Addressing Modes, Table, User Stack

1. Address Mode

2. Table

3. Stack and Subroutine Call

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1. Address Mode

2. Table

3. Stack and Subroutine Call

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Introduction

• Data (source and destination) could be in Register

Memory (data or program)

Immediate values (part of instruction)

• PIC18 provides 4 addressing modes
(The waysof accessing data is called *addressing mode*.)

Immediate

Direct

Register indirect

Indexed-ROM

Immediate and Direct Addressing mode

Add. Code 0000 0E56 MOVLW 0x56

0002 6E40 MOVWF 0x40, ACCESS

0004 C040 MOVFF 0x40, 0x50

0006 F050

0008 EF00 GOTO 0

000A F000

56 →WREG

WREG → loc 40H

(Loc 40H)→loc 50H

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SFR Registers and their addresses

- Can be access by Their name
- · Their address
- Which is easier to remember?
 MOVWF PORTB
 MOVWF 0xF81

Address	Symbol Name	Value
FE8	WREG	0x56
F80	PORTA	0x00
F81	PORTB	0x00
F83	PORTD	0x00
F82	PORTC	0x00
F89	LATA	0x00
F8A	LATB	0x00
F8B	LATC	0x00
F8C	LATD	0x00
F92	TRISA	0x7F
F93	TRISB	OxFF
F94	TRISC	0x00
F95	TRISD	OxFF
	INDFO	ed Memory
	INDF1	ed Memory
FE9	FSRO	0x0000
FE9	FSROL	0x00
FEA	FSROH	0x00
FE1	FSR1L	0x00
FE2	FSR1H	0x00
		7
		/

Immediate and Direct Addressing mode

What is the difference between

INCF fileReg, W

INCF fileReg, F

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SFR Registers and their addresses

SFR addresses are started at F80H and the last location is the address FFFH $\,$

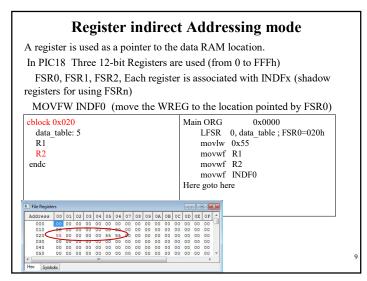
In Listing file, you will see that the SFR names are replaced with their addresses.

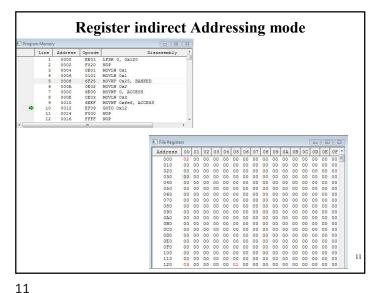
The WREG register is one of the SFR registers and has address FE8h

0000 0E56 MOVLW 0x56 ; MOVLW 56H 0002 6E81 MOVWF 0xf81, ACCESS ; MOVWF PORTB

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Register indirect Addressing mode cblock 0x000 Main ORG R2 LFSR 0, data table; FSR0=120h endc movlw 01H cblock 0x120 movlb high R1 data table: 5 movwf R1 movlw 02H R1 endc movwf R2 movlw 03H R1 address = 125H movwf INDF0 R2 address = 000HHere goto here data table address = 120H Program Memory Line Address Opcode Disassembly EE01 LFSR 0, 0x120 0002 F020 NOP 0E01 MOVLW 0x1 0101 MOVLB 0x1 6F23 MOVWF 0x25, BANKED 0E02 MOVLW 0x2 0008 000A 6E00 MOVWF 0, ACCESS OOOE OEO3 MOVLW 0x3 0010 6EEF MOVWF Oxfef, ACCESS EF09 GOTO 0x12 F000 NOP 0014

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Address Mode 1. 2. **Table** Stack and Subroutine Call 3.

Tables

Indirect addressing makes accessing data dynamic. That is why we can implement array in high level language.

Looping is possible to increment the address in FSR with

Instructions	
XXXX INDFn	After accessing fileREG pointed by FSFn, the FSRn is unchanged.
XXXX POSTINCn	After accessing fileREG pointed by FSFn, the FSRn is incremented
XXXX PREINCn	Before accessing fileREG pointed by FSFn, the FSRn is incremented.
XXXX POSTDECn	After accessing fileREG pointed by FSFn, the FSRn is decremented
XXXX PLUSWn	Address = WREG + FSRn

XXXX stays for some instructions

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Example

Solution B

MOVLW 55H

LFSR 0,0x40

MOVWF POSTINC0
MOVWF POSTINC0

MOVWF POSTINCO

 ${\color{red} \textbf{MOVWF}} \hspace{0.1cm} \texttt{POSTINC0}$

MOVWF POSTINCO

Example

Write a program to copy the value 55H into RAM locations 40h to 45h using

•Direct addressing mode

•Register indirect addressing mode

•A loop

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Example

Solution C

В1

LFSR 0, 0x040

MOVLW 0x5

MOVWF COUNT
MOVLW 0x55

MOVWF POSTINCO

DECF COUNT, F

BNZ B1

Of course C is much better if the table size is large

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Example

Write a program to clear 16 RAM location starting at location 60H using Auto increment.

COUNTREG EQU 0x10 CNTVAL EQU D'16' MOVLW CNTVAL MOVWF COUNTREG 1,0x60 LFSR CLRF POSTINC1 DECF COUNTREG, F BNZ В3

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Example

Assume that RAM locations 40-43H have the following hex data. Write a program to add them together and place the result in locations 06 and 07.

COUNTREG EQU 0x20 Address Data L BYTE EQU 0x06 040H 7D H_BYTE EQU 0x07 041H EB CNTVAL EQU 4 042H C5 MOVLW CNTVAL 043H 5B MOVWF COUNTREG 0,0x40LFSR WREG CLRF CLRF H BYTE POSTINCO, W ADDWF BNC OVER H BYTE, F INCF COUNTRÉG, F OVER DECF BNZВ5 MOVWF L BYTE

Example

Write a program to copy a block of 5 bytes of data from location starting at 30H to RAM locations starting at 60H.

COUNTREG EQU 0x10 CNTVAL EQU D'5' MOVLW CNTVAL MOVWF COUNTREG 0, 0x30LFSR LFSR 1, 0x60 в3 MOVF POSTINCO, W MOVWF POSTINC1 COUNTREG, F DECF в3 BNZ

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Example

Write a program to add the following multi-byte BCD numbers and save the result at location 60H. 12896577 +23647839

COUNTREG EQU 0x20 CNTVAL EQU D'4' MOVLW CNTVAL COUNTREG MOVWF 0.0x30LFSR Address Data 1,0x50 LFSR 030H 77 2,0x60 LFSR 031H 65 BCF STATUS, C MOVF POSTINCO, W 032H 89 ADDWFC POSTINC1,W 033H 12 DAW 050H 39 MOVWF POSTINC2 051H 78 DECF COUNTREG, F 052H 64 BNZ в3 053H

Fixed Data in Program Space

ROM has enough space to store fixed data

- •DB directive, which means **Define Byte**, is widely used to allocate ROM program memory in byte-sized chunks
- •Use single quotes (´) for a single character or double quotes (") for a string

Org 0x500

DATA1 DB 0x39

DATA2 DB 'z'

DATA3 DB "Hello All"

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Fixed Data in Program Space

• Is there any instruction to load 21 bits

MOVLW low MYDATA
MOVWF TBLPTRL
MOVLW high MYDATA
MOVWF TBLPTRH
MOVLW upper MYDATA
MOVWF TBLPTRU

 A register to store the read byte TBLAT: keeps the data byte once it is fetched into the CPU. Use instruction TBLRD* to read the data from ROM (put it in TBLLAT)

TBLRD*	Table Read	After Read, TBLPRTR stays the sam
TBLRD*+	Table Read with Post-inc	Reads and inc. TBLPTR
TBLRD*-	Table Read with Post-dec	Reads and dec TBLPTR
TBLRD+*	Table Read with pret-inc	Increments TBLPTR and then reads

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Fixed Data in Program Space

- Program conter is 21-bit, which is used to point to any location in ROM space.
- How to fetch data from the code space? Known as a table processing: register indirect ROM addressing mode.
- To read the fixed data byte. We need an address pointer: TBLPTR: Points to data to be fetched
 - 21 bits as the program counter!!
 - Divided into 3 registers:

TBLPTRL, TBLPTRH, TBLPTRU (all parts of SFR)

- Is there any instruction to load 21 bits
- A register to store the read byte TBLLAT: keeps the data byte once it is fetched into the CPU

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Example Write a program to read a byte in MYDATA, and then put it to PORTD 000000 00004 Main ORG 0x0000 000000 000000 6A95 000002 0E18 000004 6EF6 000006 0E00 000008 6EF7 00000A 0E00 00000C 6EF8 00000E 0008 00004 00005 00006 00007 CLRF TRISD low MYDATA TBLPTRL high MYDATA TBLPTRH MOVLW MOVWF 00008 00009 00010 MOVLW MOVWF MOVLW MOVWF TBLRD* upper MYDATA TBLPTRU 00010 00011 00012 00013 00014 00015 00016 TABLAT, W 000010 50F5 000012 6E83 MOVE PORTD 000014 EF0A F000 goto here here 00017 00018 MYDATA DB "A" 000018 0041 00020

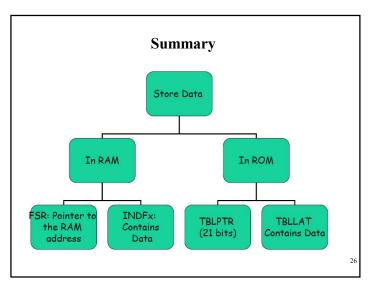
Example Assume that ROM space starting at 250H contains "Embedded System", write a program to send all characters to PORTD one byte at a time 00004 Main ORG 0x0000 CLRF 000000 6A95 000002 0E50 000004 6EF6 00005 00006 MOVLW low MYDATA 00007 MOVWF TBLPTRL 000004 6EF6 000006 0E02 000008 6EF7 00000A 0E00 00000C 6EF8 00000E 0009 000010 50F5 high YDATA MOVLW MOVWF TBĹPTRH upper MYDATA TBLPTRU MOVLW MOVWF TBLRD TABLAT, W MOVF 000012 E002 000014 6E83 000016 D7FB MOVWF PORTD BRA GOTO ORG 000018 EF0C F000 00017 EXIT 0x250 6D45 6562 6464 6465 5320 7379 6574 006D 00019 MYDATA "Embedded System", 0 00020 END

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Look-Up table

- Used to access elements of a frequently used with minimum operations
- Example: x^2
- We can use a look-up table instead of calculating the values **WHY?**
- Store the function f(x) in a table (RAM or ROM)
- We can get the base address of the table
- Input x provides the displacement

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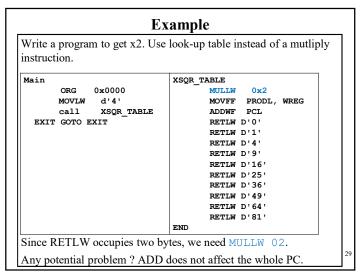
Look-Up table with RETLW

- RETLW (Return from subroutine with Literal to WREG) will provide the desired look-up table element in WREG
- Before execute RETLW, we need to a fixed value (displacement) to the PCL to index into the look-up table. So PC=PC+disp

And the corresponding RETLW is at the PC+disp.

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Look-Up table in RAM

Store data in a continue location

Using FSR as a pointer and the working register as an index

For example:

MOVFF PLUS2, PortD

Will copy data from location pointed by

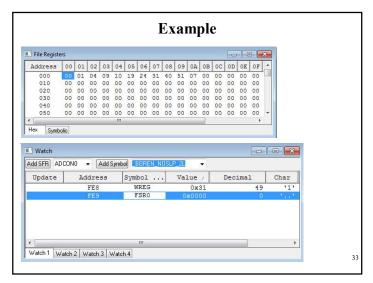
FSR2+WREG into PortD

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Example Write a program to get x2. Use look-up table instead of a mutliply instruction. Use TBLAT Input EQU 0 ORG 0x0FE ORG 0x0000 XSQR TABLE MOVLW d'9' db D'0', D'1', D'4', D'9', MOVWF input db D'16',D'25',D'36',D'49', call XSQR db D'64',D'81' EXIT GOTO EXIT END XSQR low XSQR TABLE MOVLW MOVWF TBLPTRL MOVLW high XSQR TABLE MOVWF TBLPTRH upper XSQR_TABLE MOVLW TBLPTRU MOVWF MOVF input, W TBLPTRL, F ADDWF MOVLW ADDWFC TRIPTRH ADDWFC TBLPTRU TBLRD* MOVF RETURN

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```
Example
Write a program to get x2. Use look-up table instead of a mutliply
instruction. Use FSR. FSR provide the base addr. W provide disp.
 table
        equ 0
                                         MOVLW d'7'
 Main
         ORG
               0x0000
                                         call XSQR
         LFSR 0, table
                                   EXIT GOTO EXIT
         MOVLW 0
         MOVWF POSTINCO
         MOVLW d'1'
                                    XSQR
                                         LFSR
                                                   0, table
         MOVWF POSTINCO
                                           MOVFF
                                                   PLUSWO, WREG
         MOVLW d'4'
         MOVWF POSTINCO
                                           RETURN
         MOVLW d'9'
                                           END
         MOVWF POSTINCO
         MOVLW d'16'
         MOVWF POSTINCO
         MOVLW d'25'
         MOVWF POSTINCO
         MOVLW d'36'
         MOVWF POSTINCO
         MOVLW d'49'
         MOVWF POSTINCO
         MOVLW d'64'
         MOVWF POSTINCO
         MOVWF POSTINCO
   ;for initialization of the table in RAM
```



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- 1. Address Mode
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Review: bit-addressability

- Often need to access individual bit of the port instead of the entire 8 bits.
- PIC18 provides instructions that alter individual bits without altering the rest of the bits in the port.
- Most common bit-oriented instructions:

fileReg
fileReg
e fileReg
fileReg, skip
fileReg, skip

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Stack and Subroutine Call

- In modern computers, we need a stack to store the return address during subroutine call. Besides, we need to save the register values that will be modified by the subroutine.
- In the PIC18, we only have a hardware stack for storing return address.
- How to solve this problem.

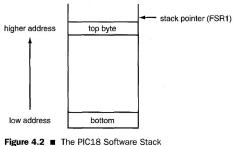
Stack and Subroutine Call

asciil EQU 0x0 ;input parameter gcU 0x1 ;input parameter cout bod EQU 0x2; output value local_var EQU 0x3 Main program CALL SUB1 CALL SUB1 CALL SUB1

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Stack and Subroutine Call

We can use FSR to implement a user stack.



Stack and Subroutine Call

Use some locations to store the affected registers

```
ascii 1 EQU 0x0 ;input parameter ascii 1 EQU 0x1 ;input parameter out bcd EQU 0x2 ;output value local var EQU 0x3 W TEMP SUB1 EQU 0x4 FTag TEMP SUB1 EQU 0x5 BSR_TEMP_SUB1 EQU 0x6
                                                                   MOVWF W_TEMP_SUB1
                                                                    MOVFF STATUS, STATUS TEMP) SUB1
                                                                    MOVFF BSR, BSR_TEMP_SUB1
                                                                    this routine may modify
                                                                    WREG, STATUS, BSR
 Main program
                                                                   MOVFF BSR_TEMP_SUB1, BSR
MOVFF STATUS_TEMP_SUB1, STATUS
 CALL SUB1
                                                                   MOVF W_TEMP_SUB1,W
                                                                   RETURN
CALL SUB1
```

Waste resource if we have many routines. Also, the routines cannot be re-entrence

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Stack and Subroutine Call

We can use FSR to implement a user stack.

Use LFSR to define a stack

LFSR FSR1, 0x500

The following instruction will push the STATUS register onto the stack:

MOVFF STATUS, PREINCL

The following instruction sequence will pull the top byte of the stack onto the STATUS register:

MOVFF POSTDEC1, STATUS

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```
Stack and Subroutine Call
          equ 0x200
         equ 0x200
equ 0
equ 1
equ 2
ORG 0x0000
LFSR 0, stack
MOVING DO
R0
R1
R2
Main
                                          ; define a user stack
          MOVWF RO
          MOVLW 0x10
          MOVWF R1
MOVLW 0x20
          MOVWF R2
           MOVLW 0x99
                                                                    FSR0=0201
FSR0=0202
FSR0=0203
          MOVFF R0, PRETINCO
                                          ; push R0 to stack
          MOVFF R1, PRETINCO
MOVFF R2, PRETINCO
MOVFF WREG, PRETINCO
                                           push R1 to stack
push R2 to stack
                                          ; push WREG to stack FSR0=0204
          MOVLW 0x71
MOVWF R0
MOVWF R1
                                           ; Now program change W, R0, R1, R2
          MOVWF R2
                                           restore the value from stack FSR0=0203
          MOVFF POSTDECO, WREG
          MOVFF POSTDECO, RO
                                            FSR0=0202
          MOVFF POSTDECO, R1
MOVFF POSTDECO, R2
                                            FSR0=0201
                                          ; FSR0=0200
          EXIT GOTO EXIT
 END
```

Stack and Subroutine Call

Use stacks store the affected registers

```
ascii 1 EQU 0x0 ;input parameter ascii h EQU 0x1 ;input parameter out bod EQU 0x2 ;output value local var EQU 0x3 stack EQU 0x200; define a stack Main program LFSR 0, stack ..... this routine may modify WREG, STATUS, BSR ..... MOVFF POSTDECO, BSR MOVFF POSTDECO, STATUS MOVFF POSTDECO, WREG RETURN
```

In general, modern computer systems use stack to pass parameters and to handle local variables. But the procedure is very complicated. If you are interested in this issue, read PIC Microcontroller: An Introduction to Software & Hardware Interfacing Chapter 4.