 : (5 : (7 : []))))$
 $=$
 $(+1) 1 : \text{map } (+1) (3 : (5 : (7 : [])))$
 $=$
 $(+1) 1 : ((+1) 3 : \text{map } (+1) (5 : (7 : [])))$
 $=$
 $(+1) 1 : ((+1) 3 : ((+1) 5 : \text{map } (+1) (7 : [])))$
 $=$
 $(+1) 1 : ((+1) 3 : ((+1) 5 : ((+1) 7 : \text{map } (+1) [])))$
 $=$
 $(+1) 1 : ((+1) 3 : ((+1) 5 : ((+1) 7 : [])))$
 $=$
 $2 : (4 : (6 : (8 : [])))$
 $=$
 $[2, 4, 6, 8]$

The higher-order library function filter selects every element from a list that satisfies a predicate.

```
filter :: (a → Bool) → [a] → [a]
```

For example:

```
> filter even [1..10]  
[2,4,6,8,10]
```


Filter can be defined using a list comprehension:

```
filter p xs = [x | x ← xs, p x]
```

Alternatively, it can be defined using recursion:

```
filter p []      = []  
filter p (x:xs)  =  
    | p x        = x : filter p xs  
    | otherwise  = filter p xs
```

A number of functions on lists can be defined using the following simple pattern of recursion:

$$\begin{aligned} f \ [] &= v \\ f \ (x:xs) &= x \oplus f \ xs \end{aligned}$$

f maps the empty list to a value v , and any non-empty list to a function \oplus applied to its head and f of its tail.

For example:

```
sum [] = 0  
sum (x:xs) = x + sum xs
```

$v = 0$
 $\oplus = +$

```
product [] = 1  
product (x:xs) = x * product xs
```

$v = 1$
 $\oplus = *$

```
and [] = True  
and (x:xs) = x && and xs
```

$v = \text{True}$
 $\oplus = \&\&$

The higher-order library function foldr (“fold right”) encapsulates this simple pattern of recursion, with the function \oplus and the value v as arguments.

For example:

```
sum      = foldr (+) 0  
  
product = foldr (*) 1  
  
and      = foldr (&&) True
```

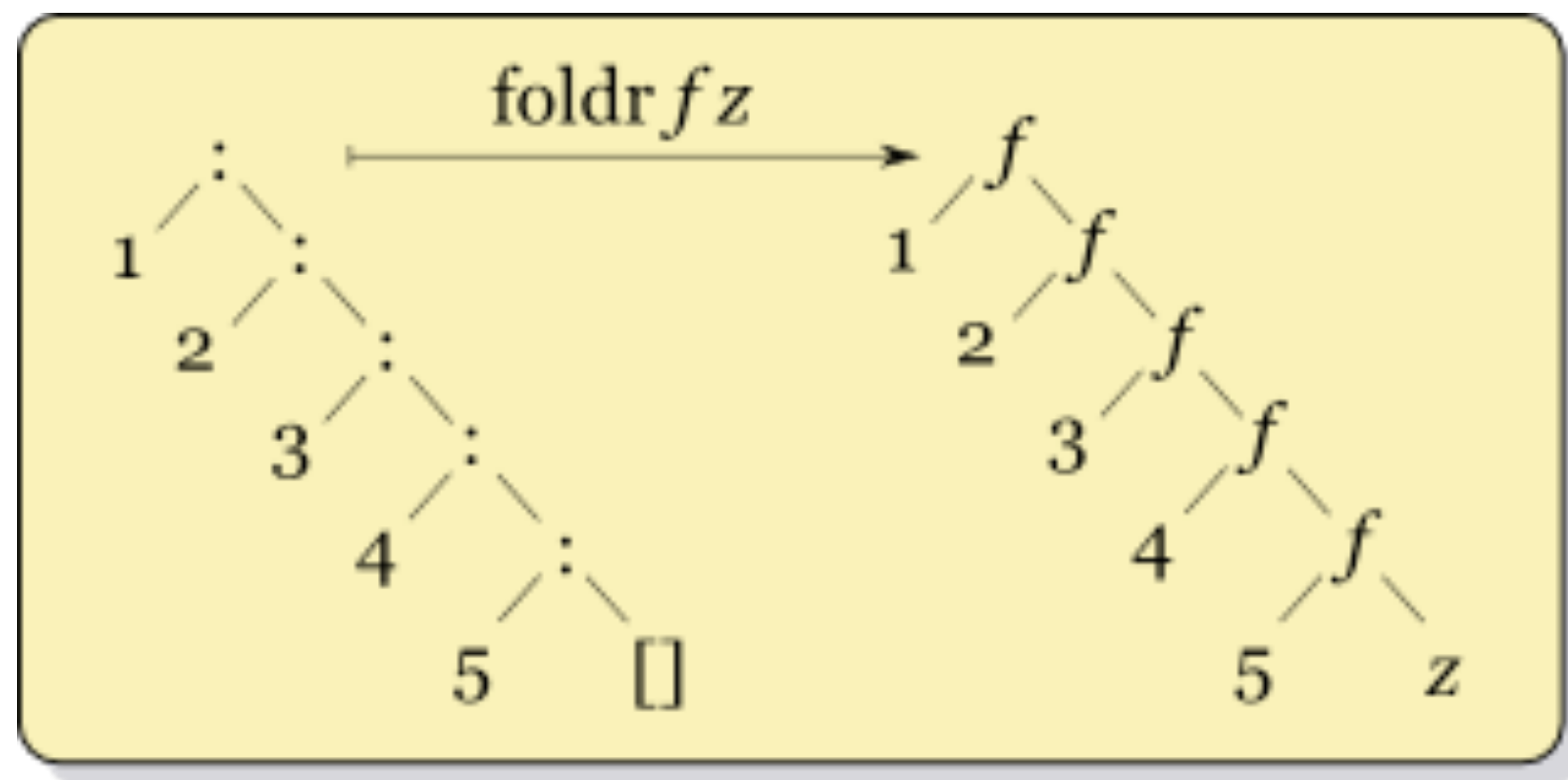
Foldr defined using recursion:

```
foldr (⊕) v []      = v
```

```
foldr (⊕) v (x:xs) =
```

```
  x ⊕ foldr (⊕) v xs
```

Think of `foldr` non-recursively, as simultaneously replacing each `cons` in a list by an infix function, and `[]` by a value.



For example:

`sum [1,2,3]`

`=`

`= foldr (+) 0`

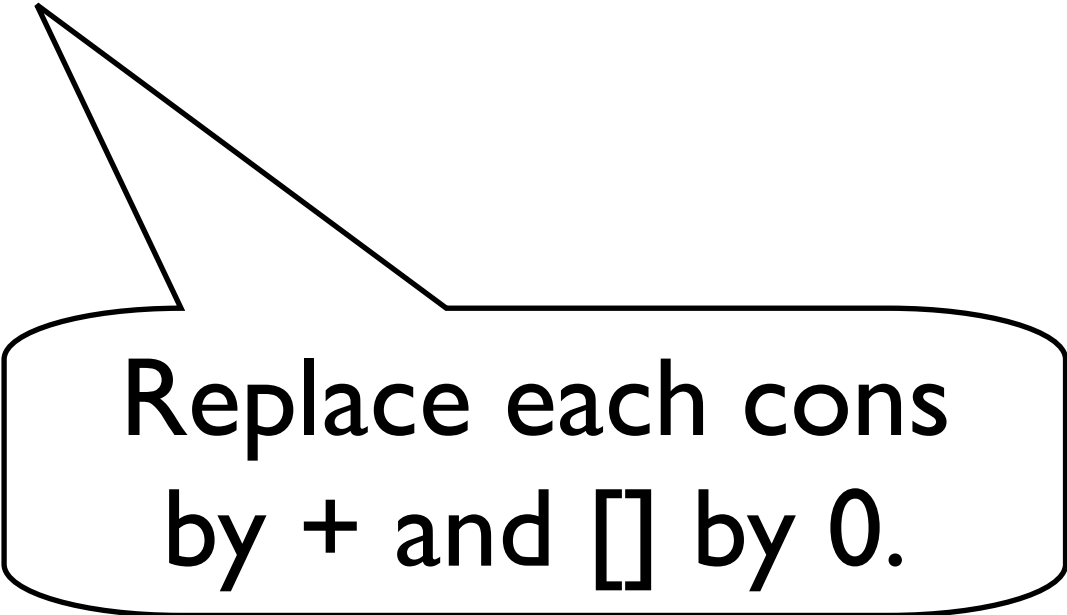
`foldr (+) 0 (1:(2:(3:[])))`

`=`

`1+(2+(3+0))`

`=`

`6`



Replace each cons
by + and [] by 0.

sum using foldr: how it works

```
sum [1,2]
=
foldr 0 [1,2]
=
foldr (+) 0 (1 : (2 : []))
=
1 + (foldr (+) 0 (2 : []))
=
1 + (2 + (foldr (+) 0 []))
=
1 + (2 + 0)
=
3
```

what is the type of foldr?

Type of foldr

```
foldr :: (a -> b -> b) -> b -> [a] -> b
```

Combining

Sum of squares of positive integers in a list

list
comprehension

```
f :: [Int] -> Int
f xs = sum [x*x | x<-xs, x>0]
```

recursion

```
f :: [Int] -> Int
f [] = 0
f (x:xs) | x>0 = (x*x) + f xs
          | otherwise = f xs
```

Higher order functions

```
f :: [Int] -> Int
f xs = foldr (+) 0 (map sqr (filter pos xs))
  where
    sqr x = x*x
    pos x = x > 0
```

*find the largest number under
100,000 that's divisible by 3829*

```
largestDivisible :: (Integral a) => a  
largestDivisible = head (filter p [100000,99999..])  
  
where p x = x `mod` 3829 == 0
```

What do these
evaluate to?

`foldr(:) [] xs`

`foldr(:) ys xs`

`foldr(:) [] xs`

Replaces “:” by “:”, and [] by [] --no change!

The result is equal to `xs`.

`foldr(:) ys(a:(b:(c:[])))`

`= a:(b:(c:ys))`

The result is `xs++ys`!

Folding an operation into a non-empty list

$$\text{foldr1} :: (a \rightarrow a \rightarrow a) \rightarrow [a] \rightarrow a$$
$$\text{foldr1 } f [x] = x$$
$$\text{foldr1 } f (x:xs) = f x (\text{foldr1 } f xs)$$

Operator: function composition

“do something, do something else”

$$(f . g) x = f (g x)$$

Note: $f . g x$ not the same as $f (g x)$

not all pairs can be composed
output type of g must be input type of f

$$(\cdot) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$$

input of f and output of g are of the same type: b
 $f \cdot g$ has input type a , same as g ,
and output type c , same as f

composition is associative

$$f \cdot (g \cdot h) = (f \cdot g) \cdot h$$

forward composition: $>.>$

order of composition is significant
 $(f . g)$ means first apply g then apply f

Can define an operator that
applies functions in the opposite order

$$(>.>) :: (a \rightarrow b) \rightarrow (b \rightarrow c) \rightarrow (a \rightarrow c)$$
$$g >.> f = f . g$$

the application operator: \$

application of function `f` to argument `e`

`f e`

can also be explicit

`f $ e`

can use as an alternative to parentheses

`flipV (flipH (rotate horse))`

becomes

`flipV $ flipH $ rotate horse`

can use as a function

`zipWith ($) [sum, product] [[1,2], [3,4]]`

Application and composition

suppose f has type $\text{Integer} \rightarrow \text{Bool}$

$f . x$ means f composed with x

so x must have type $s \rightarrow \text{Integer}$ for some type s

$f\ x$ means f applied to x so x must be of type Integer

$f\ \$\ x$ also means f applied to x so x has type Integer

Exercises

let `id` be the polymorphic identity function

`id x = x`

explain the behaviour of the following expressions:

`(id . f)` `(f . id)` `id f`

define a function `composeList` which composes
a list of functions into a single function type

give the type

explain why the function has this type

what is the effect on an empty list?

what is the type of the application operator, \$?

what is the result of
`zipWith ($) [sum, product] [[1,2], [3,4]]`?

explain the behaviour of the expressions

`(id $ f)` `(f $ id)` `id ($)`

lambda abstractions

`addOne x = x + 1`

could save the definition overhead and just use

`\x -> x + 1`

e.g. `map (\x -> x + 1) [3,4,5]`

`mapFuns fs x = map (\f -> f x) fs`

`comp2 :: (a -> b) -> (b -> b -> c) -> (a -> a -> c)`

`comp2 f g = (\x y -> g (f x) (f y))`

e.g.

`comp2 sq add 3 4`

Exercises

define a function `total`

`total :: (Integer -> Integer) -> (Integer -> Integer)`

so that `total f` is the function which at value `n` gives

$f\ \$\ 0 + f\ \$\ 1 + \dots f\ \$\ n$

given a function `f` of type `a -> b -> c`

write down a lambda abstraction that describes

a function of type `b -> a -> c` which behaves

like `f` but takes its arguments in the other order

define

`flip :: (a -> b -> c) -> (b -> a -> c)`

which reverses the order in which a function takes its arguments

partial application

```
multiply :: Int -> Int -> Int  
multiply x y = x * y
```

`multiply 2` given a number `y` returns `2 * y`

any function taking two or more arguments
can be partially applied to one or more arguments

```
doubleAll :: [Int] -> [Int]  
doubleAll = map (multiply 2)
```

partial applications can specialise general functions

Exercise

standard operators in Haskell are called operator sections
when partially applied

Find operator sections `sec1` and `sec2` so that
`map sec1 . filter sec2`
has the same effect as
`filter (>0) . map (+1)`

currying and uncurrying

```
multiply :: Int -> Int -> Int  
multiply x y = x * y
```

neater, permits partial application

```
multiplyUC :: (Int, Int) -> Int  
multiplyUC (x, y) = x * y
```

```
curry :: ((a, b) -> c) -> (a -> b -> c)  
curry g x y = g (x, y)
```

$\text{uncurry} :: (a \rightarrow b \rightarrow c) \rightarrow ((a, b) \rightarrow c)$
 $\text{uncurry } f \ (x, y) = f \ x \ y$

`curry` and `uncurry` are inverses of each other

inverse of
 $\text{unzip} :: [(a, b)] \rightarrow ([a], [b])$
is not `zip`

$\text{zip} :: [a] \rightarrow [b] \rightarrow [(a, b)]$
but

$\text{uncurry zip} :: ([a], [b]) \rightarrow [(a, b)]$

so

$\text{prop_zip } xs = \text{uncurry zip } \$ \text{unzip } xs == xs$

Exercises

what is the effect of `uncurry ($)`? what is its type?
what about `uncurry (:)`, `uncurry (.)`?

define functions

`curry3 :: ((a,b,c) -> d) -> (a -> b -> c -> d)`

`uncurry3 :: (a -> b -> c -> d) -> ((a,b,c) -> d)`