

# Lecture 5: Search with Other Agents

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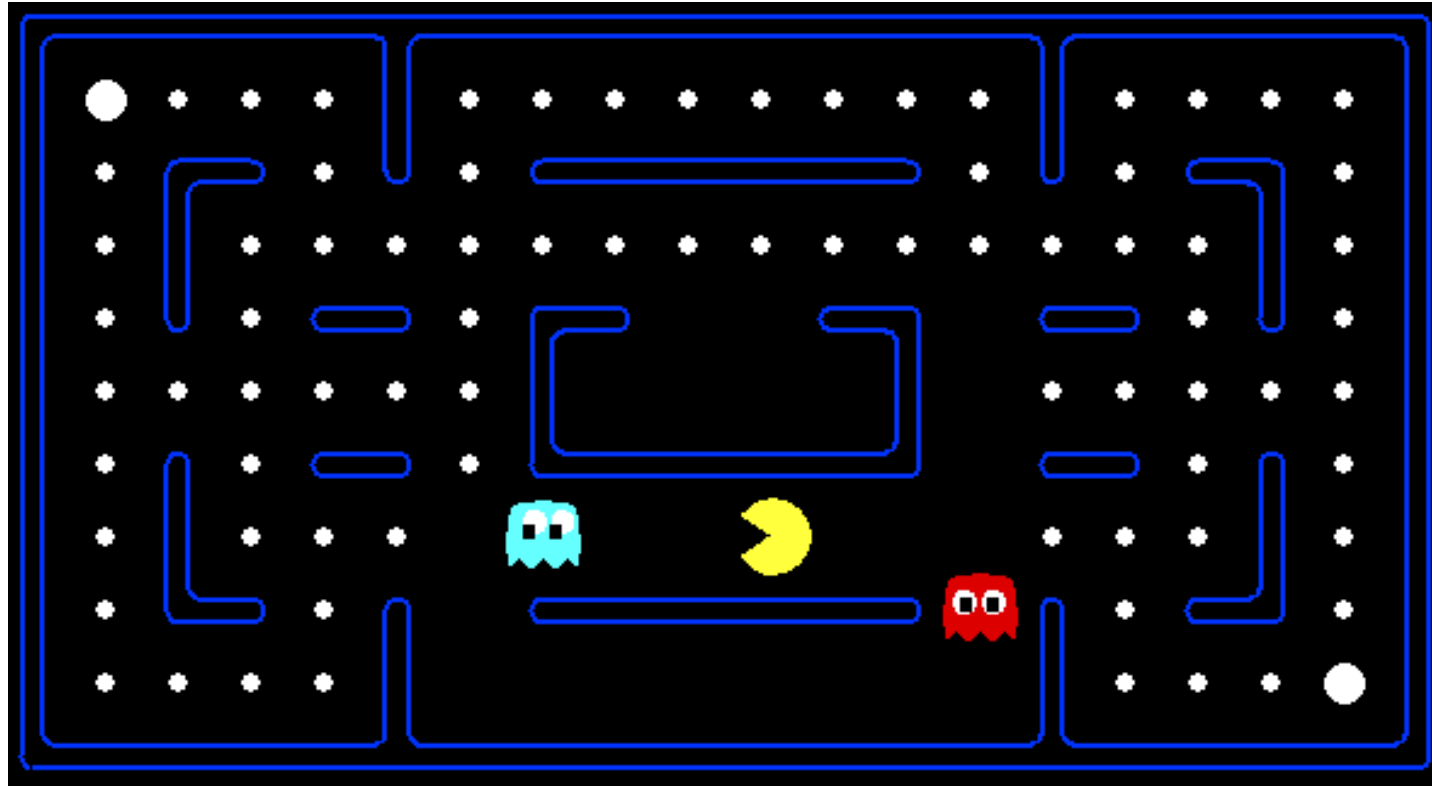
<https://shuaili8.github.io>

<https://shuaili8.github.io/Teaching/CS410/index.html>

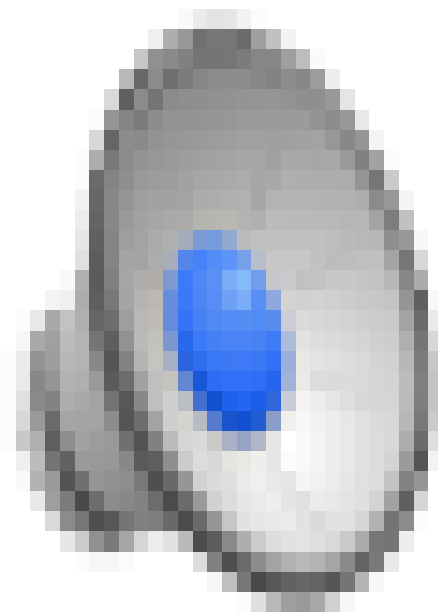
Part of slide credits: CMU AI & <http://ai.berkeley.edu>

# Game Type

# Behavior from Computation



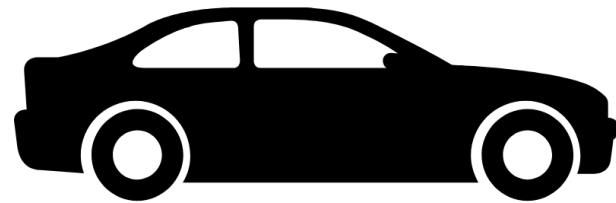
# Video of Demo Mystery Pacman



# Agents Getting Along with Other Agents

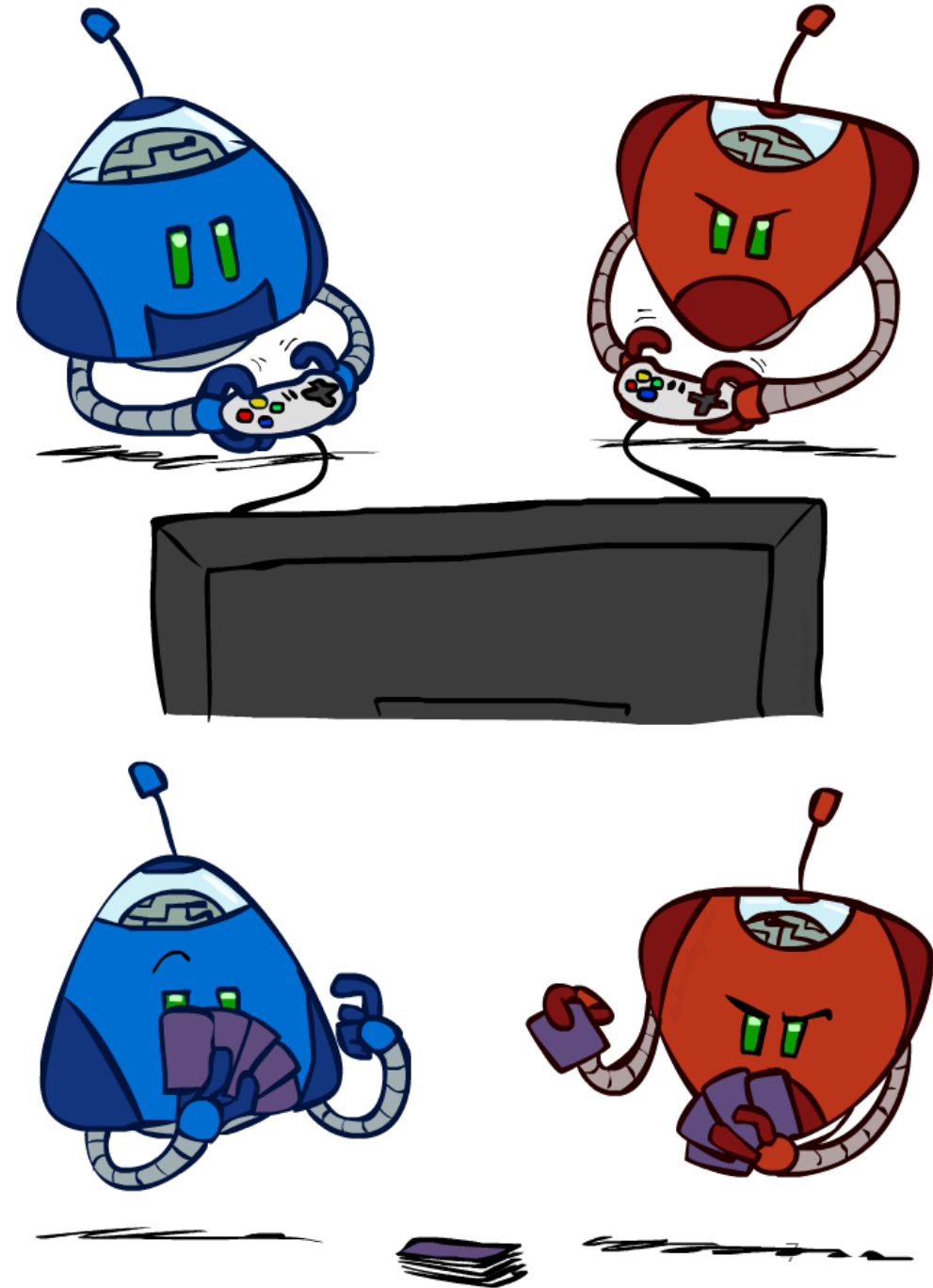


# Agents Getting Along with Humans



# Types of Games

- Many different kinds of games!
- Axes:
  - Deterministic or stochastic?
  - One, two, or more players?
  - Zero sum?
  - Perfect information (can you see the state)?
- Want algorithms for calculating a **contingent plan** (a.k.a. **strategy** or **policy**) which recommends a move for every possible eventuality



# History of Game AI

1956 checkers

1992 backgammon

1994 checkers

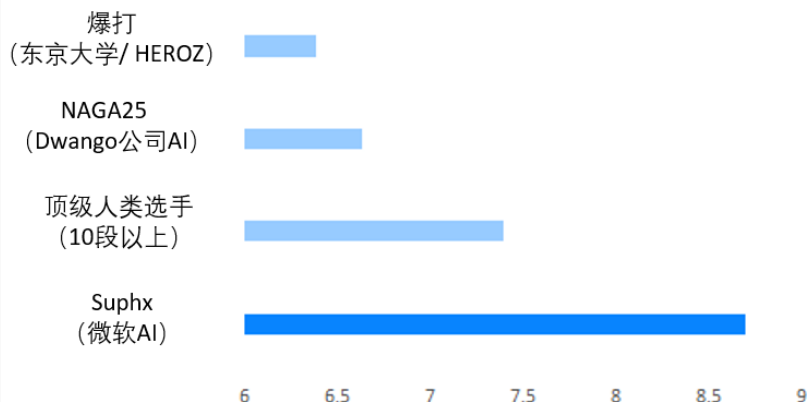
1997 chess

2016 Go

2017 Texas hold'em

2019 Majiang

天凤平台“特上房”稳定段位对比



信息集平均大小



1956年

Arthur Samuel开发的  
AI程序首先应用于国际跳棋

1992年

Hans Berliner开发的AI程序  
在双陆棋中取得突破

1994年

Jonathan Schaeffer  
团队开发的AI程序  
战胜国际跳棋  
世界冠军



许峰雄团队开发的AI程序  
击败国际象棋世界冠军

1997年

DeepMind团队  
开发的AI程序  
战胜围棋世界冠军  
2016年



Deep Blue

2017年

卡耐基梅隆大学、Facebook AI以及阿尔伯特大学开发的  
AI程序在德州扑克取得突破

下一个里程碑在哪里?



AlphaGo



Libratus

Pluribus

DeepStack

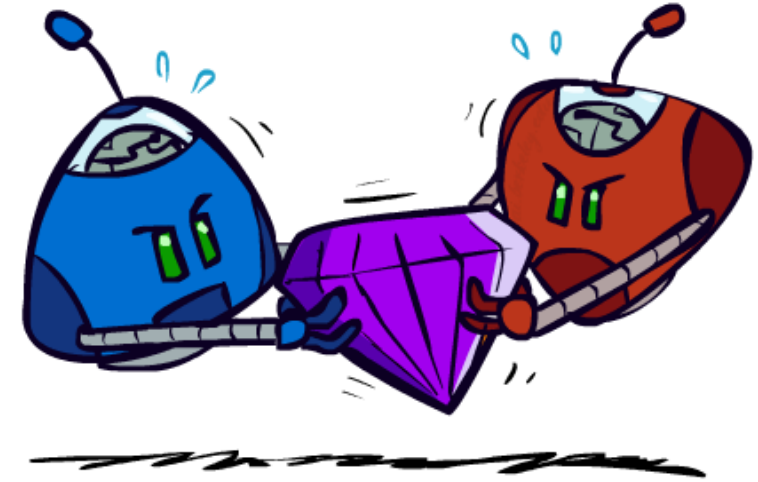


# Types of Games 2



- General Games

- Agents have **independent** utilities (values on outcomes)
- Cooperation, indifference, competition, shifting alliances, and more are all possible
  - We don't make AI to act in isolation, it should
    - a) work around people and b) help people
  - That means that every AI agent needs to solve a game



- Zero-Sum Games

- Agents have **opposite** utilities (values on outcomes)
- Lets us think of a single value that one **maximizes** and the other **minimizes**
- Adversarial, pure competition

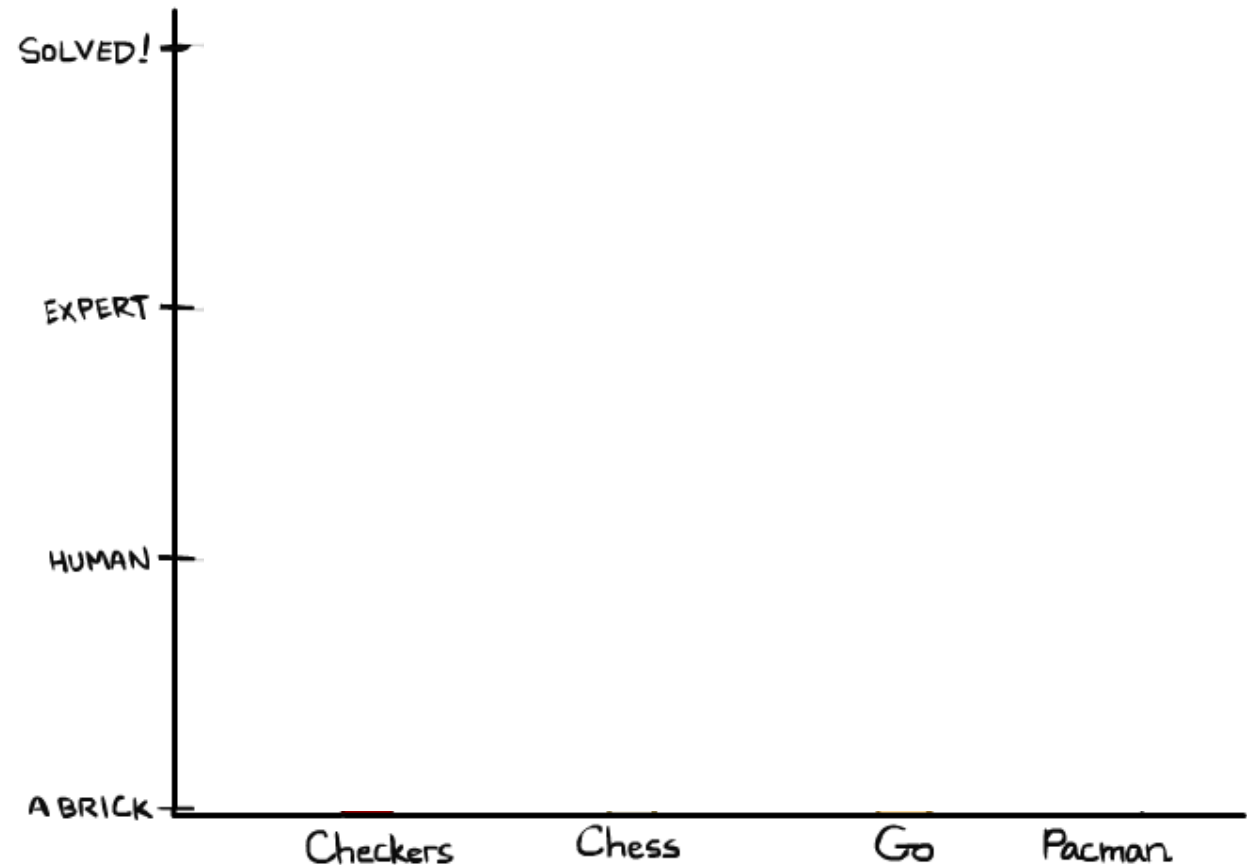
# Types of Games 3



- Common payoff games
  - Discussion: Use a technique you've learned so far to solve one!

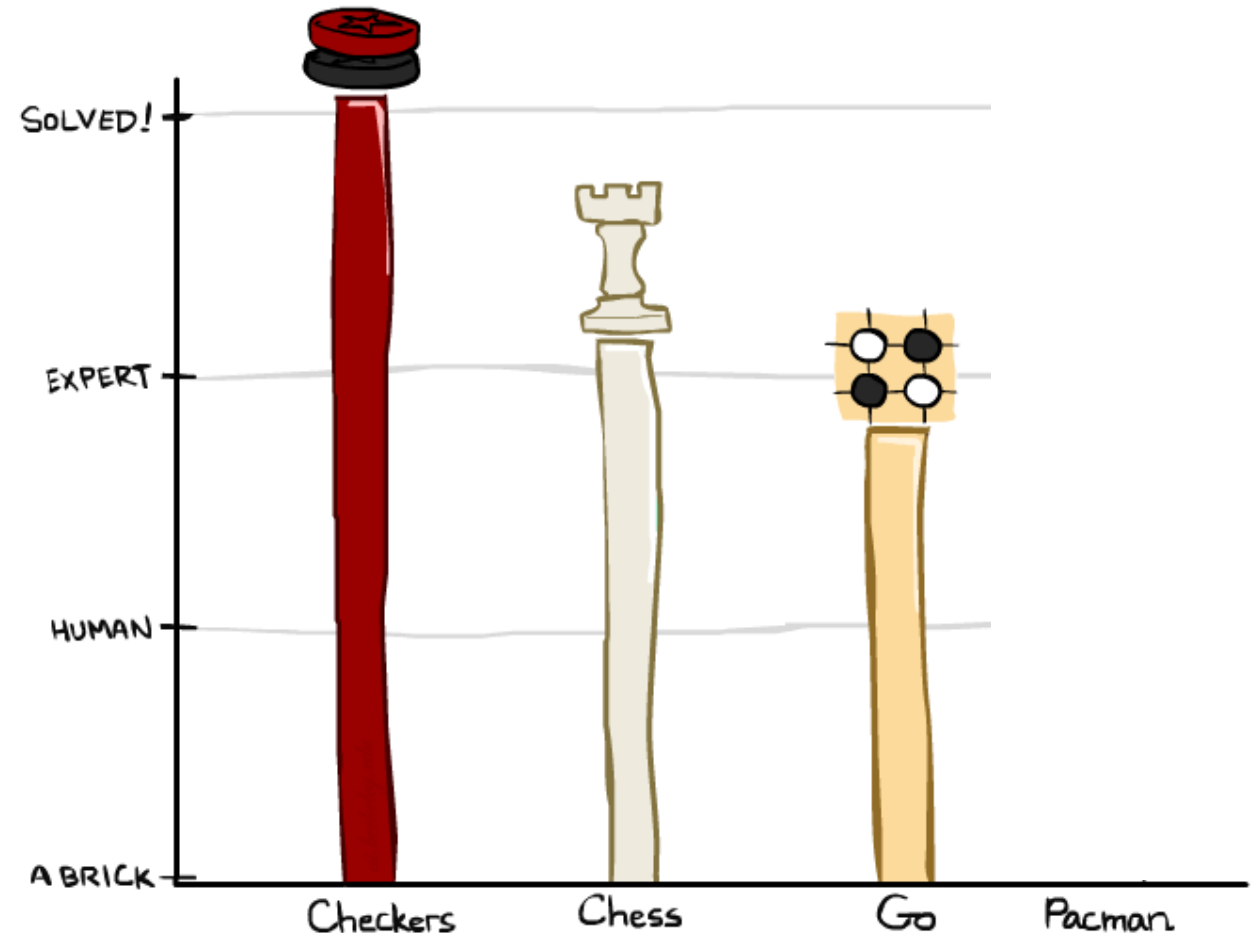
# Zero-Sum Games

- **Checkers:** 1950: First computer player. 1994: First computer champion: Chinook ended 40-year-reign of human champion Marion Tinsley using complete 8-piece endgame. 2007: Checkers solved!
- **Chess:** 1997: Deep Blue defeats human champion Gary Kasparov in a six-game match. Deep Blue examined 200M positions per second, used very sophisticated evaluation and undisclosed methods for extending some lines of search up to 40 ply. Current programs are even better, if less historic.
- **Go:** 2016: Alpha GO defeats human champion. Uses Monte Carlo Tree Search, learned evaluation function.

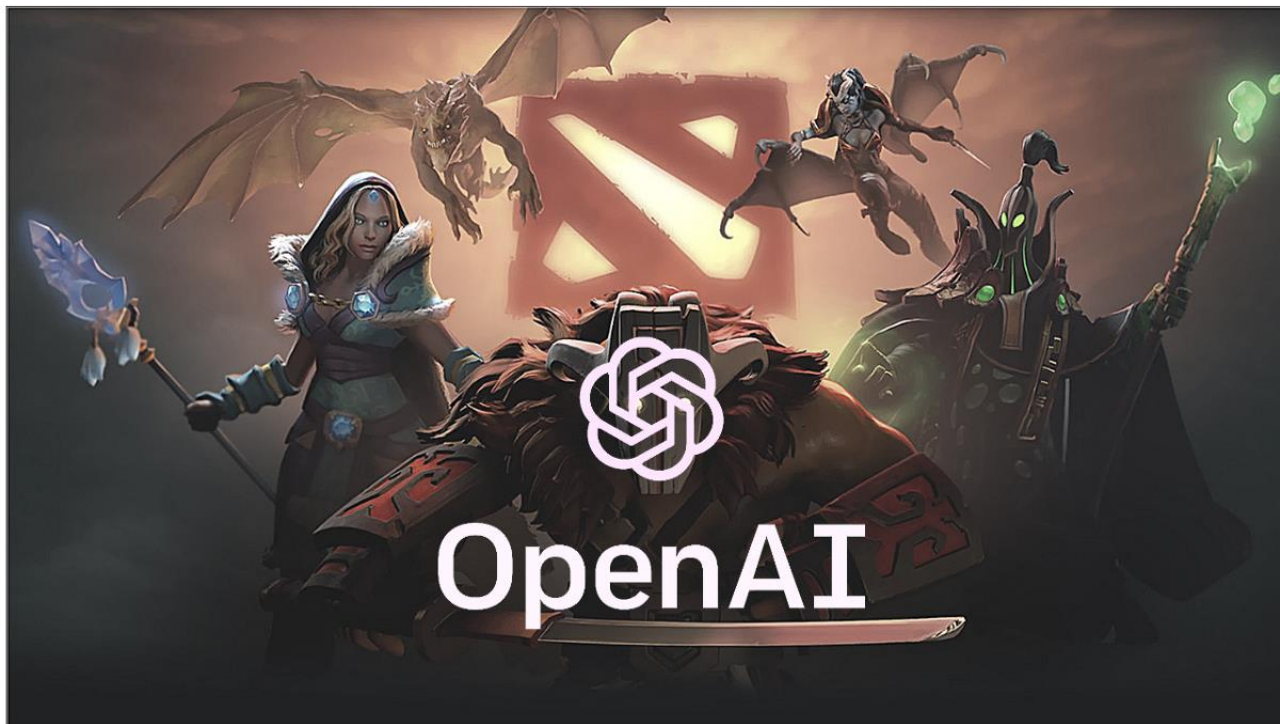


# Zero-Sum Games 2

- **Checkers:** 1950: First computer player. 1994: First computer champion: Chinook ended 40-year-reign of human champion Marion Tinsley using complete 8-piece endgame. 2007: Checkers solved!
- **Chess:** 1997: Deep Blue defeats human champion Gary Kasparov in a six-game match. Deep Blue examined 200M positions per second, used very sophisticated evaluation and undisclosed methods for extending some lines of search up to 40 ply. Current programs are even better, if less historic.
- **Go :**2016: Alpha GO defeats human champion. Uses Monte Carlo Tree Search, learned evaluation function.
- **Pacman**

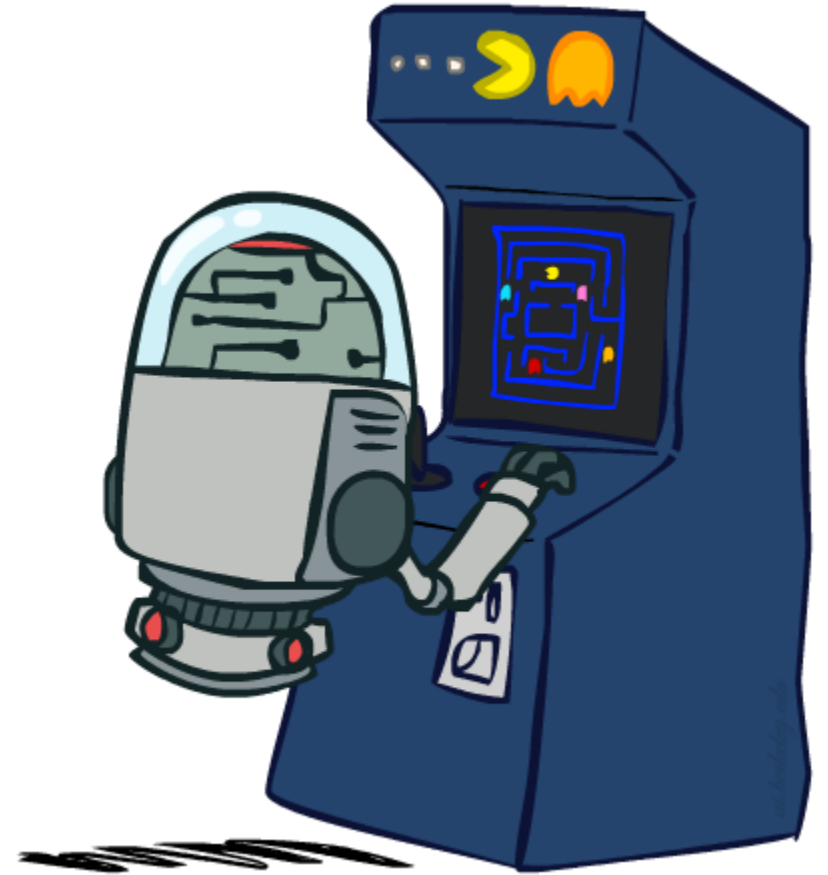


# Game playing – state of the art

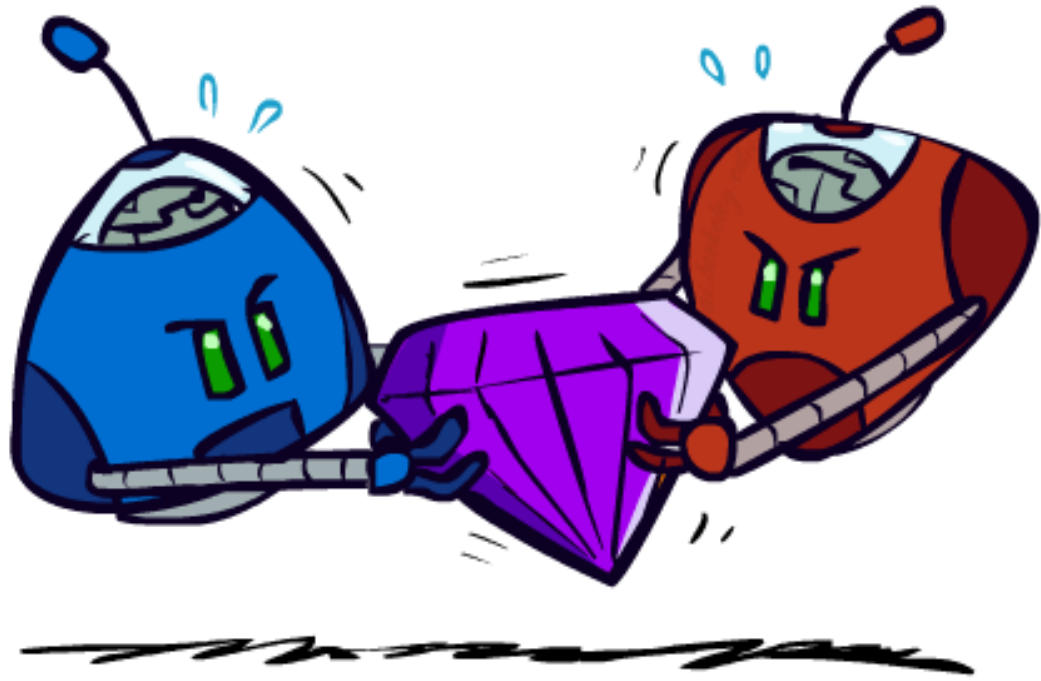


# “Standard” Games

- Standard games are **deterministic**, observable, two-player, turn-taking, zero-sum
- Game formulation:
  - States:  $S$  (start at  $s_0$ )
  - Players:  $P=\{1\dots N\}$  (usually take turns)
  - Actions:  $A$  (may depend on player / state)
  - Transition Function:  $S \times A \rightarrow S$
  - Terminal Test:  $S \rightarrow \{t, f\}$
  - Terminal Utilities:  $S \times P \rightarrow R$
- Solution for a player is a **policy**:  $S \rightarrow A$

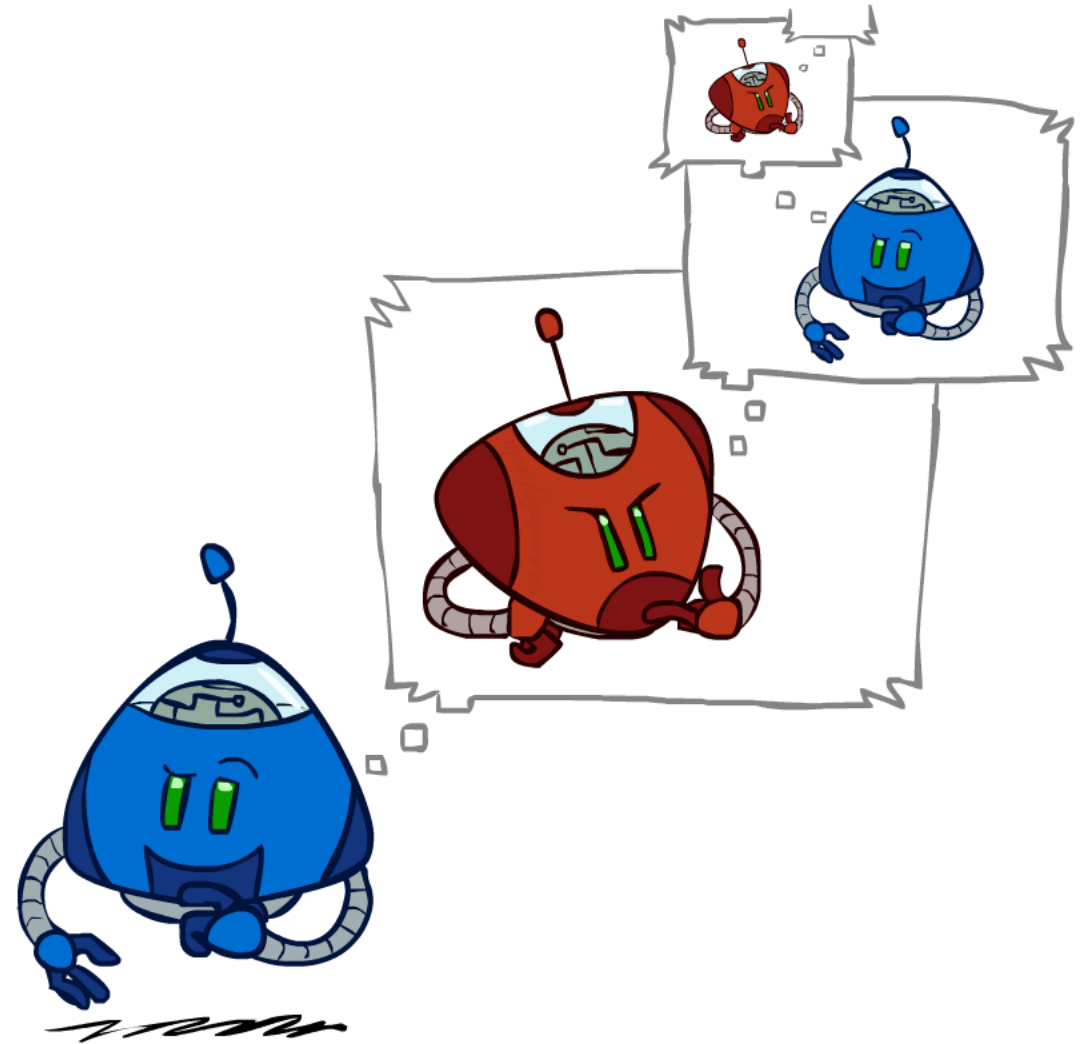




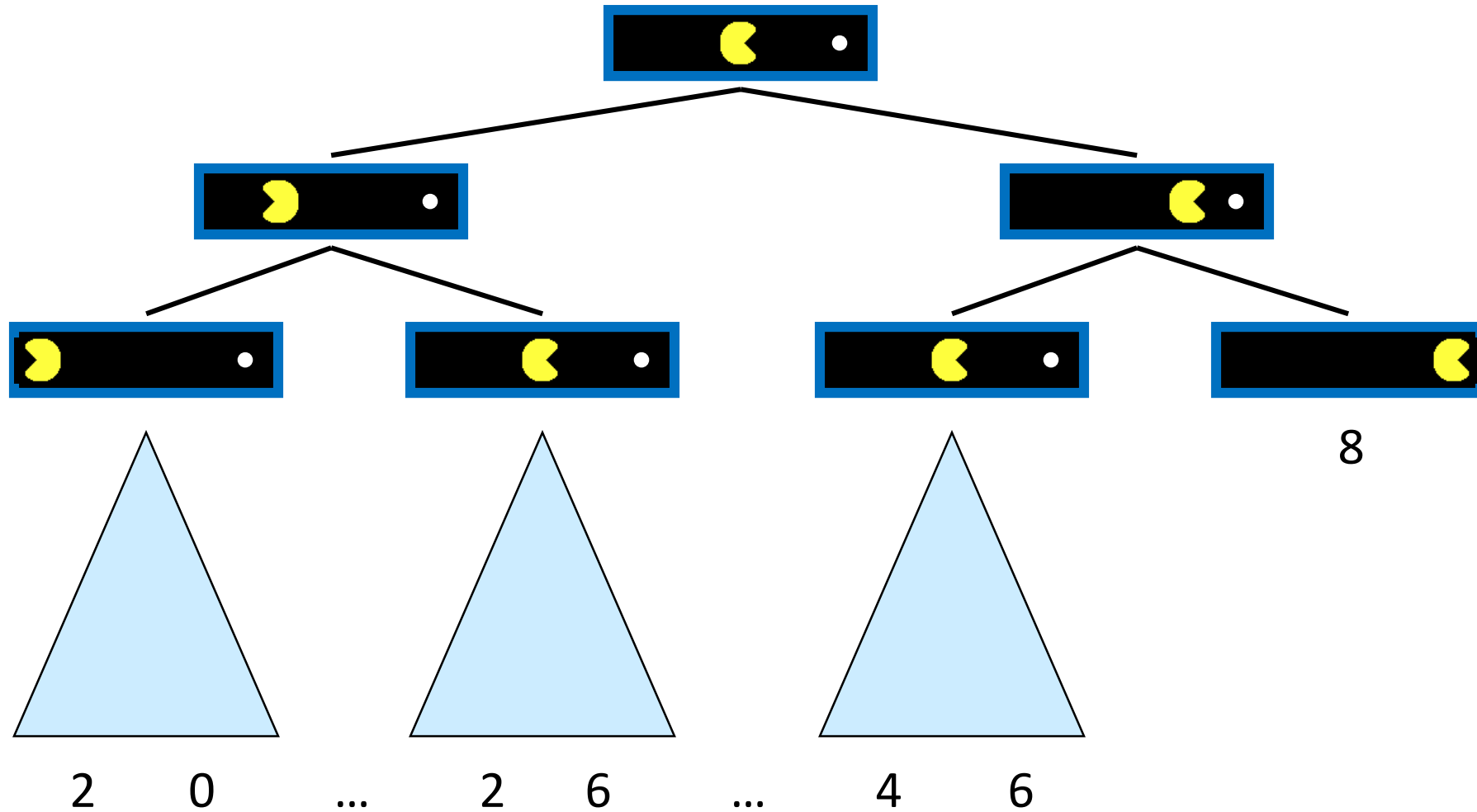


# Adversarial Search

Cost  $\rightarrow$  Utility!



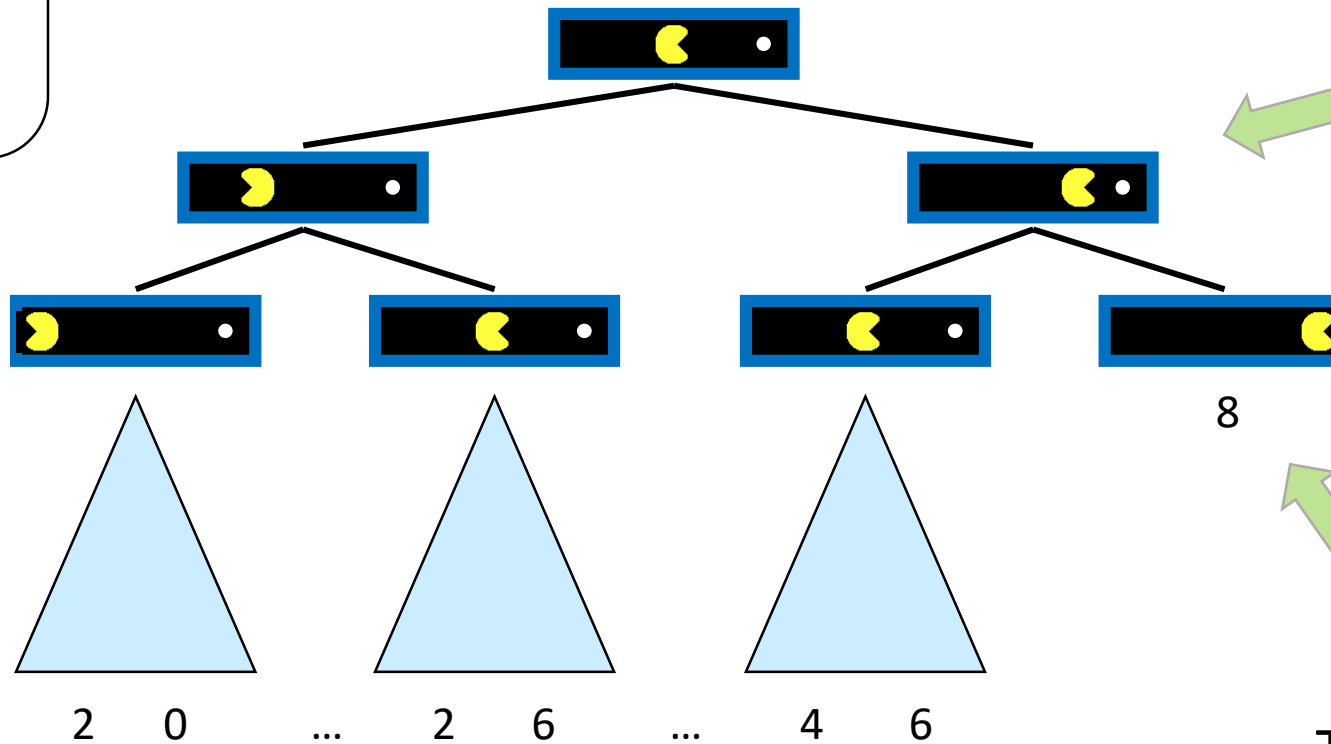
# Single-Agent Trees





# Single-Agent Trees: Value of a State

Value of a state:  
The best achievable  
outcome (utility)  
from that state



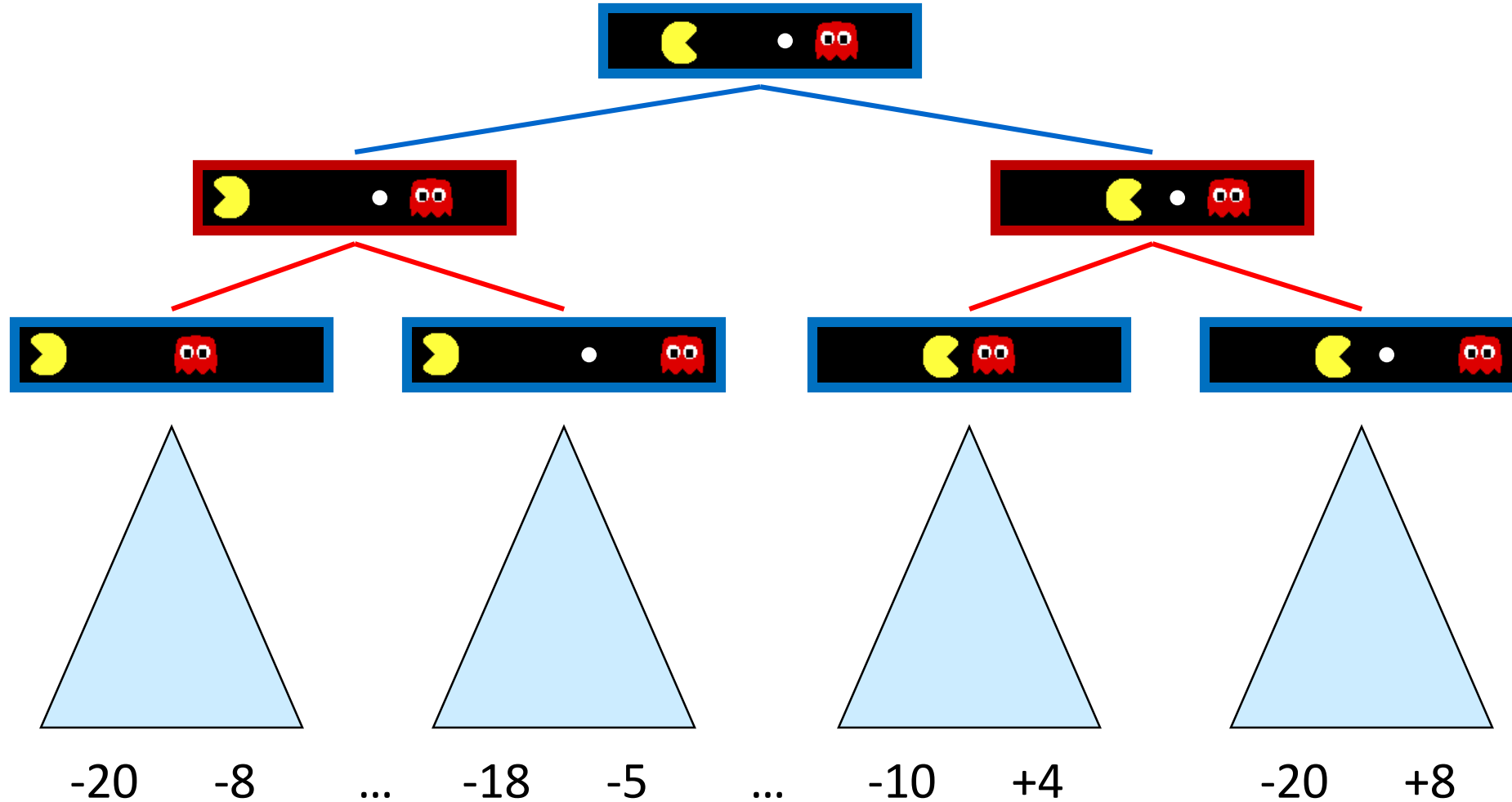
Non-Terminal States:

$$V(s) = \max_{s' \in \text{children}(s)} V(s')$$

Terminal States:

$$V(s) = \text{known } 17$$

# Adversarial Game Trees



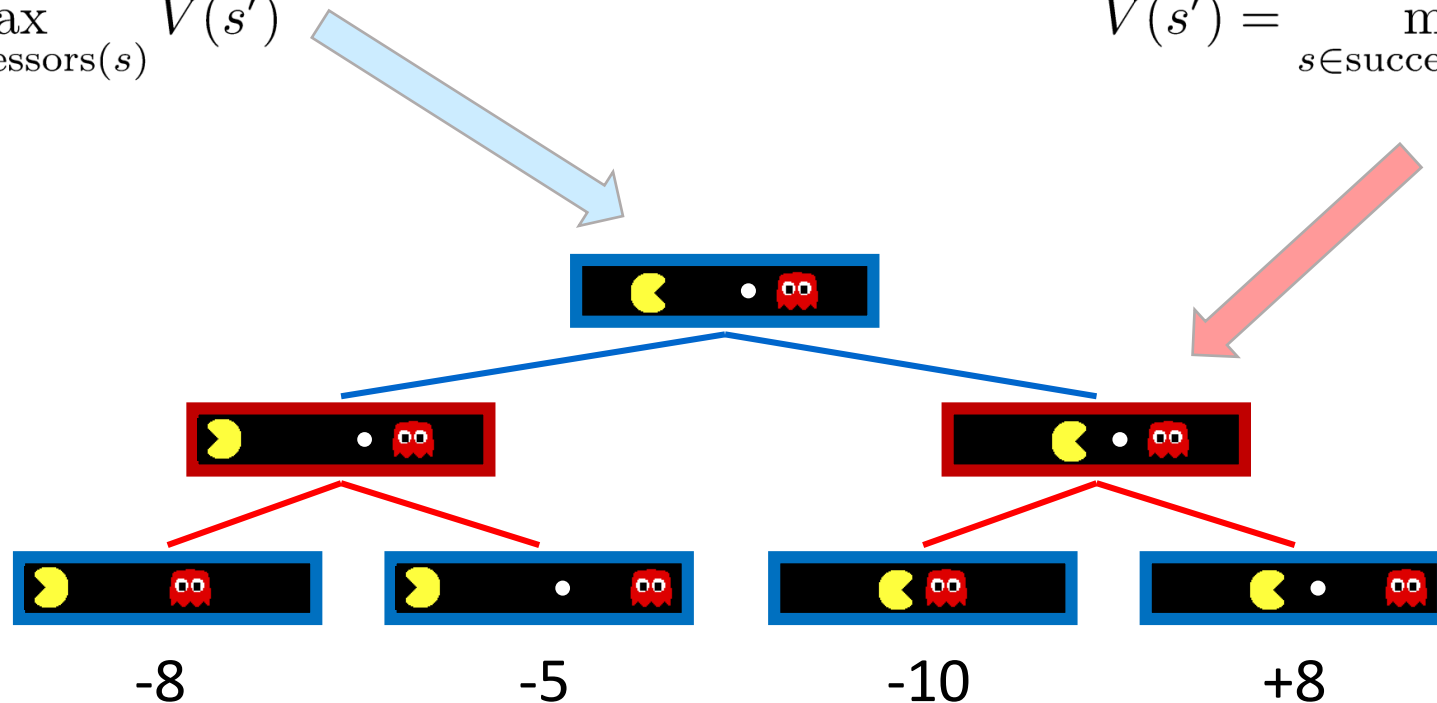
# Adversarial Game Trees: Minimax Values

States Under Agent's Control:

$$V(s) = \max_{s' \in \text{successors}(s)} V(s')$$

States Under Opponent's Control:

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$



Terminal States:

$$V(s) = \text{known}$$

# Example: Tic-Tac-Toe Game Tree

- States
- Actions
- Values



MAX (X)



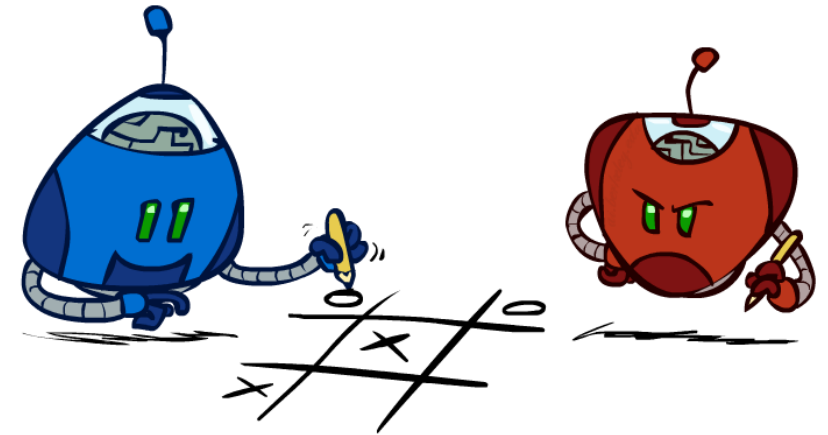
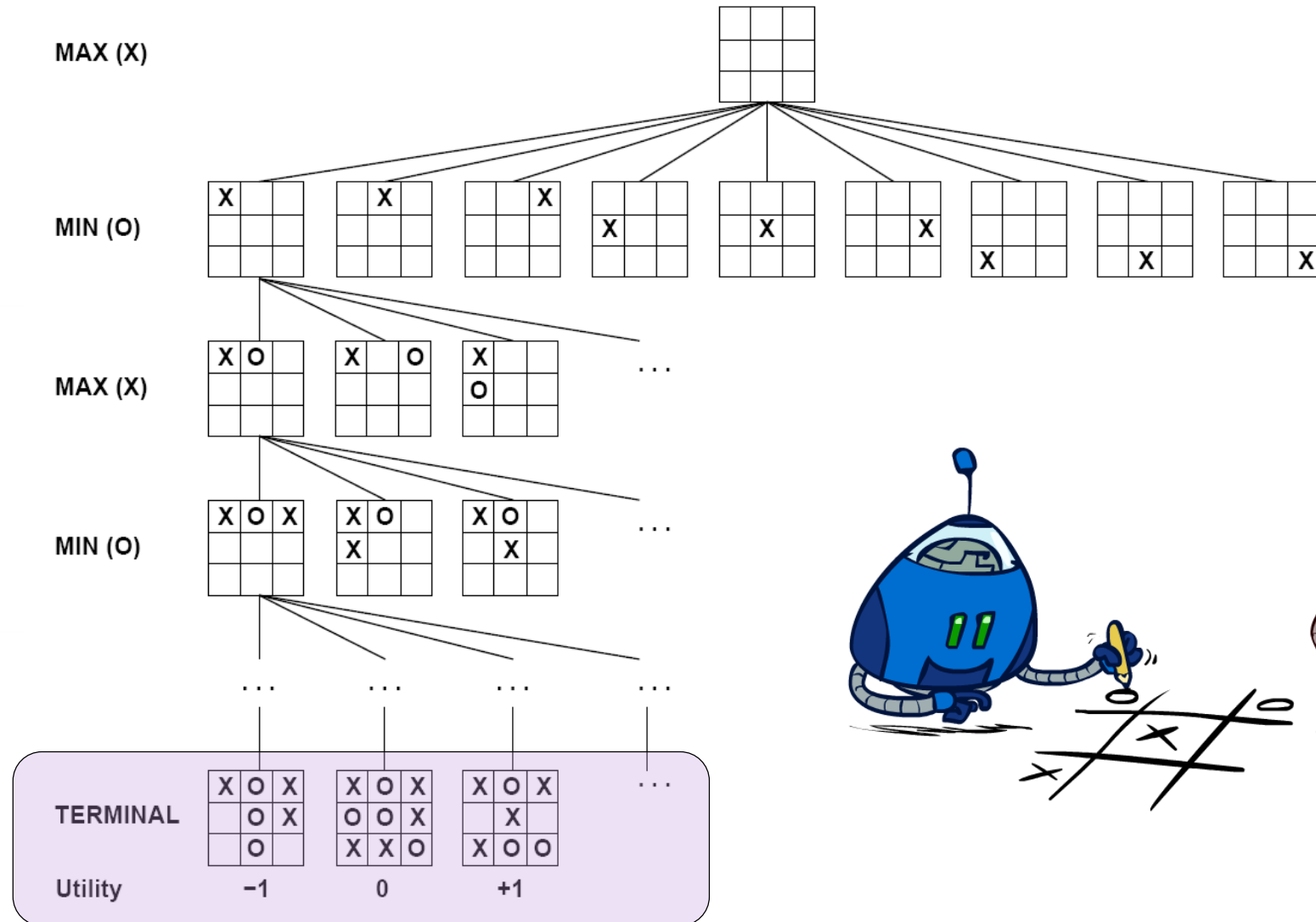
MIN (O)



MAX (X)

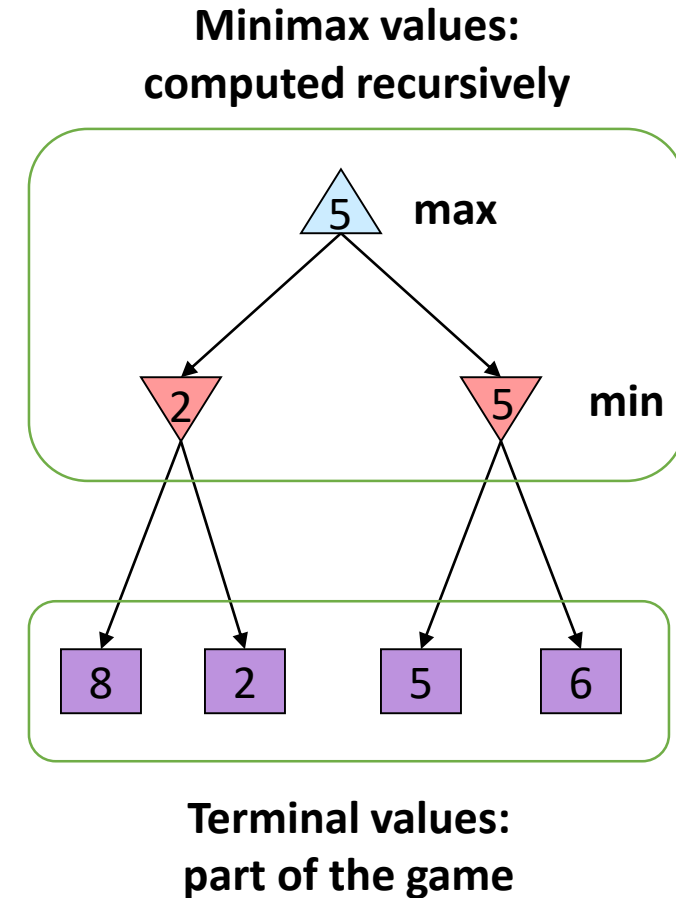


MIN (O)



# Minimax Search

- Deterministic, zero-sum games:
  - Tic-tac-toe, chess, checkers
  - One player maximizes result
  - The other minimizes result
- **Minimax** search:
  - A state-space search tree
  - Players alternate turns
  - Compute each node's **minimax value**: the best achievable utility against a rational (optimal) adversary



# Minimax Implementation

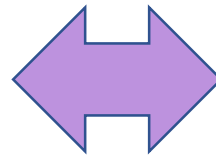
def max-value(state):

    initialize  $v = -\infty$

    for each successor of state:

$v = \max(v, \text{min-value}(\text{successor}))$

    return  $v$



def min-value(state):

    initialize  $v = +\infty$

    for each successor of state:

$v = \min(v, \text{max-value}(\text{successor}))$

    return  $v$

$$V(s) = \max_{s' \in \text{successors}(s)} V(s')$$

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$

# Minimax Implementation (Dispatch)

```
def value(state):
```

if the state is a terminal state: return the state's utility

if the next agent is MAX: return max-value(state)

if the next agent is MIN: return min-value(state)

```
def max-value(state):
```

initialize  $v = -\infty$

for each successor of state:

$v = \max(v, \text{value}(\text{successor}))$

return  $v$

```
def min-value(state):
```

initialize  $v = +\infty$

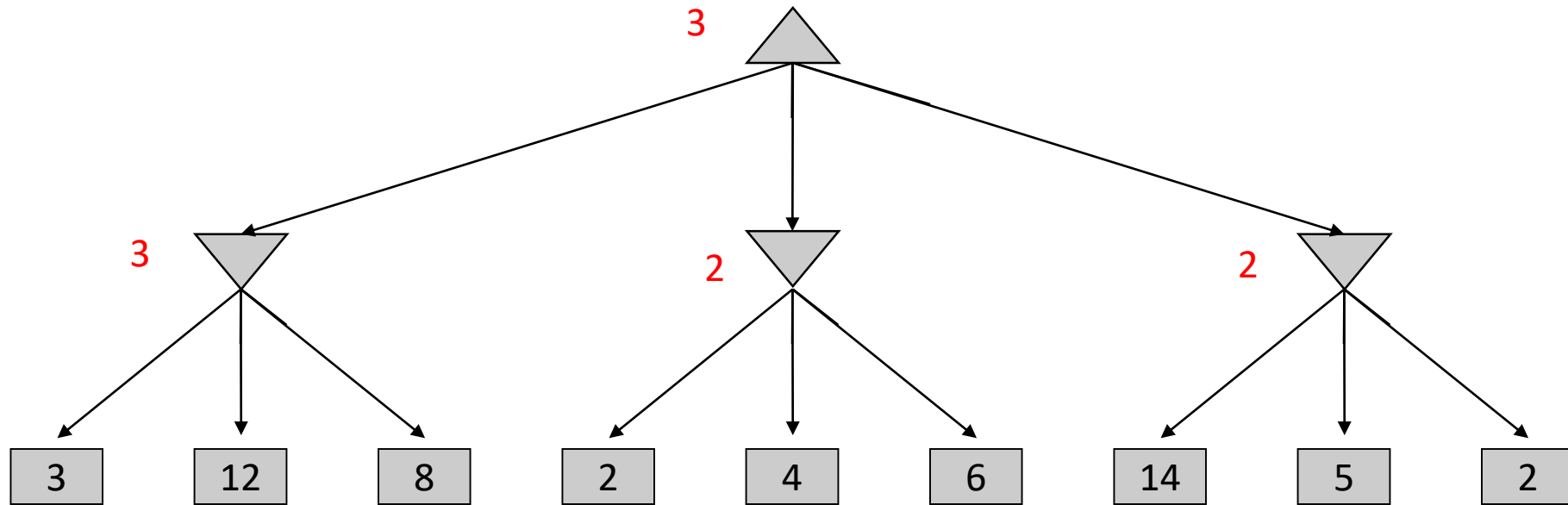
for each successor of state:

$v = \min(v, \text{value}(\text{successor}))$

return  $v$

# Example

- Actions?



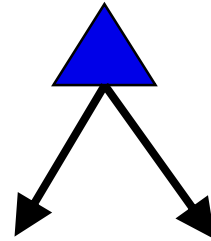


# Pseudocode for Single Agent

```
def max_value(state):  
    if state.is_leaf:  
        return state.value  
    # TODO Also handle depth limit  
  
    best_value = -10000000  
  
    for action in state.actions:  
        next_state = state.result(action)  
  
        next_value = max_value(next_state)  
  
        if next_value > best_value:  
            best_value = next_value  
  
    return best_value
```

# Pseudocode for Minimax Search

```
def max_value(state):  
    if state.is_leaf:  
        return state.value  
    # TODO Also handle depth limit  
  
    best_value = -10000000  
  
    for action in state.actions:  
        next_state = state.result(action)  
        next_value = min_value(next_state)  
  
        if next_value > best_value:  
            best_value = next_value  
  
    return best_value  
  
def min_value(state):
```

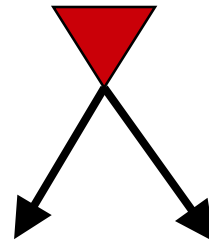


$$V(s) = \max_a V(s'),$$

where  $s' = \text{result}(s, a)$

$$\hat{a} = \operatorname{argmax}_a V(s'),$$

where  $s' = \text{result}(s, a)$



# Pseudocode for Generic Game Tree

```
function minimax_decision( state )  
    return argmaxa in state.actions value( state.result(a) )
```

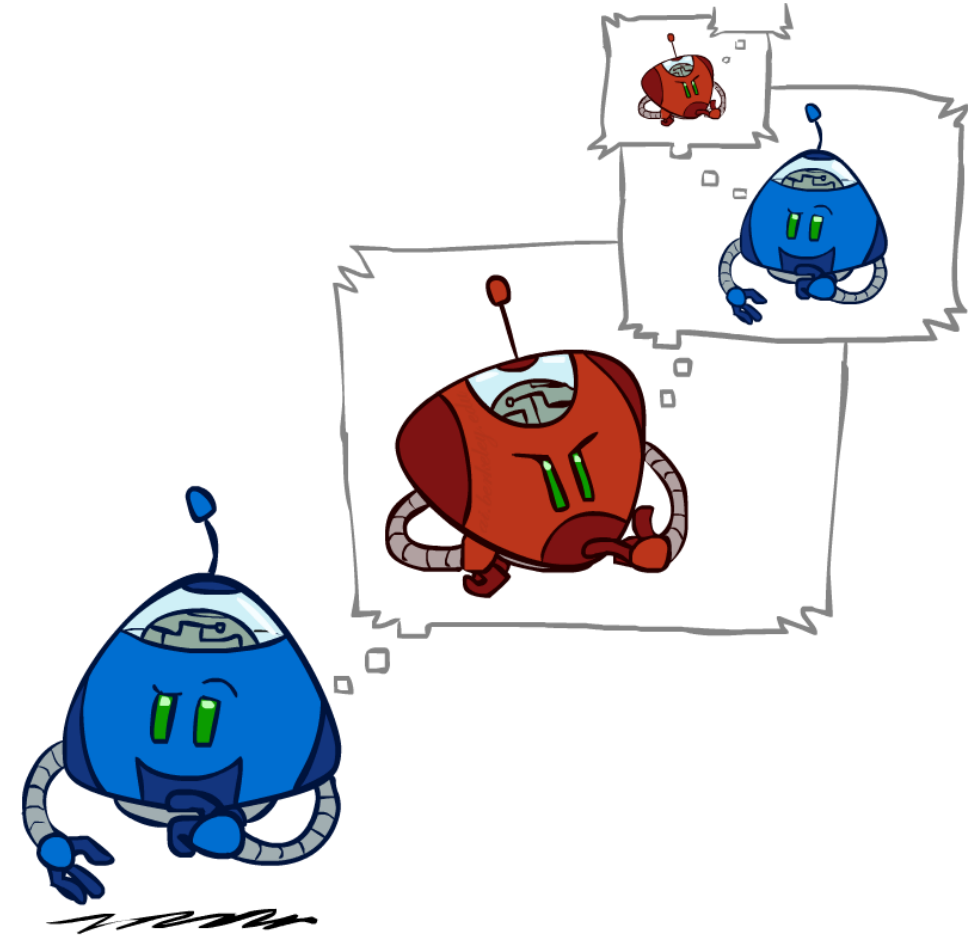
```
function value( state )  
    if state.is_leaf  
        return state.value
```

```
    if state.player is MAX  
        return maxa in state.actions value( state.result(a) )
```

```
    if state.player is MIN  
        return mina in state.actions value( state.result(a) )
```

# Minimax Efficiency

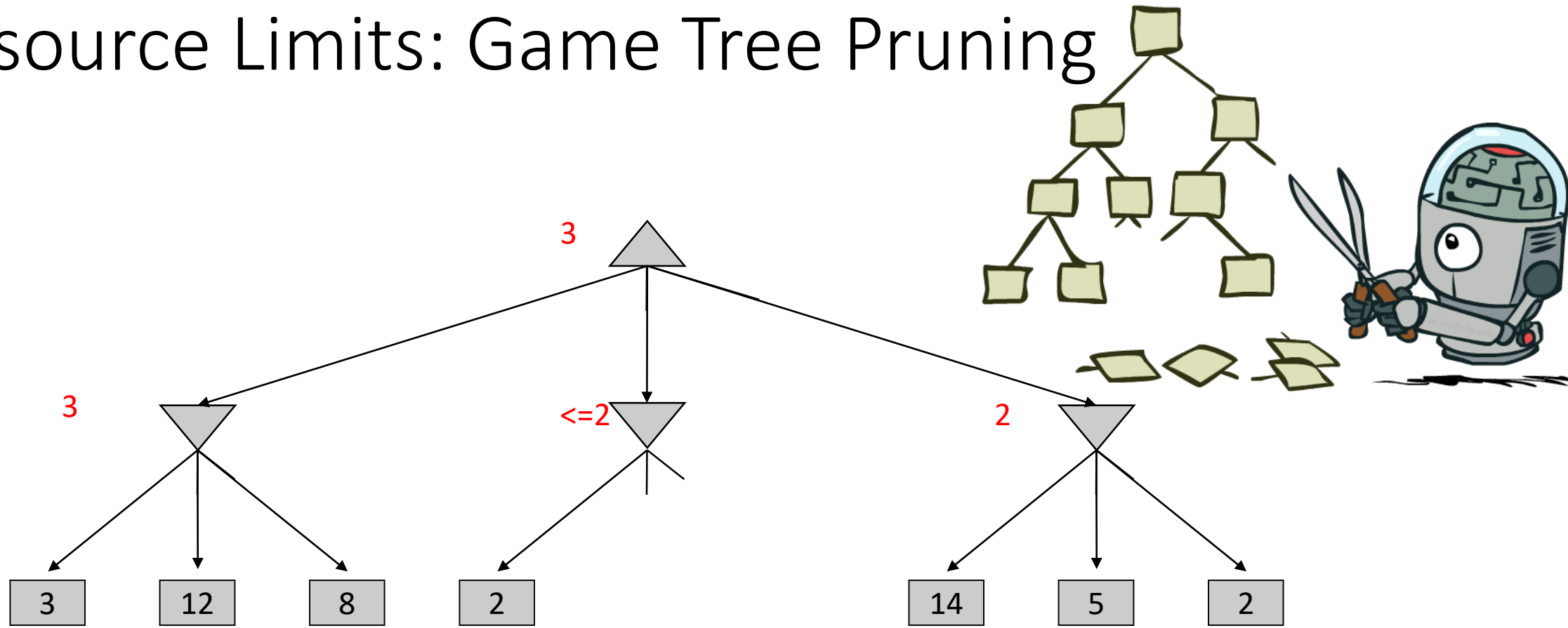
- How efficient is minimax?
  - Just like (exhaustive) DFS
  - Time:  $O(b^m)$
  - Space:  $O(bm)$
- Example: For chess,  $b \approx 35$ ,  $m \approx 100$ 
  - Exact solution is completely infeasible
  - But, do we need to explore the whole tree?
  - Humans can't do this either, so how do we play chess?
  - **Bounded rationality** – Herbert Simon



# Resource Limits



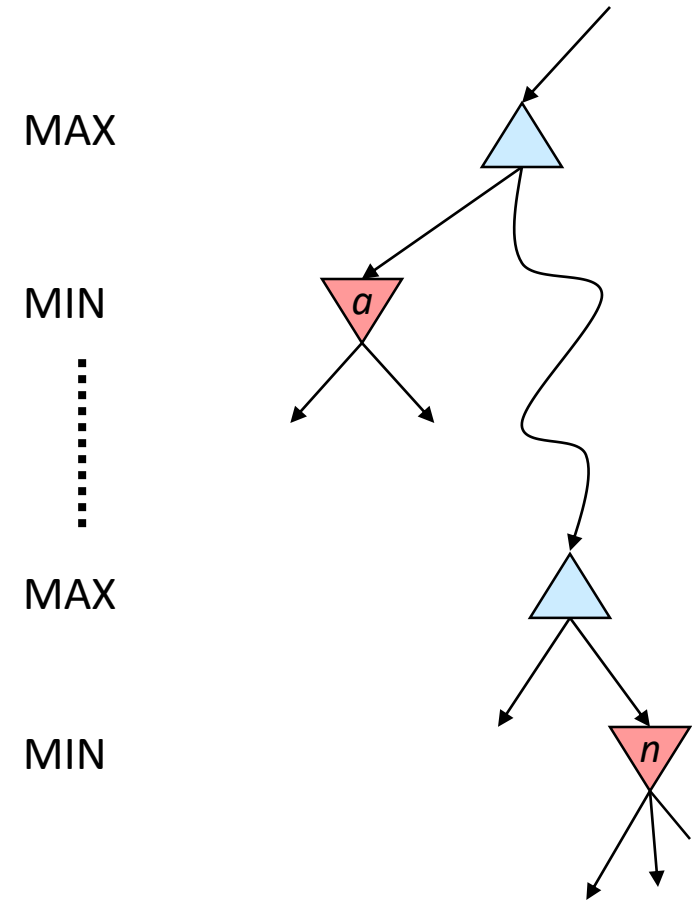
# Resource Limits: Game Tree Pruning



*The order of generation matters:* more pruning is possible if good moves come first

# Game Tree Pruning: Alpha-Beta Pruning

- General configuration (MIN version)
  - We're computing the MIN-VALUE at some node  $n$
  - We're looping over  $n$ 's children
  - $n$ 's estimate of the childrens' min is dropping
  - Who cares about  $n$ 's value? MAX
  - Let  $a$  be the best value that MAX can get at any choice point along the current path from the root
  - If  $n$  becomes worse than  $a$ , MAX will avoid it, so we can stop considering  $n$ 's other children (it's already bad enough that it won't be played)
- MAX version is symmetric



# Alpha-Beta Implementation

$\alpha$ : MAX's best option on path to root  
 $\beta$ : MIN's best option on path to root

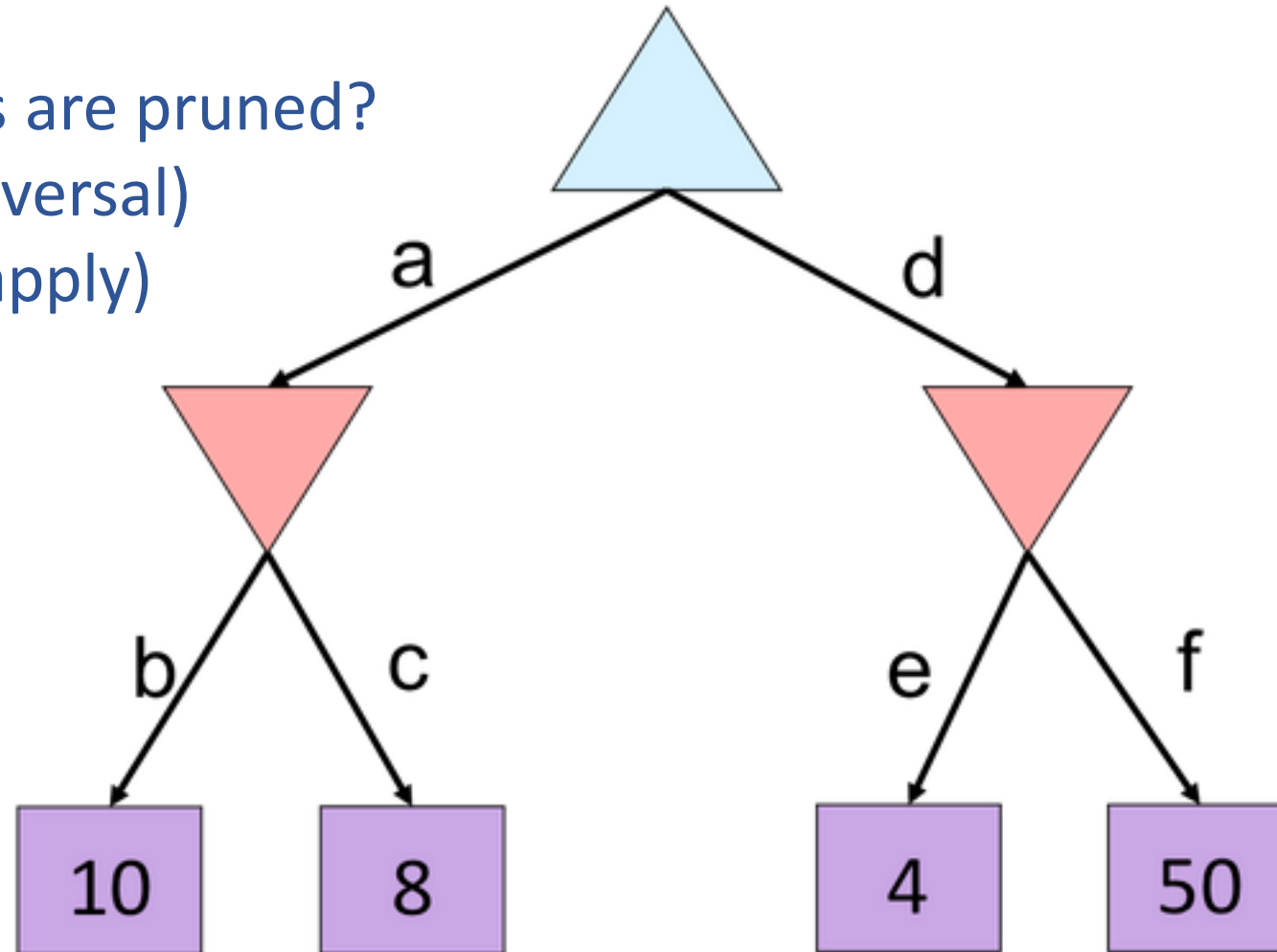
```
def max-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = -\infty$   
    for each successor of state:  
         $v = \max(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \geq \beta$  return  $v$   
         $\alpha = \max(\alpha, v)$   
    return  $v$ 
```

```
def min-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = +\infty$   
    for each successor of state:  
         $v = \min(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \leq \alpha$  return  $v$   
         $\beta = \min(\beta, v)$   
    return  $v$ 
```



# Quiz

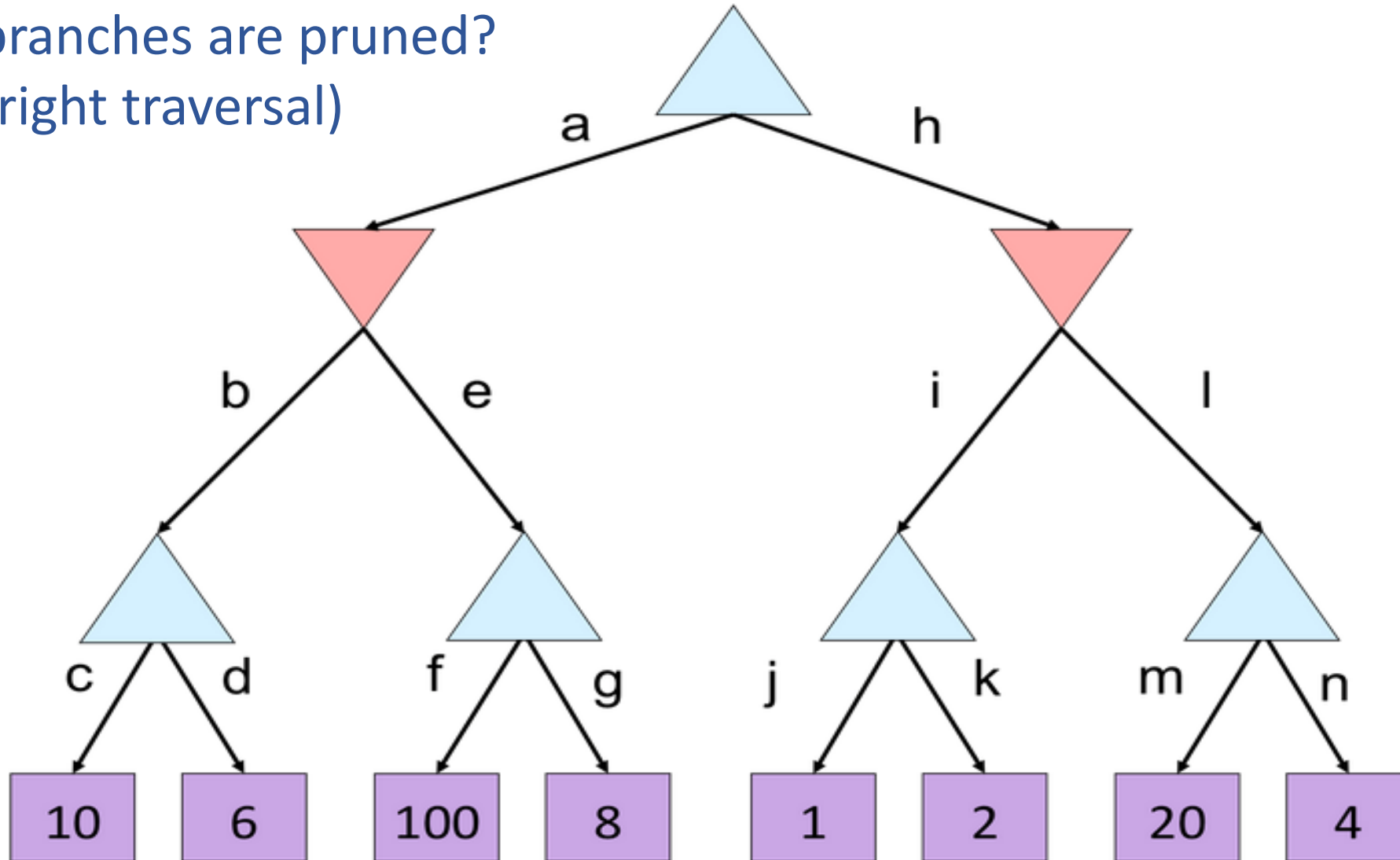
Which branches are pruned?  
(Left to right traversal)  
(Select all that apply)



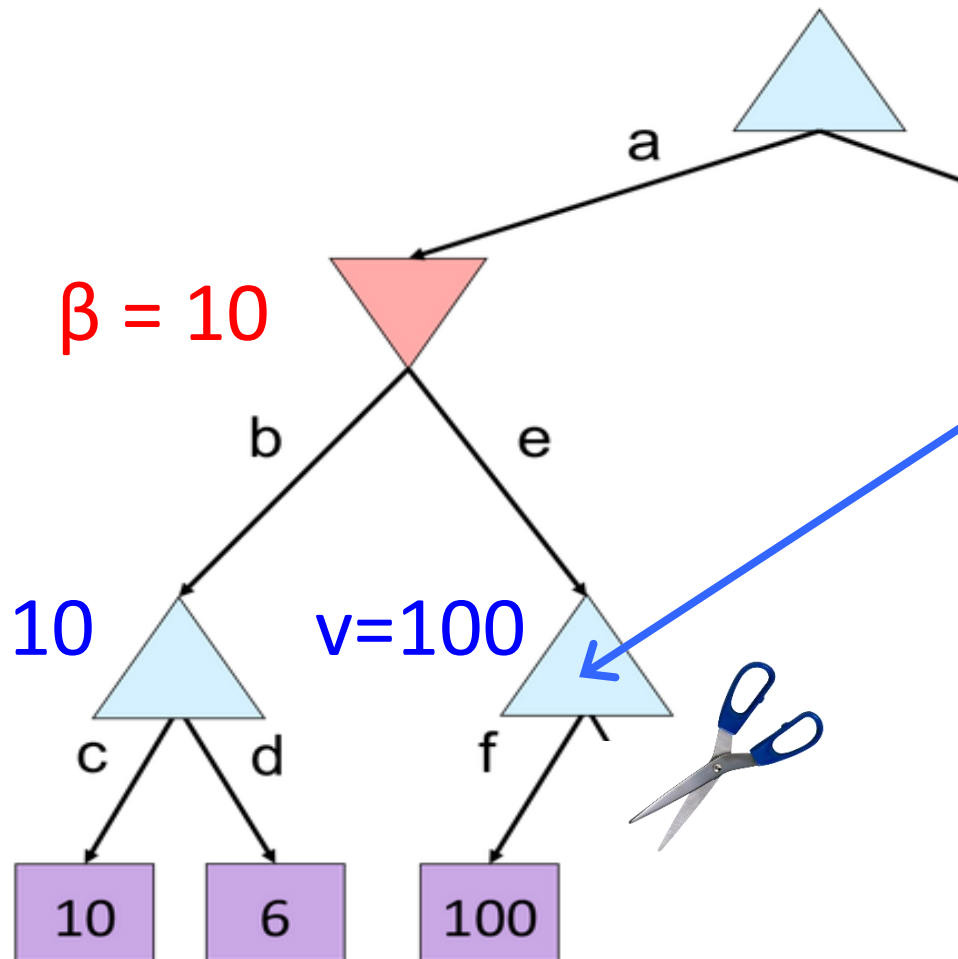
# Quiz 2

Which branches are pruned?  
(Left to right traversal)

- A) e, l
- B) g, l
- C) g, k, l
- D) g, n



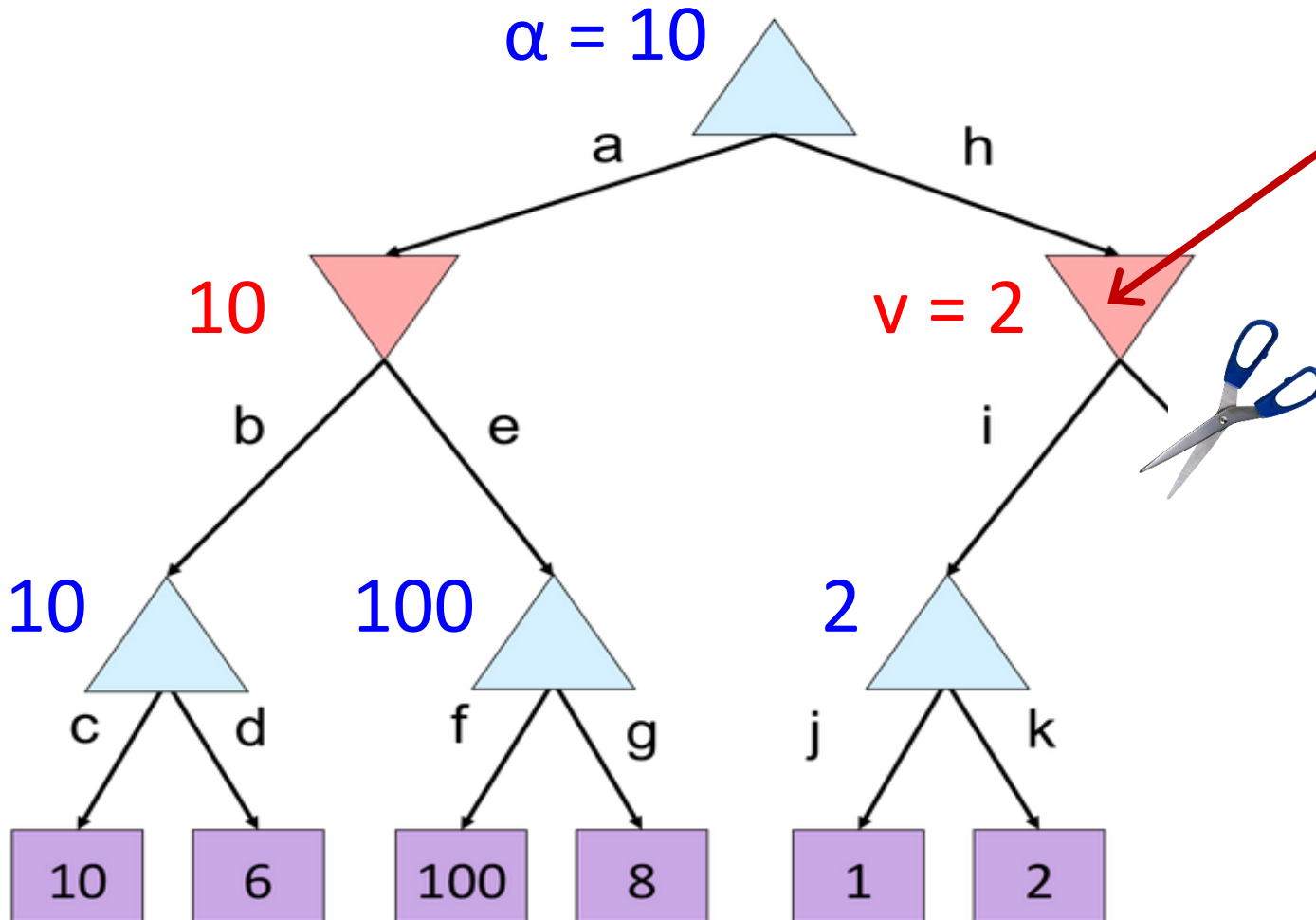
# Quiz 2 - 2



$\alpha$ : MAX's best option on path to root  
 $\beta$ : MIN's best option on path to root

```
def max-value(state,  $\alpha$ ,  $\beta$ ):  
    initialize  $v = -\infty$   
    for each successor of state:  
         $v = \max(v, \text{value}(\text{successor}, \alpha, \beta))$   
        if  $v \geq \beta$   
            return  $v$   
         $\alpha = \max(\alpha, v)$   
    return  $v$ 
```

# Quiz 2 - 3



$\alpha$ : MAX's best option on path to root

$\beta$ : MIN's best option on path to root

def min-value(state,  $\alpha$ ,  $\beta$ ):

    initialize  $v = +\infty$

    for each successor of state:

$v = \min(v, \text{value}(\text{successor}, \alpha, \beta))$

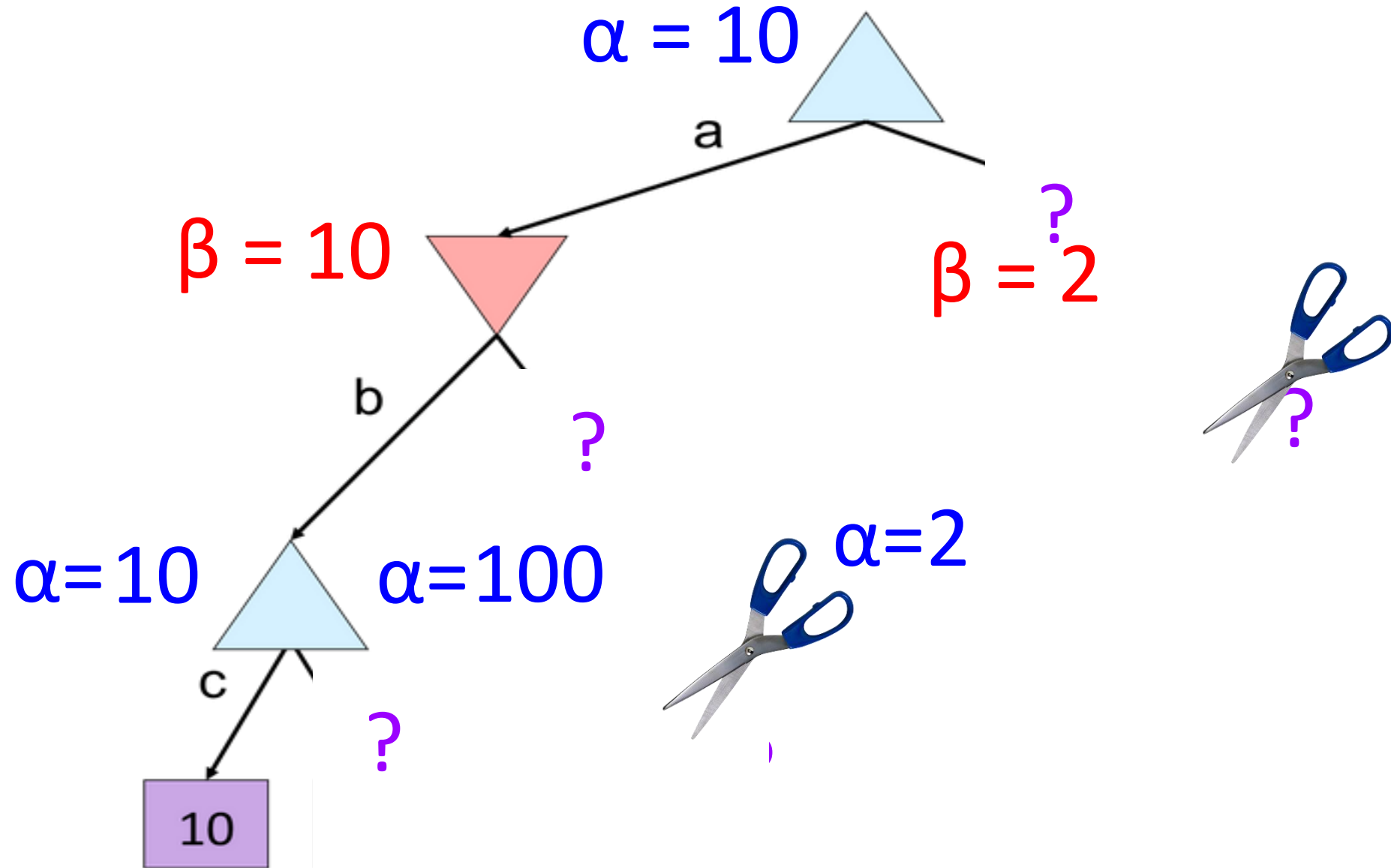
        if  $v \leq \alpha$

            return  $v$

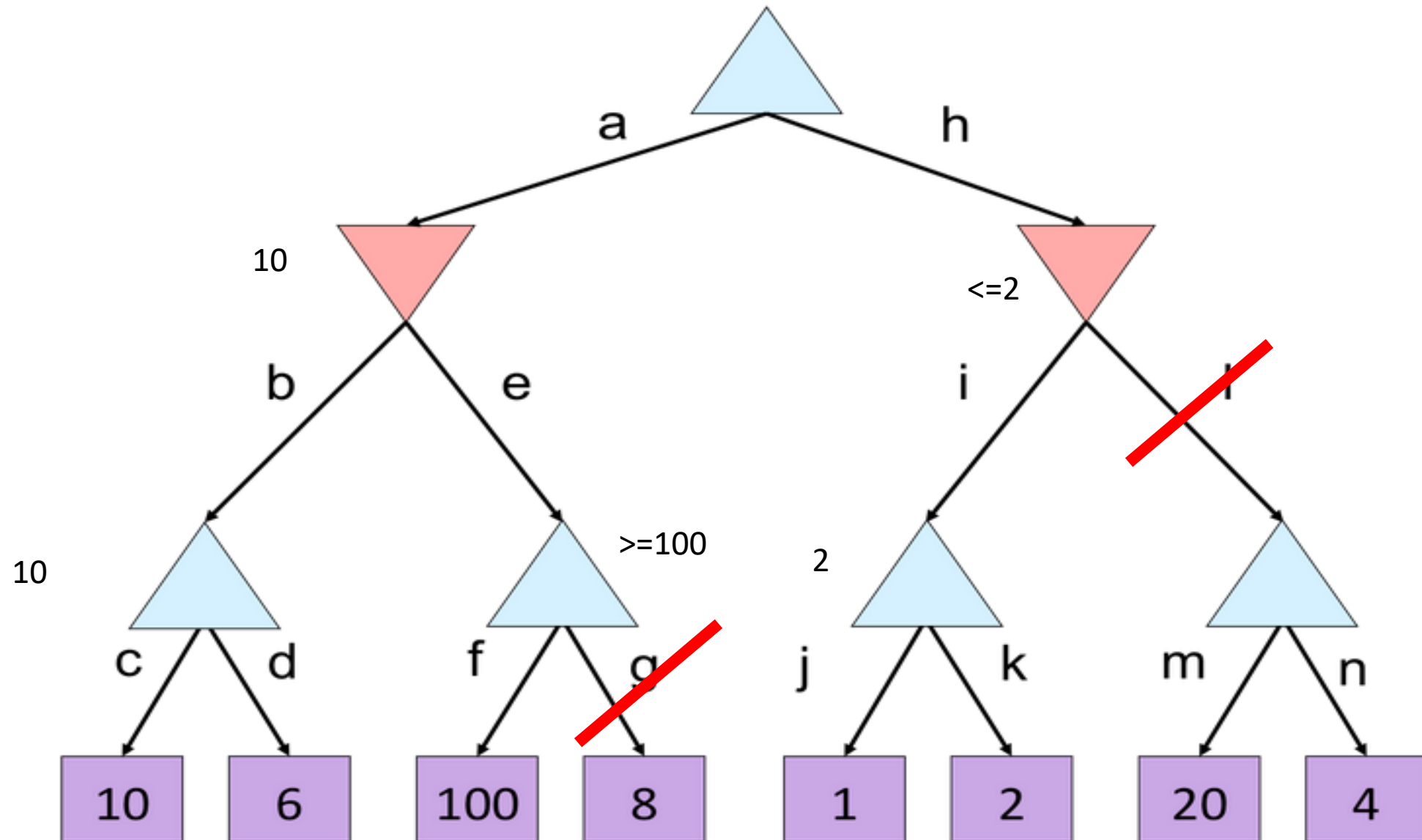
$\beta = \min(\beta, v)$

    return  $v$

# Quiz 2 - 4

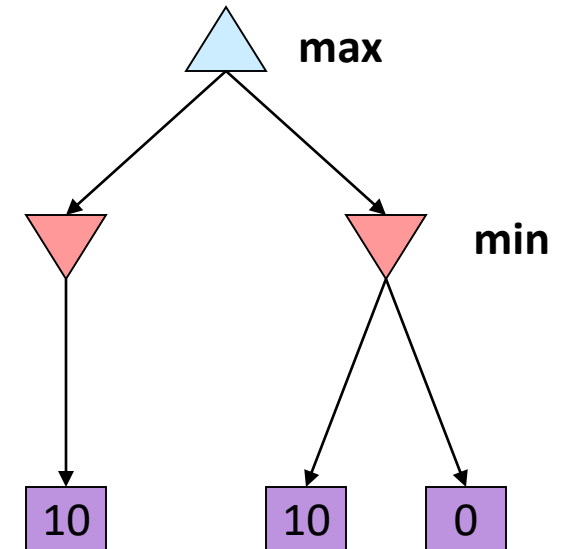


# Quiz 2 - 5



# Alpha-Beta Pruning Properties

- This pruning has **no effect** on minimax value computed for the root!
- Values of intermediate nodes might be wrong
  - Important: children of the root may have the wrong value
  - So the most naïve version won't let you do action selection
- Good child ordering improves effectiveness of pruning
- With “perfect ordering”:
  - Time complexity drops to  $O(b^{m/2})$
  - Doubles solvable depth!
  - Chess: 1M nodes/move => depth=8, respectable
  - Full search of complicated games, is still hopeless...
- This is a simple example of **metareasoning** (computing about what to compute)



# Resource Limits II

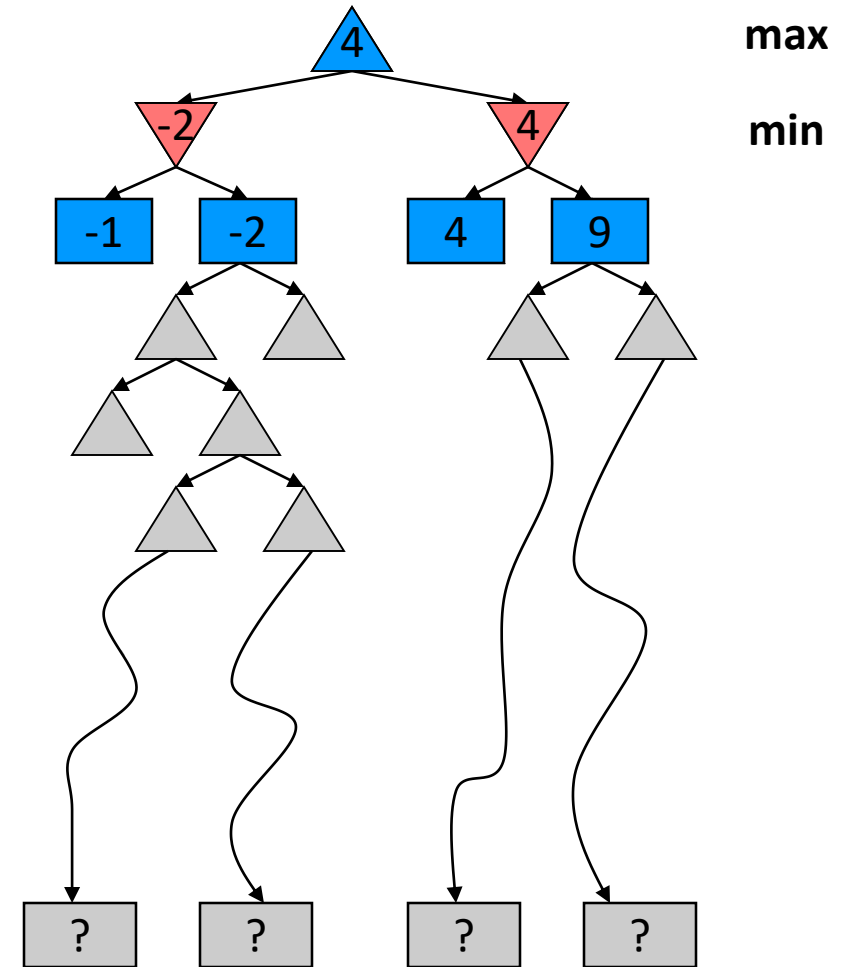
## Bounded lookahead





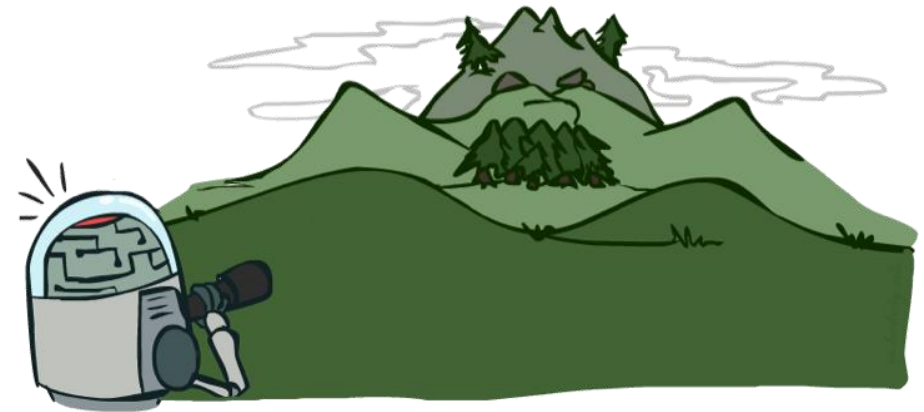
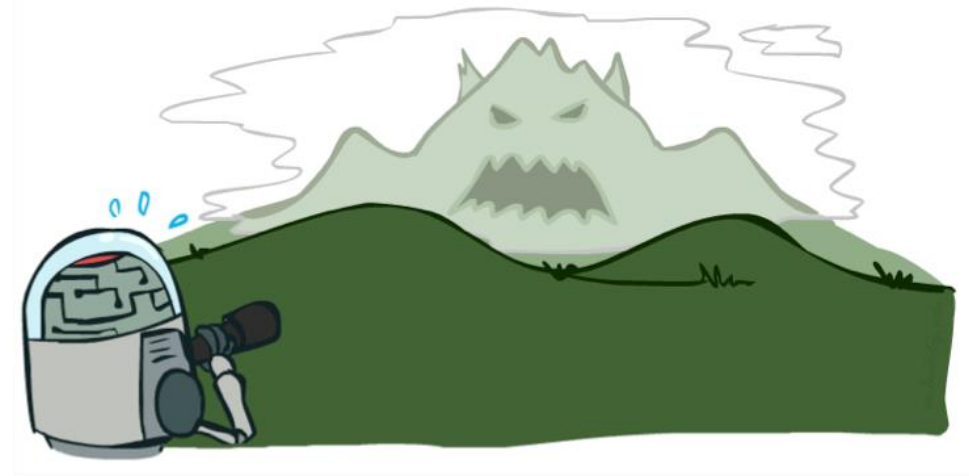
# Depth-limited search

- Problem: In realistic games, cannot search to leaves!
- Solution: Depth-limited search
  - Instead, search only to a **limited depth** in the tree
  - Replace terminal utilities with **an evaluation function** for non-terminal positions
- Example:
  - Suppose we have 100 seconds, can explore 10K nodes / sec
  - So can check 1M nodes per move
  - For chess,  $b \approx 35$  so reaches about depth 4 – not so good
  - $\alpha$ - $\beta$  reaches about depth 8 – decent chess program
- **Guarantee of optimal play is gone**
- **More plies makes a BIG difference**
- Use iterative deepening for an anytime algorithm

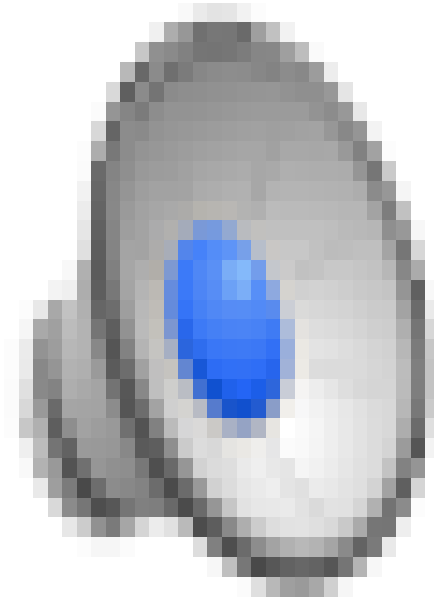


# Depth Matters

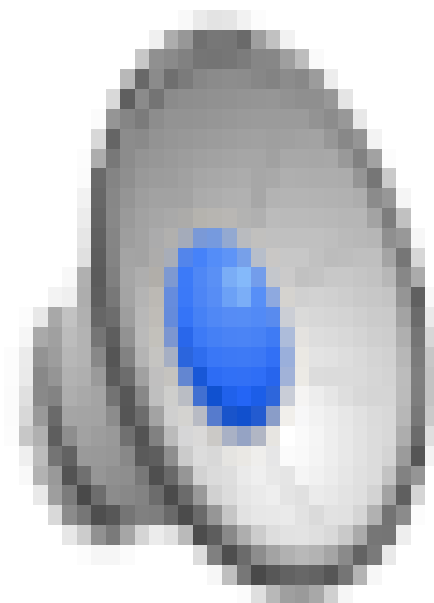
- Evaluation functions are always imperfect
- The deeper in the tree the evaluation function is buried, the less the quality of the evaluation function matters
- An important example of the **tradeoff** between complexity of features and complexity of computation



# Video of Demo Limited Depth (2)



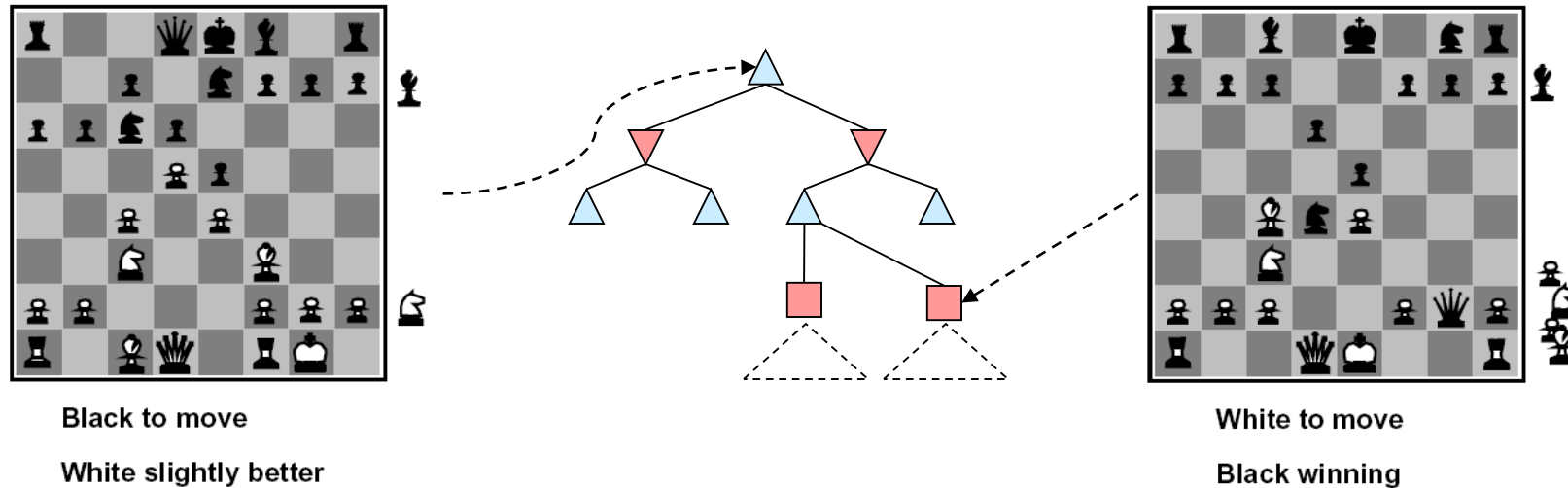
# Video of Demo Limited Depth (10)



# Evaluation Functions

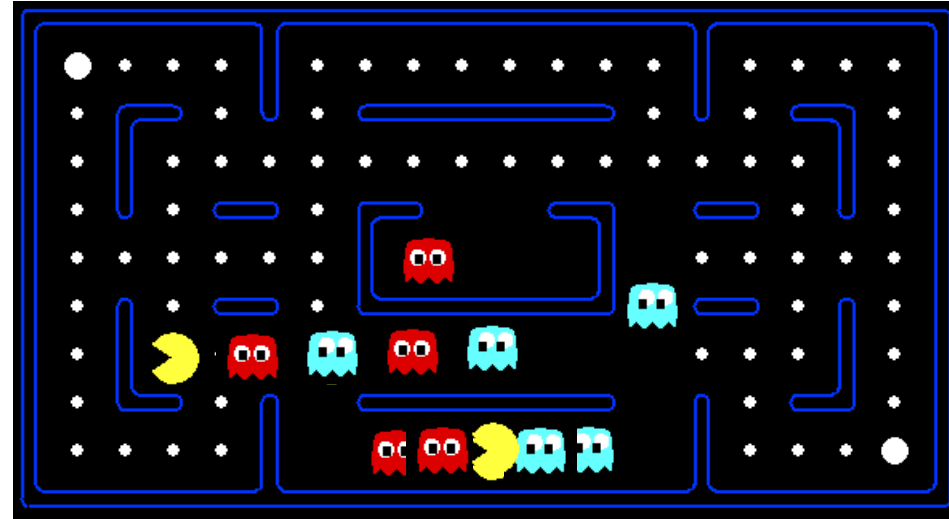


- Evaluation functions score non-terminals in depth-limited search

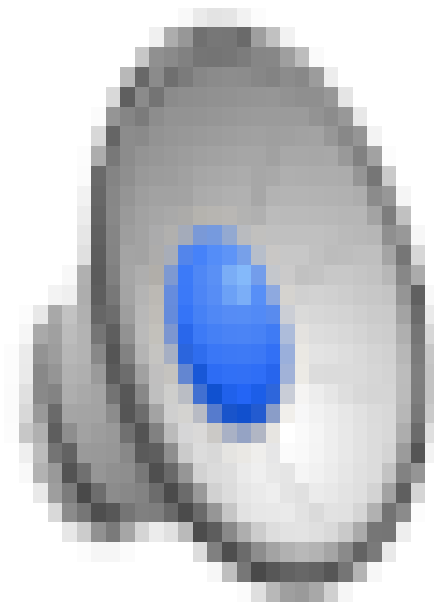


- Ideal function: returns the actual minimax value of the position
- In practice: typically weighted linear sum of features:  
$$\text{EVAL}(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$
  - e.g.  $w_1 = 9$ ,  $f_1(s) = (\text{num white queens} - \text{num black queens})$ , etc.

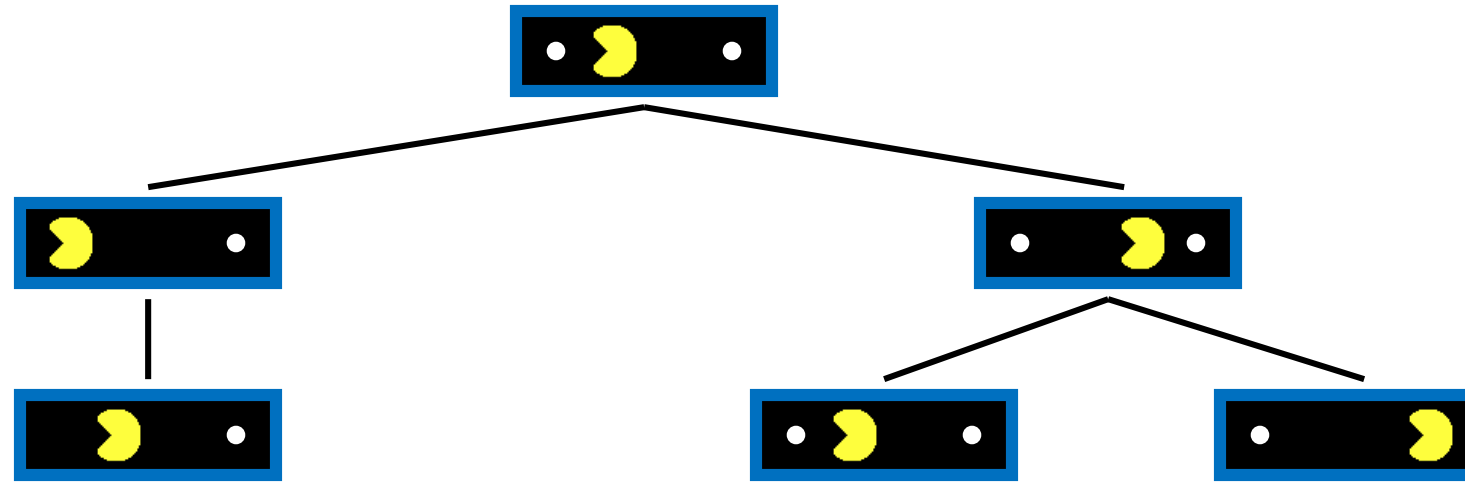
# Evaluation for Pacman



# Video of Demo Thrashing (d=2)



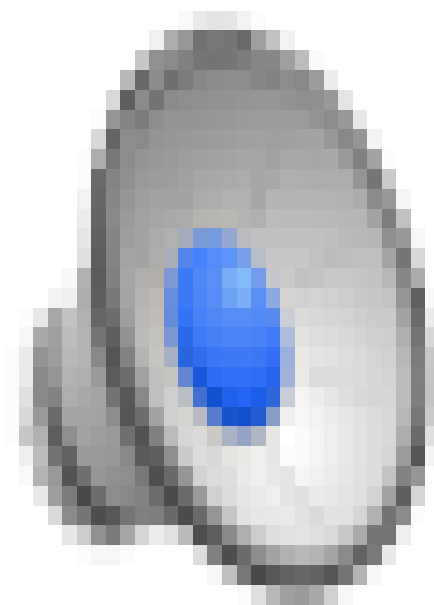
# Why Pacman Starves



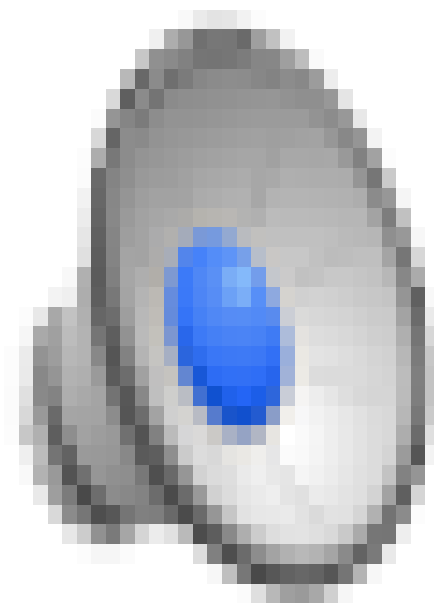
- A danger of replanning agents!
  - He knows his score will go up by eating the dot now (west, east)
  - He knows his score will go up just as much by eating the dot later (east, west)
  - There are no point-scoring opportunities after eating the dot (within the horizon, two here)
  - Therefore, waiting seems just as good as eating: he may go east, then back west in the next round of replanning!



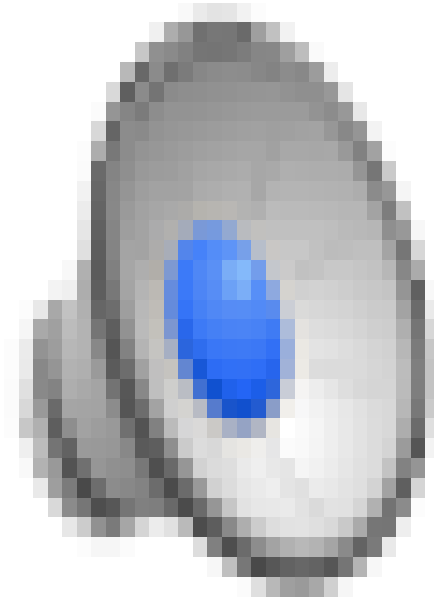
# Video of Demo Thrashing -- Fixed ( $d=2$ )



# Video of Demo Smart Ghosts (Coordination)



# Video of Demo Smart Ghosts (Coordination) – Zoomed In

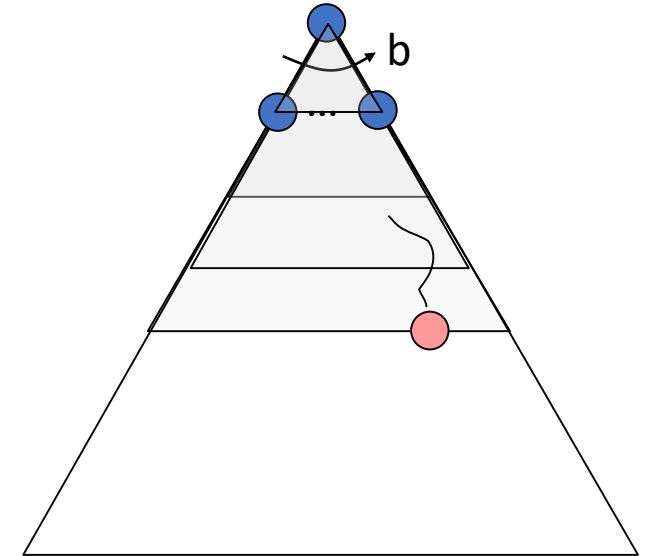


# Iterative Deepening

Iterative deepening uses DFS as a subroutine:

1. Do a DFS which only searches for paths of length 1 or less. (DFS gives up on any path of length 2)
2. If “1” failed, do a DFS which only searches paths of length 2 or less.
3. If “2” failed, do a DFS which only searches paths of length 3 or less.

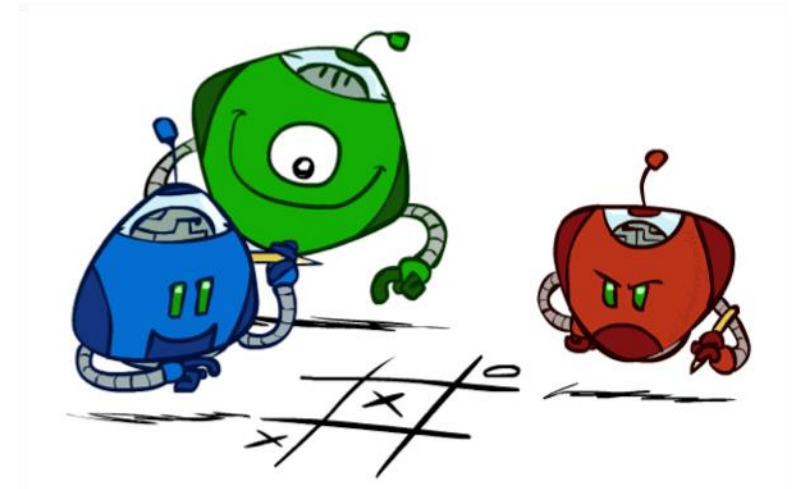
....and so on.



Why do we want to do this for multiplayer games?

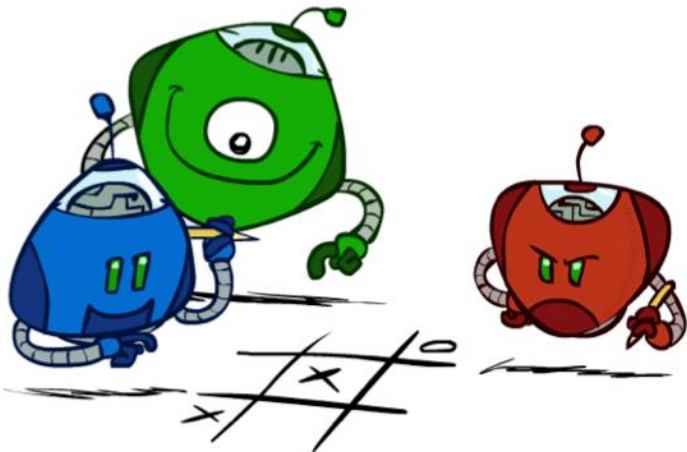
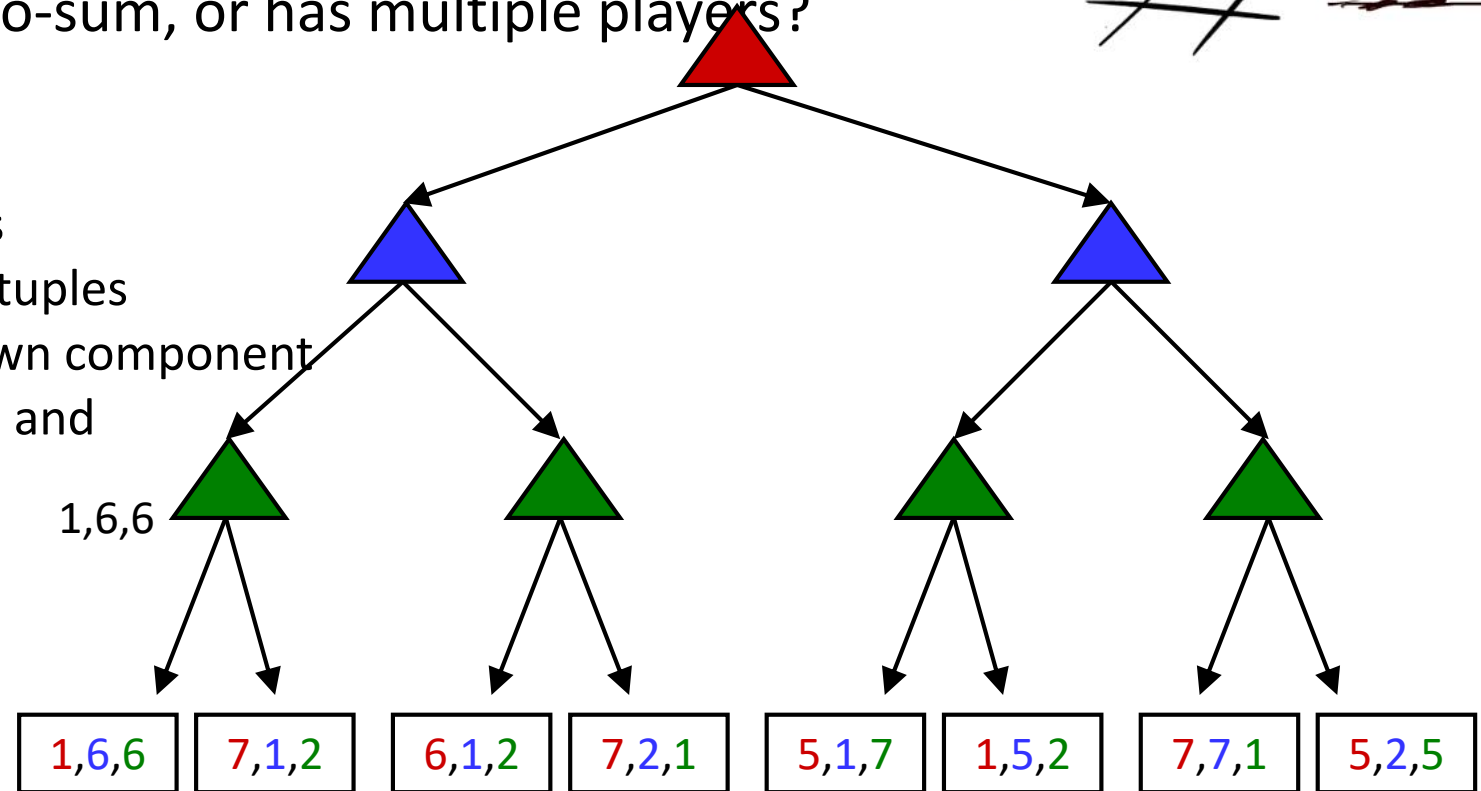
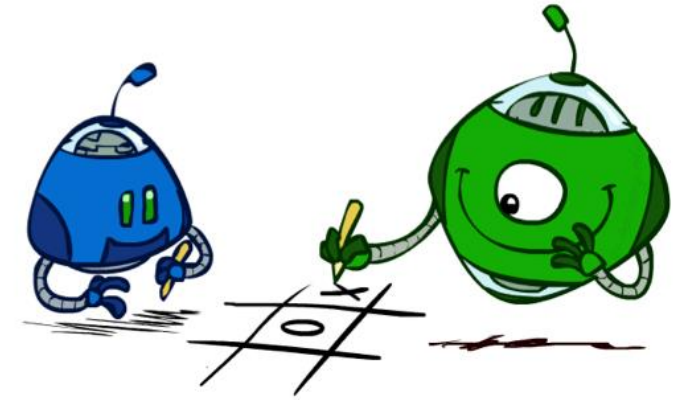
Note: wrongness of eval functions matters less and less the deeper the search goes!

# Generalized minimax



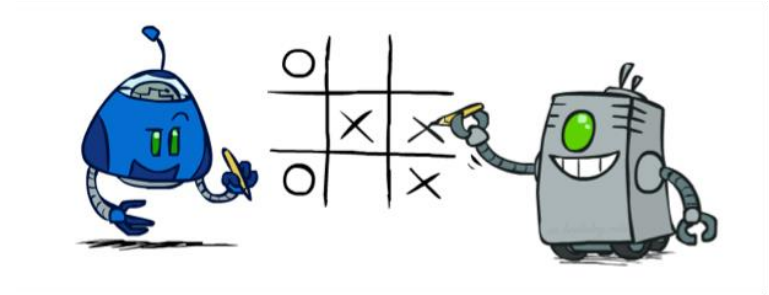
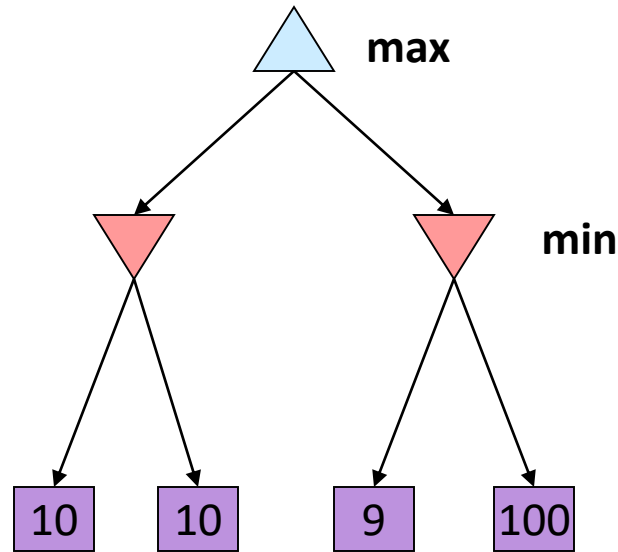
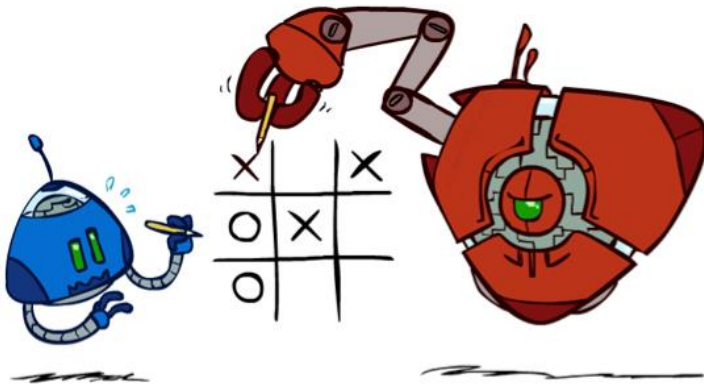
# Multi-Agent Utilities

- What if the game is not zero-sum, or has multiple players?
- Generalization of minimax:
  - Terminals have utility tuples
  - Node values are also utility tuples
  - Each player maximizes its own component
  - Can give rise to cooperation and competition dynamically...



# Modeling Assumptions

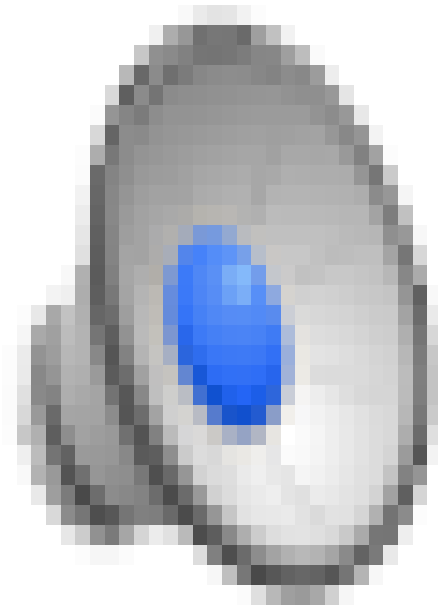
# What if your opponent isn't playing optimally?



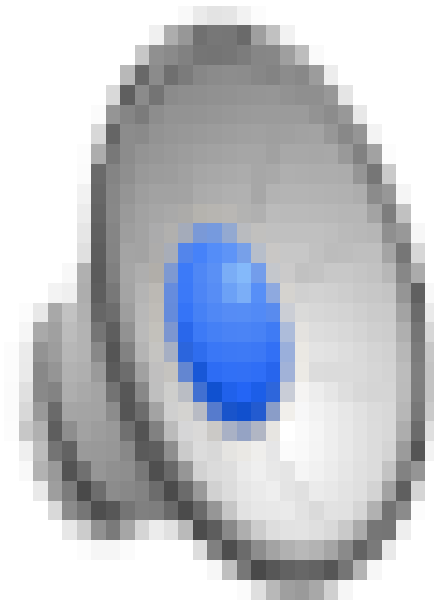
Optimal against a perfect player. Otherwise?



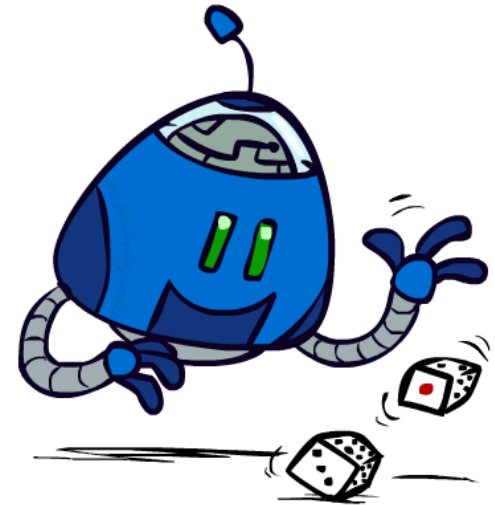
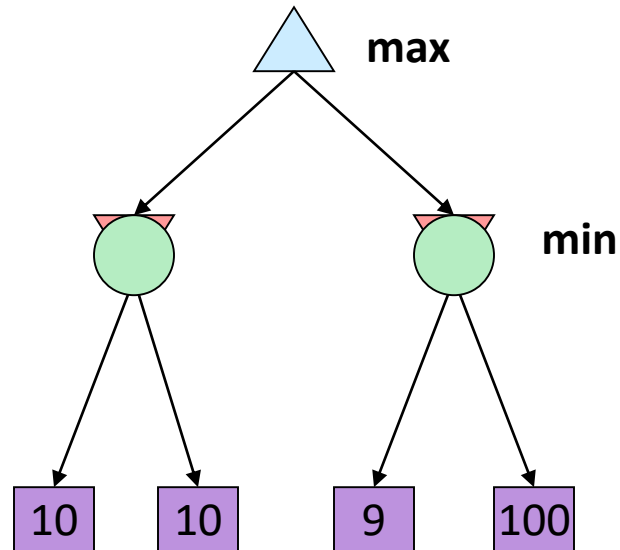
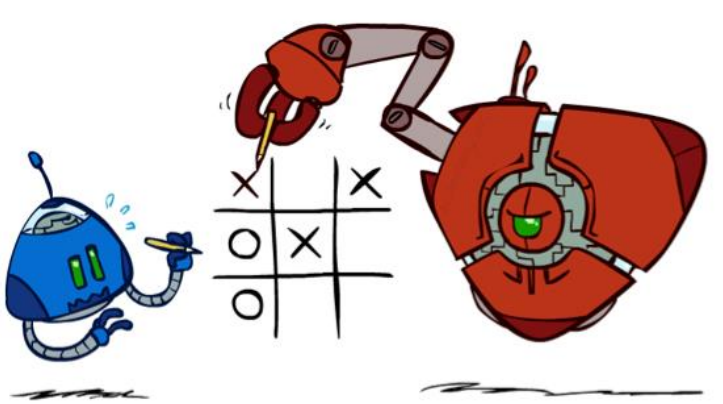
# Video of Demo Min vs. Exp (Min)



# Video of Demo Min vs. Exp (Exp)



# Worst-Case vs. Average Case



Idea: Uncertain outcomes controlled by chance, not an adversary!

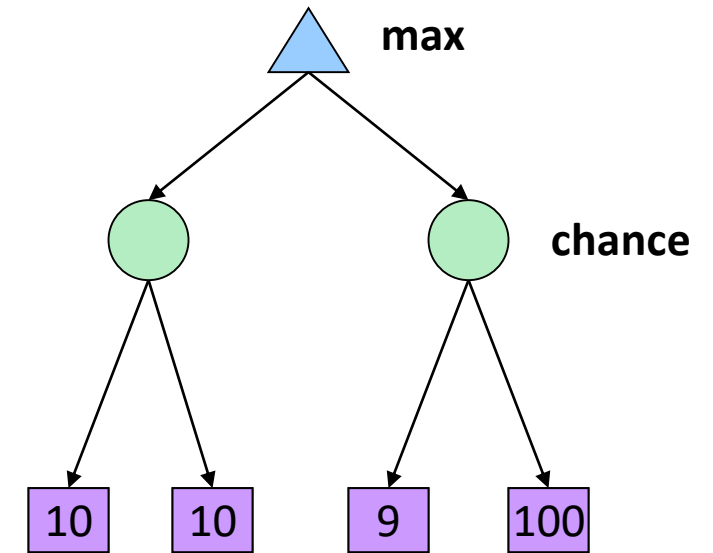
# Why not minimax?

- Worst case reasoning is too conservative
- Need average case reasoning

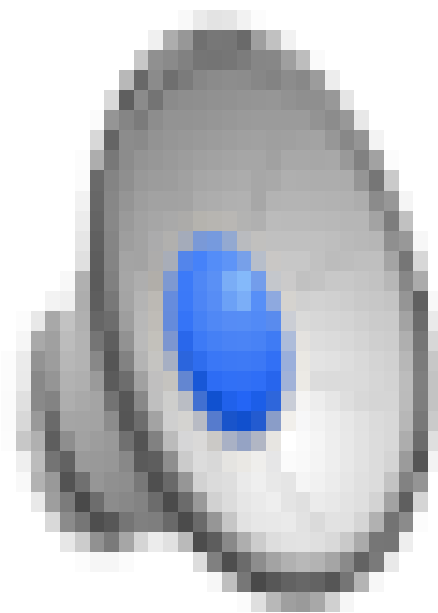


# Expectimax Search

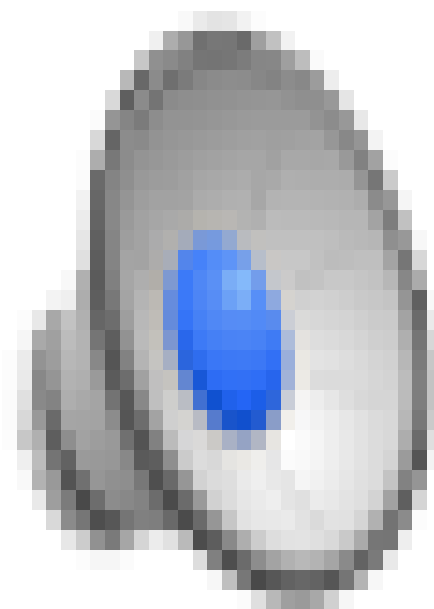
- Why wouldn't we know what the result of an action will be?
  - Explicit randomness: rolling dice
  - Unpredictable opponents: the ghosts respond randomly
  - Unpredictable humans: humans are not perfect
  - Actions can fail: when moving a robot, wheels might slip
- Values should now reflect average-case (expectimax) outcomes, not worst-case (minimax) outcomes
- **Expectimax search:** compute the average score under optimal play
  - Max nodes as in minimax search
  - Chance nodes are like min nodes but the outcome is uncertain
  - Calculate their **expected utilities**
  - I.e. take weighted average (expectation) of children
- Later, we'll learn how to formalize the underlying uncertain-result problems as **Markov Decision Processes**



# Video of Demo Minimax vs Expectimax (Min)



# Video of Demo Minimax vs Expectimax (Exp)



# Expectimax Pseudocode

```
def value(state):
```

if the state is a terminal state: return the state's utility

if the next agent is MAX: return max-value(state)

if the next agent is EXP: return exp-value(state)

```
def max-value(state):
```

initialize  $v = -\infty$

for each successor of state:

$v = \max(v, \text{value}(\text{successor}))$

return  $v$

```
def exp-value(state):
```

initialize  $v = 0$

for each successor of state:

$p = \text{probability}(\text{successor})$

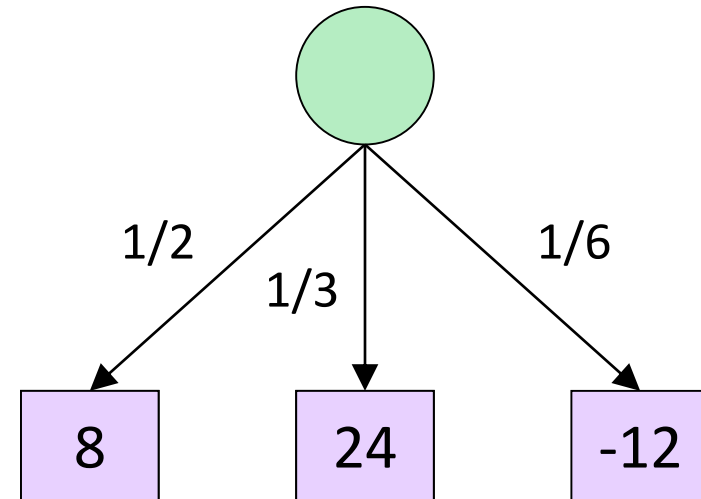
$v += p * \text{value}(\text{successor})$

return  $v$



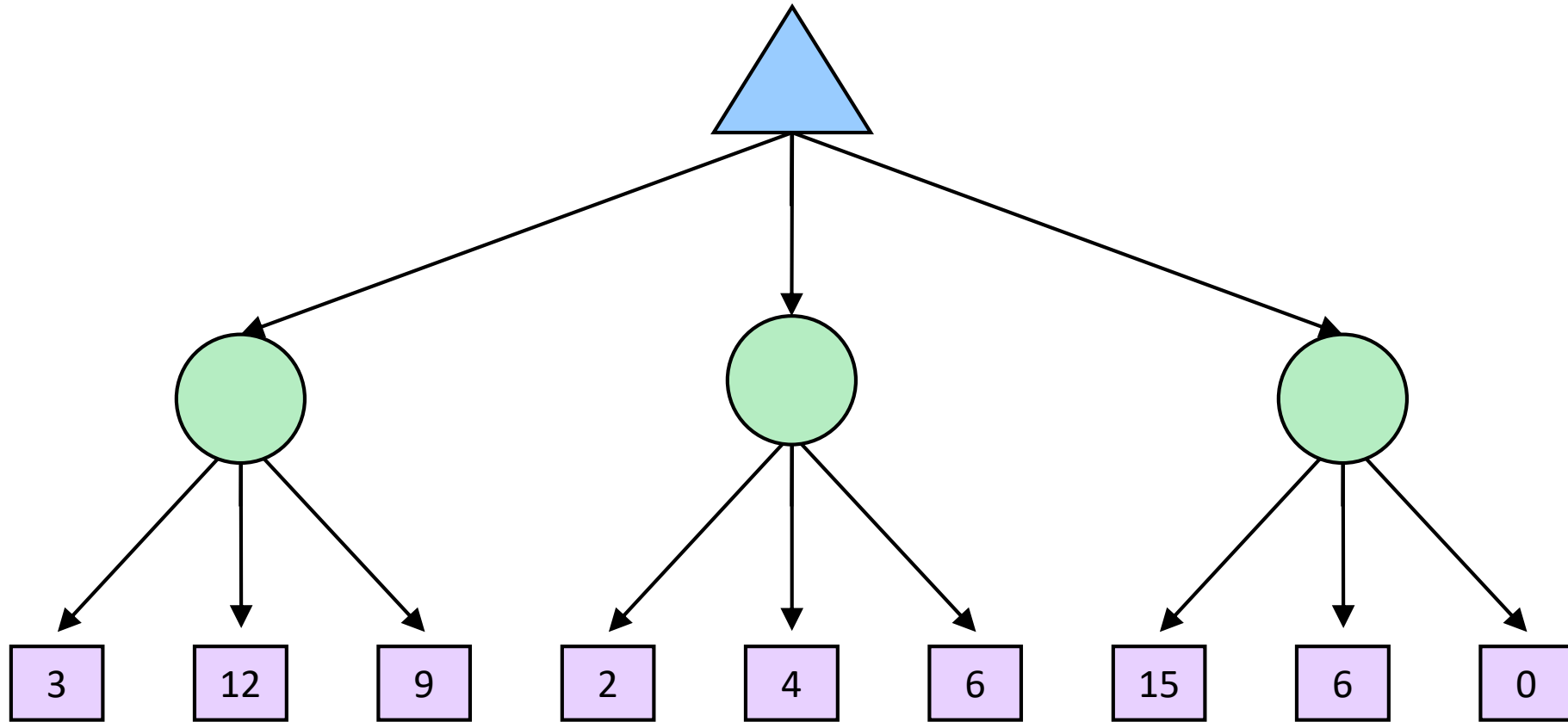
# Expectimax Pseudocode 2

```
def exp-value(state):  
    initialize v = 0  
    for each successor of state:  
        p = probability(successor)  
        v += p * value(successor)  
    return v
```

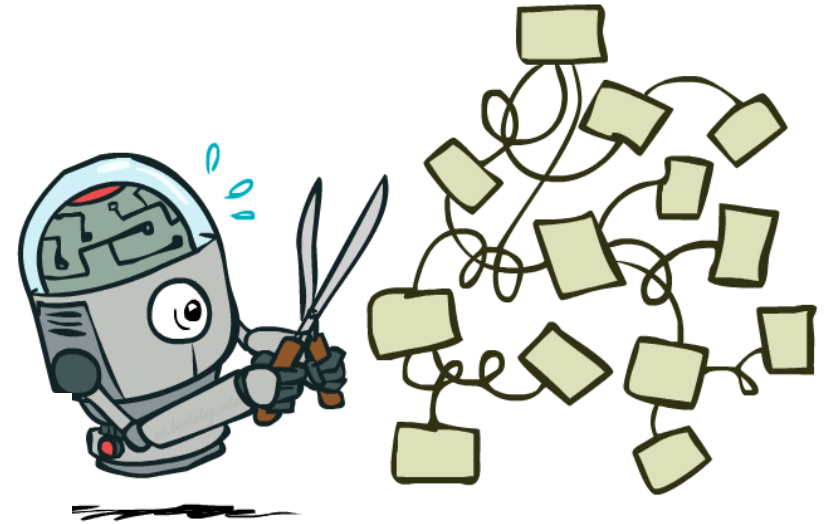
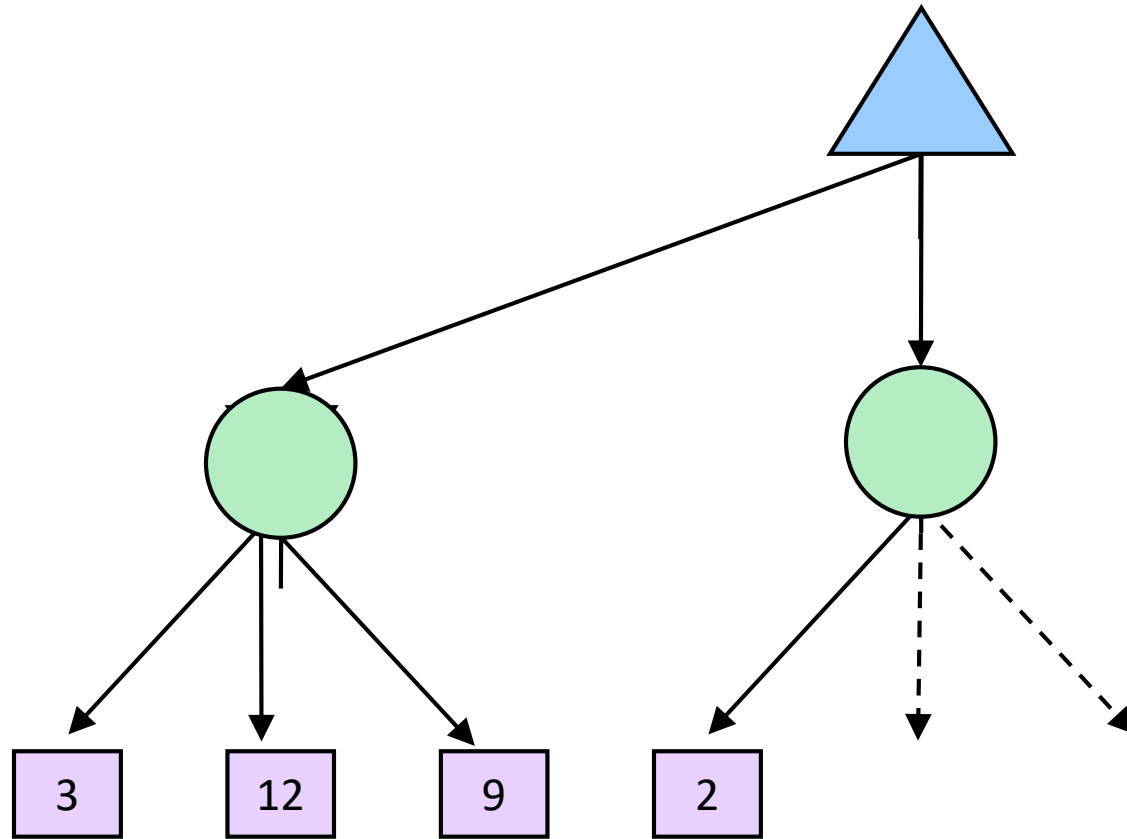


$$v = (1/2) (8) + (1/3) (24) + (1/6) (-12) = 10$$

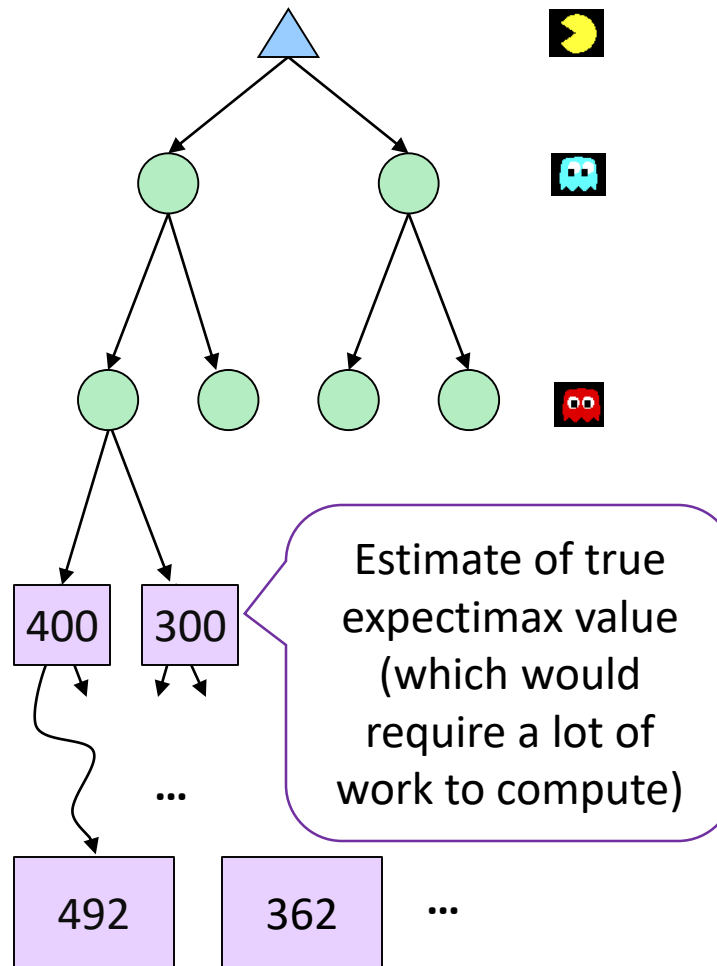
# Example



# Expectimax Pruning?

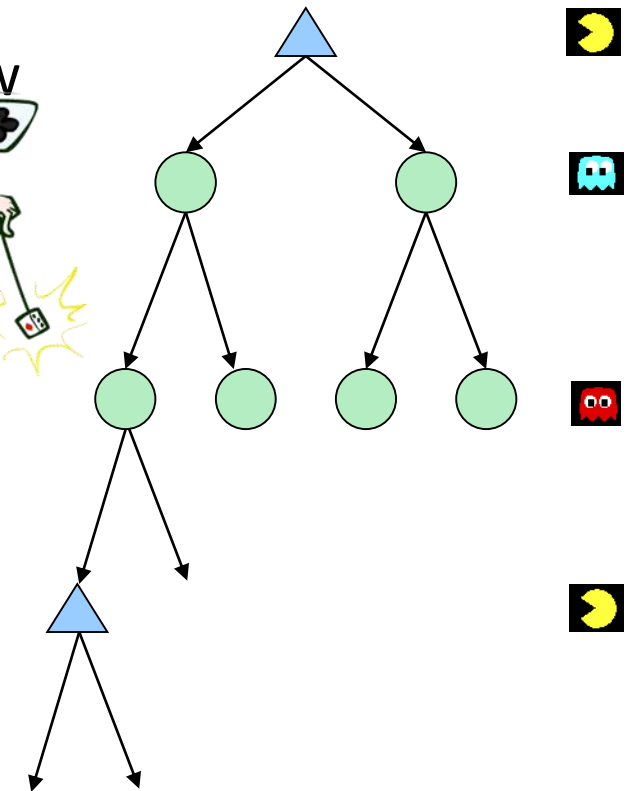


# Expectimax: Depth-Limited



# What Probabilities to Use?

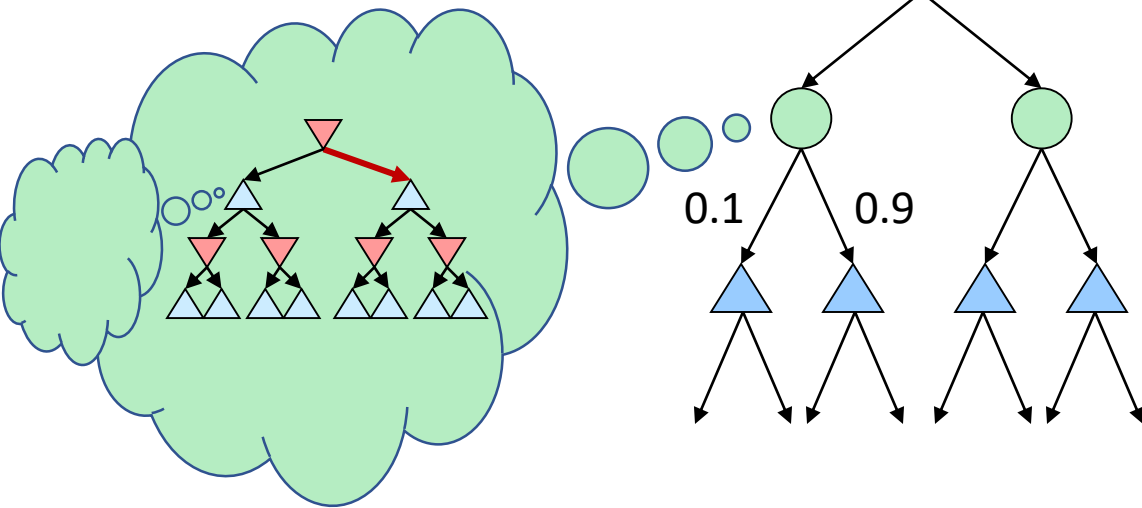
- In expectimax search, we have a probabilistic model of how the opponent (or environment) will behave in an action  
  - Model could be a simple uniform distribution (roll a die)
  - Model could be sophisticated and require a great deal of computation
  - We have a chance node for any outcome out of our control: opponent or environment
  - The model might say that adversarial actions are likely!
- For now, assume each chance node magically comes along with probabilities that specify the distribution over its outcomes



*Having a probabilistic belief about another agent's action does not mean that the agent is flipping any coins!*

# Quiz: Informed Probabilities

- Let's say you know that your opponent is actually running a depth 2 minimax, using the result 80% of the time, and moving randomly otherwise
- Question: What tree search should you use?



## ■ Answer: Expectimax!

- To figure out EACH chance node's probabilities, you have to run a simulation of your opponent
- This kind of thing gets very slow very quickly
- Even worse if you have to simulate your opponent simulating you...
- ... except for minimax and maximax, which have the nice property that it all collapses into one game tree

*This is basically how you would model a human, except for their utility: their utility might be the same as yours (i.e. you try to help them, but they are depth 2 and noisy), or they might have a slightly different utility (like another person navigating in the office)*

# Dangerous Pessimism/Optimism

## Dangerous Pessimism

Assuming the worst case when it's not likely

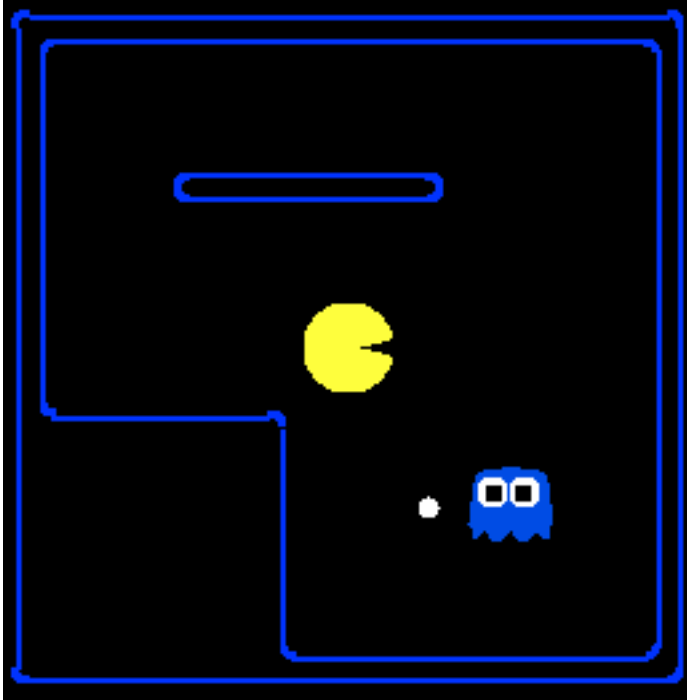


## Dangerous Optimism

Assuming chance when the world is adversarial



# Assumptions vs. Reality



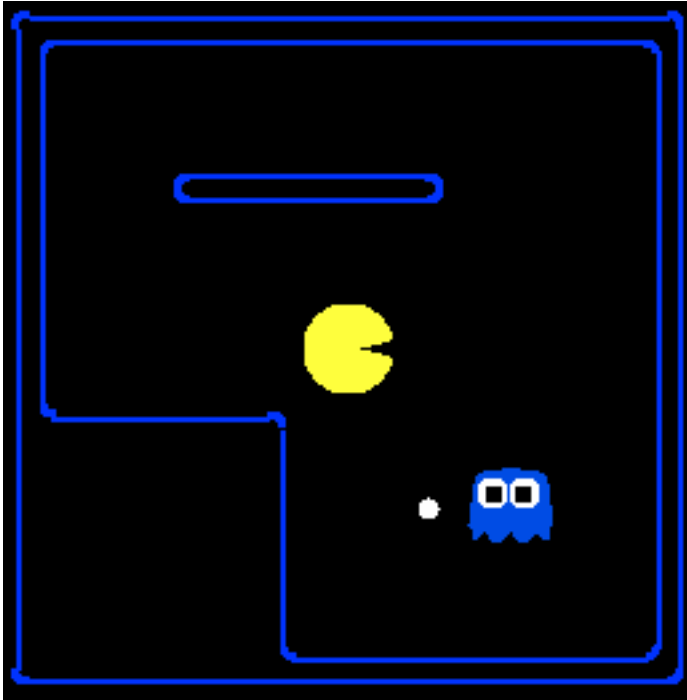
	Adversarial Ghost	Random Ghost
Minimax Pacman		
Expectimax Pacman		

Results from playing 5 games

Pacman used depth 4 search with an eval function that avoids trouble  
Ghost used depth 2 search with an eval function that seeks Pacman



# Assumptions vs. Reality 2



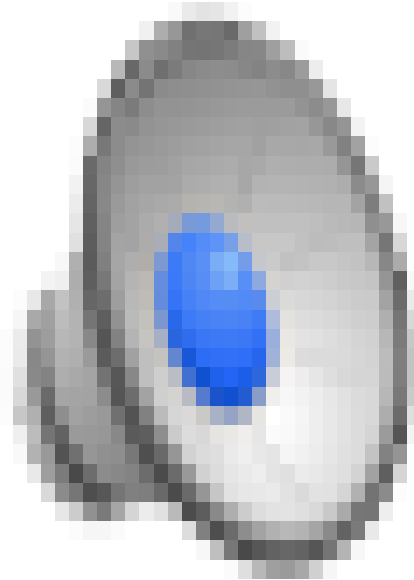
	Adversarial Ghost	Random Ghost
Minimax Pacman	Won 5/5 Avg. Score: 483	Won 5/5 Avg. Score: 493
Expectimax Pacman	Won 1/5 Avg. Score: -303	Won 5/5 Avg. Score: 503

Results from playing 5 games

Pacman used depth 4 search with an eval function that avoids trouble  
Ghost used depth 2 search with an eval function that seeks Pacman

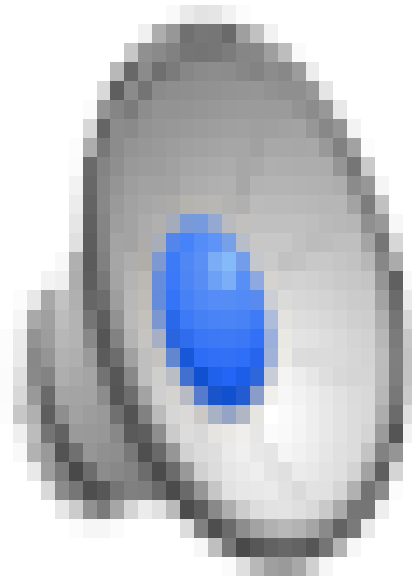
# Video of Demo World Assumptions

## Random Ghost – Expectimax Pacman



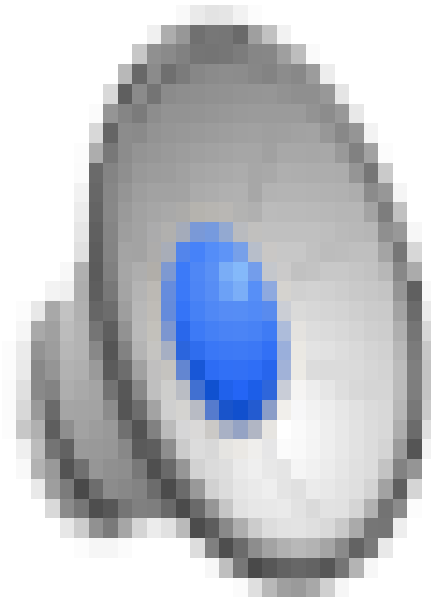
# Video of Demo World Assumptions

## Adversarial Ghost – Minimax Pacman



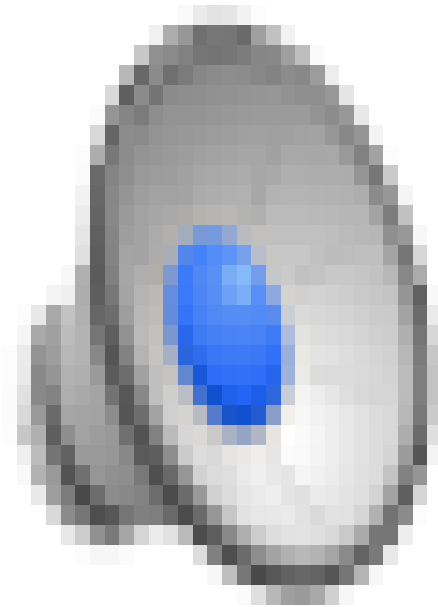
# Video of Demo World Assumptions

## Adversarial Ghost – Expectimax Pacman



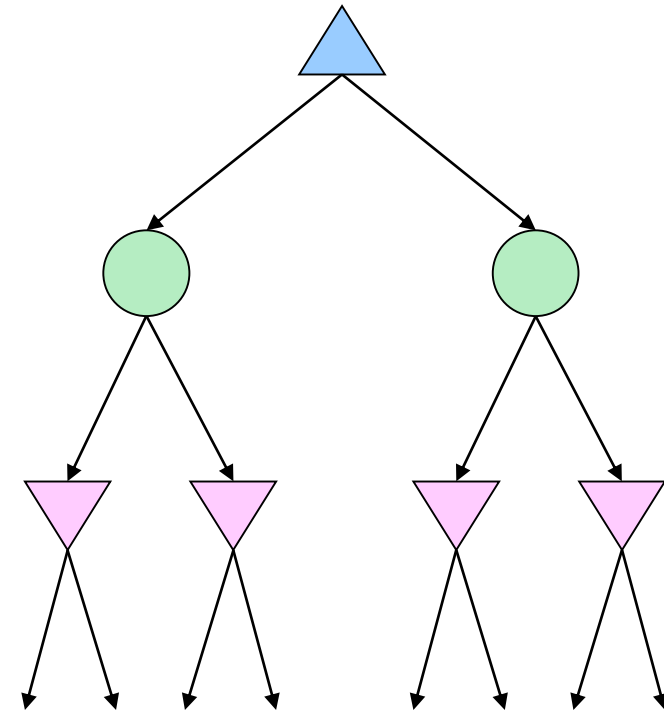
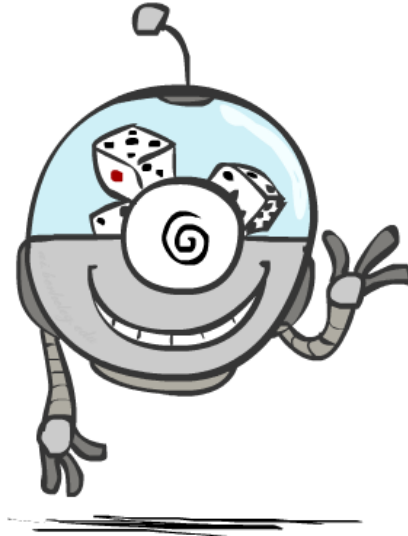
# Video of Demo World Assumptions

## Random Ghost – Minimax Pacman



# Mixed Layer Types

- E.g. Backgammon
- Expectiminimax
  - Environment is an extra “random agent” player that moves after each min/max agent
  - Each node computes the appropriate combination of its children



# Example: Backgammon

- Dice rolls increase  $b$ : 21 possible rolls with 2 dice
  - Backgammon  $\approx 20$  legal moves
  - Depth 2 =  $20 \times (21 \times 20)^3 = 1.2 \times 10^9$
- As depth increases, probability of reaching a given search node shrinks
  - So usefulness of search is diminished
  - So limiting depth is less damaging
  - But pruning is trickier...
- Historic AI: TDGammon uses depth-2 search + very good evaluation function + reinforcement learning: world-champion level play
- 1<sup>st</sup> AI world champion in any game!



# Summary

- Types of games
- Adversarial Game Trees: Minimax search
- Resource Limits I: Alpha-Beta Pruning
- Resource Limits II: Depth-limited search
  - Evaluation functions
  - Iterative deepening

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## Questions?