

From decimal expansions to reduced fractions

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```
In [1]: from math import gcd
```

A real number is rational if and only if a pattern eventually appears in its decimal expansion that repeats forever. So π , being irrational, is such that no finite sequence of consecutive digits in 3.14159265358979... eventually repeats forever. On the other hand,

- $\frac{25}{12} = 2.08333...3...$
- $\frac{97}{21} = 4.619047619047619047...619047...$
- $\frac{11941}{49950} = 0.23905905905...905...$

The decimal expansion is unique except for fractions that in reduced form, have a power of 10 as denominator: those fractions have two decimal expansions, one that ends in 0 repeating forever, another one that ends in 9 repeating forever. For instance, $\frac{1234567}{1000} = 1234.567000...0... = 1234.566999...9...$

We want to, given two nonempty strings of digits σ and τ (that we treat as strings or numbers depending on the context), find out the unique natural numbers p and q such that the decimal expansion of $\frac{p}{q}$ reads as $0.\sigma\tau\tau\tau...\tau...$ and

- either $p = 0$ and $q = 1$ (case where σ and τ consist of nothing but 0's), or
- p and q are coprime, so $\frac{p}{q}$ is in reduced form (including the case where $p = 1$ and $q = 1$ because σ and τ consist of nothing but 9's).

For instance, if $\sigma = 23$ and $\tau = 905$, then $p = 11941$ and $q = 49950$.

Writing $|\sigma|$ for the length (number of digits) in a string of digits σ , we compute:

$$\begin{aligned} 0.\sigma\tau\tau\tau...\tau... &= \sigma 10^{-|\sigma|} + \tau(10^{-|\sigma|-|\tau|} + 10^{-|\sigma|-2|\tau|} + 10^{-|\sigma|-3|\tau|} + \dots) \\ &= \sigma 10^{-|\sigma|} + \frac{\tau 10^{-|\sigma|-|\tau|}}{(1 - 10^{-|\tau|})} \\ &= \sigma 10^{-|\sigma|} + \frac{\tau 10^{-|\sigma|}}{(10^{|\tau|} - 1)} \\ &= \frac{\sigma 10^{-|\sigma|}(10^{|\tau|} - 1) + \tau 10^{-|\sigma|}}{(10^{|\tau|} - 1)} \\ &= \frac{\sigma(10^{|\tau|} - 1) + \tau}{(10^{|\tau|} - 1)10^{|\sigma|}} \end{aligned}$$

Reducing the last fraction if needed provides the desired answer.

The result of the previous computation immediately translates to the function that follows:

```
In [2]: def compute_fraction(sigma, tau):
        numerator = int(sigma) * (10 ** len(tau) - 1) + int(tau)
        denominator = (10 ** len(tau) - 1) * 10 ** len(sigma)
        return numerator, denominator

        compute_fraction('23', '905')
        compute_fraction('000', '97')
        compute_fraction('97', '000')
        compute_fraction('01234', '543210')
```

```
Out[2]: (23882, 99900)
```

```
Out[2]: (97, 99000)
```

```
Out[2]: (96903, 99900)
```

```
Out[2]: (1234541976, 99999900000)
```

To reduce a fraction, it suffices to divide its numerator and its denominator by their gcd (greatest common divisor). The **math** module has a **gcd** function:

```
In [3]: gcd(1234541976, 99999900000)
```

```
Out[3]: 24
```

Let us implement the gcd function ourselves, following Euclid's algorithm, which is based on the following reasoning. Let a and b two natural numbers with $b > 0$. Since $a = \lfloor \frac{a}{b} \rfloor b + a \bmod b$:

- if n divides both a and b then it divides both a and $\lfloor \frac{a}{b} \rfloor b$, hence it divides $a - \lfloor \frac{a}{b} \rfloor b$, hence it divides $a \bmod b$;
- conversely, if n divides both b and $a \bmod b$ then it divides $\lfloor \frac{a}{b} \rfloor b + a \bmod b$, hence it divides a .

Hence n divides both a and b iff n divides both b and $a \bmod b$. So $\text{gcd}(a, b) = \text{gcd}(b, a \bmod b)$.

Since $a \bmod b < b$, we get a sequence of equalities of the form: $\text{gcd}(a, b) = \text{gcd}(a_1, b_1) = \text{gcd}(a_2, b_2) = \dots = \text{gcd}(a_{k-1}, b_{k-1}) = \text{gcd}(a_k, 0)$ with $k \geq 1$ and $b > b_1 > b_2 > \dots > b_{k-1} > 0$; as $\text{gcd}(a_k, 0) = a_k$, a_k is the gcd of a and b .

To compute $\lfloor \frac{a}{b} \rfloor$, Python offers the `//` operator; to compute $a \bmod b$, the `%` operator:

```
In [4]: # True division.
        # The result is always a floating point number.
        8 / 2, 8.0 / 2, 8 / 2.0, 8.0 / 2.0

        # Integer division.
        # The result is an integer iff both arguments are integers.
        9 // 2, 9.0 // 2, 9 // 2.0, 9.0 // 2.0

        # Remainder.
        # The result is an integer iff both arguments are integers.
        9 % 2, 9.0 % 2, 9 % 2.0, 9.0 % 2.0
```

```
Out[4]: (4.0, 4.0, 4.0, 4.0)
```

```
Out[4]: (4, 4.0, 4.0, 4.0)
```

```
Out[4]: (1, 1.0, 1.0, 1.0)
```

If a and b are arbitrary numbers (not necessarily integers) with $b \neq 0$, then the equality $a = qb + r$ together with the conditions

- q is an integer
- $|r| < |b|$
- $r \neq 0 \rightarrow (r > 0 \leftrightarrow b > 0)$

determine q and r uniquely; // and % operate accordingly:

```
In [5]: 5 // 2, 5 % 2
        -5 // 2, -5 % 2
        5 // -2, 5 % -2
        -5 // -2, -5 % -2
        print()

        7.5 // 2, 7.5 % 2
        -7.5 // 2, -7.5 % 2
        7.5 // -2, 7.5 % -2
        -7.5 // -2, -7.5 % -2
```

```
Out[5]: (2, 1)
```

```
Out[5]: (-3, 1)
```

```
Out[5]: (-3, -1)
```

```
Out[5]: (2, -1)
```

```
Out[5]: (3.0, 1.5)
```

```
Out[5]: (-4.0, 0.5)
```

```
Out[5]: (-4.0, -0.5)
```

```
Out[5]: (3.0, -1.5)
```

The **divmod()** function offers an alternative to the previous combined use of // and %:

```
In [6]: divmod(5, 2)
        divmod(-5, 2)
        divmod(5, -2)
        divmod(-5, -2)
        print()

        divmod(7.5, 2)
        divmod(-7.5, 2)
        divmod(7.5, -2)
        divmod(-7.5, -2)
```

```
Out[6]: (2, 1)
```

```
Out[6]: (-3, 1)
```

```
Out[6]: (-3, -1)
```

```
Out[6]: (2, -1)
```

```
Out[6]: (3.0, 1.5)
```

```
Out[6]: (-4.0, 0.5)
```

```
Out[6]: (-4.0, -0.5)
```

```
Out[6]: (3.0, -1.5)
```

Let us get back to Euclid's algorithm, so assume again that a and b are two natural numbers with $b > 0$. To implement the algorithm and compute $\text{gcd}(a, b)$, it suffices to have two variables, say **a** and **b**, initialised to a and b , and then change the value of **a** to b and change the value of **b** to $a \bmod b$, and do that again and again until **b** gets the value 0. To change the value of **a** to $a \bmod b$ and change the value of **b** to b , it seems necessary to introduce a third variable:

```
In [7]: a = 30
        b = 18
        c = a % b
        a = b
        b = c
        a, b
```

```
Out[7]: (18, 12)
```

But Python makes it easier:

```
In [8]: a = 30
        b = 18
        # Evaluate the expression on the right hand side;
        # the result is the tuple (18, 12).
        # Then assign that result to the tuple on the left,
        # component by component.
        a, b = b, a % b
        a, b
```

```
Out[8]: (18, 12)
```

Note that when the value of **a** is strictly smaller than the value of **b**, then **a, b = b, a % b** exchanges the values of **a** and **b**:

```
In [9]: a = 12
        b = 18
        a, b = b, a % b
        a, b
```

```
Out[9]: (18, 12)
```

On the other hand, if the value of **a** is at least equal to the value of **b**, then this holds too after **a, b = b, a % b** has been executed. Let us trace all stages in the execution of Euclid's algorithm. The code makes use of a **while** statement whose condition is not a boolean expression. Applying **bool()** to an expression reveals which one of **True** or **False** the expression evaluates to in contexts where one or the other is expected:

```
In [10]: bool(None)
          bool(0), bool(5), bool(-3)
          bool(0.0), bool(0.1), bool(-3.14)
          bool([], bool([0]), bool([[]]))
          bool({}), bool({0: 0}), bool({0: None, 1: None})
          bool(''), bool(' '), bool('0000')
```

```
Out[10]: False
```

```
Out[10]: (False, True, True)
```

```
Out[10]: (False, True, True)
```

```
Out[10]: (False, True, True)
```

```
Out[10]: (False, True, True)
```

```
Out[10]: (False, True, True)
```

```
In [11]: def trace_our_gcd(a, b):
          while b:
              a, b = b, a % b
              print(a, b)

          for a, b in (1233, 1233), (1233, 990), (990, 1233):
              print(f'\nTracing the computation of gcd of {a} and {b}:')
              trace_our_gcd(a, b)
```

```
Tracing the computation of gcd of 1233 and 1233:
1233 0
```

```
Tracing the computation of gcd of 1233 and 990:
990 243
243 18
18 9
9 0
```

```

Tracing the computation of gcd of 990 and 1233:
1233 990
990 243
243 18
18 9
9 0

```

The gcd is the value of **a** when exiting the **while** loop:

```

In [12]: def our_gcd(a, b):
        while b:
            a, b = b, a % b
        return a

```

compute_fraction() returns the numerator and denominator of a fraction that another function, say **reduce()**, can easily reduce thanks to **our_gcd()**. It is natural to let **reduce()** take two arguments, the numerator and the denominator of the fraction to simplify, respectively. But **compute_fraction()** returns those as the first and second elements of a tuple; a function always returns a single value. Between the parentheses that surround the arguments of a function **f()**, one can insert the denotation of a tuple and precede it with the ***** symbol, which “unpacks” the members of the tuple and make them the arguments of **f()**:

```

In [13]: def f(a, b):
        return 2 * a, 2 * b

```

```

# Makes a equal to (1, 3), and provides no value to b.
f((1, 3))

```

```

TypeError                                Traceback (most recent call last)

```

```

<ipython-input-12-eae47008b792> in <module>()
    3
    4 # Makes a equal to (1, 3), and provides no value to b.
----> 5 f((1, 3))

```

```

TypeError: f() missing 1 required positional argument: 'b'

```

```

In [14]: f(1, 3)
        # f(f(1, 3)) would be f((2, 6)); f(*f(1, 3)) is f(2, 6)
        f(*f(1, 3))
        f(*f(*f(1, 3)))

```

```

Out[14]: (2, 6)

```

```
Out[14]: (4, 12)
```

```
Out[14]: (8, 24)
```

The `*` symbol can also be used in the definition of a function and precede the name of a parameter. It then has the opposite effect, namely, it makes a tuple out of all arguments that are provided to the function:

```
In [15]: # x is the tuple of all arguments passed to f().
```

```
def f(*x):  
    return x * 2
```

```
f()  
f(0)  
f(f(0))  
f(*f(0))  
f(f(f(0)))  
f(f(*f(0)))  
f(*f(*f(0)))
```

```
Out[15]: ()
```

```
Out[15]: (0, 0)
```

```
Out[15]: ((0, 0), (0, 0))
```

```
Out[15]: (0, 0, 0, 0)
```

```
Out[15]: (((0, 0), (0, 0)), ((0, 0), (0, 0)))
```

```
Out[15]: ((0, 0, 0, 0), (0, 0, 0, 0))
```

```
Out[15]: (0, 0, 0, 0, 0, 0, 0, 0)
```

Thanks to this syntax, it is possible to let **reduce()** as well as another function **output()** take two arguments **numerator** and **denominator**, and “pipe” **compute_fraction()**, **reduce()** and **output()** together so that the unpacked returned value of one function becomes the arguments of the function that follows:

```
In [16]: def reduce(numerator, denominator):
```

```
    if numerator == 0:  
        return 0, 1  
    the_gcd = our_gcd(numerator, denominator)  
    return numerator // the_gcd, denominator // the_gcd
```

```
In [17]: def output(numerator, denominator):
```

```
    print(f'{numerator}/{denominator}')
```

```
In [18]: output(*reduce(*compute_fraction('23', '905')))  
output(*reduce(*compute_fraction('000', '97')))  
output(*reduce(*compute_fraction('97', '000')))  
output(*reduce(*compute_fraction('01234', '543210')))
```

11941/49950
97/99000
97/100
51439249/4166662500