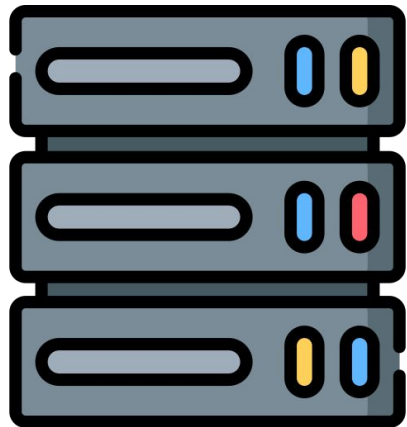
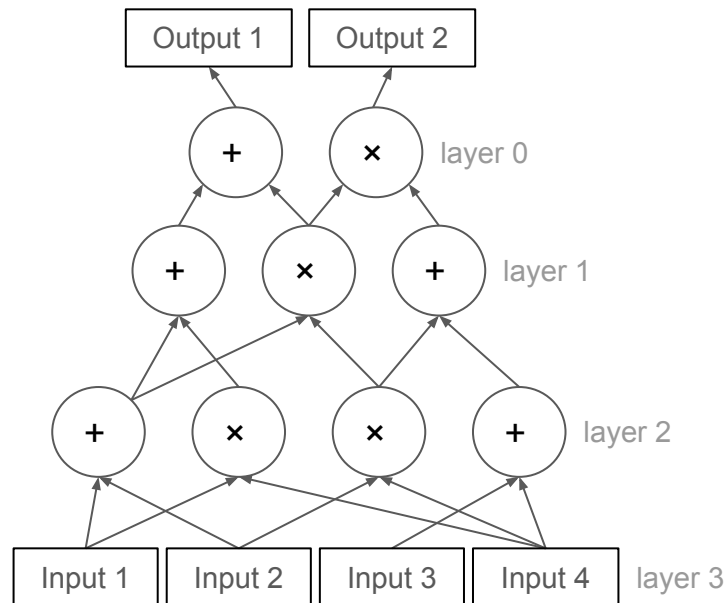


**Prover and Verifier agrees on a  
arithmetic circuit, meaning  
both parties are aware of the  
circuit**

Prover  $P$



Verifier  $V$

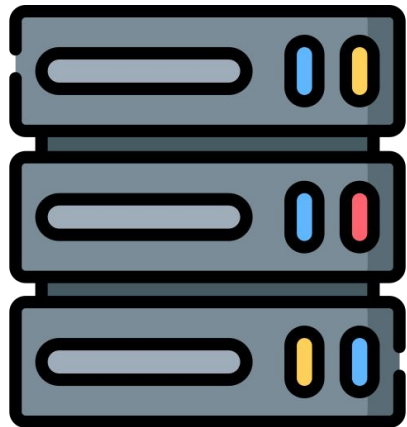


**example circuit**

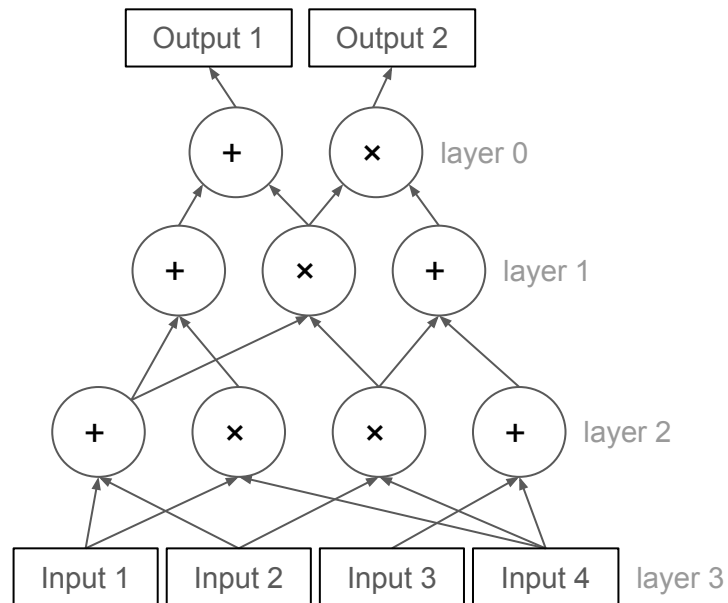
Also all relation function for each layer are  
**given**

$$\{\tilde{add}_i(g, a, b), \tilde{mult}_i(g, a, b) \mid \forall i\}$$

Prover  $P$

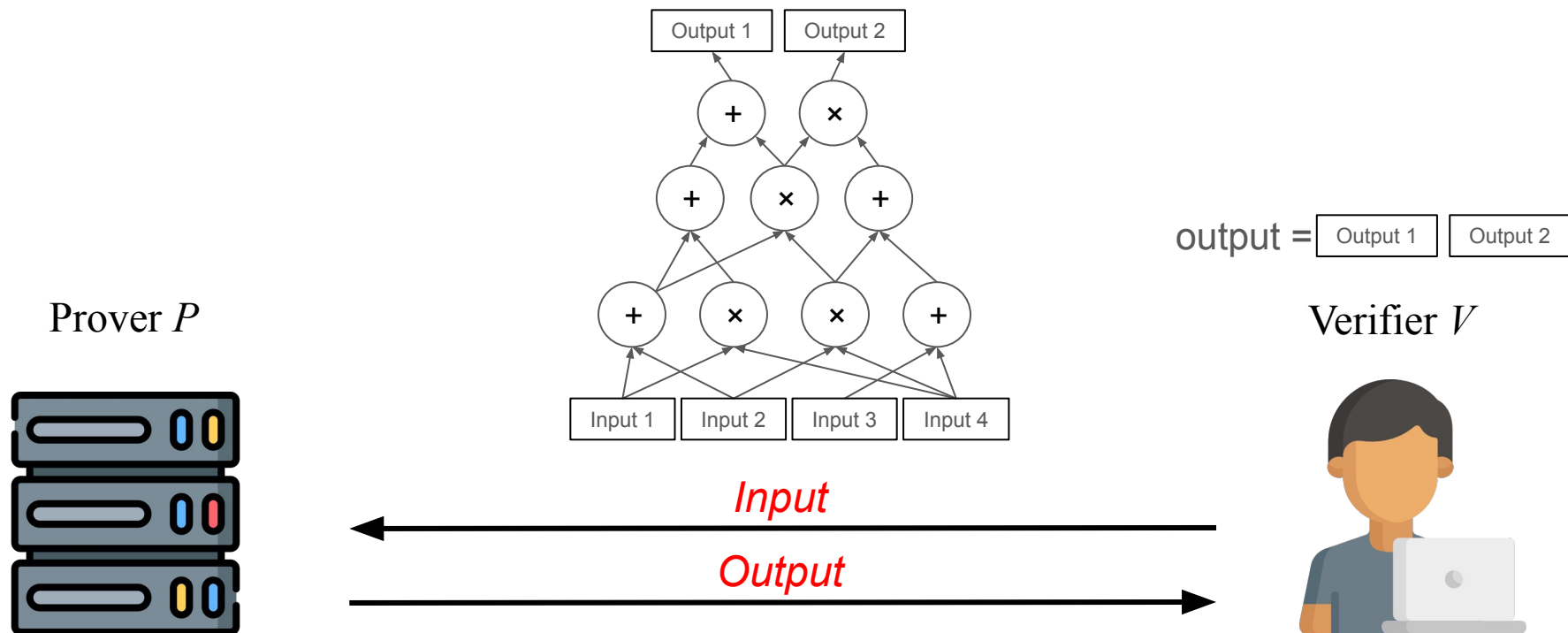


Verifier  $V$



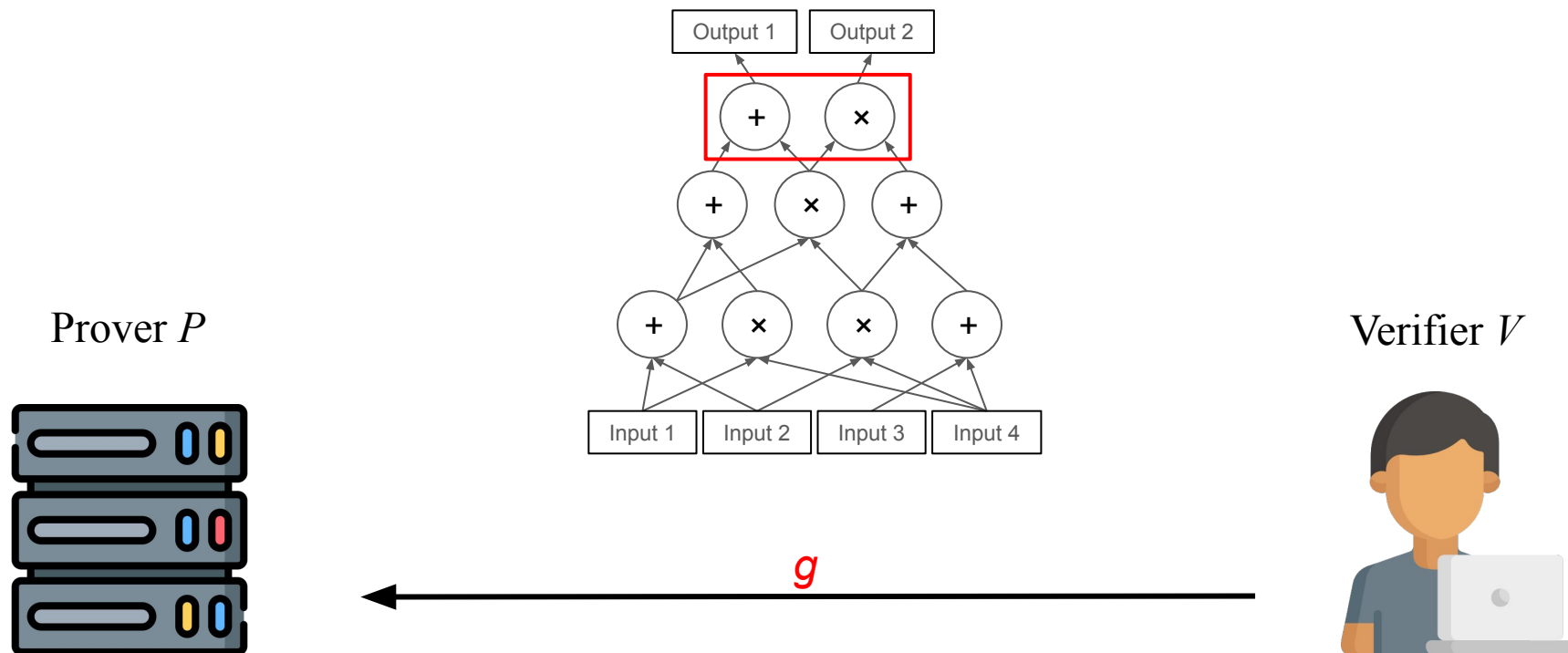
**example circuit**

# GKR Protocol



Verifier sends input and receives output

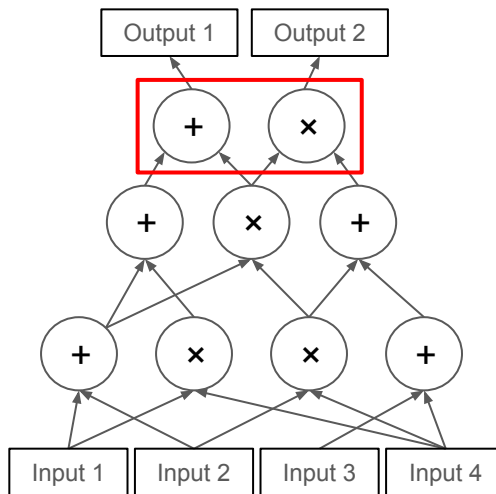
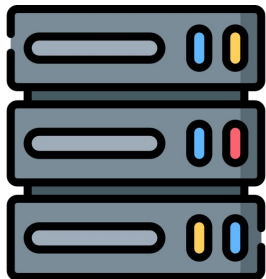
# GKR Protocol



Verifier Picks a random label from the first layer

# GKR Protocol

Prover  $P$



Verifier  $V$

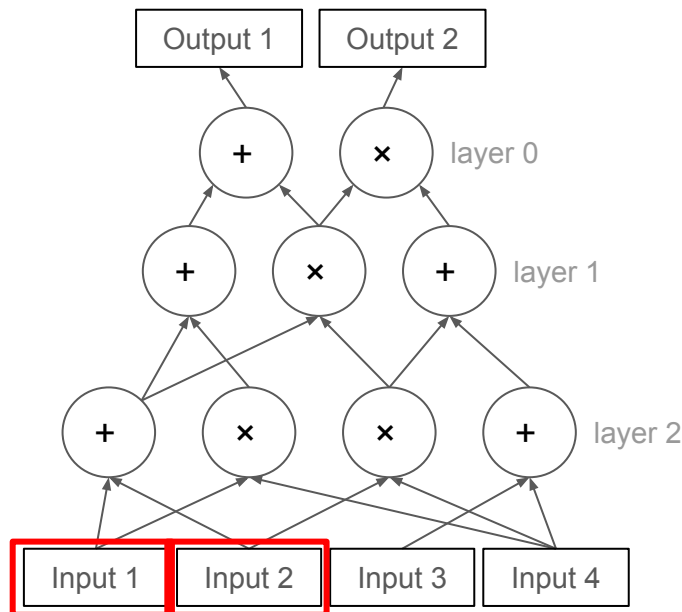


Both parties compute  $V_0(g)$  ( the  $g$ -th gate output)  
Verifier validates  $V_0(g)$  using sumcheck

# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$



# GKR Protocol (Single Layer Sumcheck)

**Sumcheck is performed on the following multivariate polynomial:**

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$

**Lagrange interpolation:**

$$\chi_h(x) = \prod_{h' \in H \setminus \{h\}} \frac{h' - x}{h' - h} \quad \left\{ \begin{array}{l} 1, \ h = x \\ 0, \ h \neq x \end{array} \right. \quad \begin{array}{l} \text{(under binary} \\ \text{inputs)} \end{array}$$

$$\chi_{h_1, \dots, h_m}(x_1, \dots, x_m) = \prod_{i=1}^m \chi_{h_i}(x_i) \quad \left\{ \begin{array}{l} 1, \ h_1 \dots h_m = x_1 \dots x_m \\ 0, \ h_1 \dots h_m \neq x_1 \dots x_m \end{array} \right. \quad \begin{array}{l} \text{(under binary} \\ \text{inputs)} \end{array}$$

# GKR Protocol (Single Layer Sumcheck)

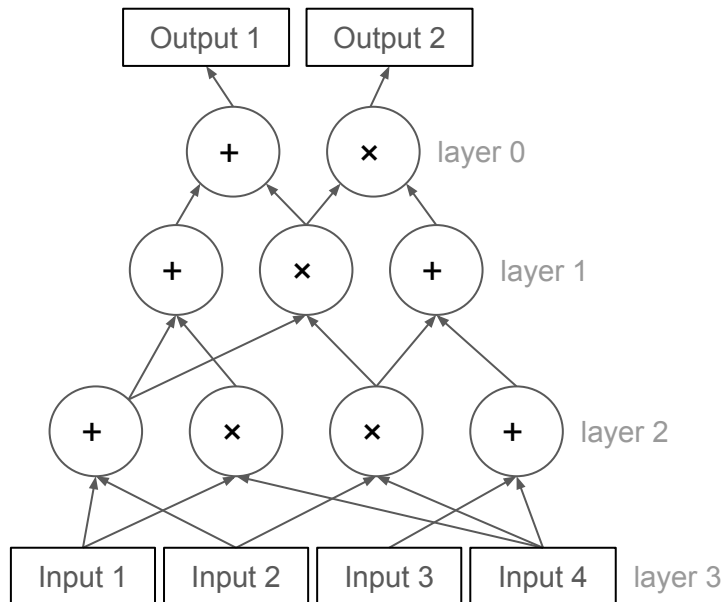
Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + \tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$

Define two functions:

$in_{i,1}(g)$  = first fan in neighbor for the  $i$ th layer

$in_{i,2}(g)$  = second fan in neighbor for the  $i$ th layer





# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( \tilde{add}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + \tilde{mult}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)) \right)$$

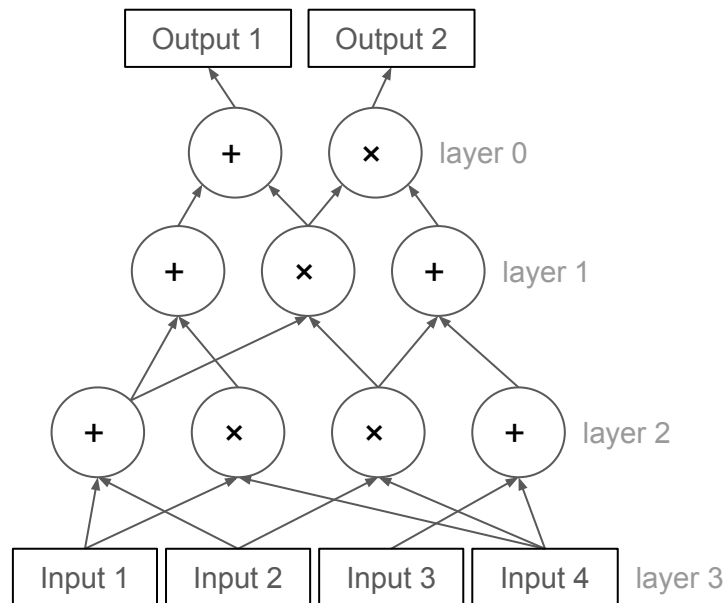
$$\tilde{add}_{i+1}(g, a, b) = \sum_{z \in \{0,1\}^{s_i}} \chi_{z, in_{i,1}(z), in_{i,2}(z)}(g, a, b) \begin{cases} 1, & \text{if } a, b == in_{i,1}(g), in_{i,2}(g) \\ 0, & \text{else} \end{cases}$$

(under binary inputs)

$$\tilde{mult}_{i+1}(g, a, b) = \sum_{z \in \{0,1\}^{s_i}} \chi_{z, in_{i,1}(z), in_{i,2}(z)}(g, a, b) \begin{cases} 1, & \text{if } a, b == in_{i,1}(g), in_{i,2}(g) \\ 0, & \text{else} \end{cases}$$

(under binary inputs)

**function for each layer is given in the beginning  
computed by verifier in the final round**



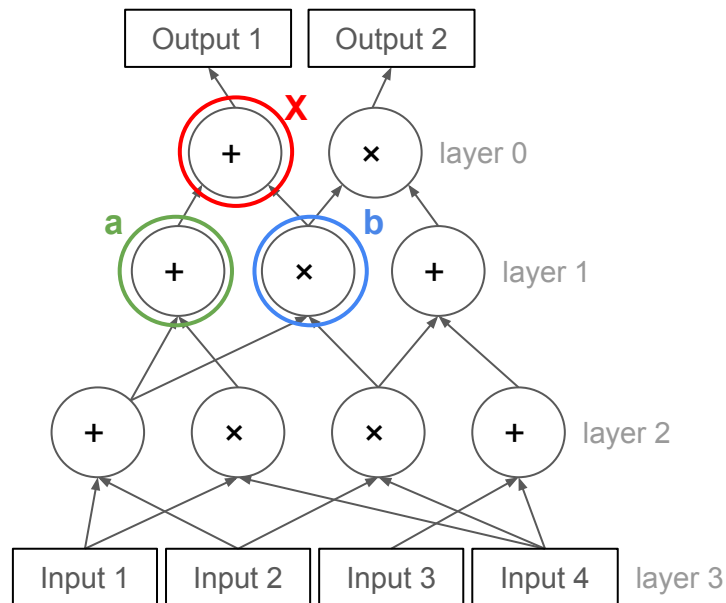
# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( \tilde{add}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + \tilde{mult}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)) \right)$$

$$\tilde{add}_{i+1}(X, a, b) = 1$$

$$\tilde{mult}_{i+1}(X, a, b) = 0$$



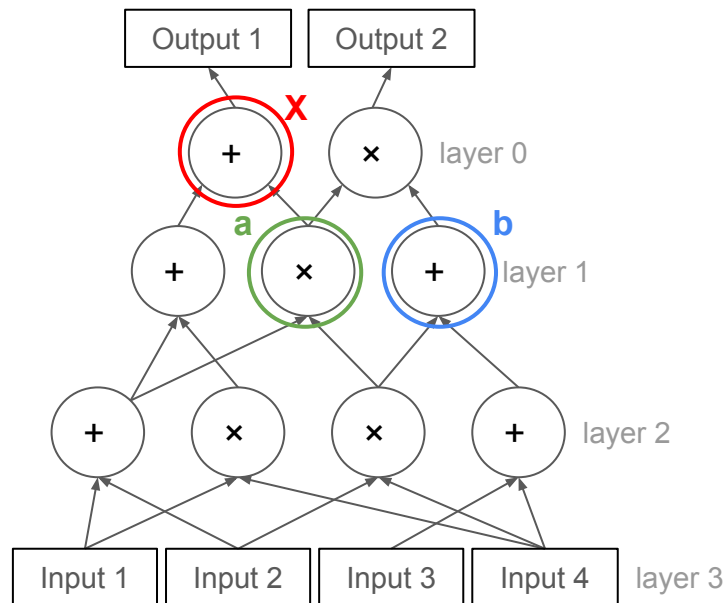
# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( \tilde{add}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + \tilde{mult}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)) \right)$$

$$\tilde{add}_{i+1}(X, a, b) = 0$$

$$\tilde{mult}_{i+1}(X, a, b) = 0$$



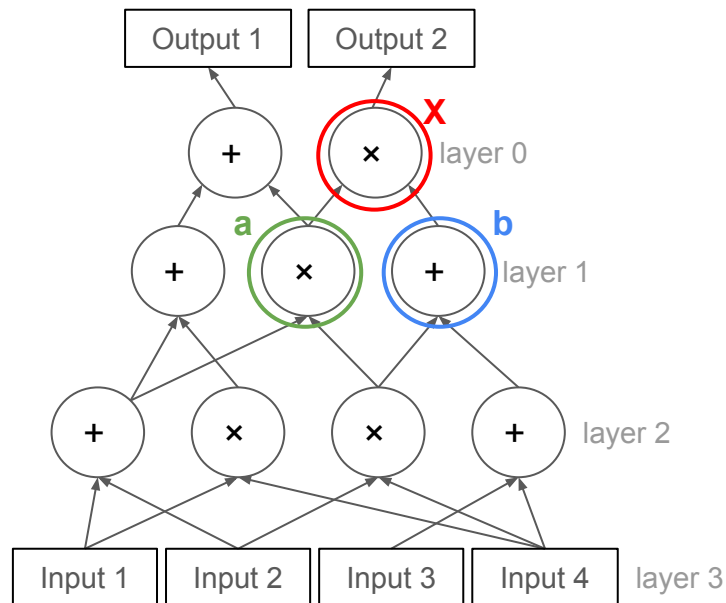
# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( \tilde{add}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + \tilde{mult}_{i+1}(g, a, b) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)) \right)$$

$$\tilde{add}_{i+1}(X, a, b) = 0$$

$$\tilde{mult}_{i+1}(X, a, b) = 1$$

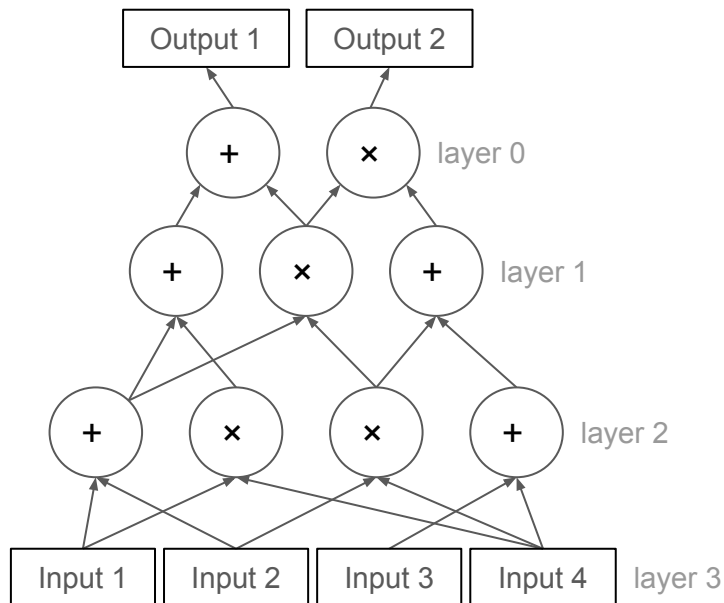


# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$

$$\tilde{V}_{i+1}(X) = \sum_{b \in \{0,1\}^{s_{i+1}}} \chi_b(X) \cdot V(b)$$

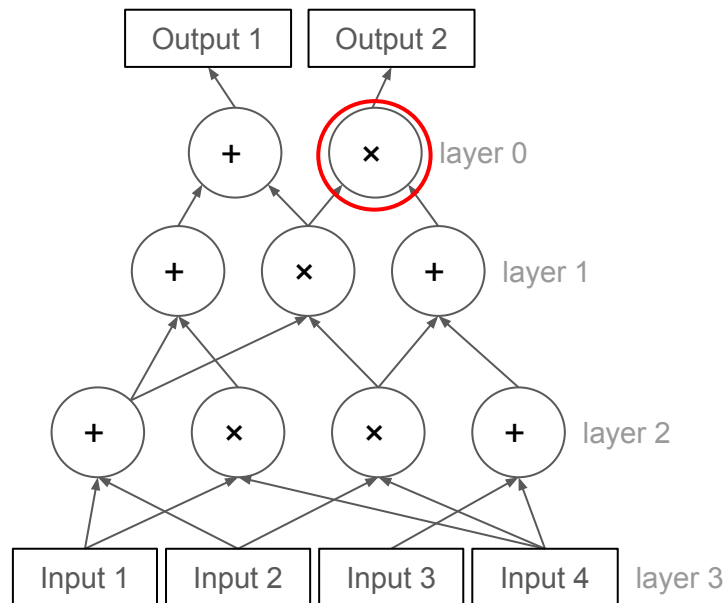


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$$\tilde{V}_i(g) =$$

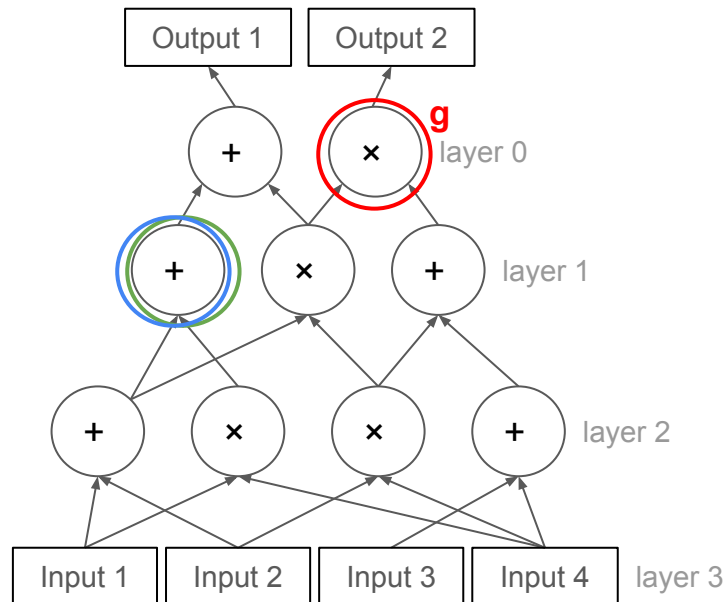


# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

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$$\tilde{V}_i(g) = \mathbf{0}$$

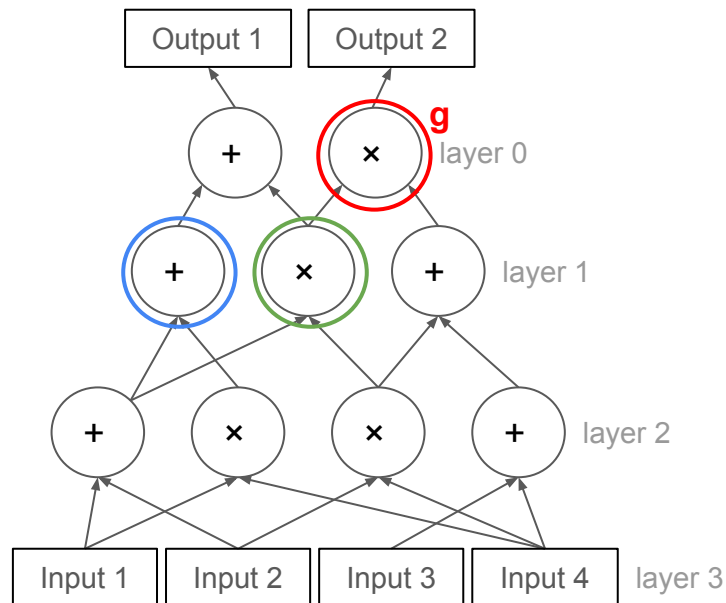


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$$\tilde{V}_i(g) = 0 + 0$$



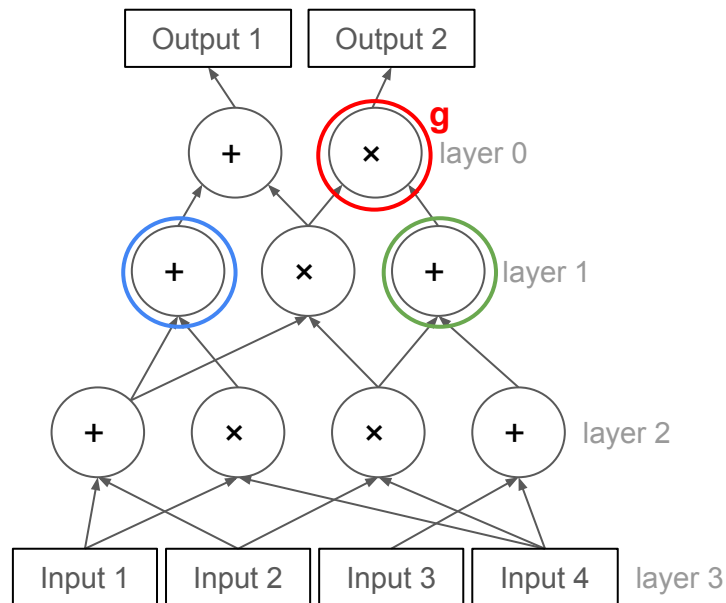


# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$

$$\tilde{V}_i(g) = 0 + 0 + 0$$

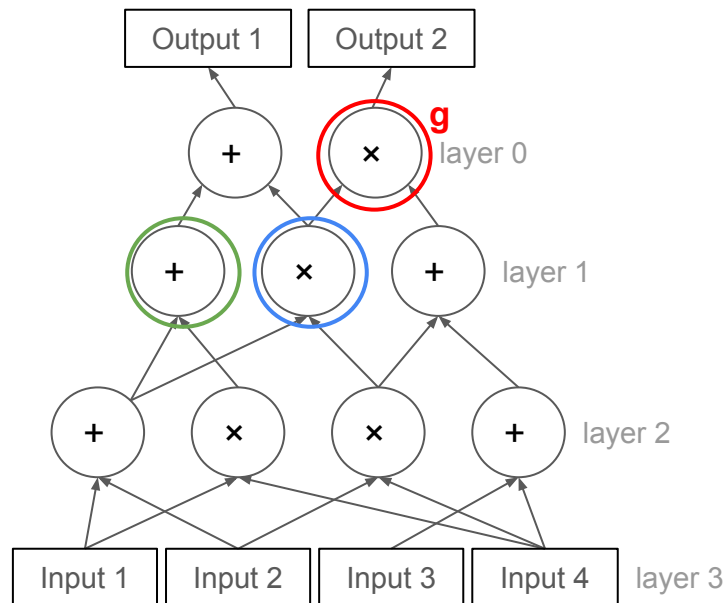


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$$\tilde{V}_i(g) = 0 + 0 + 0 + 0$$

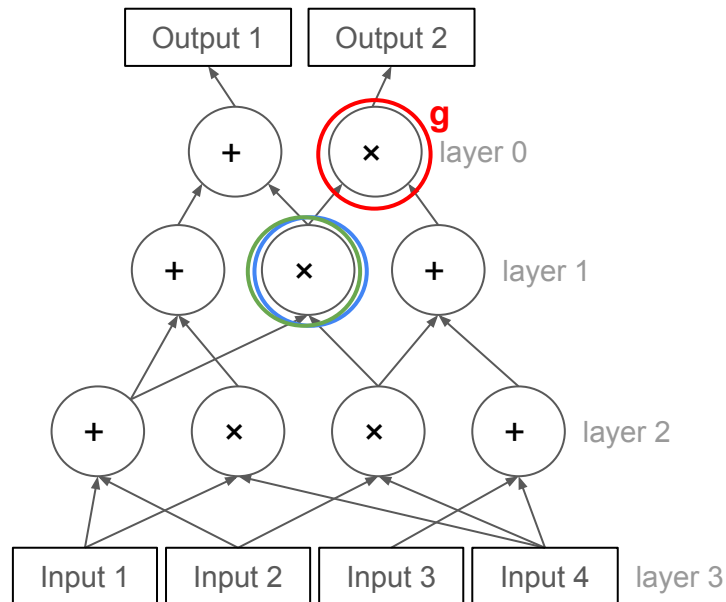


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$$\tilde{V}_i(g) = 0 + 0 + 0 + 0 + 0$$

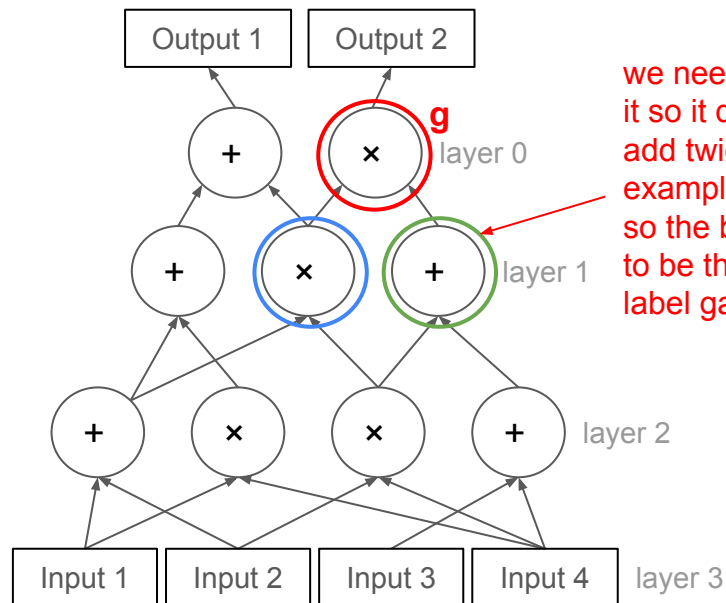


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$$\tilde{V}_i(g) = 0 + 0 + 0 + 0 + 0 + 0$$

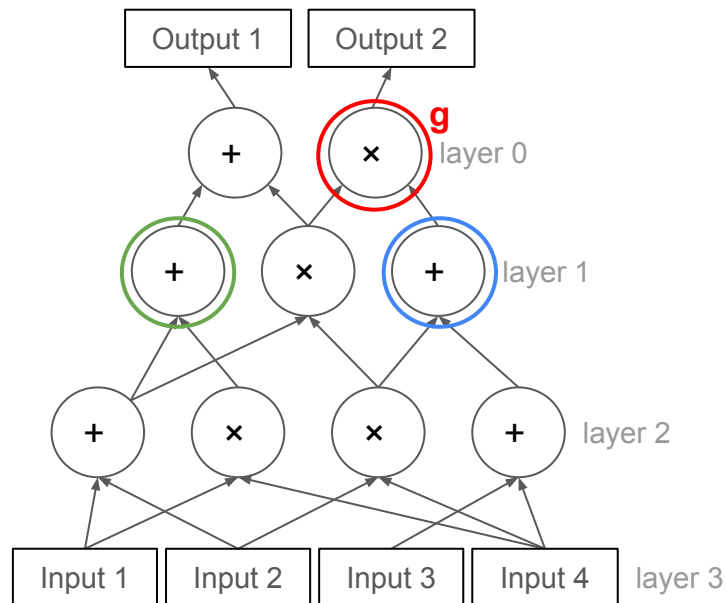


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$$\tilde{V}_i(g) = 0 + 0 + 0 + 0 + 0 + 0 + 0 + 0$$

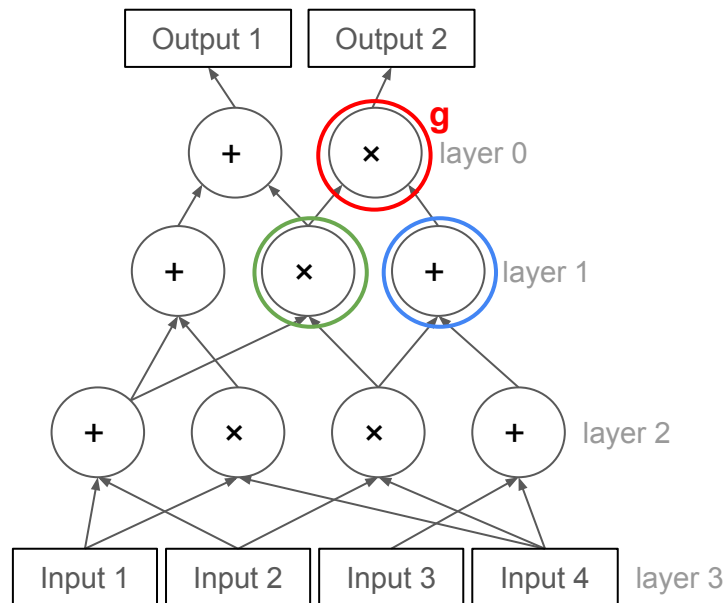


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$$\tilde{V}_i(g) = 0 + 0 + 0 + 0 + 0 + 0 + 0 + \tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)$$

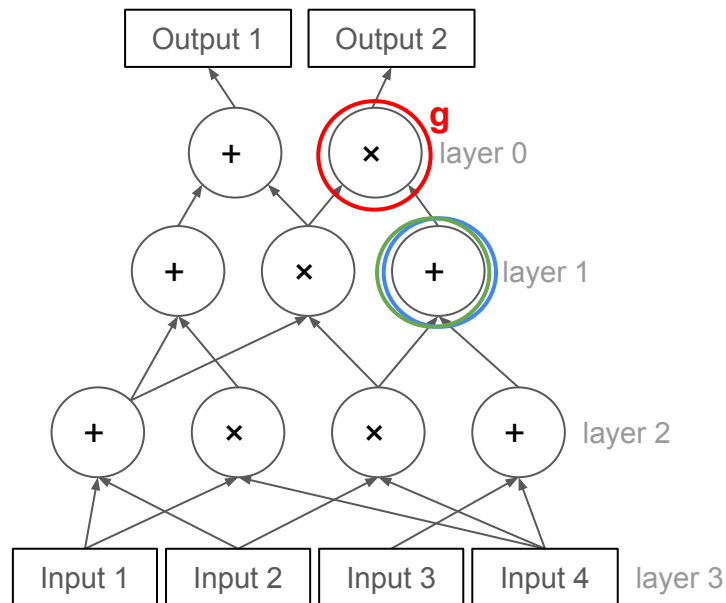


# GKR Protocol (Single Layer Sumcheck)

Sumcheck is performed on the following multivariate polynomial:

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\text{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\text{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$

$$\tilde{V}_i(g) = \mathbf{0 + 0 + 0 + 0 + 0 + 0 + 0} \\ + \tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b) + \mathbf{0}$$



# GKR Protocol (Single Layer Sumcheck)

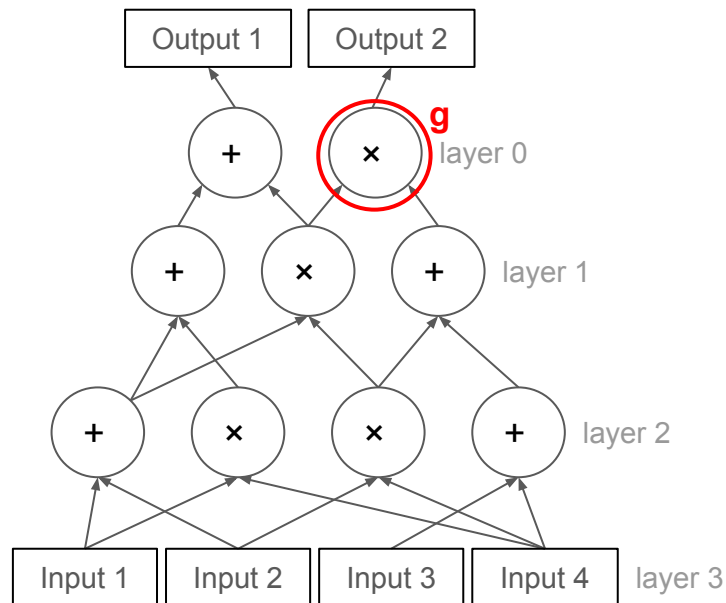
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$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$

$$\tilde{V}_i(g) = 0 + 0 + 0 + 0 + 0 + 0 + 0 + \tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b) + 0$$

$\tilde{V}_i(g)$  is a multivariate polynomial

of  $a_1, \dots, a_m$  and  $b_1, \dots, b_m$   
perform sumcheck on it!



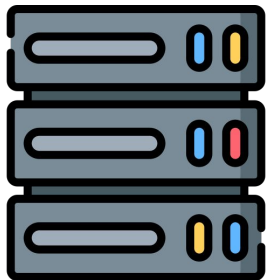


# GKR Protocol (Single Layer Sumcheck)

$$\tilde{V}_0(g, t) = \sum_{a \in \{0,1\}^{(s_{i+1}-1)}, b \in \{0,1\}^{s_{i+1}}} \left( \text{add}_{i+1}(X, t, a, b) (\tilde{V}_{i+1}(t, a) + \tilde{V}_{i+1}(b)) + (\text{mult}_{i+1}(X, t, a, b) (\tilde{V}_{i+1}(t, a) \times \tilde{V}_{i+1}(b)) \right)$$

$\underline{a_1}, \underline{a_2}, \dots, \underline{a_{s_{i+1}}}, \quad \underline{b_1}, \underline{b_2}, \dots, \underline{b_{s_{i+1}}}$

Prover  $P$



$\tilde{V}_0(g, t)$



Verifier  $V$

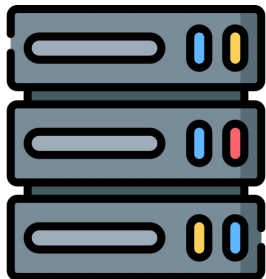


Prover sends univariate polynomial regarding the first bit of  $\mathbf{a}$  ( $a_1$ )

# GKR Protocol (Single Layer Sumcheck)

$$\underline{a_1}, \underline{a_2}, \dots, \underline{a_{s_{i+1}}}, \quad \underline{b_1}, \underline{b_2}, \dots, \underline{b_{s_{i+1}}}$$

Prover  $P$



$$\tilde{V}_0(g) == \tilde{V}_0(g, 0) + \tilde{V}_0(g, 1)$$

Verifier  $V$



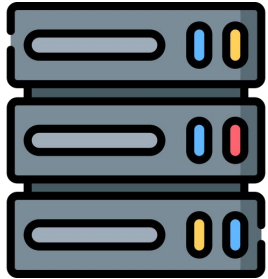
Verifier uses it to verify

$$\tilde{V}_0(g) == \tilde{V}_0(g, 0) + \tilde{V}_0(g, 1)$$

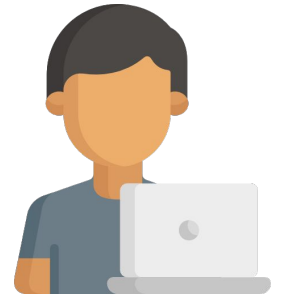
# GKR Protocol (Single Layer Sumcheck)

$$\underline{r_1}, \underline{a_2}, \dots, \underline{a_{s_{i+1}}}, \quad \underline{b_1}, \underline{b_2}, \dots, \underline{b_{s_{i+1}}}$$

Prover  $P$



Verifier  $V$



$r_1$



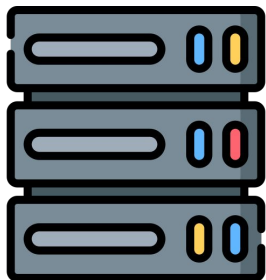
Verifier sends a random value  $r_1$

# GKR Protocol (Single Layer Sumcheck)

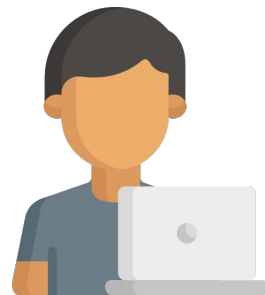
$$\tilde{V}_0(g, r_1, t) = \sum_{a \in \{0,1\}^{(s_{i+1}-2)}, b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(X, r_1, t, a, b)(\tilde{V}_{i+1}(r_1, t, a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(X, r_1, t, a, b)(\tilde{V}_{i+1}(r_1, t, a) \times \tilde{V}_{i+1}(b)))$$

$\underline{r_1}, \underline{a_2}, \dots, \underline{a_{s_{i+1}}}, \quad \underline{b_1}, \underline{b_2}, \dots, \underline{b_{s_{i+1}}}$

Prover  $P$



Verifier  $V$

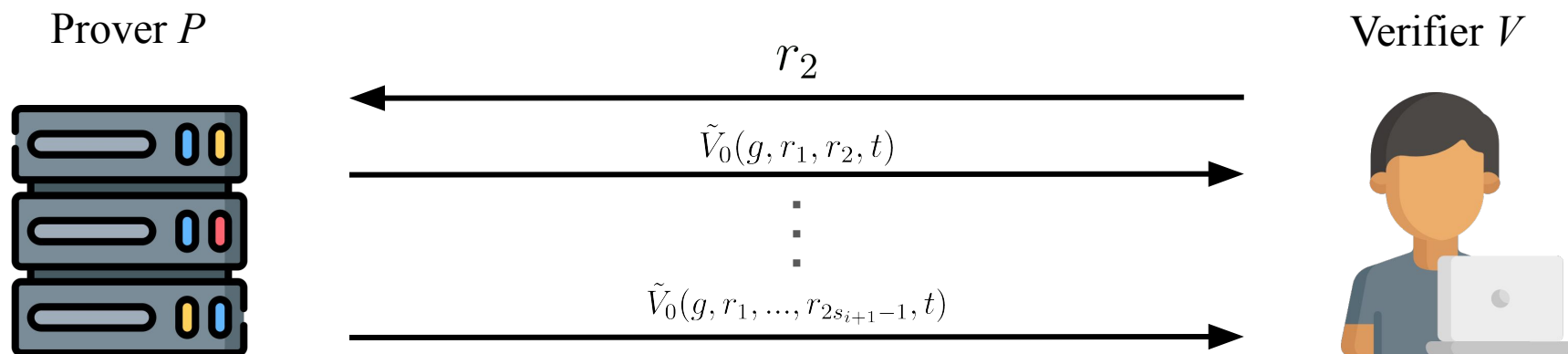


$\tilde{V}_0(g, r_1, t)$

Prover uses  $r_1$  to make the next univariate polynomial in regards to the second bit of  $\mathbf{a}$  ( $a_2$ )

# GKR Protocol (Single Layer Sumcheck)

$$\underline{r_1}, \underline{r_2}, \dots, \underline{r_{s_{i+1}}}, \quad \underline{r_{s_{i+1}+1}}, \underline{r_{s_{i+1}+2}}, \dots, \underline{r_{2s_{i+1}-1}}, \underline{b_{s_{i+1}}}$$



The protocol proceeds as in sumcheck until the last point is randomly selected

# GKR Protocol (Single Layer Sumcheck)

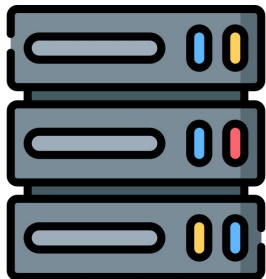
$u$

$$\underline{r_1}, \underline{r_2}, \dots, \underline{r_{s_{i+1}}},$$

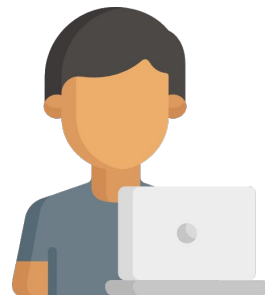
$v$

$$\underline{r_{s_{i+1}+1}}, \underline{r_{s_{i+1}+2}}, \dots, \underline{r_{2s_{i+1}}}$$

Prover  $P$



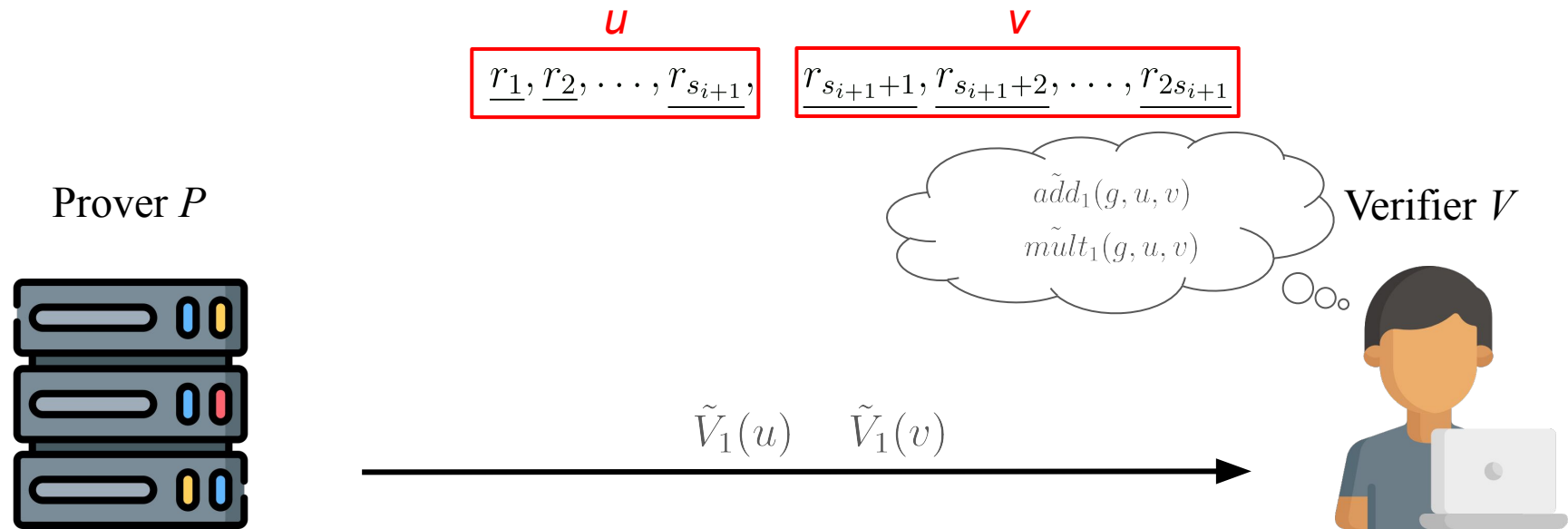
Verifier  $V$



Once the last point is selected denote the first half as  $u$  and the other half as  $v$

*then  $\tilde{V}_0(g, r_1, \dots, r_{2s_{i+1}})$  needs to be verified*

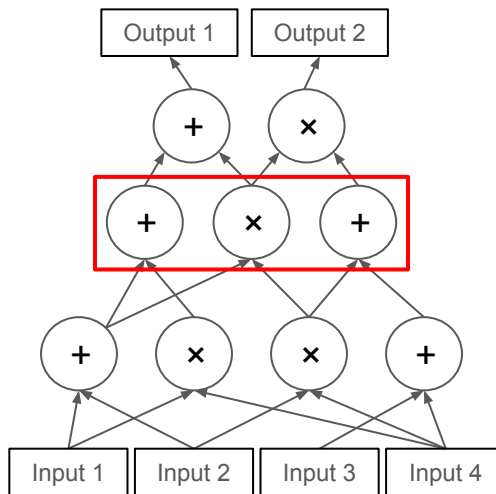
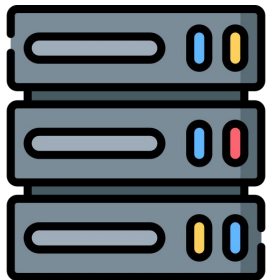
# GKR Protocol (Single Layer Sumcheck)



Prover sends  $\tilde{V}_1(u)$  and  $\tilde{V}_1(v)$   
Verifier computes locally for  $\tilde{add}_1(g, u, v)$  and  $\tilde{mult}_1(g, u, v)$   
Verifier checks  $\tilde{V}_0(g, r_1, \dots, r_{2s_{i+1}}) == \tilde{add}_1(g, u, v)(\tilde{V}_1(u) + \tilde{V}_1(v)) + \tilde{mult}_1(g, u, v)(\tilde{V}_1(u)\tilde{V}_1(v))$

# GKR Protocol

Prover  $P$



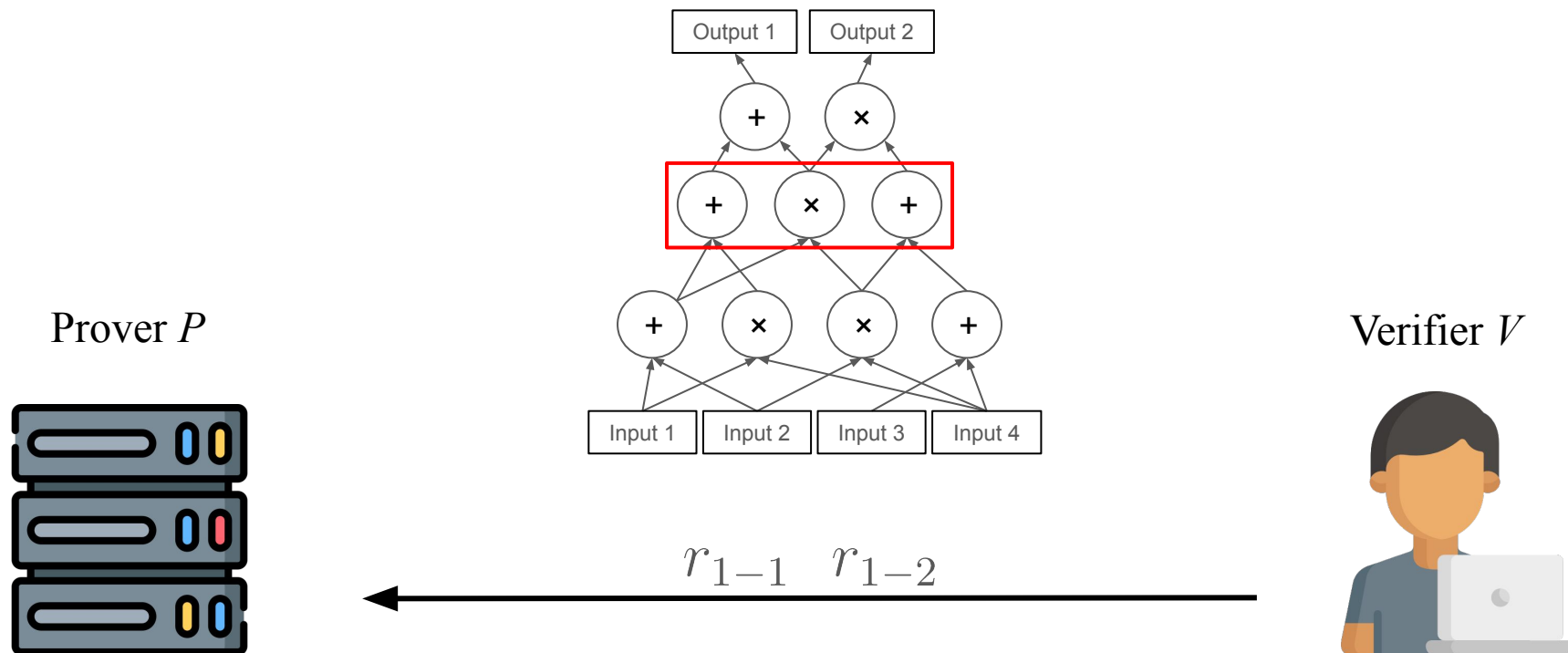
Verifier  $V$



The two claims  $\tilde{V}_1(u)$  and  $\tilde{V}_1(v)$  on the next layer now needs to be verified



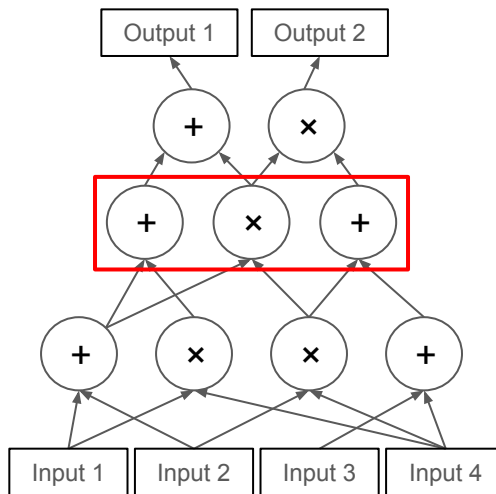
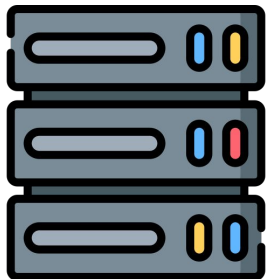
# GKR Protocol



For the new two claims Verifier randomly selects two random values  $r_{1-1}, r_{1-2}$

# GKR Protocol

Prover  $P$



Verifier  $V$

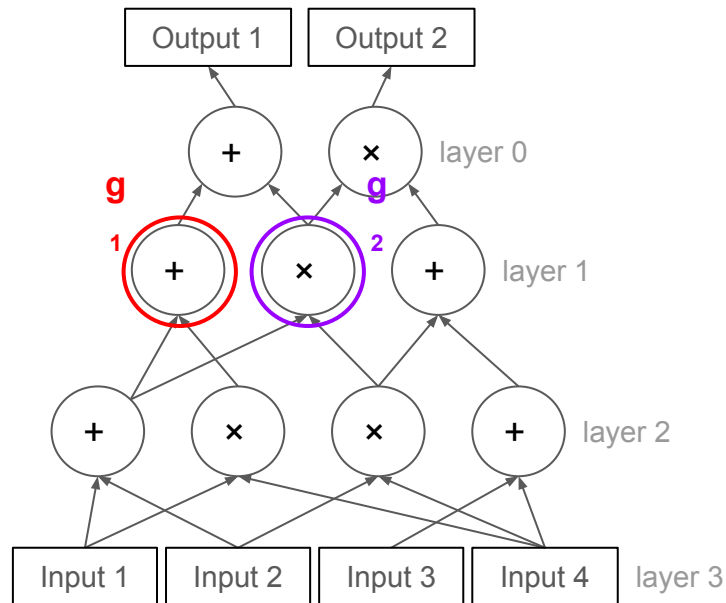


Prover uses  $r_{1-1}r_{1-2}$  to perform linear combination on  $\tilde{V}_1(u)$  and  $\tilde{V}_1(v)$

# GKR Protocol (Linear Combination)

Sumcheck is performed on the following multivariate polynomial:

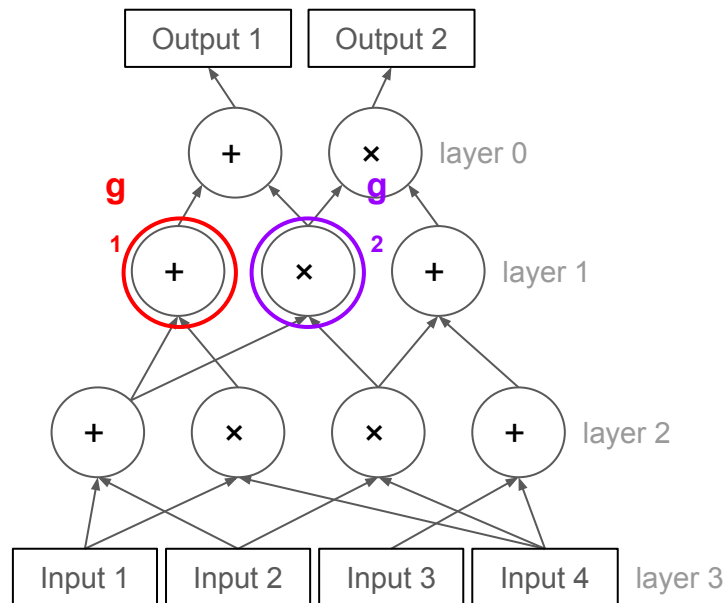
$$\begin{aligned} r_{i-1} \tilde{V}_i(g_1) + r_{i-2} \tilde{V}_i(g_2) = & r_{i-1} \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g_1, a, b)) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g_1, a, b)) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)) \\ & + r_{i-2} \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g_2, a, b)) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g_2, a, b)) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)) \end{aligned}$$



# GKR Protocol (Linear Combination)

Sumcheck is performed on the following multivariate polynomial:

$$r_{i-1}\tilde{V}_i(g_1) + r_{i-2}\tilde{V}_i(g_2) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (r_{i-1}\tilde{add}_{i+1}(g_1, a, b) + r_{i-2}\tilde{add}_{i+1}(g_2, a, b)) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) \\ + (r_{1-1}\tilde{mult}_{i+1}(g_1, a, b) + r_{1-2}\tilde{mult}_{i+1}(g_2, a, b)) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b))$$



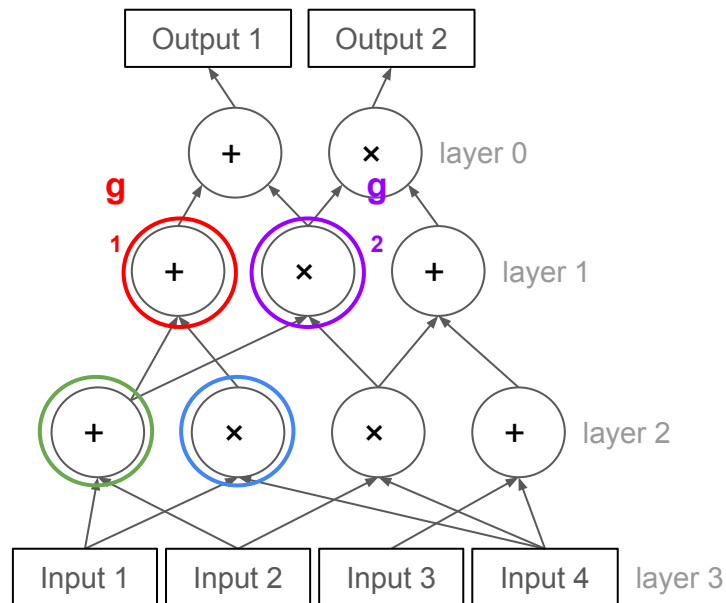
## GKR Protocol (Linear Combination)

## Sumcheck is performed on the following multivariate

**polynomial:**

$$r_{i-1}V_i(g_1) + r_{i-2}V_i(g_2) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( r_{i-1} \tilde{add}_{i+1}(g_1, a, b) + r_{i-2} \tilde{add}_{i+1}(g_2, a, b) \right) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) \\ + \left( r_{1-1} \tilde{mult}_{i+1}(g_1, a, b) + r_{1-2} \tilde{mult}_{i+1}(g_2, a, b) \right) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b))$$

**output :**  $r_{i-1}(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b))$

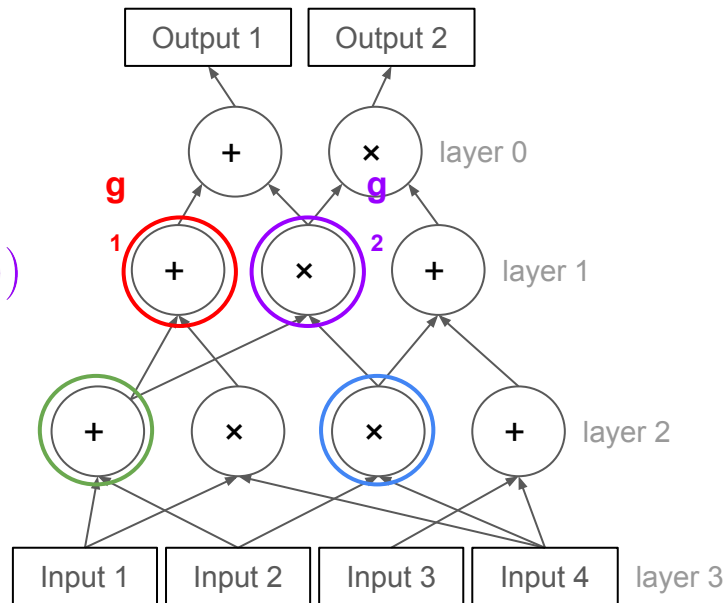


# GKR Protocol (Linear Combination)

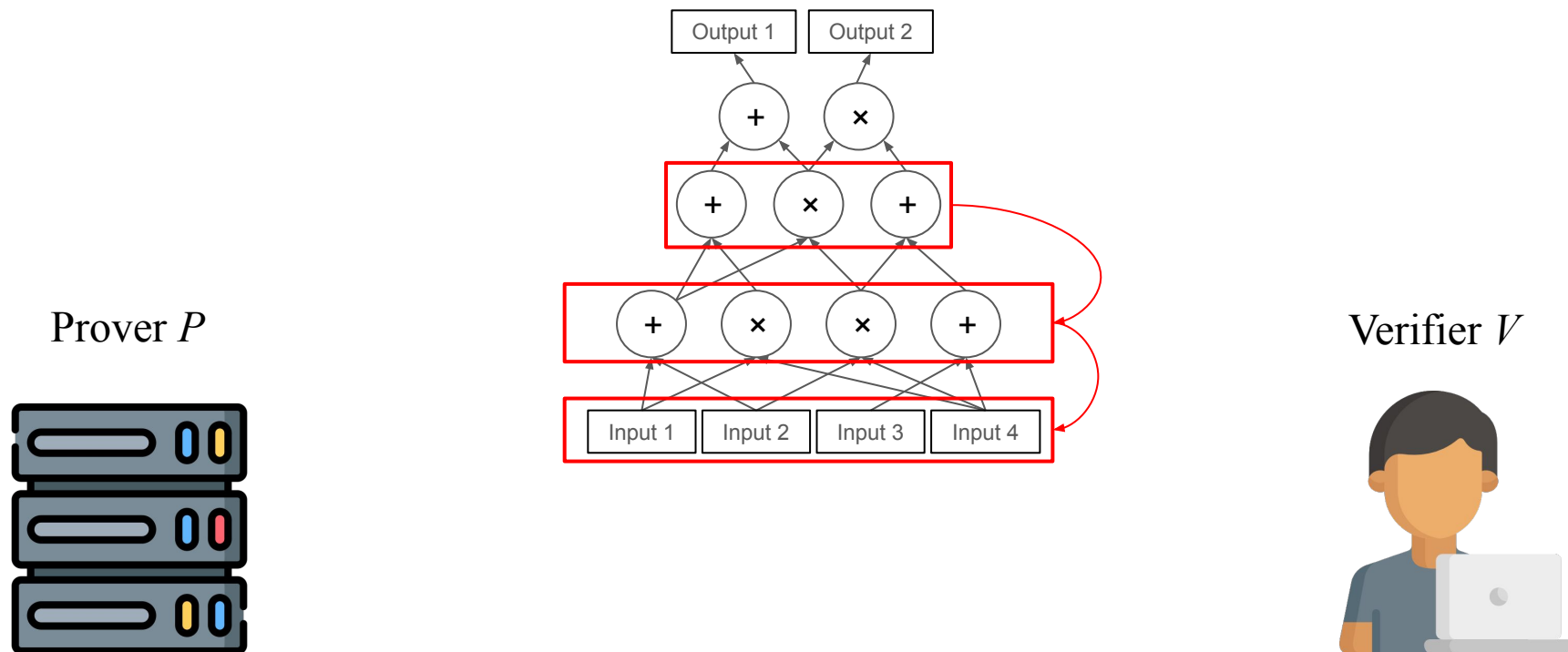
Sumcheck is performed on the following multivariate polynomial:

$$r_{i-1}\tilde{V}_i(g_1) + r_{i-2}\tilde{V}_i(g_2) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (r_{i-1}\tilde{add}_{i+1}(g_1, a, b) + r_{i-2}\tilde{add}_{i+1}(g_2, a, b)) (\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) \\ + (r_{i-1}\tilde{mult}_{i+1}(g_1, a, b) + r_{i-2}\tilde{mult}_{i+1}(g_2, a, b)) (\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b))$$

**output :**  $r_{i-2}(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b))$



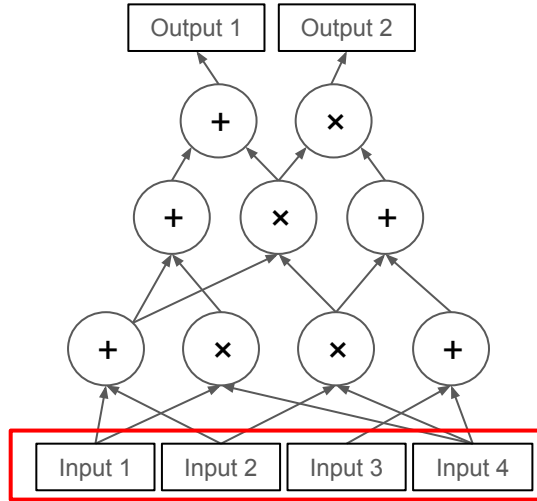
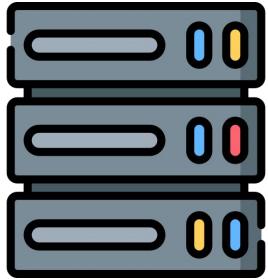
# GKR Protocol



Sumcheck is then performed on  $\tilde{V}_i(g_1) + r_{i-2}\tilde{V}_i(g_2)$ , in the end verifier selects two random gates, repeat until the input layer is reached

# GKR Protocol

Prover  $P$



Verifier  $V$



Verifier then validates the two labels values with it's own input

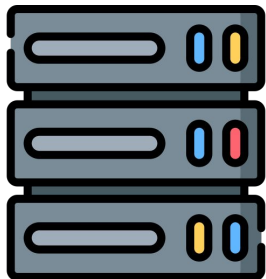


# ZK GKR Protocol

Assume there are  $d$  layers in the arithmetic circuit

declare  $d$  random polynomials  $R_1(X, w), \dots, R_d(X, w)$   
also commits to every polynomial with  $r_1, \dots, r_d$   
to get  $com_{R_1}, \dots, com_{R_d}$

Prover  $P$



Verifier  $V$



Prover defines random polynomial for each layer used to mask the gate outputs  
and commits to those polynomials using  $r_1, \dots, r_d$  to get  $com_{R_1}, \dots, com_{R_d}$

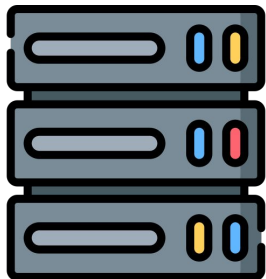
# ZK GKR Protocol

$$\tilde{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\tilde{add}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$



$$\dot{V}_i(g) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} \begin{matrix} I(\vec{0}, w) \\ + I(\vec{0}, (a, b))Z(g)R_i(g_1, w) \end{matrix} (\tilde{add}_{i+1}(g, a, b)(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + (\tilde{mult}_{i+1}(g, a, b)(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)))$$

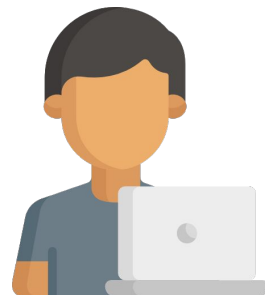
Prover  $P$



$$\chi_{\vec{0}}(a, b) = \begin{cases} 1, & \text{if } (a, b) = 0 \\ 0, & \text{else} \end{cases}$$

$$\prod_{i=1}^n z_i(1 - z_i)$$

Verifier  $V$

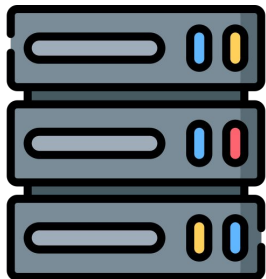


Prover defines random polynomial for each layer used to mask the gate outputs  
and commits to those polynomials using  $r_1, \dots, r_d$  to get  $com_{R_1}, \dots, com_{R_d}$

# ZK GKR Protocol

$$\dot{V}_0(g) = \sum_{a,b \in \{0,1\}^{s_1}, w \in \{0,1\}} I(\vec{0}, w) \left( \tilde{add}_1(g, a, b) (\dot{V}_1(a) + \dot{V}_1(b)) + (\tilde{mult}_1(g, a, b) (\dot{V}_1(a) \times \dot{V}_1(b)) \right. \\ \left. + I(\vec{0}, (a, b)) Z(g) R_0(g_1, w) \right)$$

Prover  $P$



Verifier  $V$

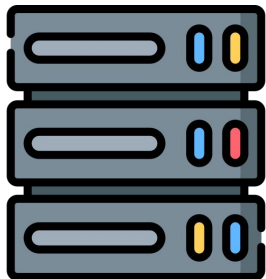


Starting from the output layer, perform ZK Sumcheck  
consider  $R_0(X, w) = 0$

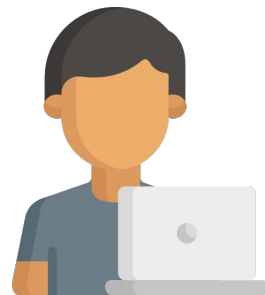
# ZK GKR Protocol

$$\dot{V}_0(g) = \sum_{a,b \in \{0,1\}^{s_1}, w \in \{0,1\}} I(\vec{0}, w) (\tilde{add}_1(g, a, b)(\dot{V}_1(a) + \dot{V}_1(b)) + (\tilde{mult}_1(g, a, b)(\dot{V}_1(a) \times \dot{V}_1(b)) + I(\vec{0}, (a, b))Z(g)R_0(g_1, w))$$

Prover  $P$



Verifier  $V$



$\dot{V}_1(u) \quad \dot{V}_1(v)$



In the last iteration for ZK Sumcheck, three points are selected  $u, v \in \mathbb{F}^{s_1}, c \in \mathbb{F}$

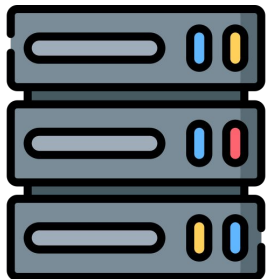
Prover aborts if  $u = v$

Prover sends  $\dot{V}_1(u) \quad \dot{V}_1(v)$

# ZK GKR Protocol

$$\dot{V}_0(g) = \sum_{a,b \in \{0,1\}^{s_1}, w \in \{0,1\}} I(\vec{0}, w) (\tilde{add}_1(g, a, b)(\dot{V}_1(a) + \dot{V}_1(b)) + (\tilde{mult}_1(g, a, b)(\dot{V}_1(a) \times \dot{V}_1(b))) \\ + I(\vec{0}, (a, b))Z(g)R_0(g_1, w)$$

Prover  $P$



Verifier  $V$



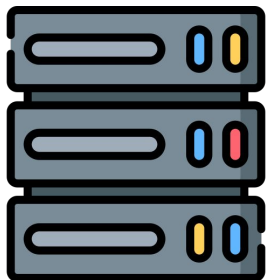
also Verifier computes  $\tilde{add}_1(g, u, v) \quad \tilde{mult}_1(g, u, v)$

Verifier checks  $h_n(r_n) - pg(r_1, \dots, r_n) == \tilde{add}_1(g, a, b)(\dot{V}_1(a) + \dot{V}_1(b)) + \tilde{mult}_1(g, a, b)(\dot{V}_1(a) \times \dot{V}_1(b))$

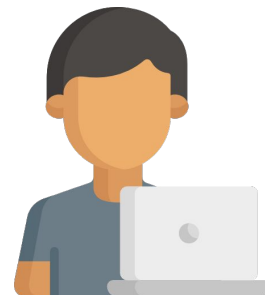
# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) \left( (r_{1-1}\tilde{add}_{i+1}(u, a, b) + r_{1-2}\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}\tilde{mult}_{i+1}(u, a, b) + r_{1-2}\tilde{mult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) \right. \\ \left. + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w)) \right)$$

Prover  $P$



Verifier  $V$



$r_{1-1} \quad r_{1-2}$

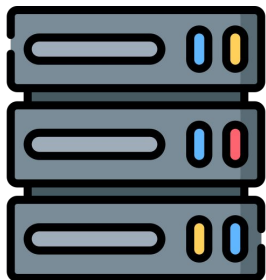


Verifier sends two randomnesses  $r_{1-1} \quad r_{1-2}$  to prover for linear combination

# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) ((r_{1-1}\tilde{add}_{i+1}(u, a, b) + r_{1-2}\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}\tilde{mult}_{i+1}(u, a, b) + r_{1-2}\tilde{mult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w))$$

Prover  $P$



Verifier  $V$



$\dot{V}_{i+1}(u^{(i+1)}) \quad \dot{V}_{i+1}(v^{(i+1)})$



At the end of ZK Sumcheck for each layer, three points are chosen

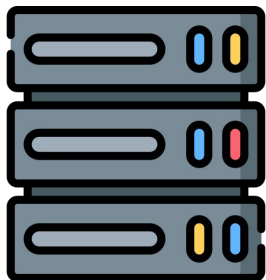
$u^{(i+1)}, v^{(i+1)} \in \mathbb{F}^{s_{i+1}}, c^{(i)} \in \mathbb{F}$ , Prover aborts if  $c^{(i+1)} = v^{(i+1)}$

Prover sends  $\dot{V}_{i+1}(u^{(i+1)}) \quad \dot{V}_{i+1}(v^{(i+1)})$

# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) ((r_{1-1}a\tilde{add}_{i+1}(u, a, b) + r_{1-2}a\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}m\tilde{ult}_{i+1}(u, a, b) + r_{1-2}m\tilde{ult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w))$$

Prover  $P$



Verifier  $V$



$\dot{V}_{i+1}(u^{(i+1)}) \quad \dot{V}_{i+1}(v^{(i+1)})$



Verifier computes:

$$r_{1-1}a\tilde{add}_{i+1}(u^{(i)}, u^{(i+1)}, v^{(i+1)}) + r_{1-2}a\tilde{add}_{i+1}(v^{(i)}, u^{(i+1)}, v^{(i+1)})$$

$$r_{1-1}m\tilde{ult}_{i+1}(u^{(i)}, u^{(i+1)}, v^{(i+1)}) + r_{1-2}m\tilde{ult}_{i+1}(v^{(i)}, u^{(i+1)}, v^{(i+1)})$$

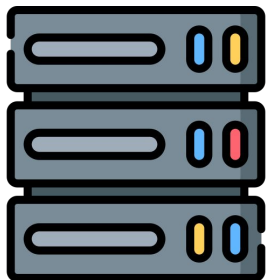
$$Z(u^{(i)}), Z(v^{(i)}), I(\vec{0}, c^{(i)}), I(\vec{0}, (u^{(i+1)}, v^{(i+1)}))$$



# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) ((r_{1-1}\tilde{add}_{i+1}(u, a, b) + r_{1-2}\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}\tilde{mult}_{i+1}(u, a, b) + r_{1-2}\tilde{mult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w))$$

Prover  $P$



Verifier  $V$



$R_i(u_1^{(i)}, c^{(i)}), \pi_{i-u}, R_i(v_1^{(i)}, c^{(i)}), \pi_{i-v}$



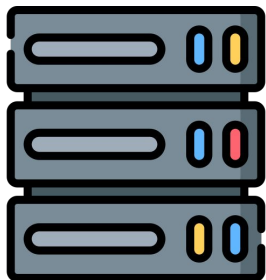
Prover opens  $R_i$  on two points and sends  
 Verifier verify  $R_i(u_1^{(i)}, c^{(i)}), R_i(v_1^{(i)}, c^{(i)})$

$(R_i(u_1^{(i)}, c^{(i)}), \pi_{i-u}) \quad (R_i(v_1^{(i)}, c^{(i)}), \pi_{i-v})$

# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) ((r_{1-1}\tilde{add}_{i+1}(u, a, b) + r_{1-2}\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}\tilde{mult}_{i+1}(u, a, b) + r_{1-2}\tilde{mult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w))$$

Prover  $P$



Verifier  $V$



$R_i(u^{(i)}, c^{(i)}), \pi_{i-u}, R_i(v^{(i)}, c^{(i)}), \pi_{i-v}$



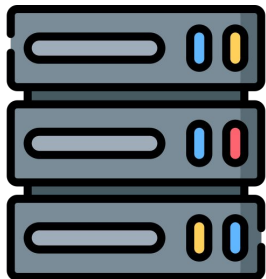
Verifier checks

$$h_n(r_n) - pg(r_1, \dots, r_n) == I(\vec{0}, c^{(i)}) ((r_{1-1}\tilde{add}_{i+1}(u^{(i)}, u^{(i+1)}, v^{(i+1)}) + r_{1-2}\tilde{add}_{i+1}(v^{(i)}, u^{(i+1)}, v^{(i+1)}))(\dot{V}_{i+1}(u^{(i+1)}) + \dot{V}_{i+1}(v^{(i+1)}))) + (r_{1-1}\tilde{mult}_{i+1}(u^{(i)}, u^{(i+1)}, v^{(i+1)}) + r_{1-2}\tilde{mult}_{i+1}(v^{(i)}, u^{(i+1)}, v^{(i+1)}))(\dot{V}_{i+1}(u^{(i+1)}) \times \dot{V}_{i+1}(v^{(i+1)}))) + I(\vec{0}, (u^{(i+1)}, v^{(i+1)}))(r_{1-1}Z(u^{(i)})R_i(u^{(i)}, c^{(i)}) + r_{1-2}Z(v^{(i)})R_i(v^{(i)}, c^{(i)}))$$

# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) \left( (r_{1-1}\tilde{add}_{i+1}(u, a, b) + r_{1-2}\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}\tilde{mult}_{i+1}(u, a, b) + r_{1-2}\tilde{mult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) \right. \\ \left. + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w)) \right)$$

Prover  $P$



Verifier  $V$

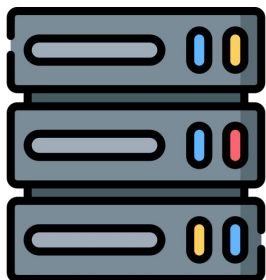


Then repeat the verifying process for the next layer

# ZK GKR Protocol

$$r_{1-1}\dot{V}_i(u) + r_{1-2}\dot{V}_i(v) = \sum_{a,b \in \{0,1\}^{s_{i+1}}, w \in \{0,1\}} I(\vec{0}, w) ((r_{1-1}\tilde{add}_{i+1}(u, a, b) + r_{1-2}\tilde{add}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) + \dot{V}_{i+1}(b)) + ((r_{1-1}\tilde{mult}_{i+1}(u, a, b) + r_{1-2}\tilde{mult}_{i+1}(v, a, b))(\dot{V}_{i+1}(a) \times \dot{V}_{i+1}(b)) + I(\vec{0}, (a, b))(r_{1-1}Z(u)R_i(u_1, w) + r_{1-2}Z(u)R_i(v_1, w))$$

Prover  $P$



$$R_d(u_1^{(d)}, 0), R_d(u_1^{(d)}, 1), R_d(v_1^{(d)}, 0), R_d(v_1^{(d)}, 1)$$

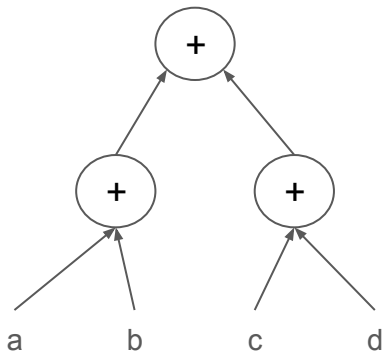


Verifier  $V$



Once the input layer is reached Verifier now has two claims on  $\dot{V}_d(v^{(d)})$   $\dot{V}_d(u^{(d)})$   
 Open  $R_d$  on four points  $R_d(u^{(d)}, 0), R_d(u^{(d)}, 1), R_d(v^{(d)}, 0), R_d(v^{(d)}, 1)$  and verify them,  
 Last Verifier checks  $\dot{V}_d(u^{(d)}) == \tilde{V}_d(u^{(d)}) + \sum_{w \in \{0,1\}} R_d(u_1^{(d)}, w)$  and  $\dot{V}_d(v^{(d)}) == \tilde{V}_d(v^{(d)}) + \sum_{w \in \{0,1\}} R_d(v_1^{(d)}, w)$

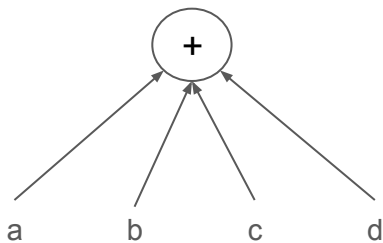
# GKR Protocol (Generalized)



$$\tilde{add}_{i+1}(X, a, b) \begin{cases} 1, & \text{if gate } a, b \text{ inputs to add gate} \\ X \\ 0, & \text{else} \end{cases}$$

two input per add gate

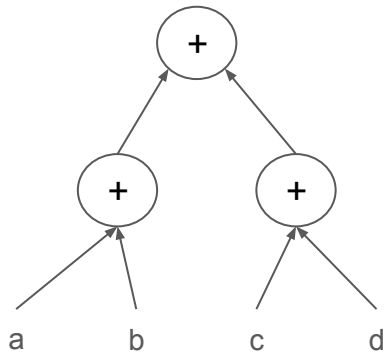
# GKR Protocol (Generalized)



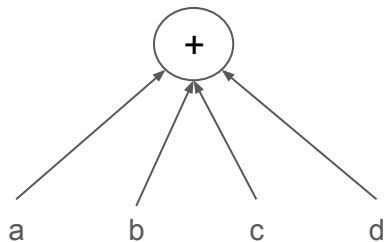
$$Xadd_{i+1}(X, a) \begin{cases} 1, & \text{if gate } a \text{ inputs to add gate } X \\ 0, & \text{else} \end{cases}$$

multiple input per add  
gate

# GKR Protocol (Generalized)

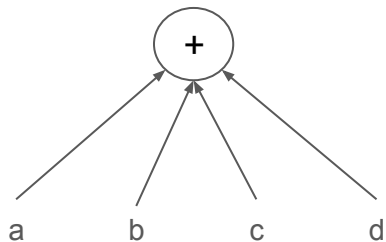


$$\tilde{V}_i(X) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( \text{add}_{i+1}(X, a, b)(\tilde{V}_{i+1}(a) + \tilde{V}_{i+1}(b)) \right. \\ \left. + (\text{mult}_{i+1}(X, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b))) \right)$$



$$\tilde{V}_i(X) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} \left( \chi_0(b) \text{add}_{i+1}(X, a)(\tilde{V}_{i+1}(a)) \right. \\ \left. + (\text{mult}_{i+1}(X, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b))) \right)$$

This is used in  
Convolution layer

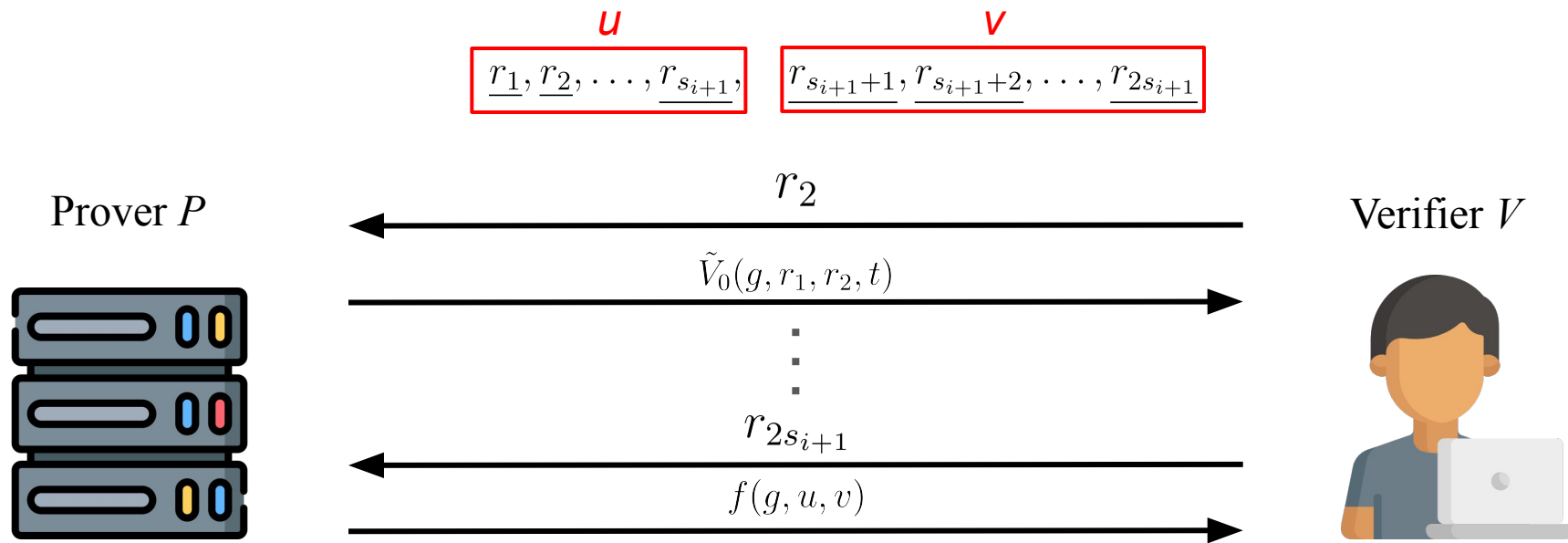


$$\tilde{V}_i(X) = \sum_{a,b \in \{0,1\}^{s_{i+1}}} (\chi_0(b) \mathit{add}_{i+1}(X, a)(\tilde{V}_{i+1}(a)) \\ + (\mathit{mult}_{i+1}(X, a, b)(\tilde{V}_{i+1}(a) \times \tilde{V}_{i+1}(b)))$$



**Instead of combining two claims using linear combination, there is an alternative method to reduce the two claim into one**

# GKR Protocol (CMT - GKR)



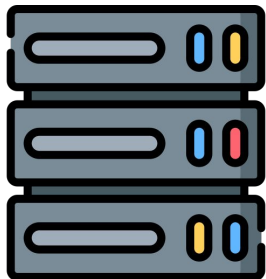
At the end of sumcheck two random labels are given, denote the first label as  $u$  and the second label as  $v$

# GKR Protocol (CMT - GKR)

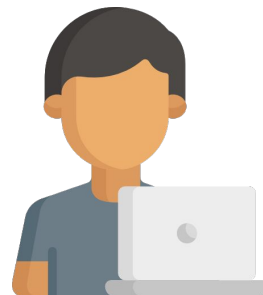
define a line  $L$  used only in the next layer, where  $L(0) = u$  and  $L(1) = v$

$$\begin{array}{cc} u & v \\ \boxed{\underline{r_1}, \underline{r_2}, \dots, \underline{r_{s_{i+1}}},} & \boxed{\underline{r_{s_{i+1}+1}}, \underline{r_{s_{i+1}+2}}, \dots, \underline{r_{2s_{i+1}}}} \end{array}$$

Prover  $P$



Verifier  $V$



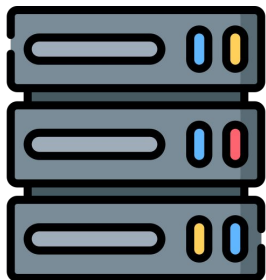
Instead of ending it right there, further define a line where  $L(0) = u$  and  $L(1) = v$

# GKR Protocol (CMT - GKR)

define a line  $L$  used only in the next layer, where  $L(0) = u$  and  $L(1) = v$

$$\begin{array}{c} u \\ \boxed{\underline{r_1}, \underline{r_2}, \dots, \underline{r_{s_{i+1}}}} \end{array} \quad \begin{array}{c} v \\ \boxed{\underline{r_{s_{i+1}+1}}, \underline{r_{s_{i+1}+2}}, \dots, \underline{r_{2s_{i+1}}}} \end{array}$$

Prover  $P$



Verifier  $V$



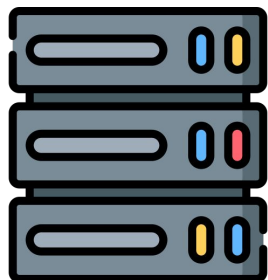
Prover then uses that line to make  $\tilde{h}_i(r') = \tilde{V}_i(L(r'))$   
In the first round of sumcheck,  $\kappa'$  is first selected

# GKR Protocol (CMT - GKR)

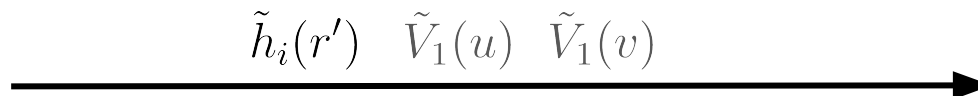
define a line  $L$  used only in the next layer, where  $L(0) = u$  and  $L(1) = v$

$$\begin{array}{c} u \qquad \qquad \qquad v \\ \boxed{\underline{r_1}, \underline{r_2}, \dots, \underline{r_{s_{i+1}}}}, \quad \boxed{\underline{r_{s_{i+1}+1}}, \underline{r_{s_{i+1}+2}}, \dots, \underline{r_{2s_{i+1}}}} \end{array}$$

Prover  $P$



Verifier  $V$



Prover sends  $\tilde{h}_i(r') \quad \tilde{V}_1(u) \quad \tilde{V}_1(v)$

Verifier uses it to verify whether  $\tilde{h}_i(0) = \tilde{V}_i(u) \quad \tilde{h}_i(1) = \tilde{V}_i(v)$