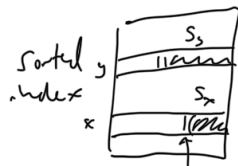


for  $x = 1$  to  $n$

for  $y = x+1$  to  $n$



$$S_x = \{z \mid d(z,x) < d(y,x)\}$$

$$S_y = \{z \mid d(z,y) < d(x,y)\}$$

compute these using sorted rows  
so  $U_{x,y} = S_x \cup S_y$

compute  $|U_{x,y}|$  by scanning through  $S_x$  and marking indices in bit vector  $b$  then scanning through  $S_y$  and marking indices in  $b$

for  $z \in S_x$

- ① update either  $C_{xz}$  or  $C_{yz}$
- ② set  $b_z$  back to zero

if  $d(z,x) < d(z,y)$   
 $C_{xz} += 1$   
else  
 $C_{yz} += 1$

for  $z \in S_y$

if  $b_z \neq 0$ ,

- ① update either  $C_{xz}$  or  $C_{yz}$
- ② set  $b_z$  back to zero

$$S_x = \{1, 3, 4\} \quad S_y = \{4, 5, 7, 8\}$$

$$b = [0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0]$$

1   2   3   4   5   6   7

$$|U_{x,y}| = |H| = 5$$

low cost for compare to  $O(n^3)$ ?

$$O(n^3) \rightarrow O(n^2 \log n + \sum_{x,y} |U_{x,y}|)$$