

Priority Queues



Priority Queue ADT

- A priority queue stores a collection of items
- Each **item** is a pair (key, value)
- Main methods of the Priority Queue ADT
 - **add** (k, x)
inserts an item with key k and value x
 - **remove_min()**
removes and returns the item with smallest key
- Additional methods
 - **min()**
returns, but does not remove, an item with smallest key
 - **len(P), is_empty()**
- Applications:
 - Standby flyers
 - Auctions
 - Stock market

Priority Queue Example

Operation	Return Value	Priority Queue
P.add(5,A)		{(5,A)}
P.add(9,C)		{(5,A), (9,C)}
P.add(3,B)		{(3,B), (5,A), (9,C)}
P.add(7,D)		{(3,B), (5,A), (7,D), (9,C)}
P.min()	(3,B)	{(3,B), (5,A), (7,D), (9,C)}
P.remove_min()	(3,B)	{(5,A), (7,D), (9,C)}
P.remove_min()	(5,A)	{(7,D), (9,C)}
len(P)	2	{(7,D), (9,C)}
P.remove_min()	(7,D)	{(9,C)}
P.remove_min()	(9,C)	{ }
P.is_empty()	True	{ }
P.remove_min()	"error"	{ }

Total Order Relations

- Keys in a priority queue can be arbitrary objects on which an order is defined
- Two distinct entries in a priority queue can have the same key
- Mathematical concept of total order relation \leq
 - Reflexive property:
 $x \leq x$
 - Antisymmetric property:
 $x \leq y \wedge y \leq x \Rightarrow x = y$
 - Transitive property:
 $x \leq y \wedge y \leq z \Rightarrow x \leq z$

Composition Design Pattern

- ❑ An **item** in a priority queue is simply a key-value pair
- ❑ Priority queues store items to allow for efficient insertion and removal based on keys

```
1 class PriorityQueueBase:
2     """ Abstract base class for a priority queue. """
3
4     class _Item:
5         """ Lightweight composite to store priority queue items. """
6         __slots__ = '_key', '_value'
7
8         def __init__(self, k, v):
9             self._key = k
10            self._value = v
11
12            def __lt__(self, other):
13                return self._key < other._key    # compare items based on their keys
14
15            def is_empty(self):                    # concrete method assuming abstract len
16                """ Return True if the priority queue is empty. """
17                return len(self) == 0
```

Sequence-based Priority Queue

- Implementation with an unsorted list



- Performance:
 - **add** takes $O(1)$ time since we can insert the item at the beginning or end of the sequence
 - **Remove_min** and **min** take $O(n)$ time since we have to traverse the entire sequence to find the smallest key

- Implementation with a sorted list



- Performance:
 - **add** takes $O(n)$ time since we have to find the place where to insert the item
 - **remove_min** and **min** take $O(1)$ time, since the smallest key is at the beginning

Unsorted List Implementation

```
1 class UnsortedPriorityQueue(PriorityQueueBase): # base class defines _Item
2     """ A min-oriented priority queue implemented with an unsorted list. """
3
4     def _find_min(self):          # nonpublic utility
5         """ Return Position of item with minimum key. """
6         if self.is_empty():      # is_empty inherited from base class
7             raise Empty('Priority queue is empty')
8         small = self._data.first()
9         walk = self._data.after(small)
10        while walk is not None:
11            if walk.element() < small.element():
12                small = walk
13            walk = self._data.after(walk)
14        return small
15
16    def __init__(self):
17        """ Create a new empty Priority Queue. """
18        self._data = PositionalList()
19
20    def __len__(self):
21        """ Return the number of items in the priority queue. """
22        return len(self._data)
```

```
23
24    def add(self, key, value):
25        """ Add a key-value pair. """
26        self._data.add_last(self._Item(key, value))
27
28    def min(self):
29        """ Return but do not remove (k,v) tuple with minimum key. """
30        p = self._find_min()
31        item = p.element()
32        return (item._key, item._value)
33
34    def remove_min(self):
35        """ Remove and return (k,v) tuple with minimum key. """
36        p = self._find_min()
37        item = self._data.delete(p)
38        return (item._key, item._value)
```

Sorted List Implementation

```
1 class SortedPriorityQueue(PriorityQueueBase): # base class defines _Item
2     """ A min-oriented priority queue implemented with a sorted list. """
3
4     def __init__(self):
5         """ Create a new empty Priority Queue. """
6         self._data = PositionalList()
7
8     def __len__(self):
9         """ Return the number of items in the priority queue. """
10        return len(self._data)
11
12    def add(self, key, value):
13        """ Add a key-value pair. """
14        newest = self._Item(key, value) # make new item instance
15        walk = self._data.last() # walk backward looking for smaller key
16        while walk is not None and newest < walk.element():
17            walk = self._data.before(walk)
18        if walk is None:
19            self._data.add_first(newest) # new key is smallest
20        else:
21            self._data.add_after(walk, newest) # newest goes after walk
22
```

```
23    def min(self):
24        """ Return but do not remove (k,v) tuple with minimum key. """
25        if self.is_empty():
26            raise Empty('Priority queue is empty.')
27        p = self._data.first()
28        item = p.element()
29        return (item._key, item._value)
30
31    def remove_min(self):
32        """ Remove and return (k,v) tuple with minimum key. """
33        if self.is_empty():
34            raise Empty('Priority queue is empty.')
35        item = self._data.delete(self._data.first())
36        return (item._key, item._value)
```


Runtime of implementing PQ using Sorted and Unsorted List

Operation	Unsorted List	Sorted List
len	$O(1)$	$O(1)$
is_empty	$O(1)$	$O(1)$
add	$O(1)$	$O(n)$
min	$O(n)$	$O(1)$
remove_min	$O(n)$	$O(1)$

Priority Queue Sorting

- We can use a priority queue to sort a set of comparable elements
 1. Insert the elements one by one with a series of **add** operations
 2. Remove the elements in sorted order with a series of **remove_min** operations
- The running time of this sorting method depends on the priority queue implementation

Algorithm *PQ-Sort*(S, C)

Input sequence S , comparator C for the elements of S

Output sequence S sorted in increasing order according to C

$P \leftarrow$ priority queue with comparator C

while $\neg S.is_empty()$

$e \leftarrow S.remove_first()$

$P.add(e, \emptyset)$

while $\neg P.is_empty()$

$e \leftarrow P.removeMin().key()$

$S.add_last(e)$

Selection-Sort

- Selection-sort is the variation of PQ-sort where the priority queue is implemented with an unsorted sequence
- Running time of Selection-sort:
 1. Inserting the elements into the priority queue with n **insert** operations takes $O(n)$ time
 2. Removing the elements in sorted order from the priority queue with n **removeMin** operations takes time proportional to

$$1 + 2 + \dots + n$$

- Selection-sort runs in $O(n^2)$ time

Selection-Sort Example

	Sequence S	Priority Queue P
Input:	(7,4,8,2,5,3,9)	()
Phase 1		
(a)	(4,8,2,5,3,9)	(7)
(b)	(8,2,5,3,9)	(7,4)
..
(g)	()	(7,4,8,2,5,3,9)
Phase 2		
(a)	(2)	(7,4,8,5,3,9)
(b)	(2,3)	(7,4,8,5,9)
(c)	(2,3,4)	(7,8,5,9)
(d)	(2,3,4,5)	(7,8,9)
(e)	(2,3,4,5,7)	(8,9)
(f)	(2,3,4,5,7,8)	(9)
(g)	(2,3,4,5,7,8,9)	()

Insertion-Sort

- Insertion-sort is the variation of PQ-sort where the priority queue is implemented with a sorted sequence
- Running time of Insertion-sort:
 1. Inserting the elements into the priority queue with n **insert** operations takes time proportional to
$$1 + 2 + \dots + n$$
 2. Removing the elements in sorted order from the priority queue with a series of n **removeMin** operations takes $O(n)$ time
- Insertion-sort runs in $O(n^2)$ time

Insertion-Sort Example

Input:

Sequence S
(7,4,8,2,5,3,9)

Priority queue P
()

Phase 1

(a)	(4,8,2,5,3,9)	(7)
(b)	(8,2,5,3,9)	(4,7)
(c)	(2,5,3,9)	(4,7,8)
(d)	(5,3,9)	(2,4,7,8)
(e)	(3,9)	(2,4,5,7,8)
(f)	(9)	(2,3,4,5,7,8)
(g)	()	(2,3,4,5,7,8,9)

Phase 2

(a)	(2)	(3,4,5,7,8,9)
(b)	(2,3)	(4,5,7,8,9)
..
(g)	(2,3,4,5,7,8,9)	()

In-place Insertion-Sort

- Instead of using an external data structure, we can implement selection-sort and insertion-sort in-place
- A portion of the input sequence itself serves as the priority queue
- For in-place insertion-sort
 - We keep sorted the initial portion of the sequence
 - We can use **swaps** instead of modifying the sequence

