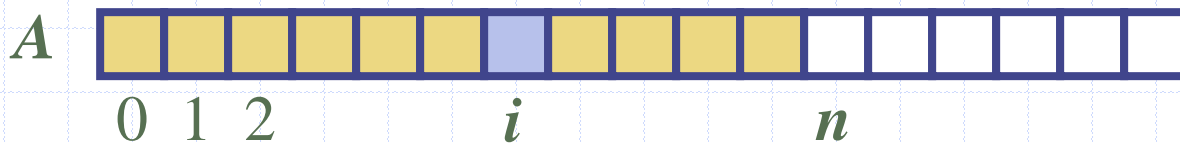


Array-Based Sequences

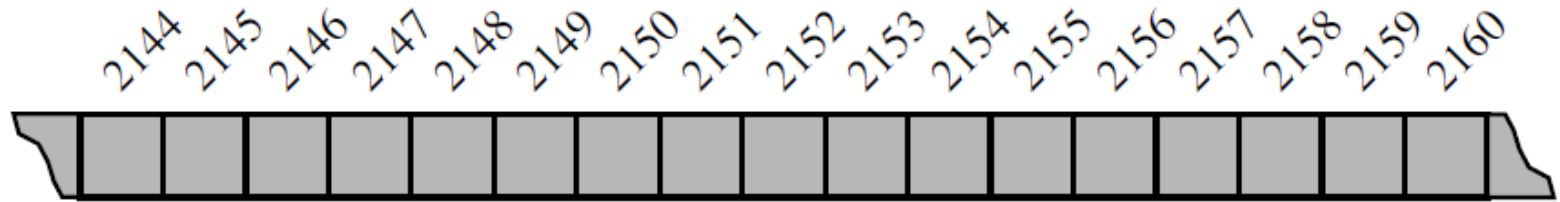


Python Sequence Classes

- Python has built-in types, **list**, **tuple**, and **str**.
- Each of these **sequence** types supports indexing to access an individual element of a sequence, using a syntax such as `A[i]`
- Each of these types uses an **array** to represent the sequence.
 - An array is a set of memory locations that can be addressed using consecutive indices, which, in Python, start with index 0.



Low-Level Arrays



A representation of a portion of a computer's memory, with individual bytes labeled with consecutive memory addresses

Random Access Memory (RAM): Theoretically we need $O(1)$ time to access any location if we already know the Address of the location.

Python String Representation ("SAMPLE")

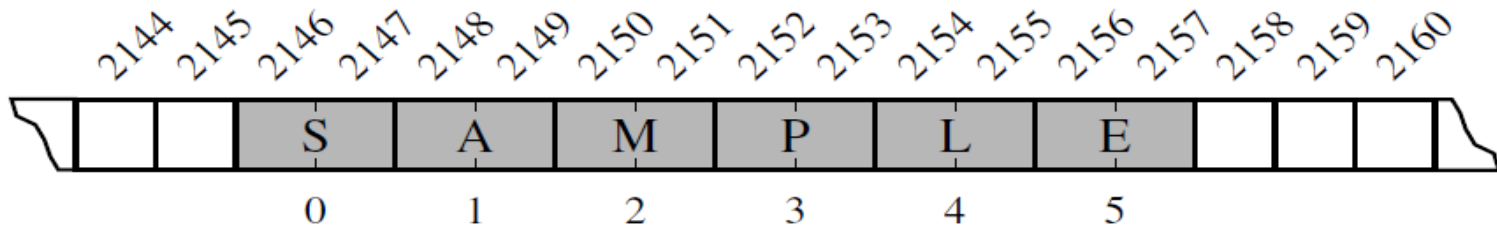


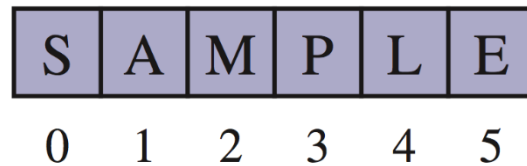
Figure 5.2: A Python string embedded as an array of characters in the computer's memory. We assume that each Unicode character of the string requires two bytes of memory. The numbers below the entries are indices into the string.

- If we know
 - memory address at which array start
 - number of bytes per element
 - index position within the array

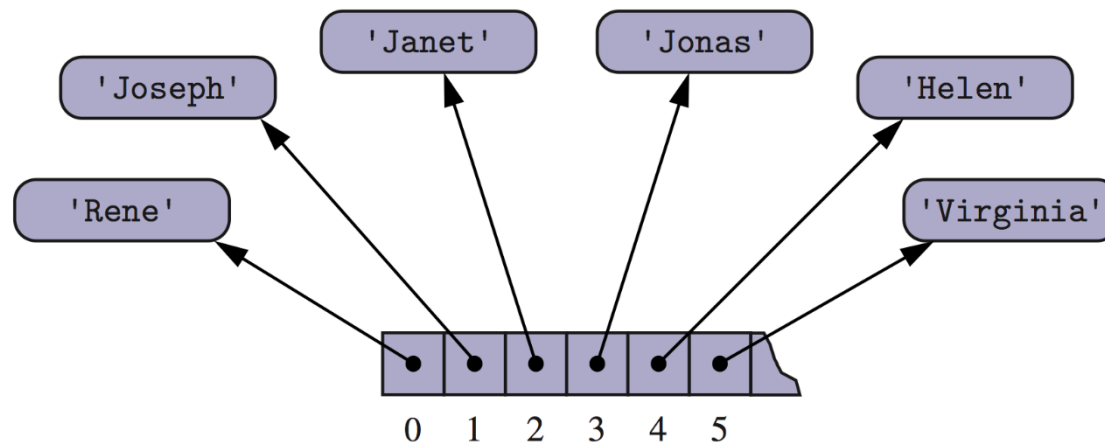
Desired memory address = start + cell_size * index

Arrays of Characters or Object References

- An array can store primitive elements, such as characters, giving us a **compact array**.



- An array can also store references to objects.



List Class Representation in Memory

- Primes=[2,3,5,7,11,13,17,19]
- Temp = Primes[3:6]

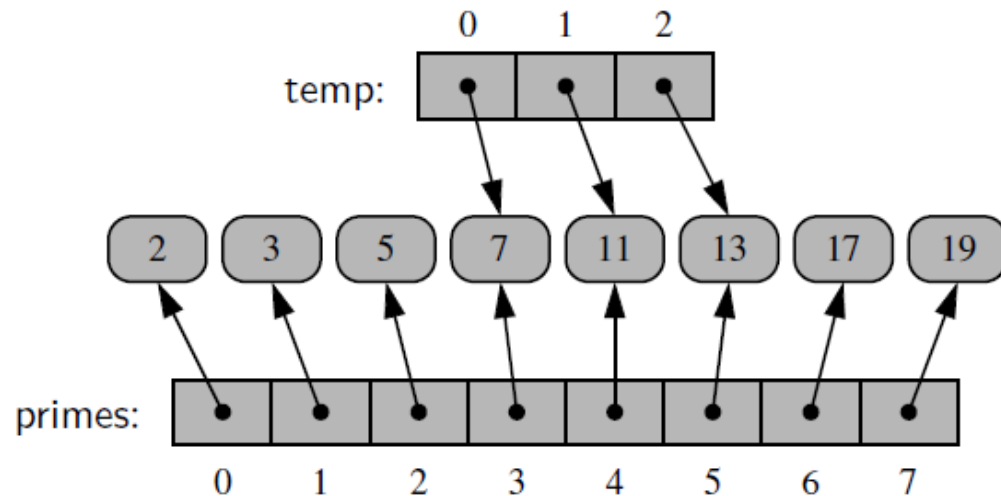
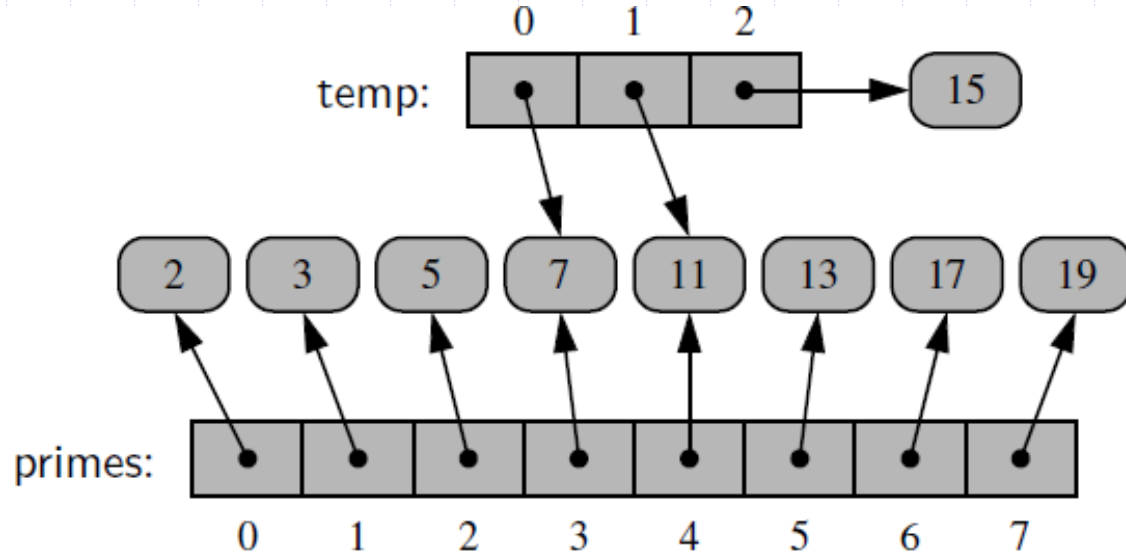


Figure 5.5: The result of the command `temp = primes[3:6]`.

List Class Representation in Memory

- ❑ Primes=[2,3,5,7,11,13,17,19]
- ❑ Temp = Primes[3:6]
- ❑ Temp[2] = 15



counters = [0] * 8

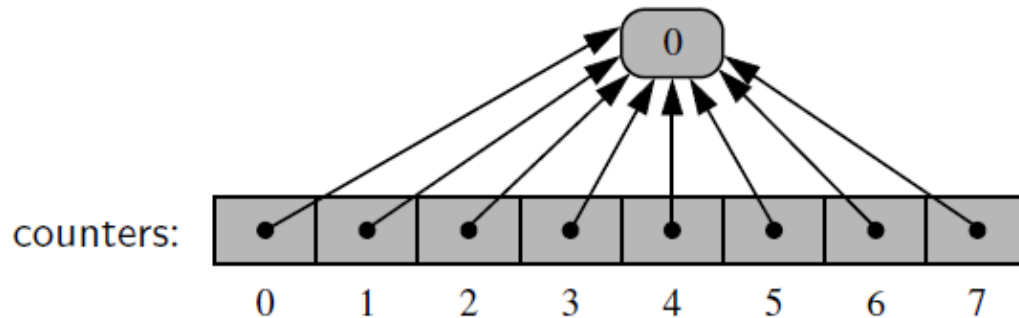
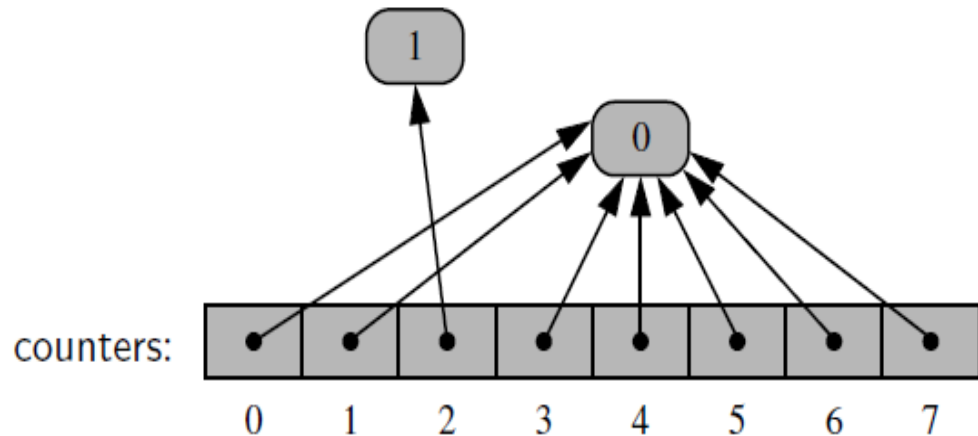


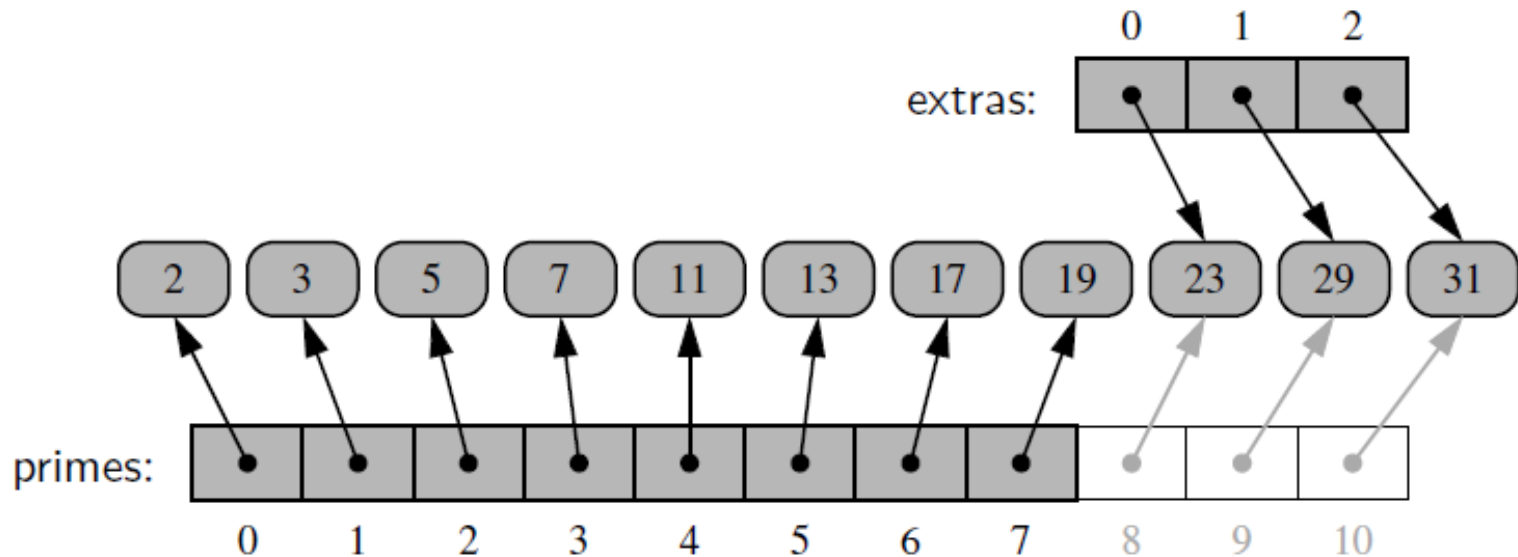
Figure 5.7: The result of the command `data = [0] * 8`.

After executing
`counters[2] += 1`



primes.extend(extras)

- primes=[2,3,5,7,11,13,17,19]
- extras = [23,29,31]



Compact Arrays

- Primary support for compact arrays is in a module named **array**.
 - That module defines a class, also named **array**, providing compact storage for arrays of primitive data types.
- The constructor for the **array** class requires a type code as a first parameter, which is a character that designates the type of data that will be stored in the array.

```
primes = array('i', [2, 3, 5, 7, 11, 13, 17, 19])
```

Type Codes in the array Class

- Python's array class has the following type codes:

| Code | C Data Type | Typical Number of Bytes |
|------|--------------------|-------------------------|
| 'b' | signed char | 1 |
| 'B' | unsigned char | 1 |
| 'u' | Unicode char | 2 or 4 |
| 'h' | signed short int | 2 |
| 'H' | unsigned short int | 2 |
| 'i' | signed int | 2 or 4 |
| 'I' | unsigned int | 2 or 4 |
| 'l' | signed long int | 4 |
| 'L' | unsigned long int | 4 |
| 'f' | float | 4 |
| 'd' | float | 8 |

Dynamic Array and Amortization

- ❑ When creating a low-level array in a computer system, the precise size of that array must be explicitly declared in order for the system to properly allocate a consecutive piece of memory for its storage.
- ❑ Otherwise system might dedicate neighboring memory locations to store other data, the capacity of an array cannot trivially be increased by expanding into subsequent cells.
- ❑ For Immutable classes like tuple or str instance, this is no problem.

Python List Class

- Dynamic Array
 - Underlying array has more capacity than current length of the list.
 - If full, create a new array with larger capacity, and then copy the elements from old array to new array.
 - Old array is no longer needed. System can claim that memory space.

Experimental Analysis (open experiment_list_size.py file)

```
import sys
```

```
try:
```

```
    n = int(sys.argv[1])
```

```
except:
```

```
    n = 100
```

```
data = []
```

```
for k in range(n):
```

```
    a = len(data)
```

```
    b = sys.getsizeof(data)
```

```
    print('Length: {0:3d}; Size in bytes:{1:4d}'.format(a,b))
```

```
    data.append(None)
```

```
# NOTE: must fix choice of n
```

```
# number of elements
```

```
# actual size in bytes
```

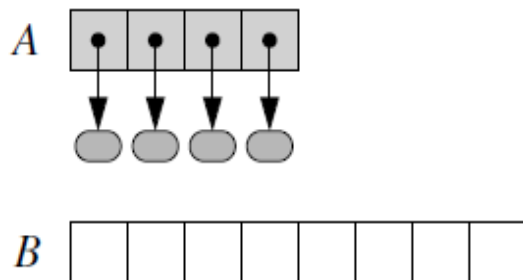
```
# increase length by one
```

Output

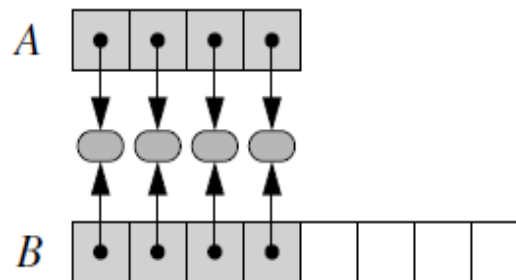
| | |
|--------------------------------|--------------------------------|
| Length: 0; Size in bytes: 72 | Length: 13; Size in bytes: 200 |
| Length: 1; Size in bytes: 104 | Length: 14; Size in bytes: 200 |
| Length: 2; Size in bytes: 104 | Length: 15; Size in bytes: 200 |
| Length: 3; Size in bytes: 104 | Length: 16; Size in bytes: 200 |
| Length: 4; Size in bytes: 104 | Length: 17; Size in bytes: 272 |
| Length: 5; Size in bytes: 136 | Length: 18; Size in bytes: 272 |
| Length: 6; Size in bytes: 136 | Length: 19; Size in bytes: 272 |
| Length: 7; Size in bytes: 136 | Length: 20; Size in bytes: 272 |
| Length: 8; Size in bytes: 136 | Length: 21; Size in bytes: 272 |
| Length: 9; Size in bytes: 200 | Length: 22; Size in bytes: 272 |
| Length: 10; Size in bytes: 200 | Length: 23; Size in bytes: 272 |
| Length: 11; Size in bytes: 200 | Length: 24; Size in bytes: 272 |
| Length: 12; Size in bytes: 200 | Length: 25; Size in bytes: 272 |
| | Length: 26; Size in bytes: 352 |

Implementation of a Dynamic Array

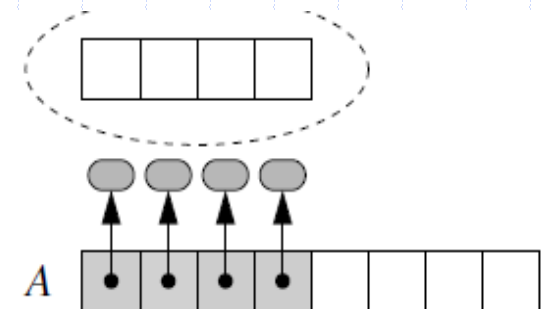
1. Allocate a new array B with larger capacity.
2. Set $B[i] = A[i]$, for $i = 0, \dots, n-1$, where n denotes current number of items.
3. Set $A = B$, that is, we henceforth use B as the array supporting the list.
4. Insert the new element in the new array.



(a)



(b)



(c)

Python Implementation

```
1 import ctypes                                # provides low-level arrays
2
3 class DynamicArray:
4     """A dynamic array class akin to a simplified Python list."""
5
6     def __init__(self):
7         """Create an empty array."""
8         self._n = 0                          # count actual elements
9         self._capacity = 1                   # default array capacity
10        self._A = self._make_array(self._capacity)  # low-level array
11
12    def __len__(self):
13        """Return number of elements stored in the array."""
14        return self._n
15
16    def __getitem__(self, k):
17        """Return element at index k."""
18        if not 0 <= k < self._n:
19            raise IndexError('invalid index')
20        return self._A[k]                    # retrieve from array
21
22    def append(self, obj):
23        """Add object to end of the array."""
24        if self._n == self._capacity:        # not enough room
25            self._resize(2 * self._capacity) # so double capacity
26        self._A[self._n] = obj
27        self._n += 1
28
29    def _resize(self, c):                     # nonpublic utility
30        """Resize internal array to capacity c."""
31        B = self._make_array(c)              # new (bigger) array
32        for k in range(self._n):              # for each existing value
33            B[k] = self._A[k]
34        self._A = B                           # use the bigger array
35        self._capacity = c
36
37    def _make_array(self, c):                 # nonpublic utility
38        """Return new array with capacity c."""
39        return (c * ctypes.py_object)()
```

Growable Array-based Array List

- In an **add(o)** operation (without an index), we could always add at the end
- When the array is full, we replace the array with a larger one
- How large should the new array be?
 - **Incremental strategy**: increase the size by a constant c
 - **Doubling strategy**: double the size

```
Algorithm add(o)  
  if  $t = S.length - 1$  then  
     $A \leftarrow$  new array of  
      size ...  
    for  $i \leftarrow 0$  to  $n-1$  do  
       $A[i] \leftarrow S[i]$   
     $S \leftarrow A$   
     $n \leftarrow n + 1$   
     $S[n-1] \leftarrow o$ 
```

Comparison of the Strategies

- We compare the incremental strategy and the doubling strategy by analyzing the total time $T(n)$ needed to perform a series of n $\text{add}(o)$ operations
- We assume that we start with an empty stack represented by an array of size 1
- We call amortized time of an add operation the average time taken by an add over the series of operations, i.e., $T(n)/n$

Incremental Strategy

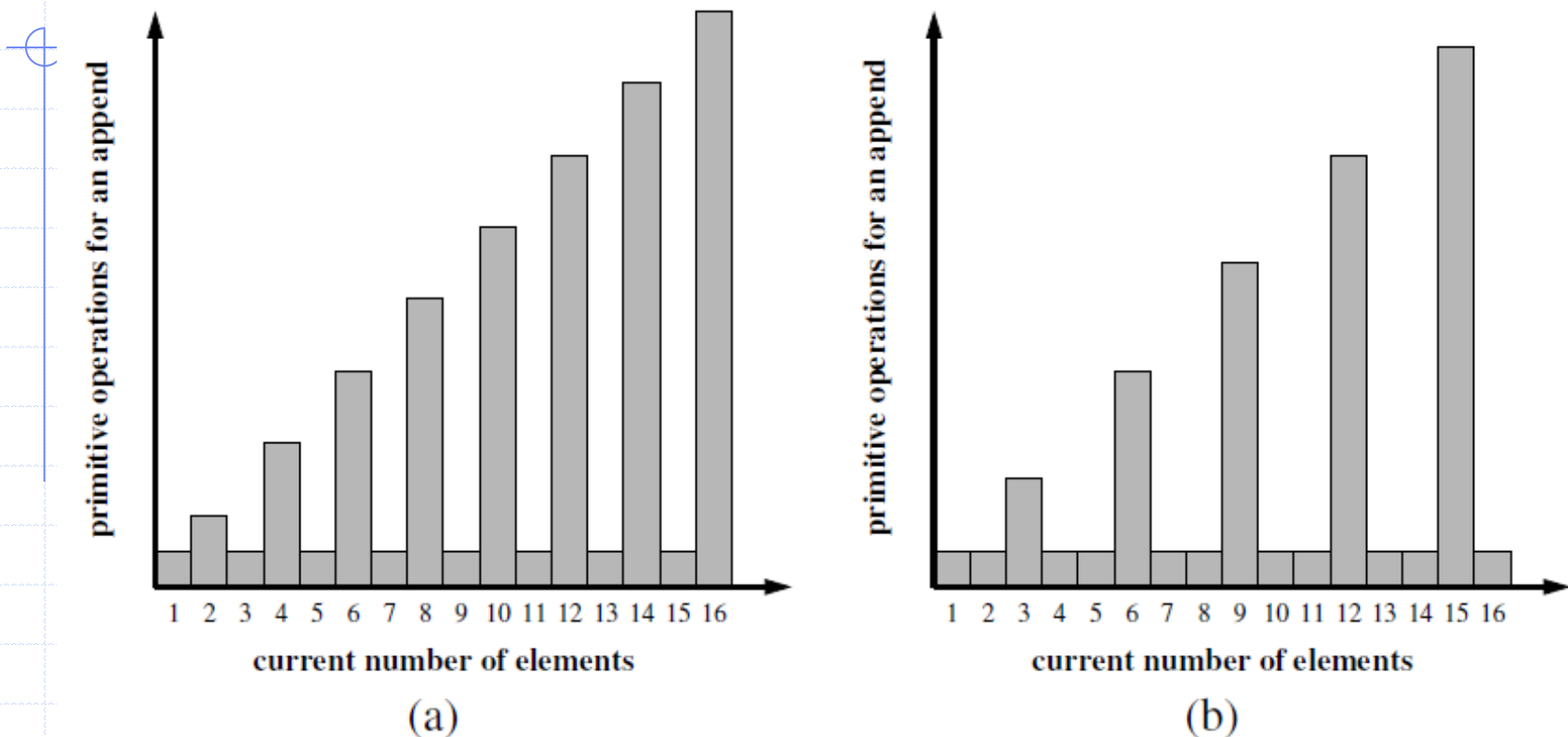


Figure 5.15: Running times of a series of append operations on a dynamic array using arithmetic progression of sizes. (a) Assumes increase of 2 in size of the array, while (b) assumes increase of 3.

Incremental Strategy Analysis

- We replace the array $k = n/c$ times
- The total time $T(n)$ of a series of n add operations is proportional to

$$n + c + 2c + 3c + 4c + \dots + kc =$$

$$n + c(1 + 2 + 3 + \dots + k) =$$

$$n + ck(k + 1)/2$$

- Since c is a constant, $T(n)$ is $O(n + k^2)$, i.e., $O(n^2)$
- The amortized time of an add operation is $O(n)$

Always Doubling Strategy

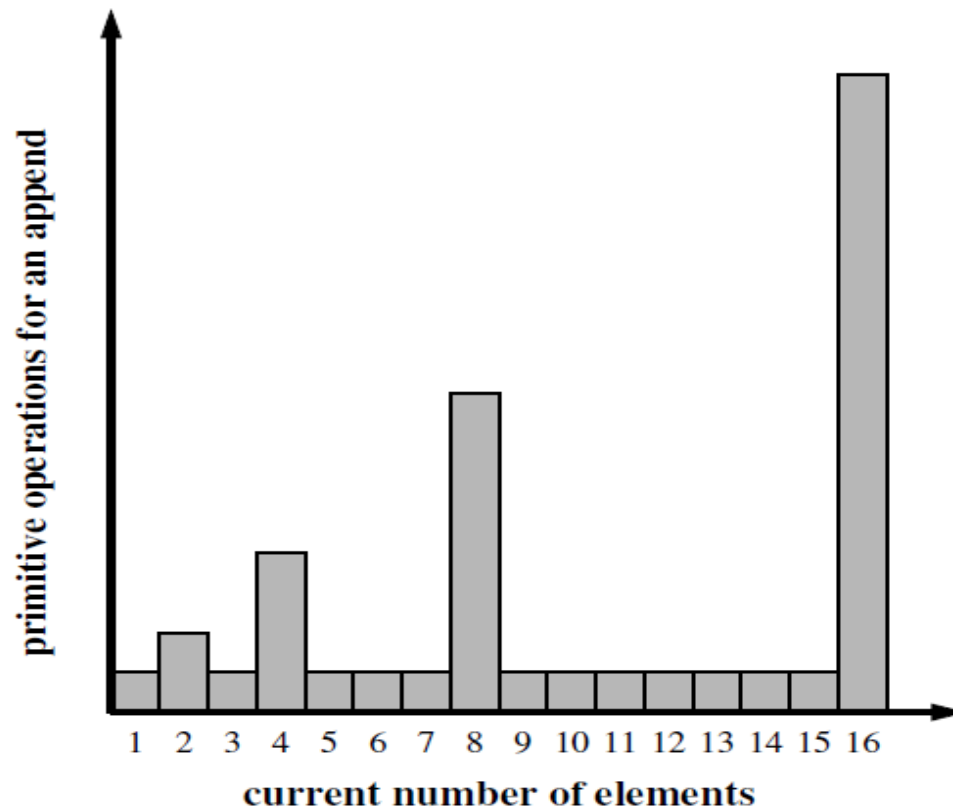


Figure 5.13: Running times of a series of append operations on a dynamic array.

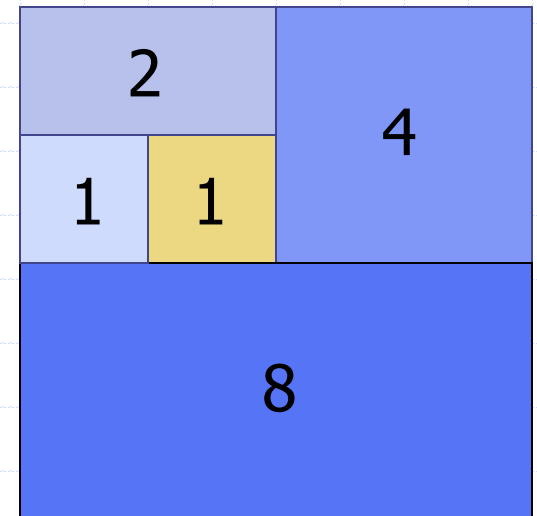
Doubling Strategy Analysis

- We replace the array $k = \log_2 n$ times
- The total time $T(n)$ of a series of n add operations is proportional to

$$\begin{aligned} n + 1 + 2 + 4 + 8 + \dots + 2^k &= \\ n + 2^{k+1} - 1 &= \\ 3n - 1 \end{aligned}$$

- $T(n)$ is $O(n)$
- The amortized time of an add operation is $O(1)$

geometric series



Asymptotic Performance of Nonmutating Behaviors

| Operation | Running Time |
|---|----------------|
| <code>len(data)</code> | $O(1)$ |
| <code>data[j]</code> | $O(1)$ |
| <code>data.count(value)</code> | $O(n)$ |
| <code>data.index(value)</code> | $O(k + 1)$ |
| <code>value in data</code> | $O(k + 1)$ |
| <code>data1 == data2</code> (similarly <code>!=</code> , <code><</code> , <code><=</code> , <code>></code> , <code>>=</code>) | $O(k + 1)$ |
| <code>data[j:k]</code> | $O(k - j + 1)$ |
| <code>data1 + data2</code> | $O(n_1 + n_2)$ |
| <code>c * data</code> | $O(cn)$ |

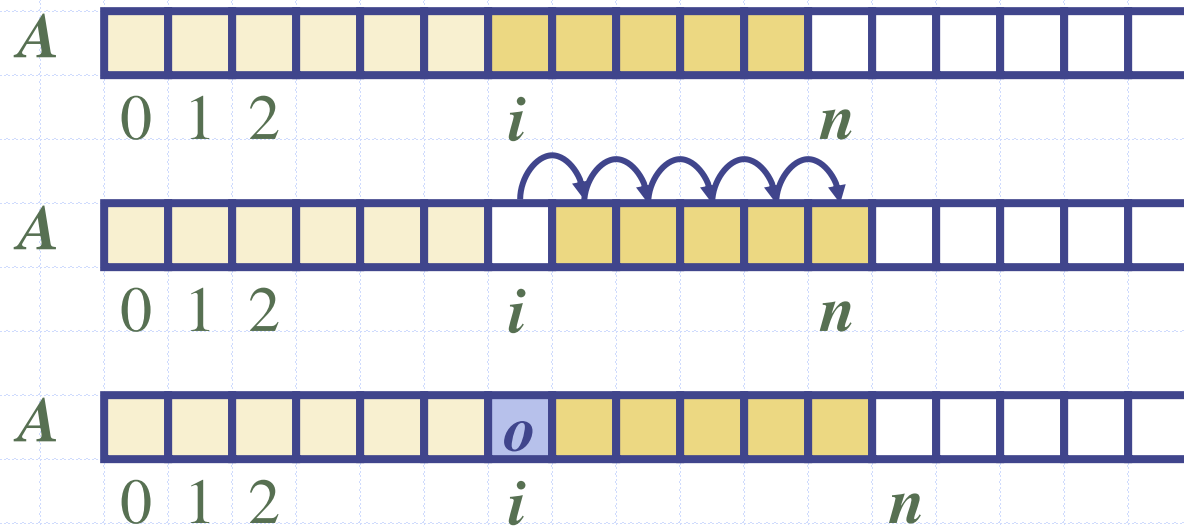
Asymptotic Performance of mutating behavior

| Operation | Running Time |
|---|------------------|
| <code>data[j] = val</code> | $O(1)$ |
| <code>data.append(value)</code> | $O(1)^*$ |
| <code>data.insert(k, value)</code> | $O(n - k + 1)^*$ |
| <code>data.pop()</code> | $O(1)^*$ |
| <code>data.pop(k)</code> <code>del data[k]</code> | $O(n - k)^*$ |
| <code>data.remove(value)</code> | $O(n)^*$ |
| <code>data1.extend(data2)</code> <code>data1 += data2</code> | $O(n_2)^*$ |
| <code>data.reverse()</code> | $O(n)$ |
| <code>data.sort()</code> | $O(n \log n)$ |

*amortized

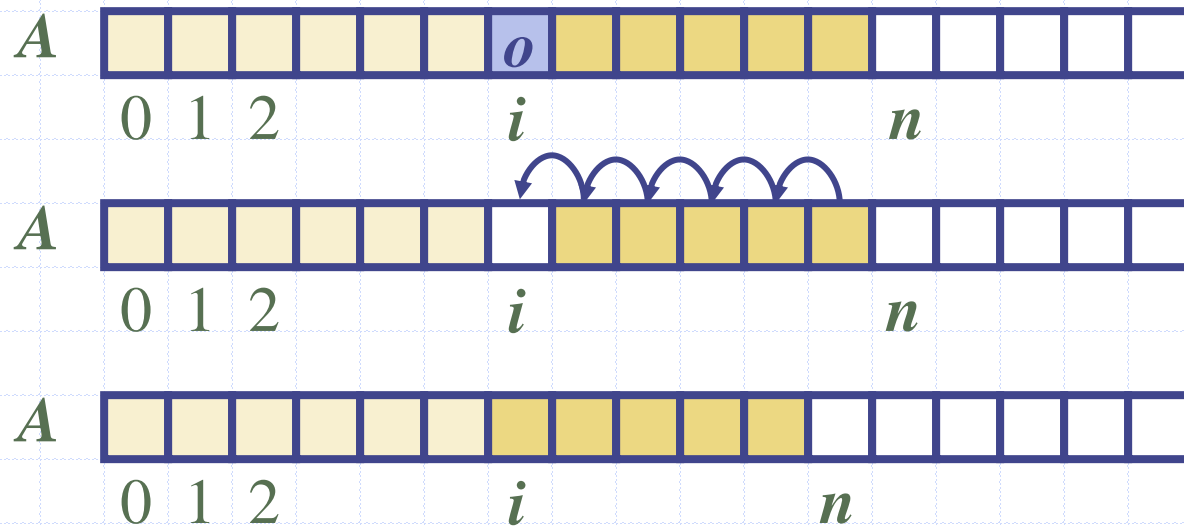
Insertion

- In an operation *add*(i, o), we need to make room for the new element by shifting forward the $n - i$ elements $A[i], \dots, A[n - 1]$
- In the worst case ($i = 0$), this takes $O(n)$ time



Element Removal

- In an operation *remove*(i), we need to fill the hole left by the removed element by shifting backward the $n - i - 1$ elements $A[i + 1], \dots, A[n - 1]$
- In the worst case ($i = 0$), this takes $O(n)$ time



Performance

- In an array based implementation of a dynamic list:
 - The space used by the data structure is $O(n)$
 - Indexing the element at I takes $O(1)$ time
 - *add* and *remove* run in $O(n)$ time in worst case
- In an *add* operation, when the array is full, instead of throwing an exception, we can replace the array with a larger one...

Creating a Multidimensional list

| | | | | |
|----|-----|-----|-----|-----|
| 22 | 18 | 709 | 5 | 33 |
| 45 | 32 | 830 | 120 | 750 |
| 4 | 880 | 45 | 66 | 61 |

```
data = [ [22, 18, 709, 5, 33], [45, 32, 830, 120, 750], [4, 880, 45, 66, 61] ]
```

Data[1][3] represents the value at row number 1 and Column number 3.

Creating a Two Dimensional List

- ❑ Two dimensional list of integers with r rows and c columns with initialize all values to 0.

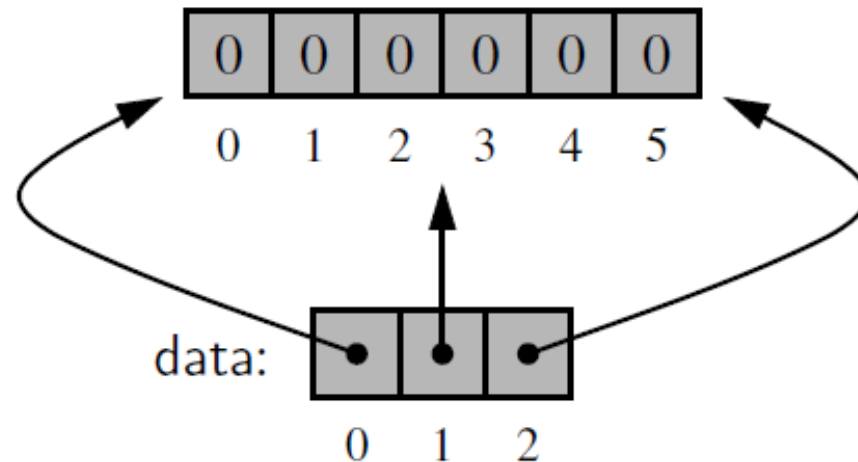
- ❑ $\text{data} = ([0] * c) * r$

Flawed Approach. Why???

- ❑ $\text{data} = [[0]*c]*r$

Flawed Approach too. Why???

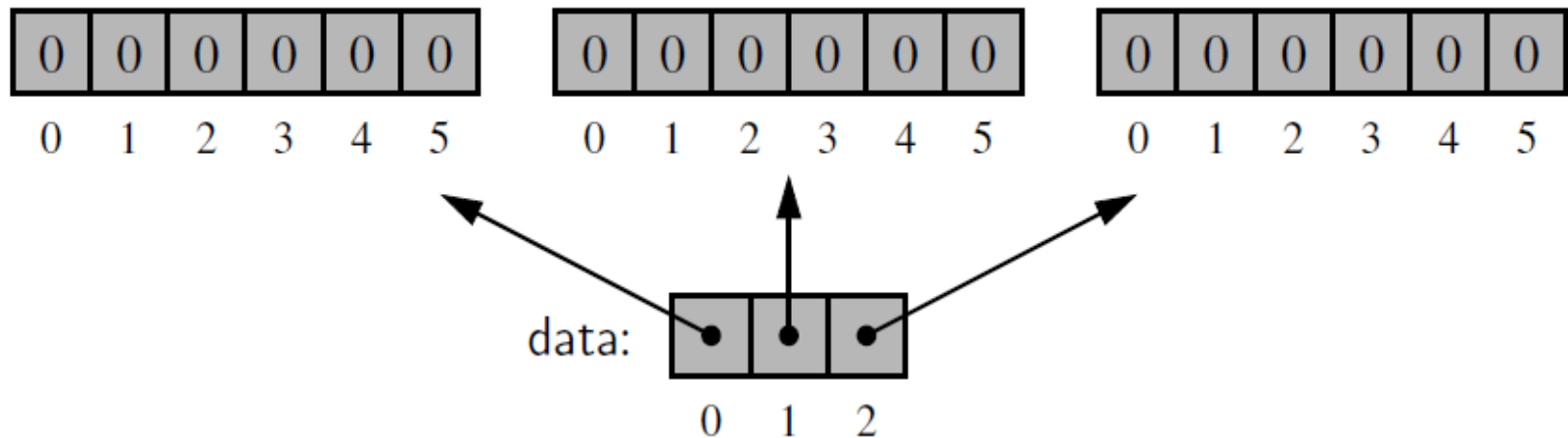
Creating a Two Dimensional List



A flawed representation of a 3×6 data set as a list of lists, created with the command `data = [[0]*6]*3`. (For simplicity, we overlook the fact that the values in the secondary list are referential.)

Creating a Two Dimensional List

```
data = [ [0]*c for j in range(r) ]
```



A valid representation of a 3×6 data set as a list of lists.

Coding Practice

- ❑ Creating our own dynamic array.
- ❑ Open the `UserDefinedDynamicArray_for_Students.py` file.
- ❑ Some functions are already implemented.
 - `Append()`, `getitem`, `delitem`, `len`, `str`, `resize`, `make_array`

Practice UserDefinedDynamic Array Class.

□ Task1:

Implement `is_empty(self)`

Implement `__iter__(self)`

□ Task2:

Implement `__setitem__(self, I, x)`

Practice UserDefinedDynamic Array Class.

- **Task3:**

Implement `extend(self,I)`

- **Task4:**

Implement `reverse(self)`

Practice UserDefinedDynamic Array Class.

□ Task5:

Implement `__contains__(self,x)`

Implement `index(self,x)`

Implement `count(self,x)`

□ Task6:

Implement `__add__(self,other)`

Implement `__mul__(self,times)`

Practice UserDefinedDynamic Array Class.

□ Task7:

Implement `__pop__(self,i=-1)`

Implement `remove(self,x)`

Using Array Based Sequences

- ❑ Storing High Scores of a game
- ❑ Simple Cryptography
- ❑ Multidimensional Array
- ❑ Tic Tac Toe game