VPR Assessment of a Novel Partitioning Algorithm

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Abstract

Field Programmable Gate Array (FPGA) systems would be well suited to space based applications except for their vulnerability to space based radiation. Various techniques for dealing with their susceptibility have been discussed in literature. This thesis aims to implement and assess a key part of a theoretical technique to protect against radiation induced Single Event Upsets (SEUs) and assess the overheads of said technique.

Contents

VPR   Versatile Place and Route

MCNC   Microelectronics Centre of North Carolina

BLE   Basic Logic Element

CLB   Configurable Logic Block

DFG   Directed Flow Graphor a DataFlow Graph

SEU   Single Event Upset

LUT   Look Up Table

VTR   Verilog To Routing

STL   Standard Template Library

FPGA   Field Programmable Gate Array

TMR   Triple Modular Redundancy

BLIF   Berkeley Logic Interchange Format

ASIC   Application Specific Integrated Circuit

LAB   Logic Array Block

I/O   Input/Output

ICAP   Internal Configuration Access Port

SRAM   Static RAM

mux   Multiplexer

CAD   Computer Aided Design

MBU   Multi Bit Upset

NRE   Non Recurring Engineering

DICE   Dual Interlock Storage Cell

VHDL   VSIC Hardware Description Language

Chapter 1

Introduction

1.1 Overview

Space plays an increasingly important role in the functioning of modern societies, being vital for fields including navigation, meteorology, and communications [17]. Field Programmable Gate Array systems (FPGAs) have many beneficial features, such as their flexibility and low Non Recurring Engineering (NRE) costs which make them highly desirable for space based applications. Unfortunately they have far greater susceptibility to space radiation. Hardened Field Programmable Gate Arrays (FPGAs) offer only a fraction of the gate counts (and hence capability of implementing large or complex circuits) of non hardened offerings prompting a search for a solution to the radiation susceptibility of FPGAs using mainstream hardware [16], one of the most popular of which is Triple Modular Redundancy (TMR). In TMR, vulnerable components are triplicated allowing for errors to be detected and mitigated. This thesis is based on the work of [8] which introduces an approach to TMR, and aims to implement a key part of the approach and assess the implementation with the aid of an open source Computer Aided Design (CAD) toolchain for FPGAs. The remainder of this chapter provides an overview of these technologies, discusses alternatives to our approach, and details why we have chosen the technique we have. Chapter 2 introduces our approach to benchmarking circuits, and presents our initial results along with a brief discussion; Chapter 3 describes our implementation and design choices made in the implementation and Chapter 4 outlines our schedule and current progress, and our final chapter presents our closing remarks.

FPGAs

Field Programmable Gate Arrays (FPGAs) are popular devices capable of implementing a wide variety of circuits. Unlike Application Specific Integrated Circuits (ASICs) which must be specially designed and manufactured for an application—a lengthy and expensive process—FPGAs are a generic off the shelf device which can be mass produced by manufacturers and then adapted for an individual user's needs. Their flexibility, low cost, and faster development time make them the most economic for a number of applications.
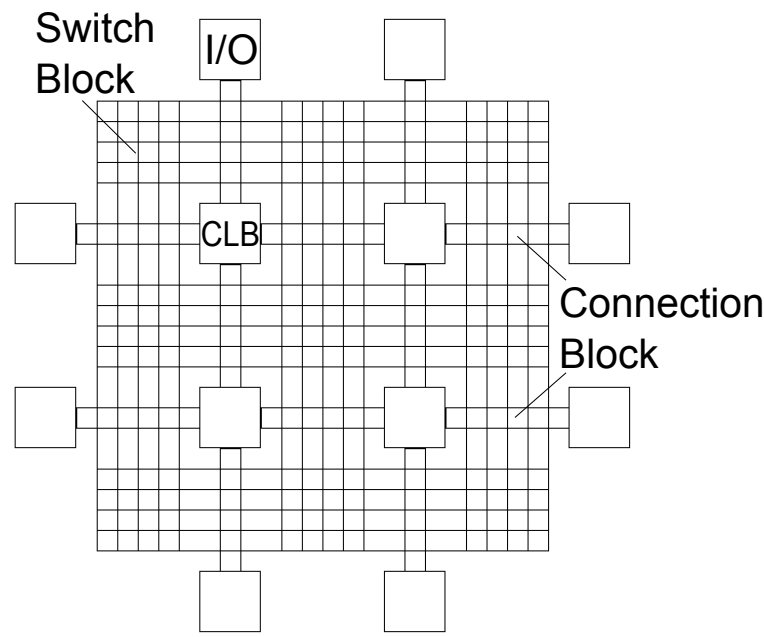
!

Figure 1.1: Island Style FPGA

TODO: Use own image. Wilton lecture notes have no license/copyright notice/etc attached, so don't know if this usage is actually allowed. Plus, doesn't look that good.

There are three main components to an FPGA: Input/Output (I/O) blocks, usually around the edge, allowing for input and output from the FPGA, Configurable Logic Blocks (CLBs) containing all the logic elements or *primitives*, and the routing between all the components. Most FPGAs also contain other structures embedded in the CLB array to provide commonly used resources such as multipliers. While they can be implemented using latches and Look Up Tables (LUTs), embedding them as discrete components allows for denser designs. The routing between components consists of channels running horizontally and vertically with a number of wires and programmable switches connecting the wires to each other and to CLBs allowing for configurable paths between arbitrary components. A typical switch or connection block has a configuration cell storing the state, and a connection can be made or unmade by writing a new value to the cell for that switch. The most common style of routing is known as island style (as the CLBs are located as islands in a sea of routing) with the routing area making up some 80%-90% of the FPGA's area [10]. Each CLB is a cluster of smaller blocks, called Basic Logic Elements (BLEs), with each BLE containing the logic primitives, typically a programmable LUT to implement combinational logic, a latch for register operations and implementing sequential logic, and a Multiplexer (mux) to switch between the two. The values for the LUT, whether the mux is selecting the latch or LUT output, and other component states are all stored in configuration memory like the routing switches and are typically implemented in SRAM.

Programming an FPGA involves loading in a bitstream which describes all the component values

(i.e. contents of the configuration memory for each cell) for a circuit, accomplished through writing the bitstream to a special configuration port on the FPGA. A number of FPGAs also allow for run time programming, or reconfiguration, of parts of a circuit through loading the bitstream for only the section of interest while the rest of the FPGA keeps running.

!

Check grammar in this next section. There are four main technologies used to implement the configuration memory in FPGAs:

- Static RAM (SRAM), which gives the highest density devices and includes the Virtex-5 family this thesis focuses on however they are volatile and must be reprogrammed every power up from an external configuration memory;

- (anti)fuse, which are only one time programmable;

- flash, which is non-volatile (thus not requiring an external configuration memory) and reprogrammable however has a lower density than SRAM based FPGAs [10].
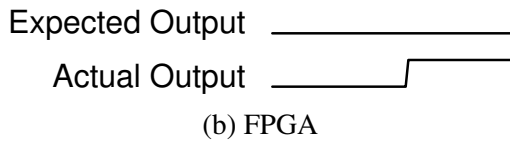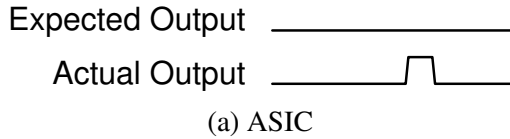
Partial Reconfiguration

Partial reconfiguration involves loading configuration information for part of a circuit during operation. Much like the complete configuration described above, it involves writing a configuration bitstream to one of the available configuration ports, in this case also including the location to reconfigure. The configuration memory of recent Virtex devices is split into frames, and one can only reconfigure entire frames. A configuration frame is 41 (32-bit) words long on a Virtex-5 device and configures a bit slice of the device that spans 20 CLB rows. As more frames are being reconfigured the larger the bitstream, and consequently the longer the time to reconfigure. The main configuration ports used are the external SelectMAP interface, or the internal Internal Configuration Access Port (ICAP), each with a bandwidth of 400MB/s in all Virtex devices [8, 22]

Space Based Applications

Space is quite different from a terrestrial environment, and FPGAs have a number of advantages due to their lower Non Recurring Engineering (NRE) costs and flexibility. As FPGAs can be reconfigured during a mission, faulty or outdated designs can be replaced remotely; however, there is a significant downside: as systems go further into space and are no longer protected by the earth's atmosphere they become increasingly likely to suffer from radiation induced errors where ionising radiation intersecting with a component causes charge build up, potentially triggering incorrect operation [20]. As outlined in Table 1.1, for higher orbits the mean time to upset is on the order of only a second, and this rate increases as technology advances and chip density further increases. Of the potential effects, which range from unnoticeable to device destruction, this thesis is concerned with mitigating Single Event Upsets (SEUs), where an incorrect signal is triggered but the underlying circuitry is not damaged. We

Orbit	SEUs per device/day	Mean time to upset (s)
LEO (560 km)	4.09	2.11×10^4
Polar (833 km)	1.49×10^4	5.81
GPS (20,200 km)	5.46×10^4	1.58
Geosynchronous (36,000 km)	6.20×10^4	1.39

Table 1.1: SEU Rate Predictions for Virtex-4 devices at various orbits [8]

Expected Output _____

Actual Output _____⊓_____

(a) ASIC

Expected Output _____

Actual Output _____⌐‾‾‾‾‾

(b) FPGA

also concern ourselves primarily with errors affecting only single bits or components rather than Multi Bit Upsets (MBUs) in which multiple components are affected at the same time. Keep in? In an ASIC, while SEUs may be picked up and latched or otherwise continue affecting the circuit in future, the component itself continues operating normally.

FPGAs on the other hand are vulnerable to configuration errors as well. When the charged particle impacts configuration memory it can flip the state of that cell changing the actual circuit. Unlike transient errors, these functional errors persist until corrected.

Additionally for SRAM devices, the off-chip configuration memory itself can be affected, so the next time the chip is reprogrammed (e.g. after power cycling), an incorrect circuit configuration will be loaded.

(Anti)fuse devices, being non reprogrammable, are immune to configuration errors, though both SRAM and flash based FPGAs are vulnerable and all three are susceptible to transient SEU [6].

How We Deal With FPGA Downsides

Clearly, in order for FPGAs to be viable in space based systems the effects of SEUs must be mitigated. A number of technologies and techniques are available, each with their own advantages and disadvantages. A number of options exist which detect errors but are unable to determine the correct result, requiring a reload of the configuration memory while the circuit is non operational until the reconfiguration completes. For many applications this downtime is impractical, thus we will be looking at options which allow the circuit to continue operating correctly. There are three main categories of SEU hardening techniques [4]:

- Charge Dissipation, which aims to keep the effect of the radiation below the level where it would

	Power (μW)	Speed (ns)	HARDNESS (e/b-d)	AREA (mm^2)
Standard	Rise – 0.7 Fall – 0.2	Rise – 0.21 Fall – 0.27	10^{-8} 1 node	360
Increased Drive Current	Rise – 1.0 Fall – 0.2	Rise – 0.16 Fall – 0.15	2×10^{-9} 1 node	460
TMR	Rise – 1.72 Fall – 1.27	Rise – 0.2 Fall – 0.27	10^{-11} 2 node	1200
DICE	Rise - 1.4 Fall - 1.1	Rise - 0.96 Fall - 0.97	1.6×10^{-10} 2 node	520

Table 1.2: Comparison of hardening techniques [4]

have an effect. This includes techniques such as increasing the drive current. These methods typically require custom hardware (increasing costs) and usually increase power usage.

- Temporal Filtering, which aims to filter out transient SEUs, includes methods such as delay-and-vote [4]. These techniques often slow down operation and are ineffective against configuration errors.

- Spatial Redundancy, which uses multiple redundant circuits to detect errors and be able to continue operating. Spatial redundancy techniques include Dual Interlock Storage Cell (DICE) [7] and Triple Modular Redundancy (TMR) and can be implemented either in hardware or at the design level not requiring any custom hardware. These methods typically increase area and power usage.

While hardened FPGAs are available, they typically lag well behind mainstream commercial offerings [16], thus solutions which can be implemented on mainstream commercial FPGA hardware are desirable. Additionally, there is very little point hardening an FPGA and not its configuration buffers and memory which take up far more surface area [10] and are thus even more vulnerable. For these reasons TMR, requiring no custom hardware and providing SEU protection against both transient and functional errors, is one of the more popular SEU hardening techniques even though it comes at the cost of more than tripling area and greatly increasing power usage. Table 1.2 details power usage, operating speed, hardness, and required area for flip flops which have been hardened using the techniques listing within the table. Explain columns One additional technique specific to SRAM based FPGAs relates to the protection of the off-chip configuration memory. As SRAM is volatile and loads the state from off chip at power up, this external configuration memory must also be protected from SEUs. This can be accomplished by incorporating error detection and correction techniques in the RAM, something already in place on a number of mainstream FPGAs such as the Virtex-4 and -5 [9].

1.2   Triple Modular Redundancy

Triple Modular Redundancy is a commonly used method for creating fault tolerant systems in which a given circuit is implemented three times with independent components, with the outputs feeding into a voter circuit to determine the majority value. Any SEU will affect the output value of at most one version, so the majority vote is still correct. For transient errors that are not in a feedback loop correcting the output is enough to fix the error; however, SEUs in feedback paths or in the configuration memory will persist, and this require some method for eliminating them. One possible approach is resetting the system but while this occurs the system is unavailable so a reset may not be a feasible solution. Instead, partial reconfiguration can reconfigure only the faulty circuit while the redundant circuits continue operating and providing output. After reconfiguration the circuit must then be resynchronised to the same state as the other two. We use the approach presented by [8] which involves running the circuit until the state converges, which is guaranteed (for acyclic circuits) to occur within a timeframe given by the number of register stages and the clock frequency. In order for this approach to always resynchronise correctly the circuit must have no feedback loops which may carry incorrect data. To solve this we simply ensure that all feedback loops are *cut*, that is, the value is voted on before being passed back into the circuit. This has the additional benefit of correcting transient errors which would otherwise be caught in a feedback cycle by ensuring the cycle data is correct.

This approach requires three times as many circuit elements (as the circuits are triplicated) plus whatever is required for voters. By minimising the number of voters, we can thus reduce the overhead of our approach.

Once an error occurs it takes up to T_{path} to reach the voter and be detected, where T_{path} is given by the clock period and number of register stages. This is called the *error detection time*. Detection of an error can then be used to trigger reconfiguration. Sending a request to the reconfiguration controller goes through a token ring network consisting of the other voters, and the reconfiguration controller. In the worst case it takes one full cycle of the network to receive the token, one full cycle to reach the reconfiguration controller, and three cycles to transmit the request, giving $5 \times TimePerHop$. Benchmarks of a sample voter indicate 50 clock cycles per hop is a good estimate. *Reconfiguration time $= T_R$* is dependent upon the circuit size. For a Virtex-5 device, each reconfiguration area consists of 160 LUTs and 160 latches, where only whole reconfiguration areas can be reconfigured. The bitstream size for this area is Fill in which takes $15.4\mu s$ to reconfigure at $100MHz$. Once the error has been detected and the circuit reconfigured it must then be resynchronised with the other partitions, which takes up to T_{path} using the previously described technique.

The error recovery time consists of the time to detect the error, send a request to the configuration controller, and then reconfigure and resynchronise the circuit, thus is a function of the circuit area, clock frequency, and number of register stages. Therefore it is required that the number of register stages and area are small enough, and frequency large enough, that our error recovery time is within a user specified

limit.

$$\text{Error Recovery Time} = \text{Error Detection Time} + \text{Reconfiguration Time} + \text{Resynchronisation Time}$$

$$\text{Error Detection Time} <= T_{path} = \text{Clock Period} \times \text{Register Stages}$$

$$\text{Communication Time} = \text{Cycles per flit per hop} \times \text{Number of flits} \times \text{NumStops} \times \text{Clock Period}$$

$$\text{Communication Time} = 50 \times 5 \times (\text{NumPartitions} + 1) \times \text{Clock Period}$$

$$\text{Reconfiguration Time} = \text{Clock Period} \times \frac{\text{Bitstream Size}}{\text{Reconfiguration Speed}}$$

$$= \text{Clock Period} \times \left\lceil \frac{\max(numLUTs, numLatches)}{160} \right\rceil \times 1.54 \times 10^{-5}$$

$$\text{Resynchronisation Time} <= T_{path} = \text{Clock Period} \times \text{Register Stages}$$

$$(1.1)$$

[8]

Additionally, as each voter circuit adds some constant overhead in terms of area, power usage and clock frequency slowdown it is desirable to have each partition as large as possible. This thesis is concerned with implementing and assessing this TMR design; a discussion of other TMR methods and our reasons for not using them is included below.

This method only works when at most one SEU occurs within the error detection and recovery time; should SEUs occur in two of the three partitions then it is impossible for the voter to determine the correct value necessitating a complete reload of the configuration memory (*scrubbing*). Therefore, we require the error detection and recovery time to be sufficiently small that the likelihood of multiple events occurring within that time period are sufficiently small.

Additionally, as mentioned earlier, it is also desirable to minimise the number of voters to reduce the overhead of this approach. To that end, having larger partitions (and hence fewer) is preferable to smaller provided we still stay within our target recovery time.

Triple Modular Redundancy Implementations

This thesis builds on the work of [8] which details a partitioning algorithm that traverses a circuit represented as a Directed Flow Graph (DFG) in a breadth first manner, creating partitions that stay within our constraints. Our goal is to create an algorithm which stays within a user specified error recovery time, doesn't require existing code to be rewritten, allows for both custom voting and reconfiguration logic to be added, can use industry standard FPGAs rather than custom hardware, and effectively protects the entire system from SEUs with as close to no downtime as achievable. Additionally, it is desirable to limit the overhead of implementing TMR through minimising the number of voters required. There are a number of existing TMR solutions, however none quite meet our requirements. Our first requirement is that standard FPGA hardware can be used, with our implementation specifically targeting Virtex-5 chips. Options with custom hardware such as [16] (with three FPGAs and an ASIC voter in a single package), are often
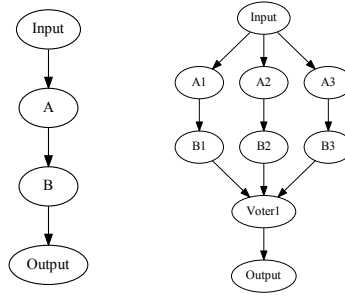
Figure 1.2: DFG before and after partitioning

prohibitively expensive, and prevent us from using our existing boards. Many FPGAs marketed specifically at space based applications are, in addition to featuring specialised hardware, only latchupexplain?
immune or only include inbuilt TMR on registers, leaving them still vulnerable to SEUs [12]. Non hardware solutions are typically implemented pre-synthesis, such as [11] (which introduces a VHDL library featuring triplicated components), and require existing code to be rewritten, or during synthesis such as [1] and [2] which support neither specifying an error recovery limit, nor for adding reconfiguration logic. Other options look at using partial TMR (such as [18]) which, while it does reduce the overhead of TMR, means the entire circuit is no longer protected, or have excessive downtimes to recover from errors such as [3], which uses idle cycles in a data path to calculate redundant results. One approach similar to ours is presented by [13] who also partition a post-synthesis netlist (represented by a DFG), but their focus is on evaluating techniques for cutting feedback loops, while we focus on partitioning circuits into smaller sub circuits.

Our Algorithm

Given a netlist description of a circuit, it is possible to represent the circuit as a DFG [10]. Our goal is to split a DFG into a number of smaller subgraphs, triplicate the components of each subgraph, and insert voting and recovery logic, with each subgraph having independent components and an error recovery time within our threshold. We can then proceed to implement our graph, made up of our new subgraphs, as normal. To do so we traverse the DFG in a depth first manner, keeping track of the number of register stages, area, and maximum frequency, extending our partition area as we do so, until our recovery time constraint would be violated. At that point we cleave off our partition and write it to a file, open a new empty partition, and repeat for all circuit elements. Once this is done, we have a set of subcircuits. We now triplicate each partition and insert our additional voting logic, then join each subcircuit back together. While doing so we must make sure that no loops exist within a partition and that all values are voted on before being reused, as otherwise the circuit may not resynchronise. This is accomplished by making sure that each node is only added once, and when inserting the voting logic that all outputs are voted on

before being used as inputs.

1.3 CAD Flow

FPGAs are typically programmed in a higher level description language such as VHDL or Verilog, and then a number of programs (collectively making up the Computer Aided Design (CAD) flow or development toolchain) turn the source into a bitstream to program a target FPGA. The design flow process can be split into a number of sub processes as illustrated in Figure 1.3 [5, 10, 15].

1. The synthesiser turns a hardware description language such as VHDL or Verilog into a netlist of basic gates and flip flops.

2. The optimiser removes redundant logic, and attempts to simplify logic.

3. The mapper maps logic elements to primitives, the basic logic elements contained on the FPGA.

4. The packer combines logic elements into CLBs.

5. The placer locates each CLB within the FPGA architecture, deciding which physical block implements which logic block.

6. The router makes the required connections between each element by deciding which switches are on or off. This includes the connections within each CLB (local routing) and in between CLBs (global routing).

For our partitioner we will insert an additional step into the design flow between mapping and packing, which operates directly on a netlist. The additional steps are described more fully in section Reference. !

How VPR Works

For this thesis we will be assessing the results of our algorithm implementation after processing by Versatile Place and Route (VPR), an open source packer, placer and router. VPR was chosen as it is open source allowing modifications to be made if necessary, and it is well documented and popular in research, making it much easier for us to determine what's happening and why, rather than relying on proprietary black box processes from commercial vendors. A brief understanding of the algorithms used in VPR and the effects of different settings is useful, though not critical, for understanding the results. [15] has a more detailed list of all the options VPR takes. Unless otherwise specified, all values are at their defaults.

Packer

VPR uses the AAPack algorithm described by [14]. This is a greedy algorithm which operates on blocks sequentially, starting with an FPGA area of 1 block by 1 block. For each block it greedily adds primitives
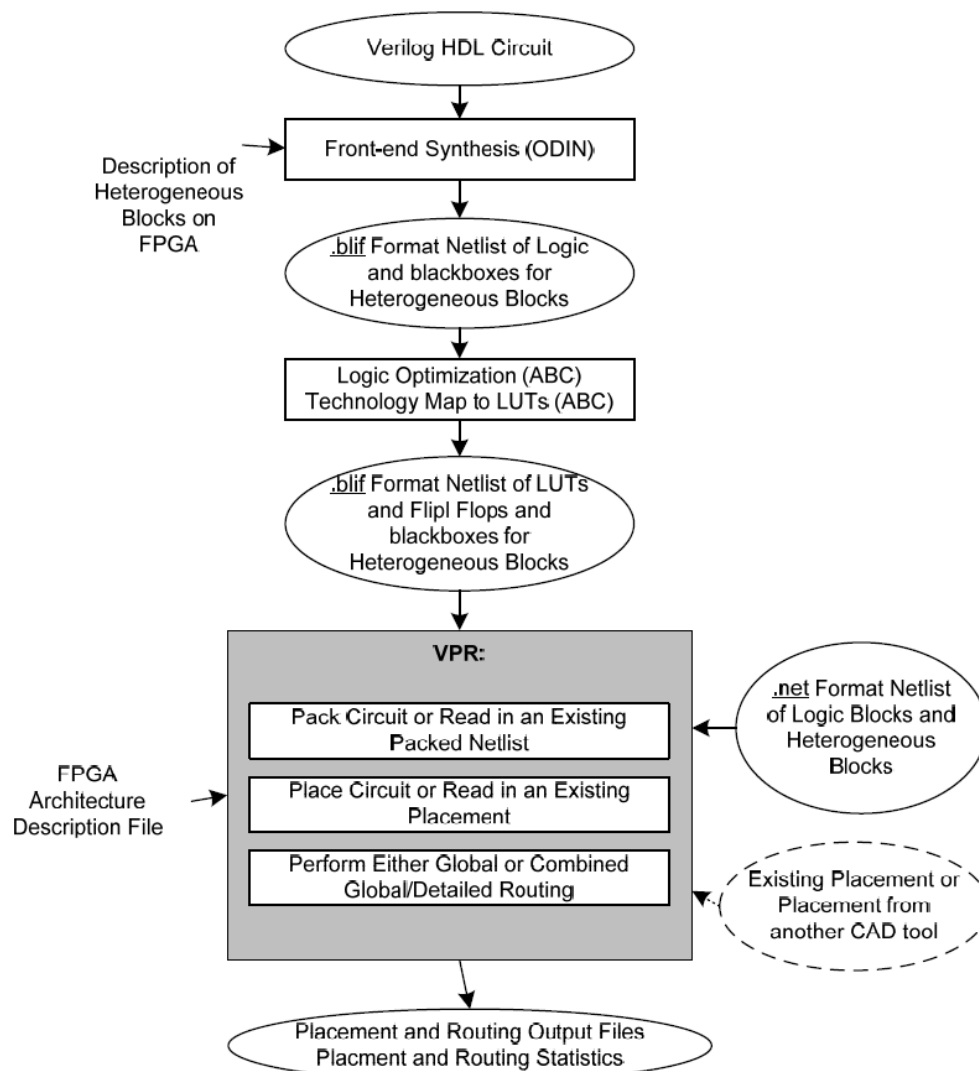
Figure 1.3: Cad Design Flow. [15]

based on a configurable cost function until no more primitives can be added. It then repeats for the next block, and the next after that, until every primitive has been packed. As it runs out of blocks in the current FPGA area it expands the FPGA area used until it reaches the physical limit specified in the architecture file (or grows indefinitely if no limit is specified). This means that even if the device is of area 40 by 40, if the packer can fit everything in a 30 by 30 area it will do so, and VPR will treat the FPGA as being only 30 by 30. The cost function can be configured through options passed to VPR, to [15]:

- prioritise optimisation of timing or area (default is prefer timing)

- prioritise absorbing nets with fewer connections over those with more (default is yes)

- when prioritising absorbing nets with fewer connections, focus more on signal sharing or absorbing smaller nets (default is greatly prefer absorbing smaller nets)

- determine the next complex block to pack based on timing or number of inputs (default is timing).

The main thing to note, as relates to our results, is that as much as possible AAPack will never leave blocks partially packed while there is still a primitive which will fit. Even when optimising timing exclusively, it will still attempt to maximally pack each cluster.

Placer

VPR's placer uses a simulated annealing algorithm where the options allow us to specify annealing schedule parameters and cost function. The default options were chosen via experimentation and are likely superior to custom options we may choose to use, and affect the quality of the result rather than materially affecting the behaviour [5, 15]. For these reasons we will be leaving them at their default. Section Reference discusses the variation in results due to the stochastic nature of the placement algorithm. !

Router

VPR's router supports three different algorithms:Awkward phrasing, fix breadth_first, which focuses solely on routing a design; timing_driven, the default, which tends to use slightly more tracks (5%) than breadth_first while providing much faster routes ($2\times-10\times$) with less CPU time; and directed_search, which like breadth_first is routability driven however uses A* to improve runtime. We will be using the default timing_driven algorithm.TODO: Reason? There are a number of options setting algorithm parameters, all of which we will leave at their defaults, though we will be changing the route_chan_width parameter as we collect results. route_chan_width specifies the width of the channels in the architecture. If omitted VPR will perform a binary search on channel capacity to determine the minimum channel width.

Chapter 2

Algorithm

For our partitioner, we operate on a netlist after optimisation and technology mapping, but before packing. Our goal is to take an input netlist and transform it into a netlist in the same format, with the same set of outputs for each set of inputs, but with redundant components.

!

Diagram of tool flow

- Partition

- Triplicate

- Join

- Flatten - Use abc to transform our heirarchical circuit into a format supported by VPR.

- Test - Use the verification abilities of abc, described more fully by reference. Essentially, given the original circuit and our transformed circuit, verify that every set of inputs produces the same set of outputs in both circuits.

!

Chapter 3

Results

3.1 Benchmarking Procedure

These results were collected by running benchmark circuits through an automated test suite written in Python by the author. For each benchmark circuit, and each target recovery time, 10 repetitions were performed, due to the variation in reported values due to the random nature of placement. The original circuit is run through VPR to collect base results, then the circuit is run through our partitioner to TMR it. The TMR'd version is then verified by abc as being functionally equivalent to the orignal, and then run through VPR to collect TMR'd results. Each run of VPR used a randomly generated seed for the placer. The mean of the reported values across the 10 repetitions was then taken. The benchmarks used were the 20 largest MCNC LGSynth93 circuits, as provided by the open source VTR project Reference and described in table Reference. The set of target recovery times used is $10^{-3}, 10^{-4}, 10^{-5}$s.

3.2 Sanity Check

The following results are for the tseng circuit at a target recovery time of 10^{-4}s. The reported values can be compared to each other as a manual sanity check. Compare Area usage. Compare Recovery Time. Compare Partition Size.

Chapter 4

Benchmarking

4.1 Overview

For benchmarking we use 20 MCNC benchmarks (detailed in Table 4.1) obtained as part of the Verilog To Routing (VTR) project[1] which we will be using to evaluate the performance of our partitioner and our TMR scheme in general. Additionally, we're looking for ways of estimating area usage and timing information from a Berkeley Logic Interchange Format (BLIF) file or DFG, without needing to actually place and route the partial circuit after each iteration, as doing so is computationally prohibitive. Here to the end of section gets rewritten To start with we made simple test circuits to compare to our benchmarks by triplicating each entire benchmark and adding in simple voter logic. As progress is made on the partitioner we can start collecting results from further partitioned circuits; however triplicating the entire circuit should be sufficient for rough approximations provided $elements_{circuit} >> elements_{voter}$.

To make the test circuits we created a small Python script which, given an input circuit and input voter circuit, triplicates the circuit and adds voting logic. It creates a hierarchical BLIF file, that is, it contains nested subcircuits, which are then passed through an external program called SIS[2] to flatten it into a format VPR can read.

Expected Results

As is well described in literature (e.g. [4]), the area usage should increase by a factor of slightly more than three. There are three copies of each component, plus additional components for the voting circuitry. Maximum frequency is expected to decrease slightly due to the additional components increasing wire length and crowding routing channels however this is likely to vary depending on circuit.

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[1]http://code.google.com/p/vtr-verilog-to-routing/

[2]Available from http://www1.cs.columbia.edu/~cs4861/sis/

Name	Number of:			
	Inputs	Outputs	Latches	LUTs
alu4	14	8	0	4574
apex2	38	3	0	5637
apex4	9	19	0	3805
bigkey	229	197	672	5294
clma	62	82	99	25177
des	256	245	0	5018
diffeq	64	39	1131	4521
dsip	229	197	672	4283
elliptic	131	114	3366	10920
ex1010	10	10	0	13804
ex5p	8	63	0	3255
frisc	20	116	2658	10733
misex3	14	14	0	4205
pdc	16	40	0	13765
s298	4	6	24	5796
s38417	29	106	4389	18232
s38584.1	38	304	3780	18835
seq	41	35	0	5285
spla	16	46	0	11116
tseng	52	122	1155	3260

Table 4.1: Benchmark circuits used

4.2   Architecture

VPR allows us to specify a custom architecture for it to run against in an XML format. Initially we are keeping the default architecture detailed in [15] consisting of a grid of CLBs each consisting of ten fully interconnected BLEs, and each BLE having a latch and 6-LUT as ilustrated in Figure 4.1. Table 4.2 details the number of primitives per CLB. Primarily of interest is that each BLE has 6 inputs and 1 output and each CLB has 33 inputs and 10 outputs.

4.3   Methodology

To start with, we wanted to collect rough estimates on the impact of partitioning a circuit to provide a baseline with which to compare our partitioning algorithm, and to develop the rough estimates needed for our partitioning algorithm. To that end we first created a simple Python script to take an arbitrary input circuit, triplicate it, and insert arbitrary voter logic. These triplicated circuits were then placed and routed by VPR, as were the original benchmarks, and the results compared.
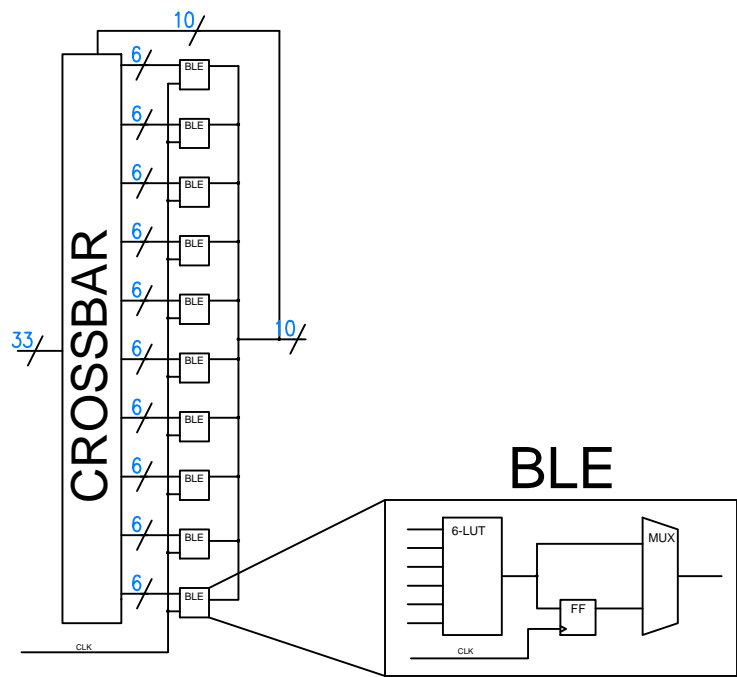
Figure 4.1: CLB Architecture

Component	Number	Notes
Flip Flop	1 per BLE	Shown as FF on Diagram
6-LUT	1 per BLE	
MUX	1 per BLE	
BLE	10 per CLB	
Crossbar	1 per CLB	
CLB	Autosized by VPR	

Table 4.2: Architecture Elements

Width	Num. Latches	Num. LUTs	FPGA Width	Channel Width	Num. Wire Segments	Used Area	Frequency	CPU Time
Auto Width	300%	301%	169%	119%	106%	301%	93%	405%
200 Width	300%	301%	169%	N/A	110%	301%	85%	385%
60 Width	300%	301%	169%	N/A	113%	302%	86%	444%

Table 4.3: Median values for specified channel widths as a percentage of non TMR equivalent

Name	Num. Latches	Num. LUTs	FPGA Width	Num. Wire Segments	Used Area	Frequency	CPU Time
pdc 200 width	N/A	300%	176%	104%	301%	75%	425%
tseng 200 width	300%	312%	169%	126%	308%	102%	388%

Table 4.4: Increase in circuits with maximum and minimum frequency slowdown

Just to fix formatting. Temporary VPR is used with our architecture file (described in Section 4.2) and
the command line options

```
1  VPR architecture.xml circuit.blif --full_stats[ --route_chan_width x]
```

where x is the width of the routing channels and - -full_stats tells VPR to be more verbose in its output. As
mentioned earlier, if - -route_chan_width is excluded then VPR determines the minimum channel width
needed to successfully route the circuit [15]. We then place and route our benchmark circuits and our
partitioned circuits. VPR itself then reports the area usage, critical path time (inverse of the frequency),
and other statistics we analysed. Unfortunately VPR does not report the number of register stages however
our partitioner will, as it needs to calculate the number of steps for its time estimation function.

4.4   Results

The results listed in Tables 4.3 and 4.4 highlight information of interest in a few key circuits, rather than
including page after page of tables. Any aggregate statistics (e.g. median) are calculated on the entire
result set, not just the results included. No collected results were considered to be outliers and the only
exclusions from the median calculation are those where the operation was not able to be completed,
so there is no data. These tables list the median characteristics of the TMR versions of our benchmark
circuits as a percentage of the original (100% is the same, 200% is double, 50% is half); additionally,
they detail the time taken to pack, place and route.

4.5 Discussion

Our simple voter circuit consists of one 3-LUT per output. Therefore we expect the number of logic elements (latches and combinational logic) to be exactly three times larger, with an additional 3-LUT per original circuit output. As shown in Table 4.3 our triplicated circuits are just slightly over three times as large. Circuit area should be roughly tripled as well, which again, matches, with the width increasing by $1.69 \approx \sqrt{3}$ and the used area increasing by just over triple. The partitioned circuits require slightly larger channels, in order to route the extra wires needed, and the additional elements and wires lead to slightly more segments per wire, and a slightly lower maximum frequency. Of note is that the time to place and route the partitioned circuits was much higher, taking around four times longer.

Circuits were 15% slower on average, with a worst case of 25% overall and a best case (in 200 width circuits) of a 2% speedup. For minimum channel routing the smallest median slowdown was observed as on average the minimum channel width needed for partitioned versions was higher, giving the router more flexibility. For 60 width circuits some were unable to be routed, therefore the results do not represent all circuits. For those reasons the reported statistics are taken from the 200 width circuits.

The speedup, while small, is unusual. It is likely due to the packer having more options, and due to the larger available area (as the packer will only increase the number of CLBs, and hence FPGA area used, when it can no longer fit new primitives in the current block).

Chapter 5

Partitioning Algorithm

!

5.1   Overview

In this section we discuss the partitioning algorithm, including how we are implementing it, progress, and the reasoning behind design choices made.

5.2   Design

We will be implementing the algorithm introduced in Section 1.2, in which we traverse a DFG in a breadth first manner splitting the DFG into subgraphs which we triplicate and insert voter logic into, while ensuring the error recovery time for each subgraph is below a threshold.

Design Choices

As much as possible, we would like our implementation to be easily extensible to multiple architectures. The actual partitioner operates on a DFG so it can be mostly architecture agnostic, only requiring the estimation functions to be architecture aware. We already have python scripts written to create our benchmark circuits which are able to manipulate BLIF files, so we opted for a toolchain incorporating them to reduce development time before we have a working implementation. Specifically, our partitioner operates on BLIF files, then generates separate BLIF files for each partition, leaving our Python scripts to perform the actual triplication, insertion of additional elements, and stitching them together. Given time we would like to combine the functionality into one program, but this is a lower priority than developing a working implementation.

Other design choices include deciding on VPR due to its open nature as discussed earlier in Section 1.3, and how we traverse our DFG. A depth first traversal tends to generate long narrow pipelines within each partition, thus increasing the number of register stages, whereas a breadth first traversal lends itself

to fewer register stages for the same number of nodes. A possible future improvement is implementing a more advanced traversal algorithm, for example A* with an appropriate heuristic could allow for more elements per partition.

Additionally, we are faced with a choice as to when in the CAD process to partition. The closer to the end of the process the more control we have, and the better our ability to estimate area and timing, but the harder it is to partition. As we are inserting new elements we want to partition before packing/placement to allow VPR to pack and place our inserted elements.

Choice of Language

We are using a combination of languages, mainly Python and C++. Language choice primarily came down to preference regarding familiarity and personal taste although a few other considerations were kept in mind. For BLIF joining and insertion of the voting logic Python was used. BLIF files are plain text and the text parsing to join and insert is computationally simple, so the primary concern was short development time while still being readable and maintainable (although Python's performance on text is still quite reasonable) [19]. For the actual partitioner C++ was chosen for a few reasons. Firstly, it was expected that the area and time estimations could be quite computationally expensive, so a lower level compiled language was chosen for performance reasons [19]. Secondly, VPR is written in C, so using C or C++ allowed for easy code reuse, or merging the partitioner and VPR. Our reason for choosing C++ over C was that we preferred an object oriented language as we felt it would be easier to maintain, and would better lend itself to our goal of extensibility, as well as its libraries making our implementation much easier.

5.3   Input File Format

The BLIF file format is a textual format which describes an arbitrary sequential or combinational network of logic functions [21]. Of the full BLIF specification, VPR only supports a subset of it, and hence our partitioner is also designed to only support that same subset.

```
1 .model voter
2 .inputs in1 in2 in3
3 .outputs out1 out 2
4 .clock clock
5 .names in1 in2 in3 out1
6 11- 1
7 1-1 1
8 -11 1
9 .latch in1 out2 re clock 1
10 ...
```

```
11 commands
12 ...
13 .end
```

Listing 5.1: BLIF file layout

Model name:	.model ⟨Name⟩	The name of the model.
Input List:	.inputs {Signal}	The model inputs.
Output List:	.outputs {Signal}	The model outputs.
Clock List:	.clock {Signal}	The model clocks.
		Commands
LUT:		.names {InputSignals} ⟨OutputSignal⟩
		{Line}
Latch:		.latch ⟨InputSignal⟩ ⟨OutputSignal⟩ [Field ClockSignal] [Field]
Optional End Marker:		.end

{Name} indicates 1 or more of Name. ⟨Name⟩ indicates a compulsory field. [Name] indicates an optional field. A combinational logic element (.name) is followed by one or more lines describing the logic function it implements. However, our partitioner only cares about node type and the signals (named with Signal above) as it builds and traverses the DFG. All other element information is stored and written back out when the node is written. Likewise for *Field*s.

VPR only supports flat BLIF files, so only one module declaration is allowed per BLIF file. SIS can be used to flatten BLIF files for use by VPR.

5.4   Implementation

Our implementation is still incomplete and so is both liable to change and is not currently match for our intended design. Most notably, we partition first, then triplicate, then join into one file; not triplicating as we partition. A simple pseudocode description is included in 5.2 (omitting code to read and write BLIF files) and is discussed below.

Reading in the BLIF file is a relatively simple process as subcircuits aren't supported. We make one pass through the input file, first reading in the list of inputs, outputs and clocks, and then creating a node for each primitive. We then iterate through the set of nodes building a list of signals, with each signal storing its sources and sinks.

```
1 Model = new BlifModel(file)
2 Queue = new Queue()
3 for each(Signal in Model->Inputs)
4     Queue.Push(Signal->Sinks)
5 Partition = new Partition()
6
```

```
 7  while(Queue.Size > 0)
 8      Node = Queue.Pop()
 9      if(AlreadyUsed(Node))
10          continue
11      if(EstimatedRecoveryTime > MAXIMUM_FAULT_RECOVERY_TIME)
12          WritePartitionToFile(FileName, Partition)
13          Partition = new Partition()
14      Partition.Add(Node)
15      for each(Signal in Node->Signals)
16          Queue.Push(Signal->Sinks)
17
18  CombinePartitions()
```

Listing 5.2: Simplified Pseudocode

As mentioned in Section 5.3, our input file format is a text file listing all the nodes. We read the file into memory and store it as a DFG, with all nodes and signals additionally stored in a hashmap to allow for quick random access. Each node contains a list of all connected signals, and each signal contains a list of all its sources and sinks making the DFG quite easy to traverse. Additionally, we store a status for each node indicating whether it is part of the current partition, a previous partition, or new, allowing us to detect feedback loops and avoid adding nodes multiple times. We then traverse the DFG in a breadth first manner while keeping running track of an estimate of the current partition's area and timing information. Once adding a new node would exceed our constraints we write the set of contained nodes to an output BLIF file, and proceed with partitioning the rest. Eventually we have one BLIF file for each partition. We then pass these to a set of existing Python scripts, written for initial benchmarking purposes and described in Section 4.3 which triplicates each partition and inserts voting logic, then connects each partition back up in a hierarchical BLIF file. This file is then passed to SIS to be flattened, at which point the partitioned circuit is ready for VPR.

5.5   Estimating Restrictions

As mentioned earlier, in order to partition our circuit we need a method of calculating the partition's (including voter logic) recovery time, which is based on circuit area (affects time to reconfigure), number of register stages (affects time to detect errors and resynchronise), and frequency (affects time to detect errors and resynchronise). Calculating the number of register stages is accomplished while traversing the DFG. Area and timing information are more difficult to determine as they rely on the placement and routing of the circuit. Placing and routing each partial partition every step as we traverse is not computationally feasible in a reasonable amount of time, as placement and routing are relatively slow processes and one of our goals is for our partitioning stage to be approximately as fast as the other
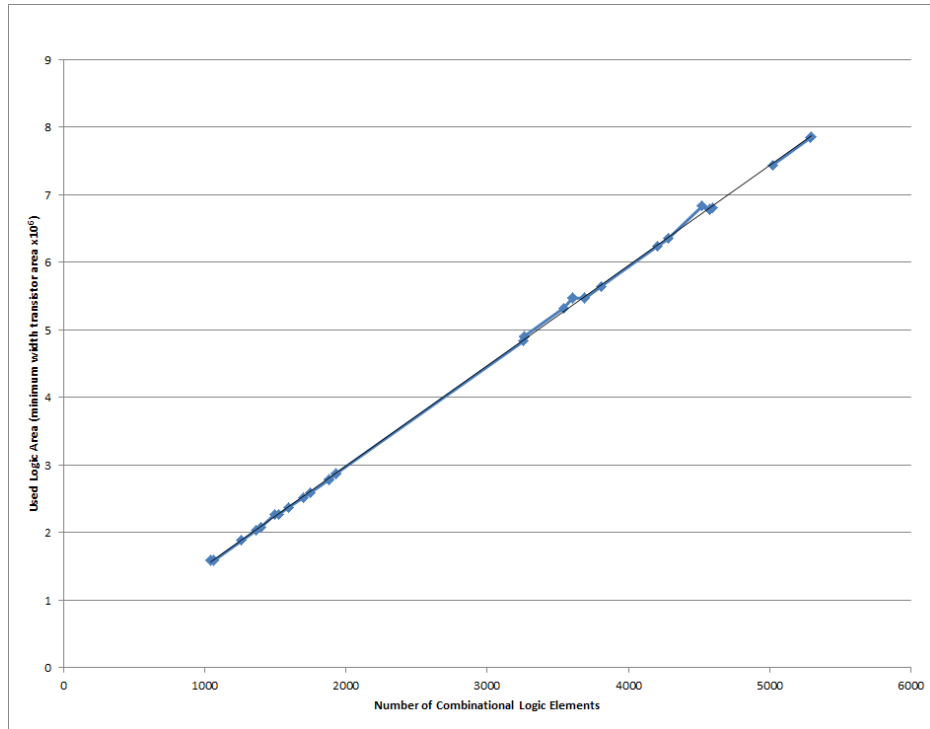
Figure 5.1: Circuit Area compared to number of Logic Elements

stages. Therefore, we need a way of estimating them. To do so we have collected preliminary benchmark information for a number of test circuits and analysed them for patterns allowing us to accurately estimate area and timing from a given circuit without placing and routing.

Results for Time and Area Estimation

Figure 5.1 illustrates the relationship between used circuit area and the number of the most common primitive. In every benchmark circuit there were more combinational logic elements than any other primitive, hence we looked at the relationship with the number of combinational logic elements. Table 5.1 duplicates what was shown in Table 4.4, and is shown again for convenience. It details the increase in circuit characteristics between the TMR and non TMR versions of the circuit, as a percentage of the original (100% is no change, 200% is double, and so on) that displayed the greatest decrease in maximum frequency.

Discussion of Time and Area Estimation

As shown in Figure 5.1 the area usage can be accurately estimated, as there is a clear relationship between the number of nodes and the area usage. The architecture we are using has one latch and one LUT per BLE, so for our supported logic elements the number of BLEs used is close to linear in $max(num_{latch}, num_{names})$. VPR's packer can be either timing or area driven. Currently we are using

Name	Num. Latches	Num. LUTs	FPGA Width	Num. Wire Segments	Used Area	Frequency	CPU Time
pdc width	200 N/A	300%	176%	104%	301%	75%	425%

Table 5.1: Circuit characteristics of benchmark with greatest frequency slowdown, expressed as a percentage of non TMR base

default settings (mostly area driven) giving us the linear relationship shown in Figure 5.1; however even when completely timing driven the packer still tries to fill every CLB [14].

After we have a basic partitioning algorithm, an area of further investigation is the impact of changing VPR's settings on the benchmark results, and the accuracy of our estimation functions.

Timing information, on the other hand, is harder to estimate with no obvious pattern, with maximum frequency appearing independent of the number of nodes. In Section 4.4 we saw that the median slowdown was 15% with a 25% worst result. Conversely, for a few rare cases the partitioned version is actually faster. Initially we will do a rough place and (optionally) route of the original circuit to determine a base time, then multiply it by an experimentally determined slowdown factor to obtain an estimate for the frequency. Initially we are using a slowdown factor of 2 (so half speed after partitioning) which easily encompasses all test circuits we've tried. We can then modify this factor by hand to examine if the impact of it on the final partition's performance warrants improving our estimation function.

Chapter 6

What next

6.1 Progress

The initial partitioner implementation is still in progress. We anticipate a basic working version by November 2012, and then will spend the next months collecting results and improving the partitioner. Done:

- Can read a BLIF file into a DFG.

- Can traverse a circuit represented as a DFG

- Have basic area and timing estimation functions.

- Can triplicate an arbitrary circuit (in a single BLIF file) and insert arbitrary voter logic (stored in another BLIF file).

- Initial benchmarks.

To Do:

- Write DFG to BLIF.

- Incorporate Python scripts into partitioning toolchain.

- Benchmark initial partitioning algorithm.

- Improve partitioner benchmarks.

- Investigate the effect of changing VPR's default parameters upon our results.

- Combine functionality of Python scripts and C++ partitioner into one program.

- Incorporate that single program into VTR's design flow, likely as part of VPR.
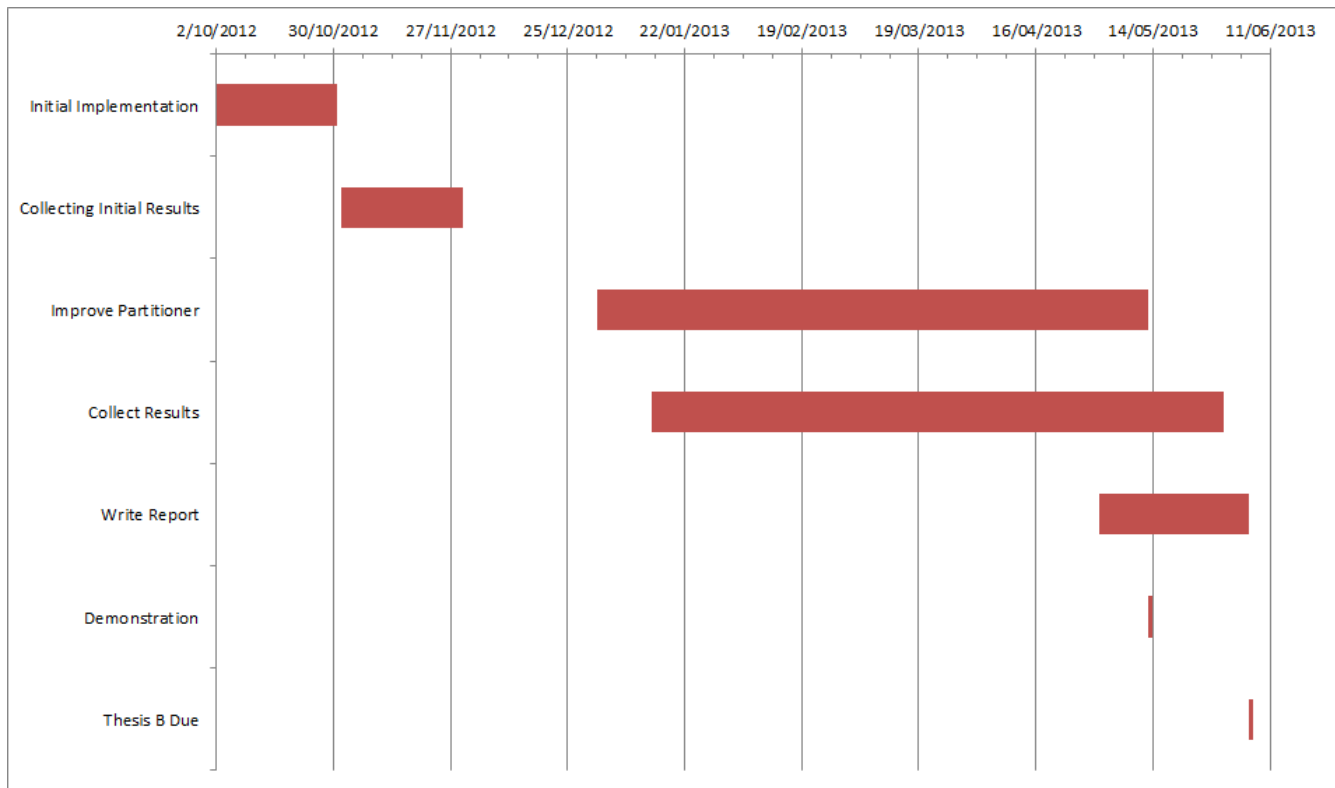
Figure 6.1: Schedule

6.2   Schedule

Figure 6.1 illustrates our anticipated schedule from the beginning of October 2011 until Thesis B is due. The rest of the year is allocated towards an initial partitioner implementation and preliminary results. December is set aside for Christmas holidays, and then most of 2012 will be used for collecting results and further refining the partitioning algorithm on the basis of the collected results.

Chapter 7

Conclusion

We are currently on track with our initial schedule, and expect to begin collecting results on our partitioner implementation, rather than the effects of generic TMR, at the end of this year. During the Christmas holidays progress will likely be quite limited however as we are very close to a working implementation it is believed to be achievable. Our results collected so far reflect only TMR in general, rather than our specific implementation, and they match those of literature and our expectations. We expect slightly worse performance from our partitioner due to additional logic being inserted, but not significantly so.

Appendix A

Results

This appendix tabulates the data used to calculate the relationships discussed in this thesis.

Name	Num. Latches	Num. LUTs	FPGA Width	Channel Width	Av. No. Wire Segments	Frequency (MHz)	CPU Time (s)
alu4	0	1522	15	200	4.76391	225	34.566
alu4TMR	0	4574	25	200	5.29148	188	132.41
apex2	0	1878	16	200	5.5124	216	47.45
apex2TMR	0	5637	29	200	6.21459	172	197.825
apex4	0	1262	14	200	5.0436	249	31.272
apex4TMR	0	3805	23	200	5.33318	204	119.002
bigkey	224	1699	16	200	3.70506	427	56.61
bigkeyTMR	672	5294	27	200	4.93838	377	193.388
clma	33	8365	34	200	5.62286	115	379.801
clmaTMR	99	25177	59	200	6.2762	103	2146.4
des	0	1591	16	200	4.07107	262	98.709
desTMR	0	5018	27	200	3.9862	229	263.883
diffeq	377	1494	15	200	3.0611	157	60.103
diffeqTMR	1131	4521	25	200	3.44041	156	204.691
dsip	224	1362	14	200	3.823	433	60.372
dsipTMR	672	4283	24	200	4.74885	377	177.405
elliptic	1122	3602	23	200	4.2743	134	123.967
ellipticTMR	3366	10920	39	200	5.15534	130	513.637
ex1010	0	4598	25	200	5.31938	176	146.81
ex1010TMR	0	13804	44	200	5.4995	134	655.814
ex5p	0	1064	13	200	4.95796	240	30.659
ex5pTMR	0	3255	22	200	5.2736	187	112.023
frisc	886	3539	23	200	5.52549	94.3	144.196
friscTMR	2658	10733	39	200	5.82234	90.1	525.145
misex3	0	1397	14	200	4.93	244	35.739
misex3TMR	0	4205	24	200	5.46202	198	136.961
pdc	0	4575	25	200	7.21792	175	167.232
pdcTMR	0	13765	44	200	7.51339	131	711.155
s298	8	1930	17	200	4.76794	124	47.789
s298TMR	24	5796	29	200	4.83777	116	185.169
s38417	1463	6042	30	200	3.53974	158	241.471
s38417TMR	4389	18232	51	200	3.78557	149	942.7
s38584.1	1260	6177	30	200	3.67662	206	263.313
s38584.1TMR	3780	18835	52	200	4.36338	157	1199.39
seq	0	1750	16	200	5.31333	253	50.624
seqTMR	0	5285	27	200	6.23725	193	184.467
spla	0	3690	23	200	6.44786	190	122.311
splaTMR	0	11116	39	200	6.75377	152	496.416
tseng	385	1046	13	200	3.04212	161	30.561
tsengTMR	1155	3260	22	200	3.82465	164	118.583

Table A.1: Results for 200 width channels

Name	Num. Latches	Num. LUTs	FPGA Width	Channel Width	Av. No. Wire Segments	Frequency (MHz)	CPU Time (s)
alu4	0	1522	15	60	4.79098	241	22.625
alu4TMR	0	4574	25	60	5.38017	193	101.828
apex2	0	1878	16	60	5.53926	212	34.256
apex2TMR	0	5637	29	60	6.34397	164	153.485
apex4	0	1262	14	60	5.13079	241	20.928
apex4TMR	0	3805	23	60	5.66967	198	90.166
bigkey	224	1699	16	60	3.79333	421	43.024
bigkeyTMR	672	5294	27	60	5.12677	363	154.662
clma				Could Not Route			
clmaTMR				Could Not Route			
des	0	1591	16	60	4.15635	257	50.68
desTMR	0	5018	27	60	4.07778	230	141.366
diffeq	377	1494	15	60	3.20978	151	28.392
diffeqTMR	1131	4521	25	60	3.60614	148	115.085
dsip	224	1362	14	60	3.84115	434	36.32
dsipTMR	672	4283	24	60	4.87989	368	118.515
elliptic	1122	3602	23	60	4.76756	137	101.398
ellipticTMR	3366	10920	39	60	5.74891	133	475.005
ex1010				Could Not Route			
ex1010TMR				Could Not Route			
ex5p	0	1064	13	60	5.29429	241	18.899
ex5pTMR				Could Not Route			
frisc	886	3539	23	60	6.17131	94.7	107.33
friscTMR				Could Not Route			
misex3	0	1397	14	60	4.99571	241	21.286
misex3TMR	0	4205	24	60	5.64103	197	91.433
pdc				Could Not Route			
pdcTMR				Could Not Route			
s298	8	1930	17	60	4.62137	122	27.868
s298TMR	24	5796	29	60	4.75444	118	120.257
s38417	1463	6042	30	60	3.66998	160	172.413
s38417TMR	4389	18232	51	60	3.87309	153	758.585
s38584.1	1260	6177	30	60	3.83166	202	195.832
s38584.1TMR	3780	18835	52	60	4.67375	159	944.364
seq	0	1750	16	60	5.30778	245	32.426
seqTMR	0	5285	27	60	6.56298	188	149.187
spla	0	3690	23	60	6.46323	188	86.03
splaTM				Could Not Route			
tseng	385	1046	13	60	3.22621	161	20.384
tsengTMR	1155	3260	22	60	3.93092	164	79.889

Table A.2: Results for 60 width channels

Name	Num. Latches	Num. LUTs	FPGA Width	Channel Width	Av. No. Wire Segments	Frequency (MHz)	CPU Time (s)
alu4	0	1522	15	34	5.61203	196	88.236
alu4TMR	0	4574	25	48	5.6293	182	245.343
apex2	0	1878	16	48	5.78099	204	80.968
apex2 TMR	0	5637	29	60	6.3913	167	885.343
apex4	0	1262	14	48	5.74387	175	55.509
apex4 TMR	0	3805	23	58	5.59716	181	357.909
bigkey	224	1699	16	42	4.00538	422	84.037
bigkey TMR	672	5294	27	48	5.46703	351	314.02
clma	33	8365	34	64	5.9249	111	826.48
clma TMR	99	25177	59	74	6.81963	94.3	6225.16
des	0	1591	16	46	4.46904	248	85.209
des TMR	0	5018	27	44	4.43131	229	329.551
diffeq	377	1494	15	38	3.67617	152	53.137
diffeq TMR	1131	4521	25	44	3.94535	156	238.817
dsip	224	1362	14	40	4.14675	446	86.123
dsip TMR	672	4283	24	38	5.10172	377	207.014
elliptic	1122	3602	23	50	5.19297	133	289.63
elliptic TMR	3366	10920	39	60	5.74707	133	1562.75
ex1010	0	4598	25	64	5.33384	172	401.434
ex1010 TMR	0	13804	44	62	5.70191	131	1399.9
ex5p	0	1064	13	50	5.64865	229	53.45
ex5p TMR	0	3255	22	62	6.03046	174	452.114
frisc	886	3539	23	56	6.66421	96.2	429.658
frisc TMR	2658	10733	39	66	6.5103	90.1	1383.44
misex3	0	1397	14	42	5.74857	233	60.019
misex3 TMR	0	4205	24	50	5.86212	190	368.648
pdc	0	4575	25	64	7.8308	156	397.958
pdc TMR	0	13765	44	72	7.74884	121	1005.6
s298	8	1930	17	30	5.16031	117	2620.11
s298 TMR	24	5796	29	34	4.97929	119	425.303
s38417	1463	6042	30	42	4.02152	158	376.671
s38417 TMR	4389	18232	51	50	3.96346	147	7381.12
s38584.1	1260	6177	30	44	4.22808	200	459.989
s38584.1 TMR	3780	18835	52	56	4.69062	156	2790.9
seq	0	1750	16	46	5.86111	207	105.526
seq TMR	0	5285	27	58	6.68624	190	448.058
spla	0	3690	23	54	7.08891	144	357.396
spla TMR	0	11116	39	66	6.92343	145	1320.57
tseng	385	1046	13	34	3.63807	163	49.618
tseng TMR	1155	3260	22	42	4.36026	163	183.167

Table A.3: Results for autosized width channels

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