

Time & Space Complexity Cheat Sheet

BIG O NOTATION HIERARCHY (Fastest to Slowest)

$O(1) < O(\log n) < O(\sqrt{n}) < O(n) < O(n \log n) < O(n^2) < O(n^3) < O(2^n) < O(n!)$

Constant < Logarithmic < Root < Linear < Linearithmic < Quadratic < Cubic < Exponential

ACCEPTABLE TIME COMPLEXITIES FOR HACKERRANK

For $n = 10^5$ (typical Citadel constraint):

Complexity	Max n	Example
$O(1)$	Any	Hash lookup, array access
$O(\log n)$	10^{18}	Binary search
$O(\sqrt{n})$	10^{14}	Prime checking
$O(n)$	10^8	Single loop
$O(n \log n)$	10^6	Sorting, heap operations
$O(n^2)$	10^4	Nested loops (DANGER for $n=10^5$)
$O(n^3)$	500	Triple nested loops
$O(2^n)$	20	Subset generation

For 75-minute test with $n=10^5$:

- **SAFE:** $O(n)$, $O(n \log n)$
- **RISKY:** $O(n^2)$ will likely timeout
- **NO GO:** $O(n^3)$ or worse

PYTHON DATA STRUCTURES COMPLEXITY

List (Array)

Operation	Average	Worst	Notes
arr[i]	O(1)	O(1)	Index access
arr.append(x)	O(1)	O(1)	Add to end
arr.insert(i, x)	O(n)	O(n)	Insert at position
arr.pop()	O(1)	O(1)	Remove last
arr.pop(i)	O(n)	O(n)	Remove at position
arr.remove(x)	O(n)	O(n)	Remove by value
x in arr	O(n)	O(n)	Search
arr.sort()	O(n log n)	O(n log n)	In-place sort
sorted(arr)	O(n log n)	O(n log n)	New sorted list
arr.reverse()	O(n)	O(n)	In-place reverse
arr[::-1]	O(n)	O(n)	New reversed list
min(arr), max(arr)	O(n)	O(n)	Find min/max
Slicing arr[a:b]	O(b-a)	O(b-a)	Create sublist

Dictionary (Hash Table)

Operation	Average	Worst	Notes
d[key]	O(1)	O(n)	Access
d[key] = value	O(1)	O(n)	Insert/Update
del d[key]	O(1)	O(n)	Delete
key in d	O(1)	O(n)	Check existence
Iteration	O(n)	O(n)	All keys/values

Set

Operation	Average	Worst	Notes
<code>s.add(x)</code>	$O(1)$	$O(n)$	Add element
<code>s.remove(x)</code>	$O(1)$	$O(n)$	Remove element
<code>x in s</code>	$O(1)$	$O(n)$	Check membership
<code>s1 s2</code>	$O(\text{len}(s1) + \text{len}(s2))$		Union
<code>s1 & s2</code>	$O(\min(\text{len}(s1), \text{len}(s2)))$		Intersection
<code>s1 - s2</code>	$O(\text{len}(s1))$		Difference

Deque (Double-ended Queue)

Operation	Complexity	Notes
<code>append(x)</code>	$O(1)$	Add to right
<code>appendleft(x)</code>	$O(1)$	Add to left
<code>pop()</code>	$O(1)$	Remove from right
<code>popleft()</code>	$O(1)$	Remove from left
<code>d[i]</code>	$O(n)$	Random access (slow!)

Heap (Priority Queue)

Operation	Complexity	Notes
<code>heappush(h, x)</code>	$O(\log n)$	Insert
<code>heappop(h)</code>	$O(\log n)$	Remove min
<code>h[0]</code>	$O(1)$	Peek min
<code>heapify(arr)</code>	$O(n)$	Build heap from list

Counter

Operation	Complexity	Notes

Creation	$O(n)$	Count all elements
count[x]	$O(1)$	Get count
most_common(k)	$O(n \log k)$	Top k elements

COMMON ALGORITHM COMPLEXITIES

Sorting Algorithms

```
# Built-in sort -  $O(n \log n)$  time,  $O(n)$  space
arr.sort() # Timsort
sorted(arr)

# Counting sort -  $O(n + k)$  where k is range
# Only for integers in limited range
def counting_sort(arr, max_val):
    count = [0] * (max_val + 1)
    for num in arr:
        count[num] += 1

    result = []
    for num, freq in enumerate(count):
        result.extend([num] * freq)
    return result
```

Search Algorithms

```
# Linear search -  $O(n)$ 
def linear_search(arr, target):
    for i, val in enumerate(arr):
        if val == target:
            return i
    return -1

# Binary search -  $O(\log n)$ 
def binary_search(arr, target):
    left, right = 0, len(arr) - 1
    while left <= right:
        mid = (left + right) // 2
        if arr[mid] == target:
            return mid
        elif arr[mid] < target:
```

```

        left = mid + 1
    else:
        right = mid - 1
    return -1

```

Graph Algorithms

```

# BFS - O(V + E) time, O(V) space
# DFS - O(V + E) time, O(V) space
# Dijkstra - O((V + E) log V) with heap
# Bellman-Ford - O(VE)
# Floyd-Warshall - O(V3)
# Kruskal's MST - O(E log E)
# Prim's MST - O(E log V) with heap
# Topological Sort - O(V + E)

```

Tree Algorithms

```

# Tree traversal - O(n) time, O(h) space
# BST search - O(h) average, O(n) worst
# BST insert - O(h) average, O(n) worst
# Balanced tree operations - O(log n)

```

String Algorithms

```

# Pattern matching (naive) - O(nm)
# KMP pattern matching - O(n + m)
# Rabin-Karp - O(n + m) average
# String comparison - O(min(n, m))
# Substring search - O(n)

```

SPACE COMPLEXITY

Common Space Patterns

```

# O(1) - Constant space
def constant_space(arr):
    result = 0
    for num in arr:
        result += num
    return result

```

```

# O(n) - Linear space
def linear_space(arr):
    return arr[:] # Copy array

# O(n) - Hash table for frequency
def frequency_count(arr):
    freq = {} # O(n) space
    for num in arr:
        freq[num] = freq.get(num, 0) + 1
    return freq

# O(h) - Recursion depth for tree
def tree_height(root):
    if not root:
        return 0
    return 1 + max(tree_height(root.left), tree_height(root.right))

# O(2^n) - All subsets
def all_subsets(arr):
    result = []

    def backtrack(index, current):
        if index == len(arr):
            result.append(current[:])
            return

        # Include current element
        current.append(arr[index])
        backtrack(index + 1, current)
        current.pop()

        # Exclude current element
        backtrack(index + 1, current)

    backtrack(0, [])
    return result

```

OPTIMIZATION TECHNIQUES

1. Use Hash Table for O(1) Lookup

```

# SLOW - O(n^2)
for num in arr1:
    if num in arr2: # O(n) search

```

```

    result.append(num)

# FAST - O(n)
set2 = set(arr2)  # O(n) to build
for num in arr1:  # O(n)
    if num in set2: # O(1) lookup
        result.append(num)

```

2. Avoid Repeated Calculations

```

# SLOW - O(n2)
for i in range(n):
    total = sum(arr[:i])  # Recalculates sum each time

# FAST - O(n)
prefix = [0]
for num in arr:
    prefix.append(prefix[-1] + num)

```

3. Two Pointers Instead of Nested Loops

```

# SLOW - O(n2)
for i in range(n):
    for j in range(i+1, n):
        if arr[i] + arr[j] == target:
            return [i, j]

# FAST - O(n) with sorted array
left, right = 0, n-1
while left < right:
    total = arr[left] + arr[right]
    if total == target:
        return [left, right]
    elif total < target:
        left += 1
    else:
        right -= 1

```

4. Sliding Window Instead of Recalculating

```

# SLOW - O(n2)
for i in range(n - k + 1):
    window_sum = sum(arr[i:i+k])  # O(k) each time

```

```

# FAST - O(n)
window_sum = sum(arr[:k])
for i in range(k, n):
    window_sum += arr[i] - arr[i-k]  # O(1) update

```

5. Use Deque for Queue Operations

```

# SLOW - O(n2) because list.pop(0) is O(n)
queue = []
queue.append(x)
x = queue.pop(0)  # O(n)

# FAST - O(1) for all operations
from collections import deque
queue = deque()
queue.append(x)
x = queue.popleft()  # O(1)

```

6. Binary Search Instead of Linear Search

```

# SLOW - O(n)
for i, val in enumerate(sorted_arr):
    if val == target:
        return i

# FAST - O(log n)
left, right = 0, len(sorted_arr) - 1
while left <= right:
    mid = (left + right) // 2
    if sorted_arr[mid] == target:
        return mid
    elif sorted_arr[mid] < target:
        left = mid + 1
    else:
        right = mid - 1

```

COMMON MISTAKES TO AVOID

1. Nested Loops on Large Input

```
# TIMEOUT for n=10^5
```

```
for i in range(n):
    for j in range(n):  # O(n2)
        process(i, j)
```

2. Sorting Inside Loop

```
# TIMEOUT - O(n2 log n)
for i in range(n):
    sorted_sub = sorted(arr[:i])  # Sort every iteration
```

3. String Concatenation in Loop

```
# SLOW - O(n2) because strings are immutable
result = ""
for char in chars:
    result += char  # Creates new string each time

# FAST - O(n)
result = ''.join(chars)
```

4. Using List for Frequent Membership Testing

```
# SLOW - O(n) per lookup
seen = []
for num in arr:
    if num not in seen:  # O(n)
        seen.append(num)

# FAST - O(1) per lookup
seen = set()
for num in arr:
    if num not in seen:  # O(1)
        seen.add(num)
```

5. Unnecessary Deep Copies

```
# SLOW - Copies entire array each time
def backtrack(arr):
    if condition:
        result.append(arr[:])  # OK - need copy here

    for i in range(len(arr)):
```

```

new_arr = arr[:] # BAD - unnecessary copy
new_arr[i] = x
backtrack(new_arr)

# FAST - Modify in place
def backtrack(arr):
    if condition:
        result.append(arr[:])

    for i in range(len(arr)):
        old_val = arr[i]
        arr[i] = x # Modify
        backtrack(arr)
        arr[i] = old_val # Restore

```

QUICK COMPLEXITY CHECKS

Before implementing, ask:

1. **What's n?** (array length, string length, etc.)
2. **How many nested loops?** Each adds $O(n)$
3. **Am I sorting?** That's $O(n \log n)$
4. **Am I using hash table?** Lookups are $O(1)$
5. **Am I searching unsorted array?** That's $O(n)$
6. **Will this timeout?**
 - $n=10^5$ and $O(n^2) \rightarrow$ YES
 - $n=10^5$ and $O(n \log n) \rightarrow$ NO
 - $n=10^3$ and $O(n^2) \rightarrow$ NO

RULE OF THUMB FOR CITADEL

- **$n \leq 10$:** $O(n!)$ is acceptable (brute force permutations)
- **$n \leq 20$:** $O(2^n)$ is acceptable (subset enumeration)
- **$n \leq 500$:** $O(n^3)$ is acceptable
- **$n \leq 10^4$:** $O(n^2)$ is acceptable
- **$n \leq 10^5$:** $O(n \log n)$ or $O(n)$ required

- $n \leq 10^6$: $O(n)$ or $O(\log n)$ required

For $n=10^5$ (most Citadel problems): $O(n)$, $O(n \log n)$ $O(n^2)$ - likely timeout $O(n^3)$ or worse - guaranteed timeout