## Chapter 4: The Well-Founded Sets

4-1. The class WF of well-founded sets

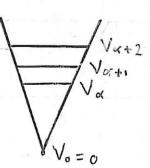
Def: We define by Wansfirste indudion on 
$$\alpha \in ON$$

$$\begin{cases} V_0 = 0 \\ V_{\alpha+1} = \mathcal{P}(V_{\alpha}) \\ V_{\alpha} = U\{V_{\zeta}: \zeta \in \alpha \ \zeta \ , \ \text{for} \ \alpha \ \text{ limit} \end{cases}$$

called the class of well-founded sets: this name will become clear later

$$V_{1} = P(0) = \{0\}$$
 $V_{2} = P(V_{1}) = \{0, \{0\}\} = \{0, 1\}$ 
 $V_{3} = P(V_{2}) = \{0, 1, 2, \{1\}\}$ 

Lemma 1: i) HXEON, Vor is transchire ii) « { B => Vx c VB.



Proof: i) by induction on 
$$\alpha$$
:

 $\alpha = 0 \implies V_0 = 0 \quad \text{bransitive}$ 
 $\alpha = \beta + 1 \quad \text{and} \quad V_{\beta} \quad \text{bransitive}$ 
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 $\alpha = \beta +$ 

· B = X+1: If x=B, then Vx = VB.

If x < B, then x < X => Vx = Vx = P(Vx) = Vx+1

since Vy+1 = VB. transitive by (i).

If a=B, Va=VB

If a < B, Va = VB by def of VB.

Remark:

If  $x \in WF$ , then the least of s.t.  $x \in Vq$  is a successor ordinal Indeed: suppose it is limit, then by def Vq = U Vz, but then X & Voy means XEVy for some < < , contradiction with the minimality of <.

If  $x \in WF$ , let  $\underline{nank(x)}$  or  $\underline{p(x)}$  be the least or  $\underline{s.t.}$   $\underline{x} \in V_{X+1} \leftarrow \underline{w} \in \underline{k}$  we know that the least ordinal  $\underline{y}$  s.t.  $\underline{x} \in V_{X}$  is successor Def:

The rank of x is the level that preceeds the "first" appearance of x.

Thus rank (x) = B implies:  $(x \in V_{\beta+1} = \mathcal{V}(V_{\beta}) \Rightarrow) \times \subseteq V_{\beta}$ x & VB by minimality of B as a rank XE Vox, Yor> B, since the Vox's are increasing with respect to inclusion. Remark: characterization of each set  $V_{\varphi}$ :  $V_{\varphi} = \begin{cases} x \in WF : \text{ vank}(x) \leqslant \varphi \end{cases}$ Froof: (a)  $x \in V_{\varphi} \Rightarrow x \in WF$  and  $v_{\varphi}(x) \leqslant \varphi \end{cases}$ Thus  $V_{\varphi+1} \subseteq V_{\varphi}$  and hence  $v_{\varphi}(x) \in V_{\varphi}$ .

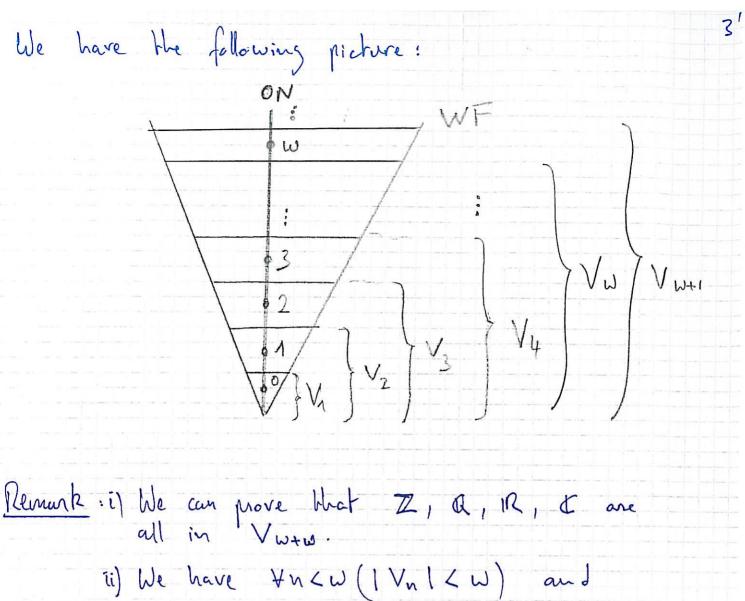
Thus  $V_{\varphi+1} \subseteq V_{\varphi}$  and hence  $v_{\varphi}(x) \in V_{\varphi}$ .

Proof: exercices

let x, y \in WF s.t. x \in y. Then 3 \langle s.t. x \in y \in V\_q => (V\_q hransihira) \times \in V\_q =) \times \in WF lemma 1.

So WF "look nice", no parhological properties.

Penna 3: Hacon (acWF ~ rank(x) = x) Proof: We show by induction on & that · If  $\alpha = 0$ , then  $\alpha \in V_1 = \{0\}$  and  $\alpha \notin V_0 = 0$ · If <= B+1= Bus B3. By I. H. we Nave BEVB+1. Hence: 19) BE VB+2 = VX+1 ( Since VB+1 = VB+2) 2°) {\begin{align\*} & \begin{align\*} & \ Then {B} = VB => B & VB, contradiction with I.H. Thus ad Va. · If of limit: If BE ∝ (i-e. B < x) Hun by I.H., B ∈ VB+1 ⊆ Va (sing B+1 & a), thus B € Va Hance & = Va => X E Vati Now assume  $\angle EVQ = UVZ$  (since of limit) Then JEZX G.t. X EVT = 5 TEVS (sing VS)
Contradiction with I.H. Thus X & Vor



| Vw | = W = 2.

iii) The cardinalities of Vox increase exponentially  $|V_{W+1}| = |P(V_{W})| = |P(W)| = 2^{20}$   $|V_{W+2}| = 2^{20}$ , etc...

More precisely, we can prove by induction on & that | Vw+x | = ] x

All mathematics take place in WF!

where I a is defined by induction on a by: 10 = W = 20 7 at1 = 2 7 a In = sup { I ?: 7 < A], for A limit

We will introduce the axiom of foundation and see that assuming it is equivalent to stating V = WF, i.e. every set is well-founded.

We need the following definitions...

Def(2F-P): A relation R is well-founded on a set A iff  $\forall X \subseteq A \left[ X \neq 0 \rightarrow \exists y \in X \left( \forall z \in X \left( \neg z R y \right) \right) \right]$ 1.2. every non-empty subset X of A has an R-minimal element.

Def (2F--P): bet A be a set, Define by induction on N: satisfications And induction on N: satisfications And N:

A=U<sup>0</sup>A U<sup>0</sup>A = A, U<sup>n+1</sup>A = U(U<sup>n</sup>A)

Existing the set  $C(A) = U(U^nA)$ Existing the we set  $C(A) = U(U^nA) = U(U^nA)$ Existing the set  $C(A) = U(U^nA) = U(U^nA) = U(U^nA)$ Existing the set  $C(A) = U(U^nA) = U(U^nA) = U(U^nA) = U(U^nA)$ Existing the set  $C(A) = U(U^nA) = U(U^n$ 

So  $cl(A) = A \circ UA \circ U^2A \circ ...$  (i.e. contains the el. of  $A + el. \circ f el. \circ f el. \circ f A + ...$ 

Lemma 4: Cl(A) is the least transitive set (ZF-P) containing A.

Proof: = If xeyecl(A), then Inst. ye U"A => x & U (U " A) = U"+1 A => x & d(A) · Thus cl (A) trunsitive. cl(A) 2 A by def · let T transhive s.t. T 2 A An indudica on u shows UMA ST Yn ∈ W. Thus d(A) = U,UMA: NEW} ⊆ T details of the induction: UOA = A S T let xe. Un+1 A = U (1) A) => => => gil. X ey & UMA and UMA GT So X Ey & T => X & T ... I.H. We have  $cl(A) = A v \{cl(x) : x \in A\}$  (exercise) let T = AU { cl(x): XEA } . Then T transitive = A = d(A) ET. · PA = cl(A) and x & A => X & cl(A) => X C Cl (A) (b) brans livity) =) cl(x) & cl(A) (by minimality

Thus Aufcliti: XEAJ = T = cl(A).

Ihm 5: Let A be a set. The following are equivalent: i) A E WF ii) d(A) & WF € is well-founded on cl (A). i) => ii): If A & WF, then by induction Proof: UOA = AEWF on n, UnAEWF, Ynew EVa, for some d (since WF closed under U). By knowstivity of WF, U"A EWF, YNEW Thus cl(A) & WF i.e. cl(A) = Vy for some or, thus cl(A) ∈ Vv+1 => cl(A) ∈ WF. ii) = ) iii): let X C cl(A) X = 0.
By (ii), CP(A) EWF = ) CP(A) EWF = ) X EWF also, = 1 we can consider

Q = min frank(y): y E X }

XEWF => Vy EX. y XEWF=> YyEX, YEWF => rankly) exist, YyEX. and let yetwith ranking) = 4. If y not E-minimal in X, By'E.X s.t. y'Ey. Thus route (y') < rank (y) = or (prop-29) contradiction with minimality of x.

iii) => i) We first show that cP(A) =WF Suppose cf(A) & WF. let X = cl(A) \ WF \ \defo let y E-min in X (exists by hyp.) If zey EX Ecl(A), Hen (2 & X (by min of y in X.) (ZE cl(A) (since cl(A) transitive.) =) 7 E WF (c.f. figure, Zecl(A) and Zd X =) ZeWF) =) y C WF => y E WF (y => y E Vx+1 CWF), contradiction with y & X = cl(A) \ WF. Thus d(A) = WF. Thus A E cl (A) & WF => A & WF ⇒ A € WF (A € Vq =) A € Vq+1 => A € WF)

Remark: the def of WF uses the power set axiom, but the equivalent formulation via trans: live closure and well-foundation of E doesn't!

Since all mathematics take place in WF, it is "reasonable" to adopt an axiom stating V = WF (:.e. every set is well-founded).

It restricts our domain of discourse to "not public losical" sets, but still permits to do all mathematics...

Since the statement V=WF is highly nonelementary to write, we give an equivalent formulation. In the pure language of out theory

Axiom 2: Foundation

Vx[∂y(yex) → ∂y(yex ∧ ¬∂2(2ex ∧ 2ey))]

i.e. if X = 0, then 3y = x (xny = 0)

i.e. every set X ≠ 0 has an E-minimal element.

The axiom of foundation is equivalent to V=WF.

Thun 6 (ZF-) The following one equivalent:

- i) the Axiom of Foundation
- ii) YA (E is well-founded on A)
- 6ci) V = W F

Proof: i) (> ii) inmediate by def of AF. ii sonpose iii (iii (= /iii We have obviously WFGV let  $A \in V$ , then by hyp,  $\epsilon$  is w-f on cl(A). By hum 5, AEWF. Thus VEWF (ii) ⇒ (i) Suppose V= WF. let AEV and XE As.t. X = 0 Then AEWF by hyp. Thus A = WF >X = WF This ensures the existence of ~= min { runk(y): y e X } Let y EX st. rank(y) = x. Then y E-min in X (otherwise ZEX s.t. ZEy implies rank(z) < rank(y) = of contradition with minimality of of ) So under AF, the universe Looks like this V=WF  $\left\{\begin{array}{c} 1 \\ 1 \\ 1 \\ 1 \end{array}\right\}$