



Introduction to Algorithms

CELEN086

Seminar 7
(w/c 25/11/2024)



Outline

In this seminar, we will study and review on following topics:

- Basic tree concepts
- Commands for binary trees
- Designing recursive algorithms on binary trees
- Binary tree with minimal height/depth

You will also learn useful Math/CS concepts and vocabularies.

Tree basics

Size of tree: 10

Height (depth) of tree: 3

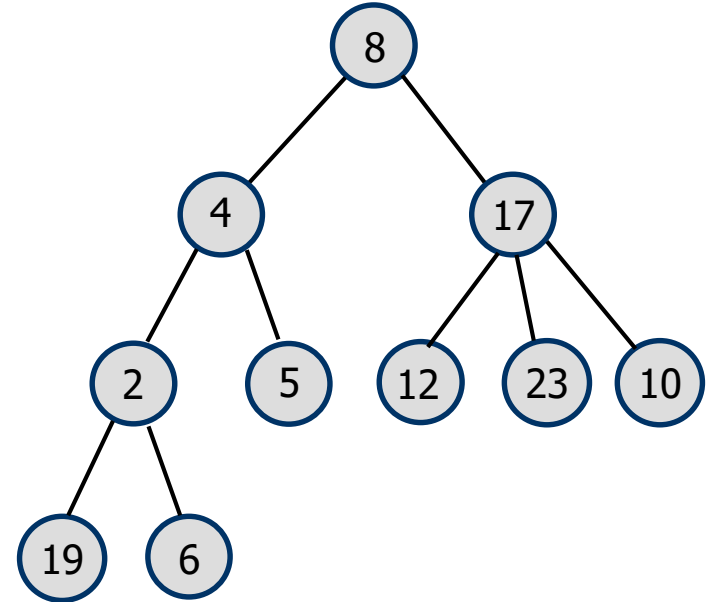
Depth of node with value 4: 1

Height of node with value 4: 2

Leaf nodes: 19 6 5 12 23 10

Binary tree? Why?

No. The node with value 17 has degree 3.





Binary tree (vs. list) commands

[tree commands]

- leaf
- node(leaf, x, leaf)
- node(left-subtree, x, right-subtree)
- isLeaf(tree)
- root(tree)
- left(tree), right(tree)
- isLeaf(left(tree))&&isLeaf(right(tree))

to check single-element tree
(leaf node)

[list commands]

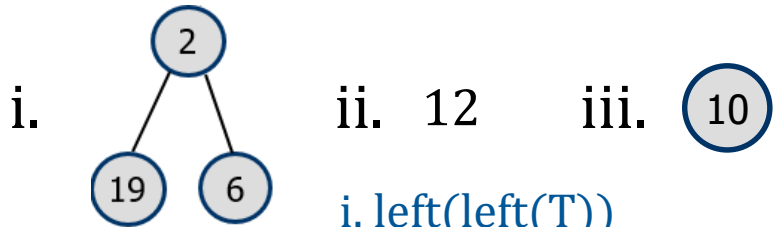
- nil
- cons(x, nil)
- cons(x, list)
- isEmpty(list)
- value(list)
- tail(list)
- isEmpty(tail(list))

to check single-element list

Binary tree

Let T be the given binary tree.

Write the pseudo code for:



i. `left(left(T))`

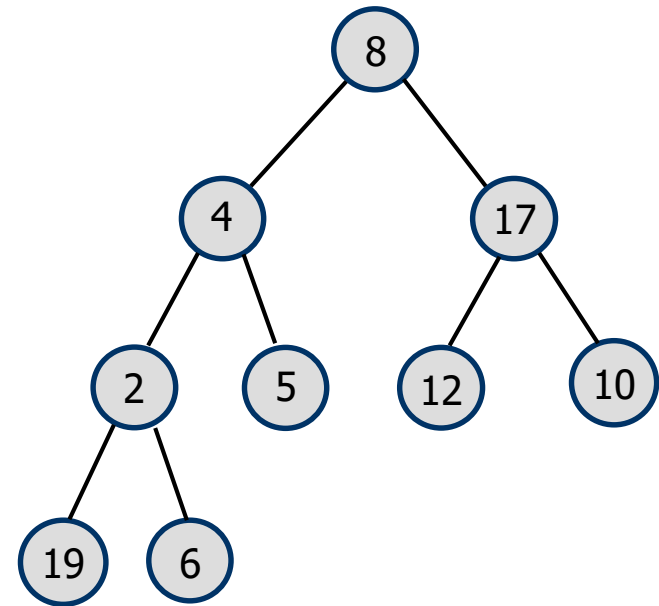
ii. `root(left(right(T)))`

iii. `right(right(T))`

vi. Draw the following tree:

`node(node(node(leaf,19,leaf),2,node(leaf,4,node(leaf,6,leaf))),5,leaf)`

Is it a subtree of T ?



Algorithm: tree size

Is the following alternative algorithm for finding the tree size correct?

```
1. if isLeaf(left(T))&& isLeaf(right(T))
2.   return 1
3. else
4.   return size(left(T))+size(right(T))+1
5. endif
```

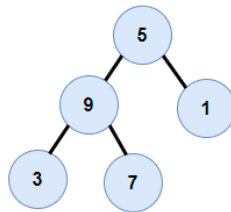
Algorithm: **size**(T)

Requires: a binary tree T

Returns: total number of nodes in T (size of T)

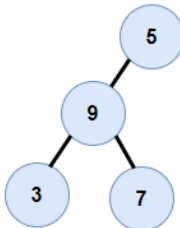
```
1. if isLeaf(T)
2.   return 0
3. else
4.   return size(left(T))+size(right(T))+1
5. endif
```

Trace it with T=



Working well.

Trace it with T=



Missing base case.

Algorithm: tree size (alternative ver.)

Algorithm: **size**(T)

Requires: a binary tree T

Returns: total number of nodes in T (size of T)

```
1. if isLeaf(left(T))&& isLeaf(right(T))
2.   return 1 // base case
3. elseif isLeaf(left(T))
4.   return 1+size(right(T)) // recursive call case#1
5. elseif isLeaf(right(T))
6.   return 1+size(left(T)) // recursive call case#2
7. else
8.   return 1+size(left(T))+size(right(T)) // recursive call case#3
9. endif
```

Practice: tree height

Write a recursive algorithm called `height(T)` that compute the height of a non-empty binary tree.

You can call the `max()` function to compute the maximum of two values.

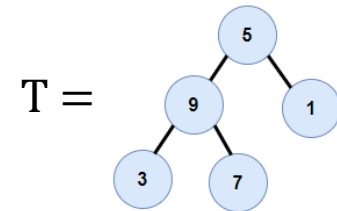
Algorithm: **height**(T)

Requires: a **non-empty** binary tree T

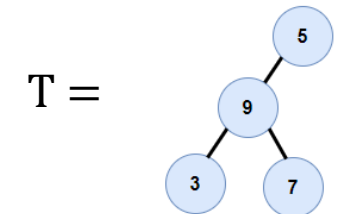
Returns: height of T

```
1. if isLeaf(left(T)) && isLeaf(right(T))
2.   return 0
3. elseif isLeaf(left(T))
4.   return 1+height(right(T))
5. elseif isLeaf(right(T))
6.   return 1+height(left(T))
7. else
8.   return 1+ max(height(left(T)), height(right(T)))
9. endif
```

Trace it with



and



Practice: delete leaf node

Write a recursive algorithm called `delete(T)` that deletes all the leaf nodes in a binary tree.

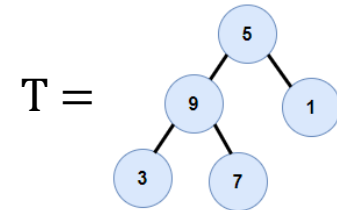
Algorithm: `delete(T)`

Requires: a `non-empty` binary tree `T`

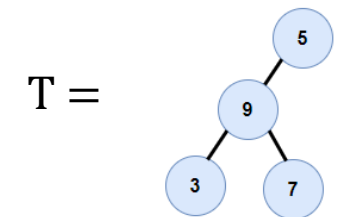
Returns: a tree after deletions of all leaf nodes in `T`

```
1. if isLeaf(left(T)) && isLeaf(right(T))
2.   return leaf
3. elseif isLeaf(left(T))
4.   return node(leaf, root(T), delete(right(T)))
5. elseif isLeaf(right(T))
6.   return node(delete(left(T)), root(T), leaf)
7. else
8.   return node(delete(left(T), root(T), delete(right(T)))
9. endif
```

Trace it with



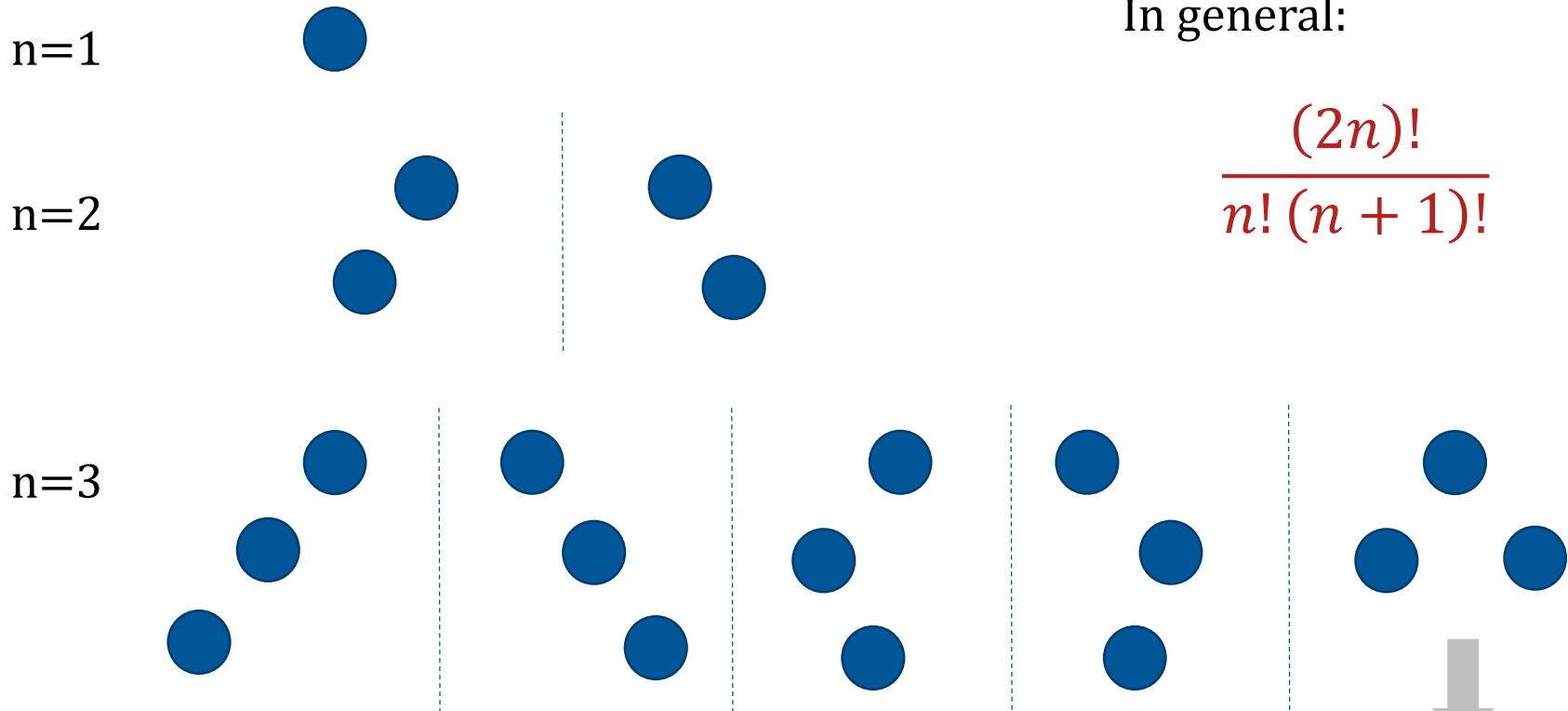
and





Binary tree with minimal height

How many binary trees can be made with n (same) nodes?

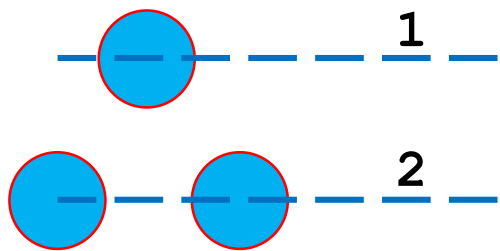


Binary tree with minimal height/depth.

Binary tree: size vs. height

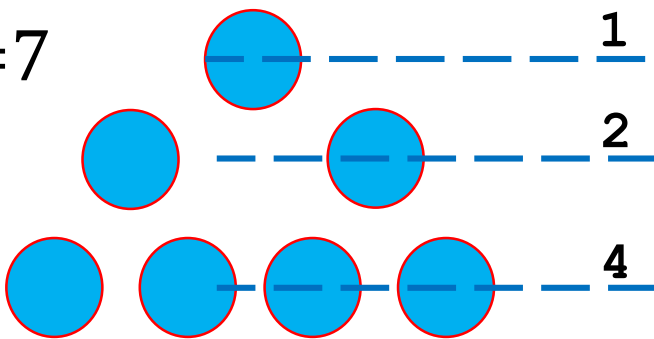
What is the minimal height of a binary tree with n nodes? $h = \lfloor \log_2 n \rfloor$

n=3



$3 = 2^0 + 2^1$
 $h = 1$

n=7



$7 = 2^0 + 2^1 + 2^2$
 $h = 2$

For a binary tree with height h , maximum number of nodes can be stored is:

$$2^0 + 2^1 + 2^2 + \dots + 2^h = 2^{h+1} - 1$$