

1.

Algorithm: isEven(n) Requires: one positive integer number n Returns: whether the number is even (True/False)
1: if n==0 2: return True 3: elseif n==1 4: return False 5: else 6: return isEven(n-2) 9: endif

Trace isEven(9):

1. isEven(9): return isEven(7)
2. isEven(7): return isEven(5)
3. isEven(5): return isEven(3)
4. isEven(3): return isEven(1)
5. isEven(1): return False.

2.

Algorithm: sumDigits(n) Requires: one positive integer number n Returns: the sum of all digits of n
1: if n/10==0 2: return n%10 3: else 4: return n%10+sumDigits(n/10) 5: endif

Trace sumDigits(1942):

1. $1942/10 \neq 0$, False, return $2 + \text{sumDigits}(194)$: $2 + 14 = 16$
2. $194/10 \neq 0$, False, return $4 + \text{sumDigits}(19)$: $4 + 10 = 14$
3. $19/10 \neq 0$, False, return $9 + \text{sumDigits}(1)$: $9 + 1 = 10$
4. $1/10 = 0$, True, return 1 (backtrack)

3.

Algorithm: numDigits(n) Requires: one positive integer number n Returns: number of digits of n
1: if n/10==0 2: return 1 3: else 4: return 1+numDigits(n/10) 5: endif

Trace numDigits(1942):

1. $1942/10 \neq 0$, False, return $1 + \text{numDigits}(194)$: $1 + 3 = 4$

2. $194/10 \neq 0$, False, return $1 + \text{numDigits}(19)$: $1+2=3$
 3. $19/10 \neq 0$, False, return $1 + \text{numDigits}(1)$: $1+1=2$
 4. $1/10 \neq 0$, True, return 1 (backtrack)
4. $\text{gcd}(2022, 12345) = 3$; $\text{gcd}(924, 198) = 6$; $\text{gcd}(234, 385) = 1$

Examples of tracing

$\text{gcd}(234, 385)$:

1. Line 1: $234 \neq 385$, False; Line 3: $234 < 385$, True; return $\text{gcd}(234, 151)$
2. Line 1: $234 \neq 151$, False; Line 3: $234 < 151$, False; return $\text{gcd}(83, 151)$
3. Line 1: $83 \neq 151$, False; Line 3: $83 < 151$, True; return $\text{gcd}(83, 68)$
4. Line 1: $83 \neq 68$, False; Line 3: $83 < 68$, False; return $\text{gcd}(15, 68)$
5. Line 1: $15 \neq 68$, False; Line 3: $15 < 68$, True; return $\text{gcd}(15, 53)$
6. Line 1: $15 \neq 53$, False; Line 3: $15 < 53$, True; return $\text{gcd}(15, 38)$
7. Line 1: $15 \neq 38$, False; Line 3: $15 < 38$, True; return $\text{gcd}(15, 23)$
8. Line 1: $15 \neq 23$, False; Line 3: $15 < 23$, True; return $\text{gcd}(15, 8)$
9. Line 1: $15 \neq 8$, False; Line 3: $15 < 8$, False; return $\text{gcd}(7, 8)$
10. Line 1: $7 \neq 8$, False; Line 3: $7 < 8$, True; return $\text{gcd}(7, 1)$
11. Line 1: $7 \neq 1$, False; Line 3: $7 < 1$, False; return $\text{gcd}(6, 1)$
12. Line 1: $6 \neq 1$, False; Line 3: $6 < 1$, False; return $\text{gcd}(5, 1)$
13. Line 1: $5 \neq 1$, False; Line 3: $5 < 1$, False; return $\text{gcd}(4, 1)$
14. Line 1: $4 \neq 1$, False; Line 3: $4 < 1$, False; return $\text{gcd}(3, 1)$
15. Line 1: $3 \neq 1$, False; Line 3: $3 < 1$, False; return $\text{gcd}(2, 1)$
16. Line 1: $2 \neq 1$, False; Line 3: $2 < 1$, False; return $\text{gcd}(1, 1)$
17. Line 1: $1 = 1$, True; return 1

$\text{Euclid}(234, 385)$

1. Line 1: $385 \neq 0$, False; return $\text{Euclid}(385, 234)$
2. Line 1: $234 \neq 0$, False; return $\text{Euclid}(234, 151)$
3. Line 1: $151 \neq 0$, False; return $\text{Euclid}(151, 83)$
4. Line 1: $83 \neq 0$, False; return $\text{Euclid}(83, 68)$
5. Line 1: $68 \neq 0$, False; return $\text{Euclid}(68, 15)$
6. Line 1: $15 \neq 0$, False; return $\text{Euclid}(15, 8)$
7. Line 1: $8 \neq 0$, False; return $\text{Euclid}(8, 7)$
8. Line 1: $7 \neq 0$, False; return $\text{Euclid}(7, 1)$
9. Line 1: $1 \neq 0$, False; return $\text{Euclid}(1, 0)$
10. Line 1: $0 = 0$, False; return 1

Note: no need for backtracking, as both are tail recursion. We can also see the improved Euclid algorithm is more efficient (there are less steps) compared to the other.

5. Refer to Seminar 3 example.

6.

Algorithm: fakeLog(x, y) Requires: two positive integers x, y Returns: largest integer k such that $x^k \leq y$
1: return fakeLogHelper(x, y, 1)

Algorithm: fakeLogHelper(x, y, k) Requires: three positive integers x, y, k Returns: largest integer k such that $x^k \leq y$
1: if $x^k > y$ 2: return k-1 3: else 4: return fakeLogHelper(x, y, k+1) 5: endif

Example of tracing

fakeLog(3, 28):

1. fakeLogHelper(3, 28, 1): Line 1: $3^1 > 28$, False; Line 4: return fakeLogHelper(3, 28, 2)
2. fakeLogHelper(3, 28, 2): Line 1: $3^2 > 28$, False; Line 4: return fakeLogHelper(3, 28, 3)
3. fakeLogHelper(3, 28, 3): Line 1: $3^3 > 28$, False; Line 4: return fakeLogHelper(3, 28, 4)
4. fakeLogHelper(3, 28, 4): Line 1: $3^4 > 28$, True; Line 2: return 3

Alternative solution without calling helper function:

Header will be the same as the main algorithm above.

1. if $x > y$
2. return 0
3. else
4. return 1+fakeLog(x, y/x) // here y/x is the integer division
5. endif

Trace this algorithm to see how it works.

7. (234, 385) are co-prime, because $\text{gcd}(234, 385)=1$.

8.

Algorithm: twinPrime(p, q) Requires: two positive integers p, q Returns: True if p, q are twin prime; False if p, q are not
1: if $ p-q \neq 2$ 2: return False 3: elseif isPrime(p) && isPrime(q) 4: return True 5: endif

Note: isPrime(n) algorithm can be referred to as the examples in Lecture 3 and Seminar 3.

i: (59, 61): True.

ii: (127, 129): False. Because $129=3 \cdot 43$, not a prime number.

9. Refer to Lecture 4 example. Change the name $F(n)$ to $\text{fib}(n)$ for this question.

10.

Algorithm: $\text{fibSum}(n)$ Requires: one positive integer n Returns: sum of first n terms in Fibonacci sequence
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1: if $n==1$ 2: return $\text{fib}(n)$ 3: else 4: return $\text{fib}(n)+\text{fibSum}(n-1)$ 5: endif
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Trace it yourself. Just focus on procedures of this algorithm.

You can call $\text{fib}(n)$ for computing the Fibonacci number directly without showing the procedures there.

11.

Algorithm: $\text{fakeSqrt}(n)$ Requires: a positive integer n Returns: returns the largest positive integer k such that $k \times k \leq n$.
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1: return $\text{fakeSqrtHelper}(n, 1)$

Algorithm: $\text{fakeSqrtHelper}(n, k)$ Requires: two positive integer n, k and $n \geq k$ Returns: returns the largest positive integer k such that $k \times k \leq n$.

1: if $k*k > n$ 2: return $k-1$ 3: else 4: return $\text{fakeSqrtHelper}(n, k+1)$ 5: endif
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Trace $\text{fakeSqrt}(10)$:

1. Line 1: return $\text{fakeSqrtHelper}(10, 1)$;
2. Line 1: $1*1 > 10$, False; Line 4: return $\text{fakeSqrtHelper}(10, 2)$;
3. Line 1: $2*2 > 10$, False; Line 4: return $\text{fakeSqrtHelper}(10, 3)$;
4. Line 1: $3*3 > 10$, False; Line 4: return $\text{fakeSqrtHelper}(10, 4)$;
5. Line 1: $4*4 > 10$, True; Line 2: return 3

Note: a tail recursion in the helper function, so no need for backtracking after obtaining 3 as answer.

12. Idea of designing this algorithm:

Given a and b , taking multiple of b ($b, 2b, 3b, 4b \dots$) until we find a multiple of b that is also divisible by a . This process does have an end, because we know at least $a*b$ is divisible by a .

Algorithm: LCM (a, b) [main] Requires: two positive integers a and b Returns: the least common multiple of a and b
1: return LCMHelper(a, b, b)

Algorithm: LCMHelper (a, b, k) [helper] Requires: three positive integers a, b and k Returns: returns the least common multiple
1: if $k\%a==0$ 2: return k 3: else 4: return LCMHelper($a, b, k+b$) 5: endif

Trace LCM(6, 8)

1. Line 1: return LCFHelper(6, 8, 8).
2. Line 1: $8\%6==0$, False; Line 4: return LCMHelper(6, 8, 16).
3. Line 1: $16\%6==0$, False; Line 4: return LCMHelper(6, 8, 24).
4. Line 1: $24\%6==0$, True; Line 2: return 24

Alternative solution (not a recursive algorithm):

Header will be the same as the main algorithm above.

1. let $m = \text{gcd}(a,b)$
2. return $a*b/m$

This method requires the mathematical fact that $a*b = \text{gcd}(a,b)*\text{lcm}(a,b)$.

And of course we are calling the $\text{gcd}()$ function as a sub-algorithm here.

Comments:

- I. So far, you should be familiar with tracing algorithms for understanding the ideas. In the future, we may exclude the tracing process in the solution sheet and leave it to you for checking through.
- II. In particular, if there is a call of helper function, think about why we need it and how it handles those extra variables in recursive calls.