## CS145 Midterm "Cheat Sheet"

## October 25, 2015

## 1 Possibly Useful Information

• Canonical SQL Statement:

SELECT <attributes>
FROM <tables>
WHERE <conditions>
GROUP BY <attributes>
HAVING <conditions>

- Functional Dependency (FD): For a relation R, and sets of attributes X and Y, the functional dependency  $X \to Y$  holds if for any  $t_1, t_2 \in R$ ,  $t_1[X] = t_2[X] \implies t_1[Y] = t_2[Y]$ .
- Armstrong's Axioms: Let the  $A_i$ s,  $B_i$ s, and  $C_k$ s be attributes:
  - 1. Split/Combine: If  $A_1, ..., A_n \to B_j$  for j = 1, ..., m, then this is equivalent to  $A_1, ..., A_n \to B_1, ..., B_m$  and vice-versa
  - 2. Reduction/Trivial:  $A_1,...,A_n \rightarrow A_i$  for any i=1,...,n
  - 3. Transitive Closure: If  $A_1,...,A_n \to B_1,...,B_m$  and  $B_1,...,B_m \to C_1,...,C_p$  then  $A_1,...,A_n \to C_1,...,C_p$
- Closure: Given a set of attributes X and a set of FDs F, the closure  $X^+$  is the set of all attributes y such that  $X \to y$ .
- Superkey: Given a relation R, a superkey is a set of attributes X such that  $X^+$  is equal to the full set of attributes of R.
- **Key:** A key is a minimal superkey, i.e. a superkey where no subset of it is also a superkey.
- Boyce-Codd Normal Form (BCNF): A relation R is in BCNF if for all sets of attributes X, either  $X^+ = X$  (X is trivial) or  $X^+ =$  the set of all attributes (X is a superkey).
- Conflicts: Two actions conflict if they are part of different TXNs, involve the same variable, and at least one of them is a write.
- Serializable: A schedule is serializable if it is equivalent to some serial ordering.
- Multi-Value Dependency (MVD): Given a relation R with a set of attributes A, two sets of attributes  $X, Y \subseteq A$ , we say that the MVD  $X \rightarrow Y$  holds if for any tuples  $t_1, t_2 \in R$  such that  $t_1[X] = t_2[Y]$ , there is a tuple  $t_3$  such that:
  - $-t_3[X] = t_1[X]$
  - $-t_3[Y] = t_1[Y]$
  - $-t_3[A \setminus Y] = t_2[A \setminus Y]$