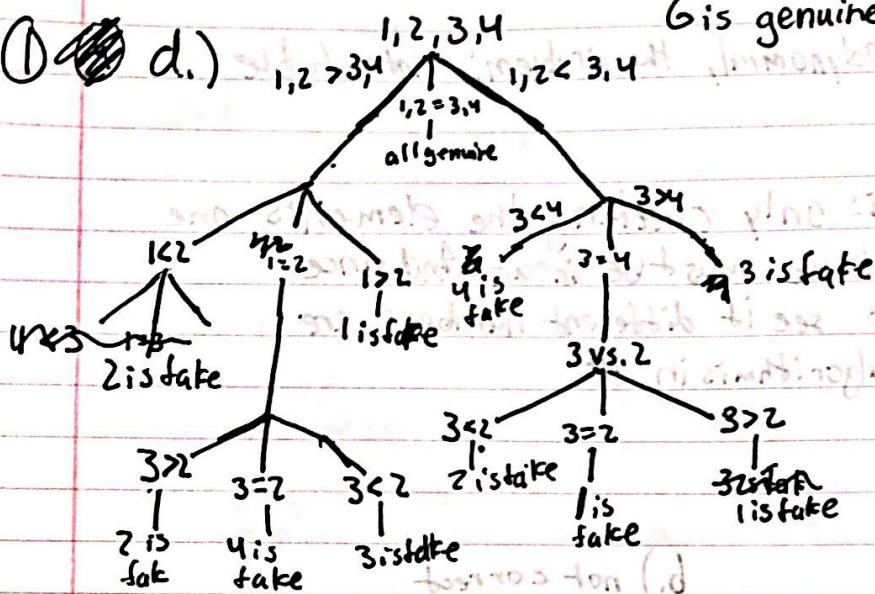


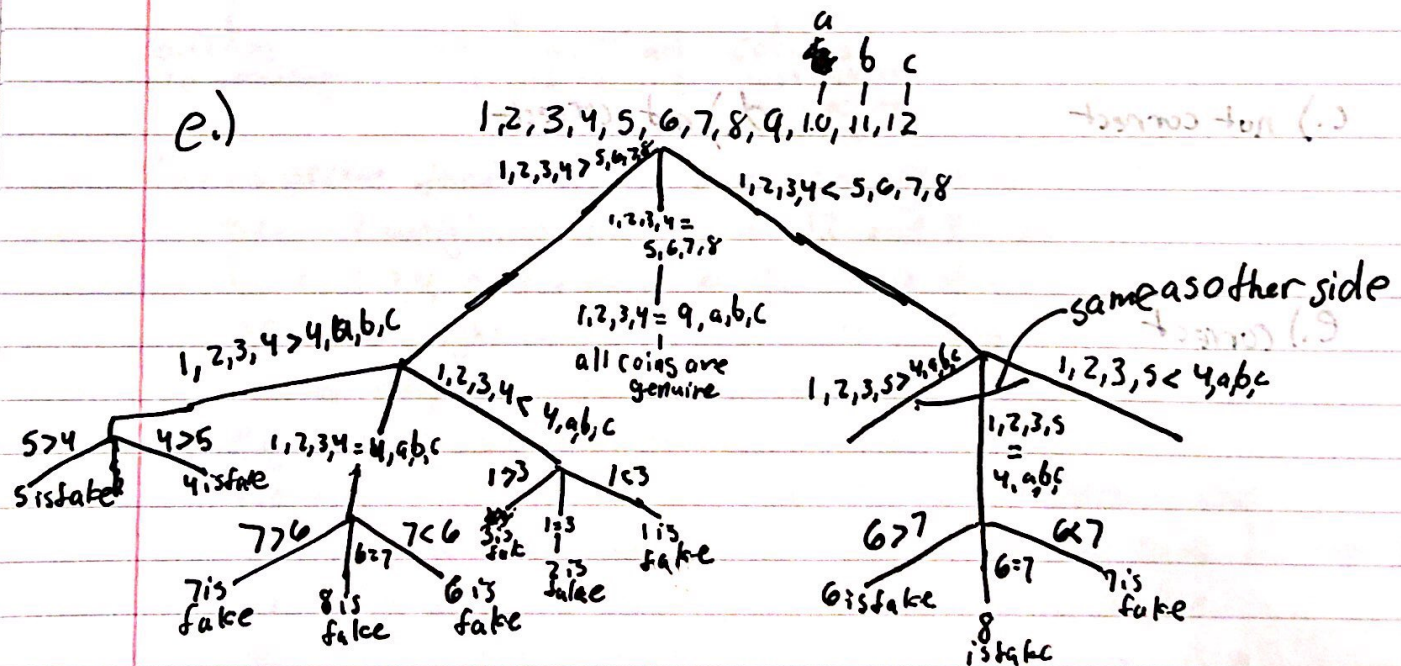
c.) to prove that all the coins are genuine would only take 2 weighings, ex. weigh 1,2 and 3,4 then weigh 1,3 and 2,4 and if these are the weight at 1,2 = weight at 3,4 and the weight of 1,3 = weight 2,4 and we know there can only be 1 fake coin, then all coin weights are equal, therefore they are all genuine.

However, ~~fit~~ to isolate the fake coin if there is one you would need 3 weighings. if you group the coins 1,2 and 3,4 again, the ~~only~~ only info information you would be able to get from the first weighing is that there is a fake coin. ~~After then if you changed the groups around to 1,3 and 2,4, the weighing of these two groups to now get new information you need to weigh 1 vs 2 and 3 vs 4 and if two coins being weighed are equal they are both genuine. if a weighing two coins are not equal, one coin needs to be weighed against one of the genuine coins to see if it is genuine or not, therefore, you need 3 weighings~~

① d.) 6 is genuine coin



e.)



$$\textcircled{2} 2^h \geq (k!)^{n/k}$$

$$h \geq \log((k!)^{n/k}) = \cancel{\log(k!)} \log(k!) \frac{n}{k}$$

$$h \geq \frac{n}{k} \log((k/2)^{k/2})$$

$$h = \frac{n}{k} \left(\frac{k}{2}\right) \log\left(\frac{k}{2}\right)$$

$$h = \frac{n}{2} \log(k) - \frac{n}{2}$$

lower bound: $n \log k$

③ since $n^{\log_2 n}$ is not polynomial, the problem is intractable

④ Since the algorithm is only checking the elements one by one, the running time must be linear. And since the algorithm checks to see if different numbers are divisors, then ~~the~~ the algorithm is in NP.

⑤ a.) not correct

b.) not correct

c.) not correct

d.) not correct

e.) correct