



# Introduction to Algorithms

Module Code: CELEN086

Lecture 10

(12/12/24)

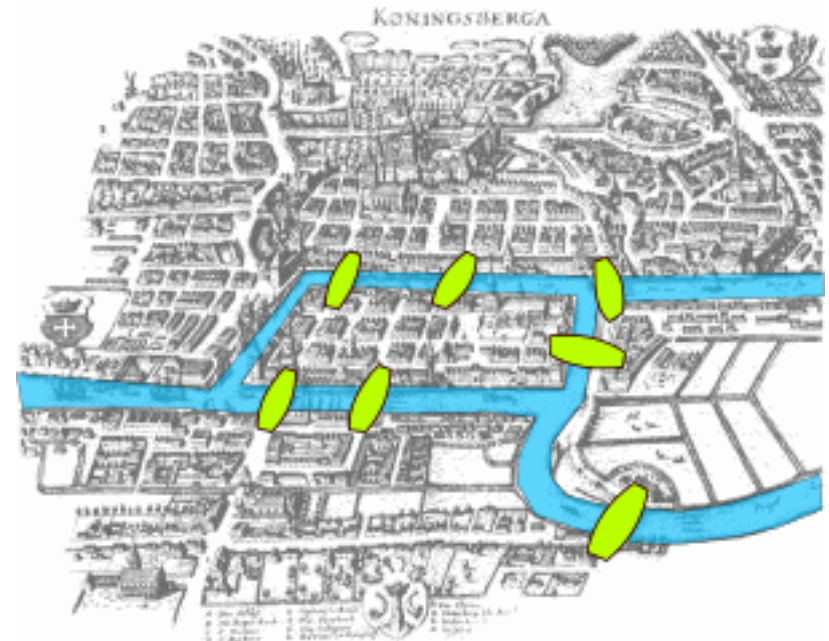
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# 7 bridge problem

In the former city of Königsberg, East Prussia (now in Russia), the river Pregel had **7 bridges**.

- Is there a path following which we can go cross all the 7 bridges exactly once?
- If such a path exists, is it possible that we end up at the same place as where we started?

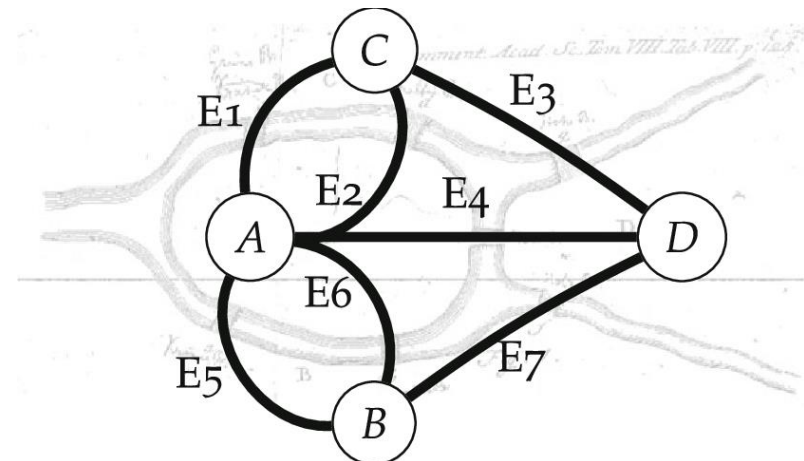
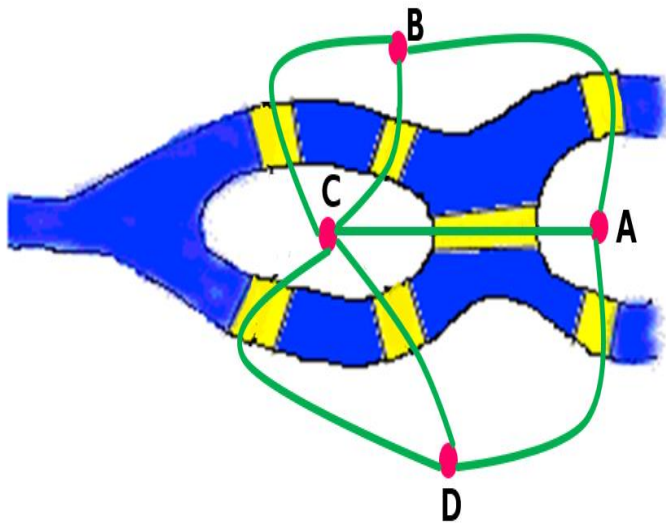


# 7 bridge problem

In 1736 Leonhard Euler explained why the problem of crossing the 7 bridges of Königsberg is impossible.



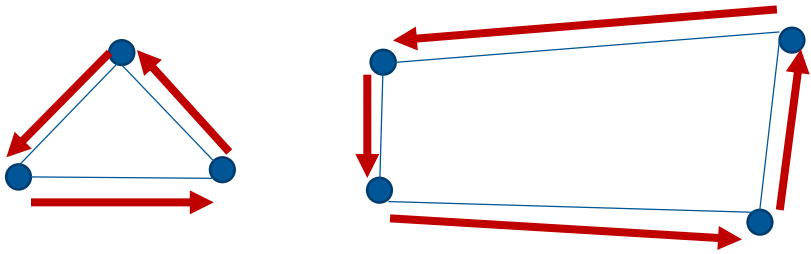
Euler laid the foundations of graph theory and prefigured the idea of topology (a mathematical branch).



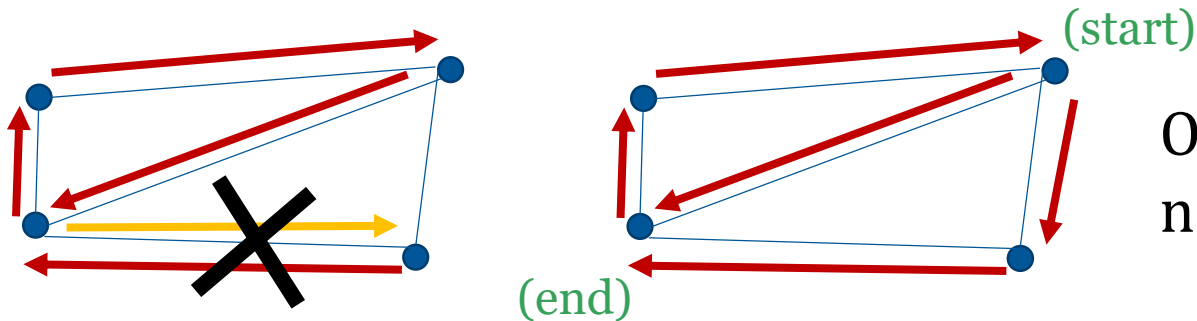
# Euler path and Euler circuit

**Euler path** (Eulerian path, Euler trail) is a path that passes through **every** edge of the graph exactly once.

If the Euler path **returns to the origin**, forming a closed path, then the path is called **Euler circuit** (Euler cycle, Euler tour).



Both have Euler circuits.

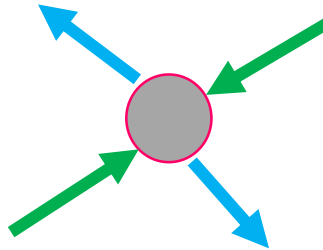


Only has Euler path,  
no Euler circuit.

# Theorem

- If a graph has Euler **circuit**, then the degree of **every** vertex must be an **even** number.

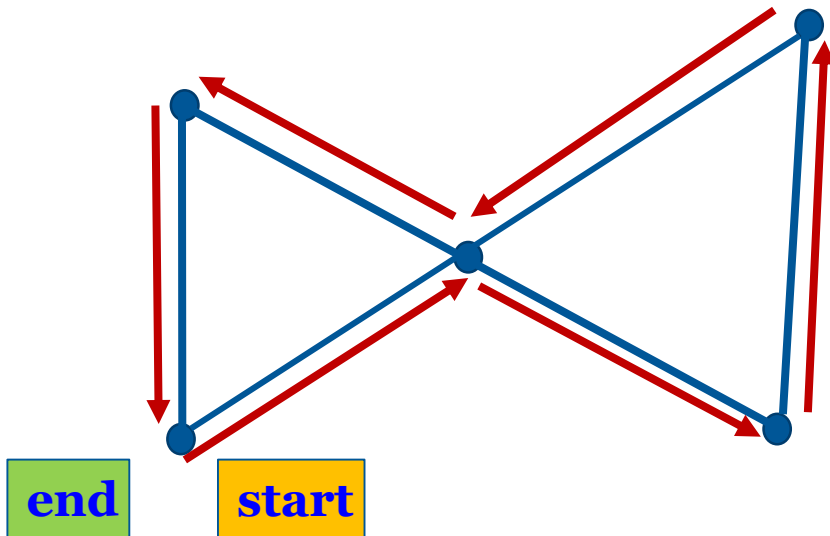
For each vertex that we enter during the tour via an edge, we must also be able to exit via another edge.



The number of edges entering each vertex, must be equal to the number of edges exiting that vertex.

- If a graph has Euler **path**, then it can have **at most** two vertices (starting/ending vertices) with **odd** degrees.

# Example



Degrees of all vertices are even.

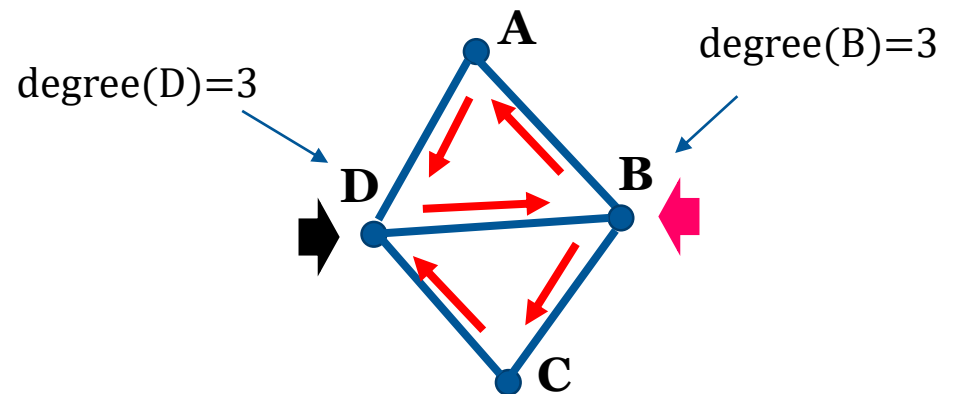
Euler circuit Yes

Euler path Yes

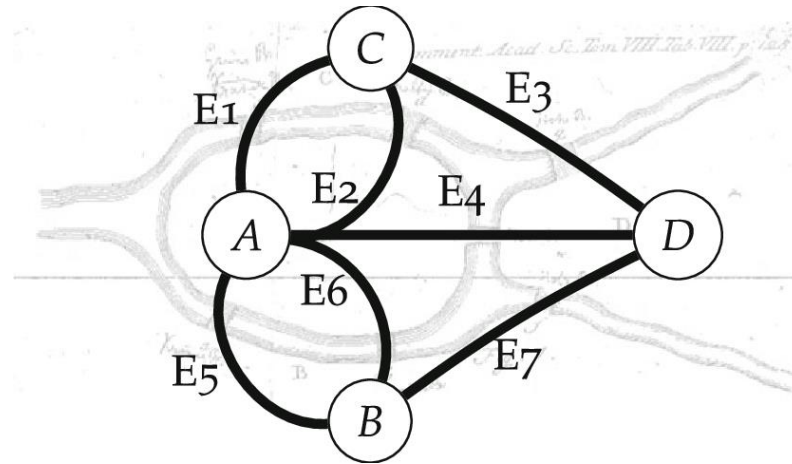
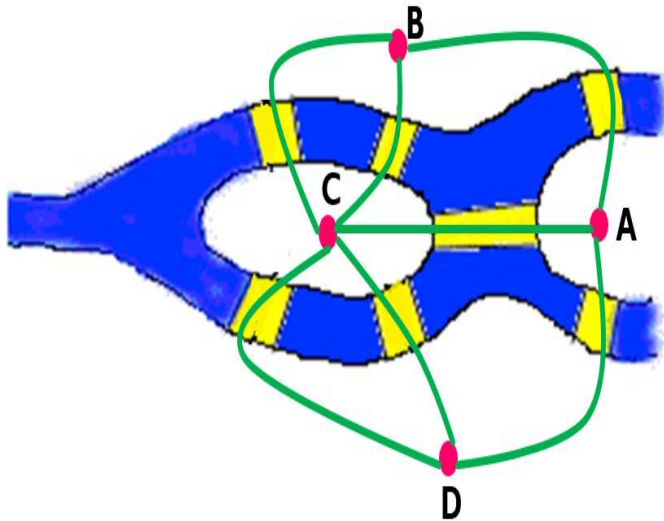
Two vertices are of odd degrees.

Euler circuit No

Euler path Yes



# Example



Vertex	Degree
A	3
B	3
C	5
D	3

More than two vertices have odd degrees.

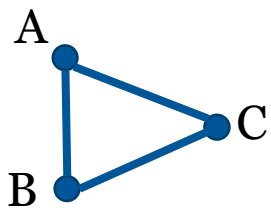
Euler circuit      No

Euler path      No

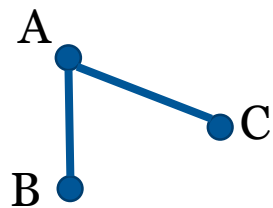
# Spanning tree

Tree is a special connected and undirected graph with no cycles.

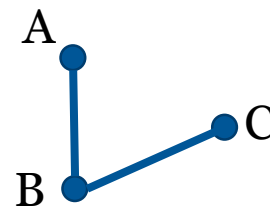
A **spanning tree** is a subset of a connected and undirected graph, which has all vertices covered with minimum possible number of edges.



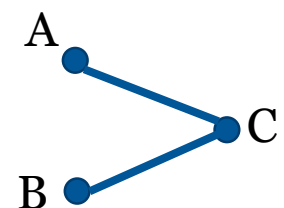
*graph*



*spanning tree*



*spanning tree*

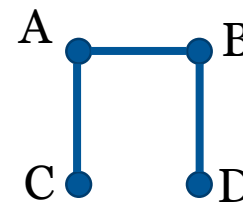
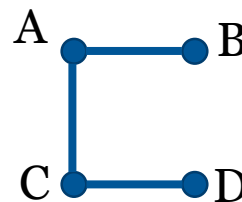
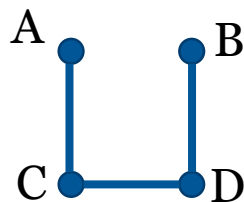
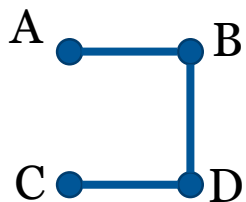
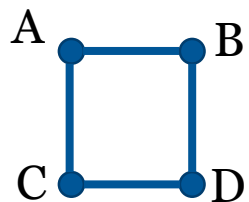


*spanning tree*

$$|V| = 3, |E| = 3, |E_s| = 2$$



# Number of spanning trees



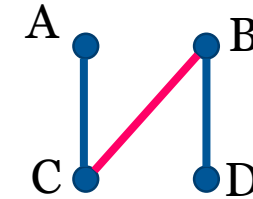
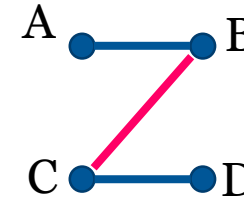
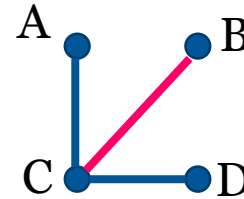
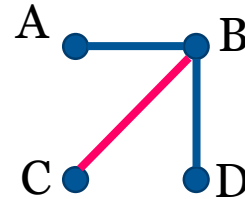
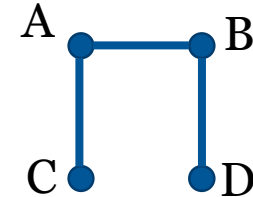
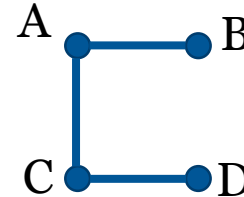
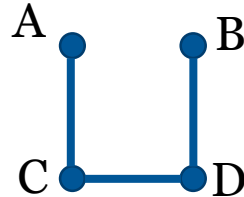
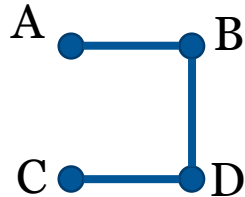
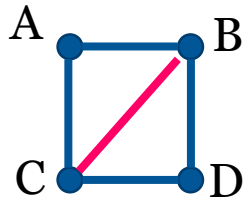
$$|V| = 4, |E| = 4, |E_s| = 3$$

$$\Rightarrow \binom{4}{3} = 4$$

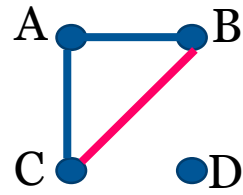
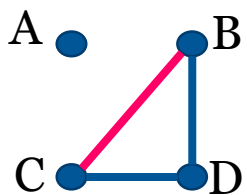
Property:  $|E_s| = |V| - 1$

This can help us identifying total number of edges needed in the spanning tree of a connected graph.

# Number of spanning trees



$$|V| = 4, |E| = 5, |E_s| = 3 \Rightarrow \binom{5}{3} = 10$$

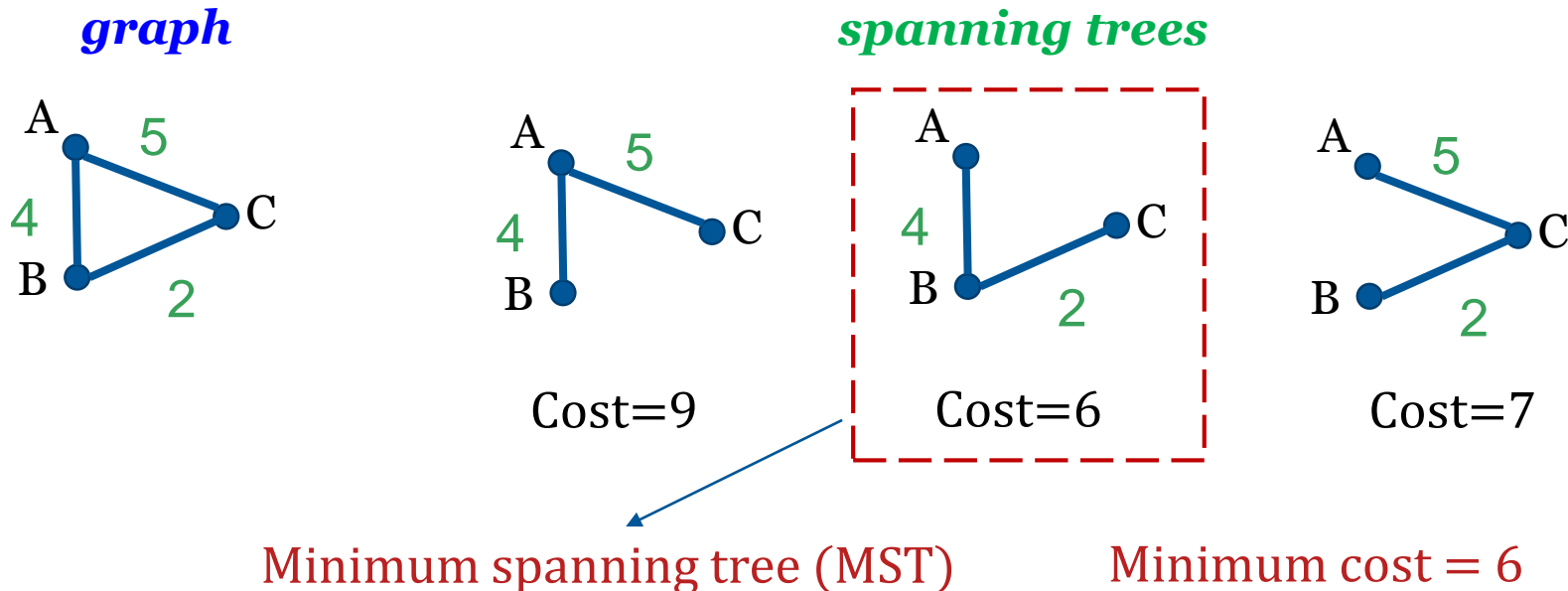


Number of spanning tree  
=  $10 - 2 = 8$

Exclude 2 cases with cycles!

# Minimum spanning tree

Minimum spanning tree is a sub graph of an undirected graph such that the subgraph spans (includes) all nodes, is connected, is acyclic, and has total minimum edge weight.

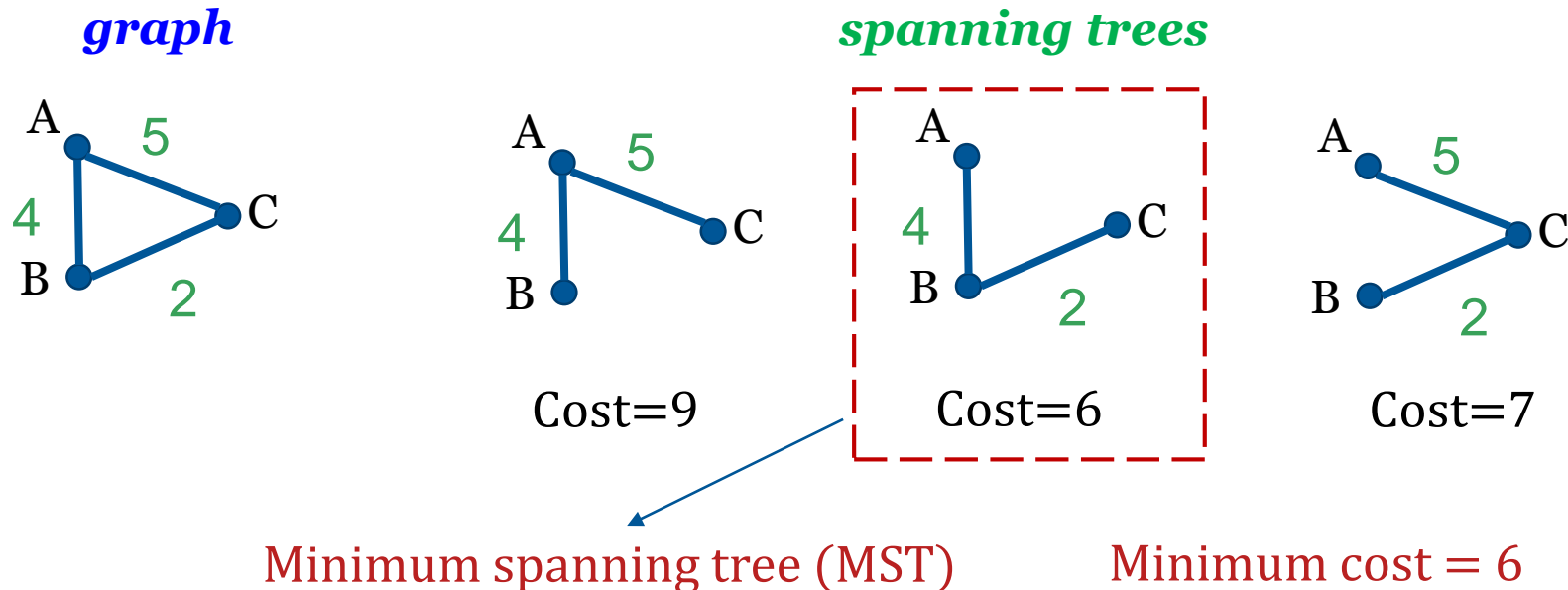


# Minimum spanning tree

For a weighted graph, we are interested in finding its spanning tree with minimum sum of its weighted edges.

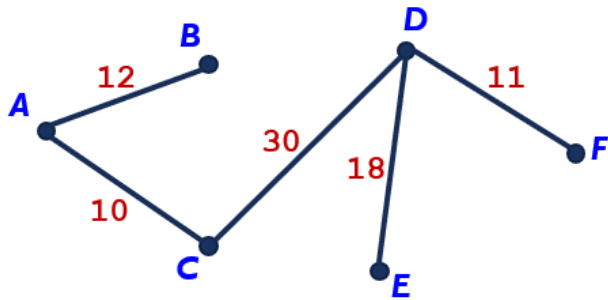
Such spanning tree is called **minimum spanning tree** of the graph.

This sum is called the **minimum cost**.

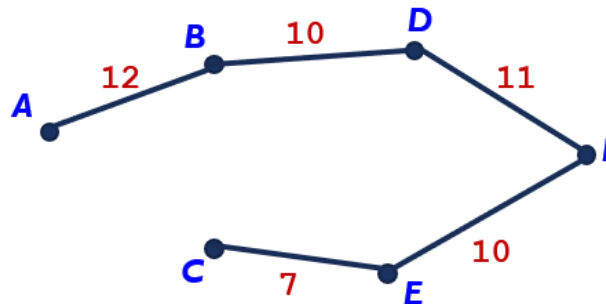


# Example

If the weighted and connected graph is complex, it is not a simple task to identify the minimum spanning tree and its minimum cost.



$$\text{Cost} = 12 + 10 + 30 + 18 + 11$$



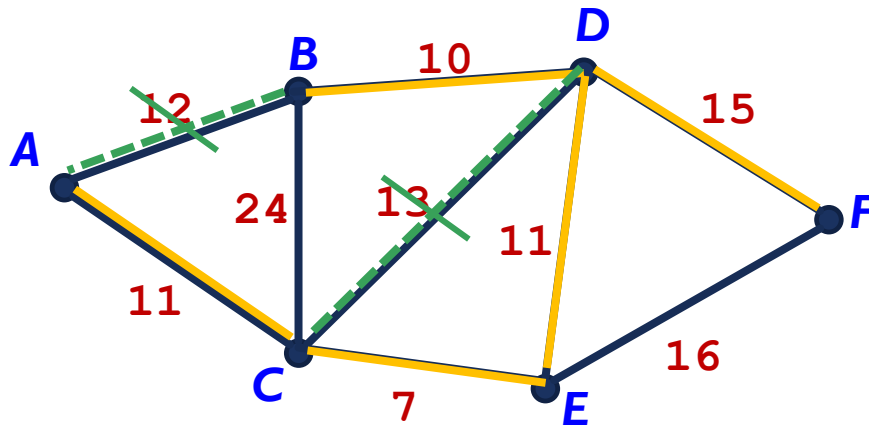
$$\text{Cost} = 12 + 10 + 11 + 10 + 7$$

We can apply two classical algorithms for solving such a task:

Kruskal's algorithm

Prim's algorithm

# Kruskal's algorithm



$$|E_s| = |V| - 1 = 6 - 1 = 5$$

Edge selected

CE

BD

AC (same weight as DE,

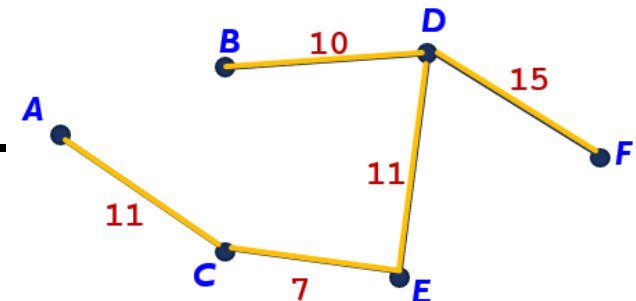
DE choose either first)

DF (avoid making a cycle)

5 edges selected, done.

- At each step, select the edge with **smallest weight**.
- Repeat this process. **Avoid forming cycles**.
- Stop when all vertices are connected.

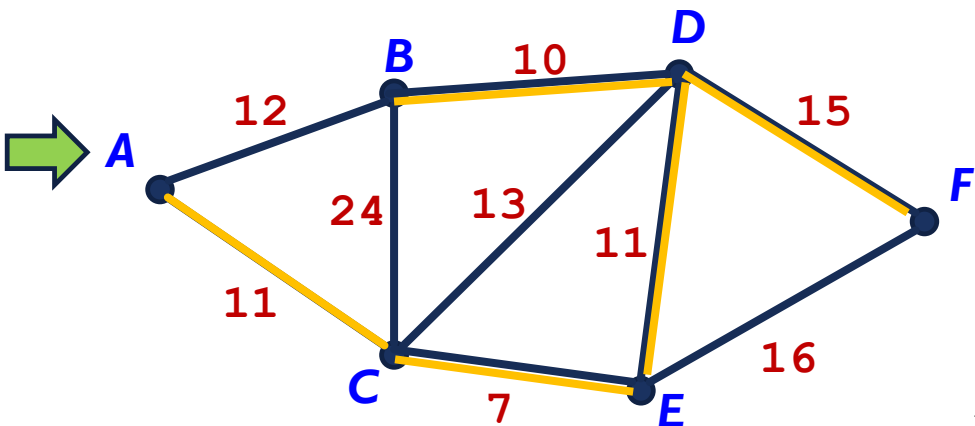
Minimum spanning tree:



Minimum cost

$$7 + 10 + 11 + 11 + 15 = 54$$

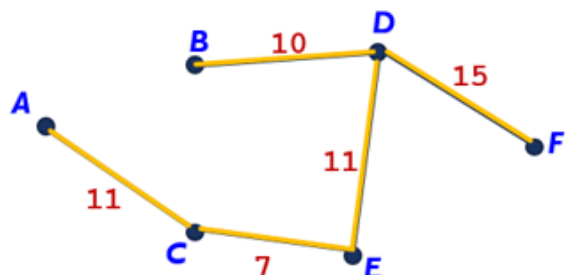
# Prim's algorithm



Visited	Unvisited	Edge selected
A	B,C,D,E,F	AC
A,C	B,D,E,F	CE
A,C,E	B,D,F	ED
A,C,E,D	B,F	DB
A,C,E,D,B	F	DF
A,C,E,D,B,F		5 edges selected, done.

- Start with any vertex (mark as visited).
- Choose the edge with **smallest weight connecting visited and unvisited vertices**.
- Mark the connected vertex as visited, repeat the previous step.
- Stop when all vertices are connected.

Minimum spanning tree:



Minimum cost  
 $7 + 10 + 11 + 11 + 15 = 54$

# Shortest path algorithms (summary)

Dijkstra's algorithm

Kruskal's/Prim's algorithms

They are all **greedy algorithms** on weighted graphs.

The optimized solution may not be unique.

To find shortest paths  
among vertices.

Work on undirected graph  
and directed graph.

To find the minimum  
spanning tree.

Work on undirected graph.





# Review

Question1: Which of the following graphs has an Euler Tour?

- (a) A graph with vertices of degrees 2, 2, 2, 2
- (b) A graph with vertices of degrees 3, 3, 3, 3
- (c) A graph with vertices of degrees 4, 4, 4, 4
- (d) A graph with vertices of degrees 1, 1, 1, 1

Question 2: What is the primary purpose of a Minimum Spanning Tree (MST)?

- (a) To find the longest path in a graph
- (b) To connect all vertices with the minimum total edge weight
- (c) To find the shortest path between two vertices
- (d) To detect cycles in a graph