

Set Theory - Supplementary

COMP9020 Tutorial

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1 Tutorial Outline

1.1 Definition

Ask yourself these questions below to test your understanding.

- What is the definition of the *subset* and *equality* relations of sets?
 - What is the *basic method* for proving them (respectively)?
 - What other methods do we have for proving or analysing them?
- What are the *empty set* and the *universe*?
 - What is the relationship between them?
- What operations can be performed on sets?
 - What are the definitions of them?
 - (How to do when we encounter *set difference*?)
- What do the following Set Theory Laws entail?
 - **Five basic laws:** Commutative, Associative, Distributive, Identity, and Complement Law
 - **Derive laws:** Idempotence, Double Complement, Annihilation, Absorption, and De Morgan's Law
 - What are **two useful results** mentioned in our lecture?
- How to calculate the *cardinality* of a *finite* set?

1.2 Brainstorming

The following questions are open and have no standard answers.

- What can the elements of a set be? Are there any restrictions?
- What is the connection between *element proof* and *proof by laws*?
- Why aren't Venn Diagrams rigorous for proofs?
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2 Explanation for Some Exercises

This section is intended to supplement the parts I didn't explain clearly or didn't have time to cover in class. If there are any proofs you provided that are more concise than mine, I will include them as well.

Exercise 7: For sets A and B , prove or disprove $\mathcal{P}(A \cup B) = \mathcal{P}(A) \cup \mathcal{P}(B)$.

Analysis¹:

1. To show $\mathcal{P}(A \cup B) \subseteq \mathcal{P}(A) \cup \mathcal{P}(B)$:

For any element $x \in \mathcal{P}(A \cup B)$

$$\Rightarrow x \subseteq A \cup B$$

$$(?) \Rightarrow x \subseteq A \text{ or } x \subseteq B$$

$$\Rightarrow x \in \mathcal{P}(A) \text{ or } x \in \mathcal{P}(B)$$

$$\Rightarrow x \in \mathcal{P}(A) \cup \mathcal{P}(B).$$

2. To show $\mathcal{P}(A) \cup \mathcal{P}(B) \subseteq \mathcal{P}(A \cup B)$:

For any element $x \in \mathcal{P}(A) \cup \mathcal{P}(B)$

$$\Rightarrow x \in \mathcal{P}(A) \text{ or } x \in \mathcal{P}(B)$$

$$\Rightarrow x \subseteq A \text{ or } x \subseteq B$$

$$(?) \Rightarrow x \subseteq A \cup B$$

$$\Rightarrow x \in \mathcal{P}(A \cup B).$$

For the content in these two boxes, we don't have a direct rule to apply. Therefore, we need to consider whether they are correct. (Please note that the essence of x is a set.)

The second box is **correct**.

For any element $x' \in x$

$$\Rightarrow x' \in A \text{ or } x' \in B$$

$$\Rightarrow x' \in A \cup B.$$

But when we consider the first box, it **cannot be directly proven**.

For any element $x' \in x$

$$\Rightarrow x' \in A \cup B$$

$$\Rightarrow x' \in A \text{ or } x' \in B.$$

This is because x' is a particular but arbitrary element of x , and we have shown that it may be in A or in B . This means that some elements of the set x are in A , while others are in B . Therefore, $x \subseteq A$ or $x \subseteq B$ may both not hold simultaneously.

Additionally, we have also identified the characteristics of a counterexample—some elements in $A \cup B$ belong **only** to A , while others belong **only** to B .

¹Please note that if you can directly find a counterexample, such thinking is unnecessary. This analysis is only for students who totally have no idea how to approach the problem.

Disproof:

Counterexample: Let $A = \{1\}$, $B = \{2\}$. Then

$$\mathcal{P}(A \cup B) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\} \neq \{\emptyset, \{1\}, \{2\}\} = \mathcal{P}(A) \cup \mathcal{P}(B). \quad \square$$

Exercise 12: Let A, B, C be languages with alphabet $\{a, b\}$ where $AC = BC$. Prove or disprove $A = B$.

Analysis:

This is a problem that requires constructing a counterexample, so it is quite challenging. Additionally, we don't have a general approach to thinking about this type of problem. Here, I will just share my own analysis.

I initially tried to **prove** that this was a *correct* conclusion. Starting with a simple case, I assume C is a finite set. Then, I followed with an *element proof*, but the condition $AC = BC$ made it difficult to handle the condition $x \in A$ or $x \in B$.

Therefore, I sought to improve this approach. Fortunately, I quickly discovered a useful property: if $AC = BC$, then the shortest words in A and B must be the same.² Thus, I **refined** the element analysis into this algorithm:

Given that $AC = BC$, assume $C = \{c_1, c_2, \dots, c_n\}$.

- The shortest word w in A and B is the same;
- remove w , obtaining A' and B' ;
- prove $A'C = B'C$;
- the shortest word w' in A' and B' is the same;
- remove w' , obtaining A'' and B'' ;
- ... and so on.

This provides a method to reduce A and B while maintaining $AC = BC$. Therefore, if A and B are also finite sets, this algorithm comes to an end in a finite step and implies A and B must contain the same element.

However, I found it difficult to prove $A'C = B'C$, because when we remove w , we cannot ensure that wc_1, wc_2, \dots, wc_n are simultaneously removed or retained in both AC and BC .

After a long time, I realized that this indeed cannot be held. Thus, I **shifted** from proving to disproving and attempted to construct a counterexample through this breakthrough³, which led to the following disproof.

Disproof:

Counterexample:

Let $A = \{\lambda, aa\}$, $B = \{\lambda, a, aa\}$, and $C = \{a, aa\}$. Then $A \neq B$ and

$$AC = \{a, aa, aaa, aaaa\} = BC. \quad \square$$

²This proposition is not obvious, and A , B , and C don't have to be finite sets. I will leave it to you, and if you're interested in exploring, you can try to prove it:)

³When constructing, please note that the fact that the shortest words in A and B are the same still holds. Through the proof of this proposition, we can also deduce that if A , B , and C are finite sets, the longest words in A and B are also the same.

Exercise 13: Let L_1 and L_2 be languages over Σ . Prove that

$$(L_1 \cap L_2)^* \subseteq L_1^* \cap L_2^*.$$

Analysis:

As we analyzed in our tutorial, I will directly list the process here.

For any element $x \in (L_1 \cap L_2)^*$

$$\begin{aligned} \Rightarrow x &\in (L_1 \cap L_2)^0 \cup (L_1 \cap L_2)^1 \cup \dots && \text{(def of Kleene Star)} \\ \Rightarrow x &\in (L_1 \cap L_2)^0 \text{ or } (L_1 \cap L_2)^1 \text{ or } \dots && \text{(def of } \cup) \\ \Rightarrow x &\in (L_1 \cap L_2)^k \text{ for some integer } k && \text{(def of } or) \\ \Rightarrow x &= w_1 w_2 \dots w_k, \text{ where each } w_i \in L_1 \cap L_2 && \text{(prop of Concatenation)} \\ \Rightarrow w_i &\in L_1 \text{ and } w_i \in L_2 && \text{(def of } \cap) \\ \Rightarrow x &\in L_1^k \text{ and } x \in L_2^k && \text{(prop of Concatenation)} \\ \Rightarrow (x &\in L_1^0 \text{ or } L_1^1 \text{ or } \dots) \text{ and } (x \in L_2^0 \text{ or } L_2^1 \text{ or } \dots) && \text{(def of } or) \\ \Rightarrow (x &\in L_1^0 \cup L_1^1 \cup \dots) \text{ and } (x \in L_2^0 \cup L_2^1 \cup \dots) && \text{(def of } \cup) \\ \Rightarrow x &\in L_1^* \text{ and } x \in L_2^* && \text{(def of Kleene Star)} \\ \Rightarrow x &\in L_1^* \cap L_2^* && \text{(def of } \cap) \end{aligned}$$

The content of *prop of Concatenation*⁴ is:

Given that A is a language and n is an integer.

- For any word $w \in A^n$, w can be expressed as $w_1 w_2 \dots w_n$, where $w_i \in A$.
- Conversely, if any word w can be expressed as $w_1 w_2 \dots w_n$, where $w_i \in A$, then $w \in A^n$.

Proof:

Please refer to the solution provided in the materials.

⁴We briefly discussed its proof in our tutorial, and I hope you still remember it;) In fact, after learning **mathematical induction**, we can prove this property in a more rigorous and formal way.