

**Math 724, Fall 2013**  
**Homework #7**

**Instructions:** Write up your solutions in LaTeX and hand in a hard copy in class on **Monday, December 9**. Collaboration is allowed (and in fact encouraged), but each student must write up his or her solutions independently and acknowledge all collaborators.

**Problem #1 Bogart, Chapter 5, Supplementary Problem #2:** There are  $m$  students attending a seminar in a room with  $n$  seats. The seminar is a long one, and in the middle the group takes a break. In how many ways may the students return to the room and sit down so that nobody is in the same seat as before?

For a set  $S$  of students, let  $i = |S|$  and let  $A_S$  be the set of reseatings so that everyone in  $S$  (and possibly others) gets his/her own seat. An element of  $A_S$  is given by a way to seat the  $m - i$  students not in  $S$  in the remaining  $n - i$  other seats, so  $|A_S| = (n - i)!/(n - m)!$ . By inclusion, the answer is

$$\sum_{i=0}^m (-1)^i \binom{m}{i} \frac{(n-i)!}{(n-m)!} = \sum_{i=0}^m (-1)^i \binom{m}{i} (n-i)^{\overline{m-i}}.$$

**Problem #2 Bogart, Chapter 5, Supplementary Problem #5.** Suppose that  $n$  children join hands in a circle for a game at nursery school. The game involves everyone falling down (and letting go). In how many ways may they join hands in a circle again so that nobody has the same person immediately to the right both times the group joins hands?

Label the children  $1, \dots, n$ . For a set  $C$  of children, let  $S_C$  be the set of rearrangements in which each child in  $C$  has the same right-hand neighbor (as though they had never let go). Then  $|S_C|$  is the number of circular rearrangements of  $n - |C|$  objects (where the “objects” are connected chains of children), so  $|S_C| = (n - i - 1)!$  where  $i = |C|$  (unless  $C = [n]$ , in which case  $|S_C| = 1$  rather than “ $(-1)!$ ”). By inclusion/exclusion, the answer is

$$(-1)^n + \sum_{i=0}^{n-1} (-1)^i \binom{n}{i} (n - i - 1)!.$$

**Problem #3 Bogart, Chapter 5, Supplementary Problem #8.** Suppose we take two graphs  $G_1$  and  $G_2$  with disjoint vertex sets, we choose one vertex on each graph, and connect these two vertices by an edge  $e$  to get a graph  $G_{12}$ . How does the chromatic polynomial of  $G_{12}$  relate to those of  $G_1$  and  $G_2$ ?

Call the vertices in question  $v_1$  and  $v_2$ . Given a proper  $k$ -coloring of  $G_1$  and a proper  $k$ -coloring of  $G_2$ , they together comprise a proper  $k$ -coloring of  $G_{12}$  iff  $v_1$  and  $v_2$  are assigned different colors. This will happen exactly  $(k - 1)/k$  of the time (since any particular color is equally likely to be assigned to any vertex). Therefore

$$\chi(G_{12}, k) = \frac{k-1}{k} \chi(G_1, k) \chi(G_2, k).$$

**Problem #4** Let  $C_n$  be the cycle graph of length  $n$  — that is, the undirected graph with vertices  $v_1, v_2, \dots, v_n$  and edges  $v_1 v_2, v_2 v_3, \dots, v_{n-1} v_n, v_n v_1$ . Find a formula for the chromatic polynomial of  $C_n$ .

Recall the inclusion/exclusion formula for the chromatic polynomial of a graph  $G = (V, E)$ :

$$\chi(G, k) = \sum_{A \subseteq E} (-1)^{|A|} k^{c(A)}$$

where  $c(A)$  means the number of components of the graph  $(V, A)$ . If  $G = C_n$ , then  $c(A) = n - |A|$  for all  $A \subsetneq E$  (for then  $A$  is acyclic) and  $c(E) = 1$ . So the formula gives

$$\begin{aligned}\chi(C_n, k) &= \left( \sum_{a=0}^{n-1} \binom{n}{a} (-1)^a k^{n-a} \right) + (-1)^n k \\ &= \left( \sum_{a=0}^n \binom{n}{a} (-1)^a k^{n-a} \right) + (-1)^n (k-1) \\ &= \boxed{(k-1)^n + (-1)^n (k-1)} \quad \text{by the binomial theorem.}\end{aligned}$$

**Problem #5** Let  $Y_n$  be the graph  $K_n$  with one edge removed. Find a formula for the chromatic polynomial of  $Y_n$ .

Let the missing edge be  $e$ . Deletion-contraction says that  $\chi(K_n, e) = \chi(Y, e) - \chi(K_{n-1}, e)$ . (The graph  $K_n/e$  by contracting an edge is not simple, but when calculating the chromatic polynomial we can ignore multiple edges, and the underlying simple graph of  $K_n/e$  is  $K_{n-1}$ . Therefore

$$\begin{aligned}\chi(Y, e) &= \chi(K_n, e) + \chi(K_{n-1}, e) = k(k-1) \cdots (k-n+1) + k(k-1) \cdots (k-n+2) \\ &= \boxed{(k(k-1) \cdots (k-n+3))(k-n+2)^2}.\end{aligned}$$

**Problem #6** ... Let  $A(G)$  denote the set of acyclic orientations of  $G$ , and let  $\alpha(G) = |A(G)|$ .  
(a) What is  $\alpha(G)$  if  $G \dots$

- is a forest?
- is  $K_n$  (the complete graph with  $n$  vertices)?
- is  $C_n$  (the cycle graph)?

If  $G$  is a forest with  $e$  edges then  $\alpha(G) = 2^e$ , because every orientation is automatically acyclic. There are exactly two orientations of  $C_n$  that are non-acyclic, so  $\alpha(C_n) = 2^n - 2$

Let  $\mathcal{O}$  be an acyclic orientation of  $K_n$ . I claim that  $\mathcal{O}$  has a unique **sink** (a vertex with no out-edges). To see this, construct a walk by starting at any vertex and following edges in the direction of the arrows. The condition that there is no sink means that the walk can be extended indefinitely — but that means that eventually some vertex will have to be repeated, which means we have traversed a cycle, which is impossible. On the other hand, there cannot be more than one sink (if there were, how could the edge between them be oriented?) Therefore, the orientations on  $K_n$  are in bijection with the total orderings of the vertices (find the unique sink, delete it, find the unique sink of the  $K_{n-1}$  remaining, delete it, ...) It follows that  $\alpha(K_n) = n!$ .

(b) Show that  $\alpha(G) = \alpha(G - e) + \alpha(G/e)$  for any edge  $e$ .

Let  $e = vw$ . Consider the map  $\pi : A(G) \rightarrow A(G - e)$  given by forgetting the orientation of  $e$ . Given any  $\mathcal{O} \in A(G - e)$ , we can add an edge  $v \rightarrow w$  (resp.,  $w \rightarrow v$ ) and obtain an acyclic orientation of  $G$  provided that  $\mathcal{O}$  contains no path from  $w$  to  $v$  (resp.,  $v$  to  $w$ ). But  $\mathcal{O}$  cannot contain both such paths — otherwise it would have a cycle! Therefore it is always possible to extend  $\mathcal{O}$  to an orientation of  $G$ , i.e.,  $\pi$  is surjective.

Call an orientation  $\mathcal{O} \in A(G - e)$  *one-way* or *two-way* depending on the cardinality of  $\pi^{-1}(\mathcal{O})$ . The two-way orientations correspond bijectively to the acyclic orientations of  $G/e$  (since if there is no path in either direction between  $v$  and  $w$  then identifying them will not produce a cycle). Therefore,

$$\begin{aligned}\alpha(G) &= \#\{\text{one-way } \mathcal{O} \in A(G - e)\} + 2\#\{\text{two-way } \mathcal{O} \in A(G - e)\} \\ &= |A(G - e)| + \#\{\text{two-way } \mathcal{O} \in A(G - e)\} \\ &= \alpha(G - e) + \alpha(G/e).\end{aligned}$$

(c) Find a formula for  $\alpha(G)$  in terms of the chromatic polynomial  $\chi_G(k)$ .

The formula is  $\alpha(G) = (-1)^n \chi_G(-1)$ , where  $n = n(G)$  means the number of vertices. We'll prove this by induction on the number of *edges*. If  $G$  has no edges, then  $\alpha(G) = 1$  and  $\chi(G, k) = k^n$ , so

$$(-1)^n \chi_G(-1) = (-1^n)(-1^n) = 1.$$

Meanwhile, if  $G$  has a loop then  $\chi(G, k)$  and  $\alpha(G)$  are both zero. Finally, if  $e$  is a nonloop edge, then by induction we have

$$\begin{aligned} \alpha(G) &= \alpha(G - e) + \alpha(G/e) && \text{(previous problem)} \\ &= (-1)^n \chi(G - e, -1) + (-1)^{n-1} \chi(G/e, -1) && \text{(induction; note that } n(G) = n(G - e) = n(G/e) - 1) \\ &= (-1)^n (\chi(G - e, -1) - \chi(G/e, -1)) \\ &= (-1)^n (\chi(G, -1)) && \text{(by the chromatic recurrence).} \end{aligned}$$

as desired.

**Problem #7** Let  $V = [n]$  and let  $G$  be a graph with vertex set  $V$ . The *chromatic symmetric function*  $X(G)$  of  $G$  is the formal power series defined by

$$X(G) = \sum_f \prod_{i=1}^n x_{f(i)}$$

where the sum ranges over all proper colorings  $f$ .

(a,b) What is  $X(G)$  if  $G$  has no edges or if  $G = K_n$ ?

If  $G$  has  $n$  vertices and no edges, then every coloring is proper, so

$$X(G) = \left( \sum_{i=1}^{\infty} x_i \right)^n.$$

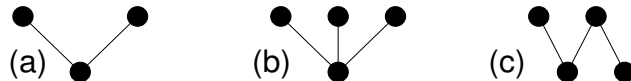
Meanwhile, if  $G = K_n$  then the only colorings possible are those in which every vertex gets a distinct color. Given a set of  $n$  distinct colors, there are  $n!$  ways to use them to properly color  $K_n$ , so

$$X(K_n) = n! \sum_{i_1 < \dots < i_n} \prod_{j=1}^n x_{i_j}.$$

(c) Let  $k$  be a positive integer. Explain how to derive the number of proper  $k$ -colorings from  $X(G)$ .

Plug in  $x_1 = x_2 = \dots = x_k = 1$  and  $x_m = 0$  for all  $m > k$ . The result is the number of colorings that use only the colors  $1, \dots, k$  — that is, the number of  $k$ -colorings.

(d) Show that the two trees shown in (b) and (c) below do *not* have the same chromatic symmetric function.



It is possible to properly color tree (b) by coloring the top three vertices blue and the bottom one red. OTOH, every red/blue coloring of tree (c) uses each color twice. Therefore the coefficients of  $x_1^2 x_2^2$  in the CSFs of the two trees are unequal. More specifically,

$$X(\text{Tree (b)}) = 24m_{1111} + 3m_{211} + m_{31}, \quad X(\text{Tree (c)}) = 24m_{1111} + 3m_{211} + 2m_{22}$$

where  $m_\lambda$  is the monomial symmetric function corresponding to  $\lambda \vdash n$  (that is, the sum of all monomials whose exponents are the parts of  $\lambda$ ).

**Extra credit.** Let  $n$  be a positive integer. For  $1 \leq i < j \leq n$ , define a hyperplane  $H_{ij} \subset \mathbb{R}^n$  by

$$H_{ij} = \{(x_1, x_2, \dots, x_n) \in \mathbb{R}^n : x_i = x_j\}.$$

Now let  $G$  be a simple graph with vertex set  $[n]$  and edge set  $E$ , and let

$$\mathcal{A}_G = \bigcup_{ij \in E} H_{ij}.$$

Thus  $\mathcal{A}_G$  is a subset of  $\mathbb{R}^n$  (it is called a *graphical hyperplane arrangement*). In terms of  $G$ , how many connected components does  $\mathbb{R}^n \setminus \mathcal{A}_G$  have?

Each component of  $\mathbb{R}^n \setminus \mathcal{A}_G$  is defined by a collection of inequalities: for each edge  $ij \in E(G)$ , either  $x_i < x_j$  or  $x_j < x_i$ . Think of these inequalities as corresponding to orienting the edge  $ij$  as either  $i \rightarrow j$  or  $j \rightarrow i$  respectively. This orientation is acyclic, because a cycle would correspond to an infeasible list of inequalities. Furthermore, every acyclic orientation gives rise to at least one point in  $\mathbb{R}^n \setminus \mathcal{A}_G$ . (For example, assign the value 0 to every sink. Delete all the sinks and assign the value  $-1$  to all sinks of the resulting graph. Delete all the sinks and assign the value  $-2 \dots$ ) So there is a bijection between connected components of  $\mathbb{R}^n \setminus \mathcal{A}_G$  and acyclic orientations!