



## Informatics II

### Exercise 11

May 05, 2020

### Dynamic programming

**Task 1.** Imagine several houses built in a circle. Festive illumination needs to be arranged for Christmas and each house has a given number of light bulbs to use for the event. In order to increase the effect of the illuminations, **only houses which are not next to each other (non-neighbor) shall participate.**

Given an array with the number of light bulbs per house, calculate the highest possible number of light bulbs which can be simultaneously lit up in the night. **Example:** Given the array `lightBulbs[ ] = {17, 5, 10, 18, 22}`, it would be optimal if the house 0 (with 17 light bulbs) and house 3 (with 18 light bulbs) would participate. Houses 0 and 4 cannot be combined since they are neighbors due to the circular structure.

Solve the illumination problem using (a) recursion, (b) memoization and (c) dynamic programming. Create a C program including the following functions:

- `int maxIlluminationRecursive(int lightBulbs[ ], int n)` which solves the problem recursively.
- `int maxIlluminationMemoized(int lightBulbs[ ], int n, int m[ ])` which solves the problem using memoization. The array `m[n]` is used in order to store the intermediate results.
- `int maxIlluminationDynamic(int lightBulbs[ ], int n)` which solves the problem with the help of dynamic programming.

Print the results of your functions for each of the following light bulbs-arrays: [ 11, 4, 3, 6, 8, 9 ], [ 4, 4, 4, 4, 4 ] and [ 13, 15, 31, 21, 9, 12, 44, 32, 12, 43, 22, 9, 11, 32, 26, 22, 21, 3, 4, 29 ].

**Task 2.** Write in C a function `int lenOfLongestGP(int set[ ], int n)` that uses dynamic programming to calculate and return **Length of the Longest Geometric Progression (LLGP)** in a given set of numbers. The common ratio of the Geometric Progression must be an integer. For example in series [5, 7, 10, 15, 20, 29], the LLGP is 3 corresponding to subseries [5, 10, 20] with common ratio of 2.



**Task 3.** Assume a rectangle of size  $n \times m$ . Determine the minimum number of squares that is prepared to tile the rectangle. The side length of rectangles must be an integer. Figure show as rectangle of  $6 \times 5$  that requires tiles to tile it with the minimum number of tiles.

The idea for dynamic programming solution is to recursively split the rectangle of size  $(n \times m)$  into either two horizontal rectangles  $((n \times h)$  and  $(n \times (m - h))$  or two vertical rectangles  $((v \times m)$  and  $((n - v) \times m)$  where  $h = 1, 2, \dots, m - 1$  and  $v = 1, 2, \dots, n - 1$ . Then finding the minimum squares that fill these splits.

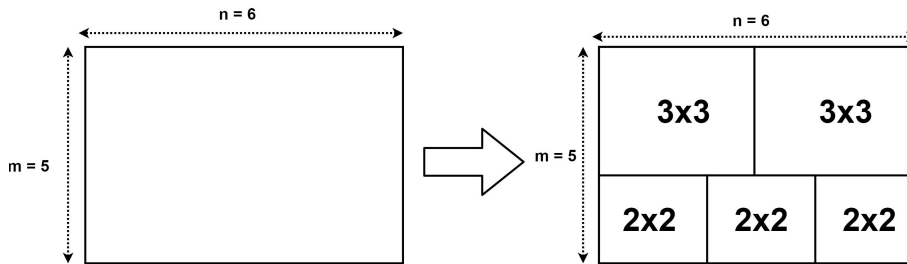


Figure 1: Minimum number of square tiles for a rectangle of size  $5 \times 6$

- Assume  $\text{MinSquare}(n, m)$  denotes the minimum squares that is required to tile a rectangle of size  $n \times m$ . State a recursive definition of  $\text{MinSquare}(n, m)$ .
- To efficiently compute the minimum squares that is required to tile a rectangle of  $n \times m$ , a dynamic programming solution with 2D array  $\text{dp}[0 \dots n][0 \dots m]$  can be used. Element  $\text{dp}[i][j]$  is the minimum number of squares that is required to tile a rectangle of size  $i \times j$ . Complete the 2D array below with the minimum squares required to tile rectangles up to size  $4 \times 3$ .

| $\begin{matrix} n \\ m \end{matrix}$ | 0 | 1 | 2 | 3 | 4 |
|--------------------------------------|---|---|---|---|---|
| 0                                    |   |   |   |   |   |
| 1                                    |   |   |   |   |   |
| 2                                    |   |   |   |   |   |
| 3                                    |   |   |   |   |   |
|                                      |   |   |   |   |   |
|                                      |   |   |   |   |   |

- For a rectangle of size  $n \times m$ , write a C program that uses a dynamic programming solution to compute the minimum number of squares that is required to tile the rectangle.