

CS 106X, Lecture 10

Recursive Backtracking

reading:

Programming Abstractions in C++, Chapter 9

Plan For Today

- Ex. Printing Binary
- What is Recursive Backtracking?
- Ex. Dice Sums
- **Announcements**
- Ex. Subsets

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Exhaustive search

- **exhaustive search:** Exploring every possible combination from a set of choices or values.
 - often implemented recursively

Applications:

- producing all permutations of a set of values
- enumerating all possible names, passwords, etc.
- combinatorics and logic programming

- Often the search space consists of many ***decisions***, each of which has several available ***choices***.
 - Example: When enumerating all 5-letter strings, each of the 5 letters is a *decision*, and each of those decisions has 26 possible *choices*.

Exhaustive search



A general pseudo-code algorithm for exhaustive search:

Search(decisions, storage). Specifically: decisions:left
decisions I decide to do sth....

function **Search** (*decisions*):

each call makes a choice

- If there are decisions left to make:

// Let's handle one decision ourselves, and the rest by recursion.

- For each available choice *C* for this decision:

- Choose *C*. loop maybe to go over each choice you can make.
 - Search the remaining decisions that could follow *C*.

- Otherwise, if there are no more decisions to make: Stop.

- *Observation:* The "base case" no longer represents a simple case of the algorithm; rather, it is the case where the algorithm is finished working and has no more choices left to make. **stop base case**

Backtracking



- **backtracking:** Finding solution(s) by trying partial solutions and then abandoning them if they are not suitable.
 - a "brute force" algorithmic technique (tries all paths)
 - often implemented recursively

Applications:

- producing all permutations of a set of values
- parsing languages
- games: anagrams, crosswords, word jumbles, 8 queens
- combinatorics and logic programming
- escaping from a maze

Backtracking



A general pseudo-code algorithm for backtracking:

function Search (decisions):

Note: understanding **what choices you have** each step before starting is vitally important.

- If there are decisions left to make:
 - // Let's handle one decision ourselves, and the rest by recursion.
 - For each available choice C for this decision:
 - Choose C .
 - Search the remaining decisions that could follow C .
 - Un-choose C . (*backtrack!*)
 - Otherwise, if there are no more decisions to make: Stop.

- Key tasks:
 - Figure out appropriate smallest unit of work (decision).
 - Figure out how to enumerate all possible choices/options for it.



Exercise: printAllBinary

- Write a recursive function **printAllBinary** that accepts an integer number of digits and prints all binary numbers that have exactly that many digits, in ascending order, one per line.

- `printAllBinary(2);`

00
01
10
11

`printAllBinary(3);`

000
001
010
011
100
101
110
111

Wrapper?

Base Case?

Recursive case?

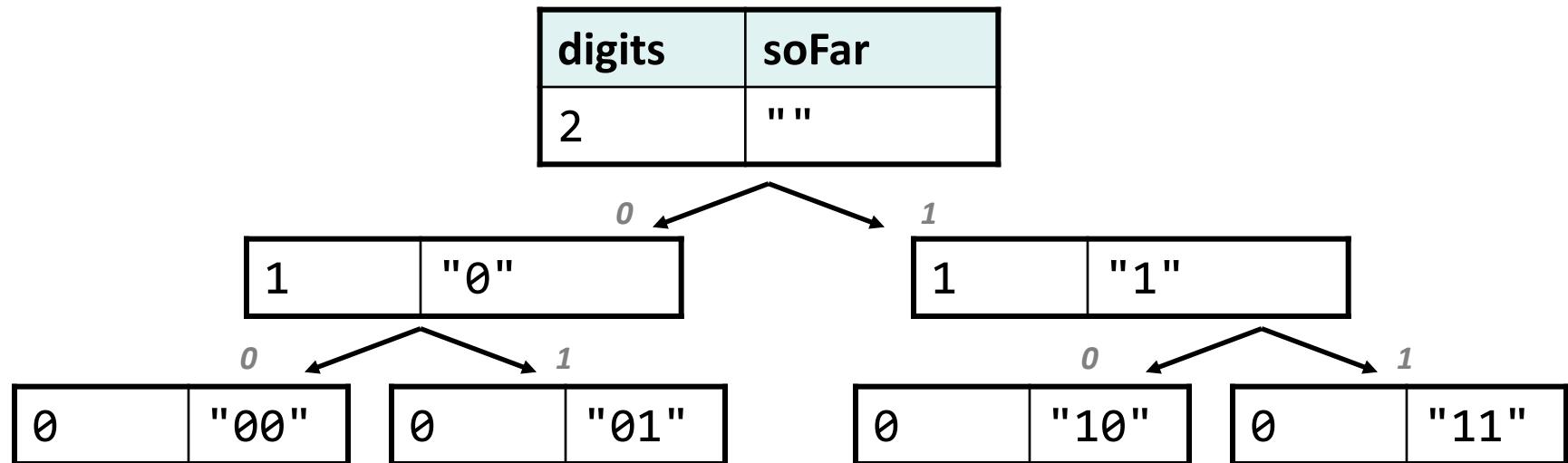
All inputs?

printAllBinary solution

```
void printAllBinary(int numDigits) {  
    printAllBinaryHelper(numDigits, "");  
}  
  
void printAllBinaryHelper(int digits, string soFar) {  
    if (digits == 0) {  
        cout << soFar << endl;  
    } else {  
        printAllBinaryHelper(digits - 1, soFar + "0");  
        printAllBinaryHelper(digits - 1, soFar + "1");  
    }  
}
```

A tree of calls

- `printAllBinary(2);`



- This kind of diagram is called a *call tree* or *decision tree*.
- Think of each call as a choice or decision made by the algorithm:
 - Should I choose 0 as the next digit?
 - Should I choose 1 as the next digit?

The base case

```
void printAllBinaryHelper(int digits, string soFar) {  
    if (digits == 0) {  
        cout << soFar << endl;  
    } else {  
        printAllBinaryHelper(digits - 1, soFar + "0");  
        printAllBinaryHelper(digits - 1, soFar + "1");  
    }  
}
```

- The **base case** is where the code stops after doing its work.
 - pAB(3) → pAB(2) → pAB(1) → pAB(0)
- Each call should keep track of the work it has done.
- Base case should print the result of the work done by prior calls.
 - Work is kept track of in some variable(s) - in this case, `string soFar`.

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Recursive Backtracking

- **Recursive Backtracking:** using recursion to explore solutions to a problem and abandoning them if they are not suitable.
 - Determine whether a solution exists
 - Find a solution
 - Find the best solution
 - Count the number of solutions
 - Print/find all the solutions
- Applications:
 - Puzzle solving (Sudoku, Crosswords, etc.)
 - Game playing (Chess, Solitaire, etc.)
 - Constraint satisfaction problems (scheduling, matching, etc.)

The Recursion Checklist

- ❑ Find what information we need to keep track of.**
What inputs/outputs are needed to solve the problem at each step? Do we need a wrapper function?
- ❑ Find our base case(s).** What are the simplest (non-recursive) instance(s) of this problem?
- ❑ Find our recursive step.** How can this problem be solved in terms of one or more simpler instances of the same problem that lead to a base case?
- ❑ Ensure every input is handled.** Do we cover all possible cases? Do we need to handle errors?

The Recursion Checklist

- ❑ **Find what information we need to keep track of.**
What inputs/outputs are needed to solve the problem at each step? Do we need a wrapper function?
- ❑ **Find our base case(s).** What are the simplest (non-recursive) instance(s) of this problem?
- ❑ **Find our recursive step.** How can this problem be solved in terms of one or more simpler instances of the same problem that lead to a base case?
- ❑ **Ensure every input is handled.** Do we cover all possible cases? Do we need to handle errors?

The Backtracking Checklist

- ❑ **Find what choice(s) we have at each step.** What different options are there for the next step?

For each valid choice:

- ❑ **Make it and explore recursively.** Pass the information for a choice to the next recursive call(s).
- ❑ **Undo it after exploring.** Restore everything to the way it was before making this choice.
- ❑ **Find our base case(s).** What should we do when we are out of decisions?

printAllBinary solution

```
void printAllBinary(int numDigits) {  
    printAllBinaryHelper(numDigits, "");  
}  
  
void printAllBinaryHelper(int digits, string soFar) {  
    if (digits == 0) {  
        cout << soFar << endl;  
    } else {  
        printAllBinaryHelper(digits - 1, soFar + "0");  
        printAllBinaryHelper(digits - 1, soFar + "1");  
    }  
}
```

Choose

Explore

Un-choose

Base case

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Exercise: Dice rolls

`diceRoll`

- Write a recursive function `diceRoll` that accepts an integer representing a number of 6-sided dice to roll, and output all possible combinations of values that could appear on the dice.

`diceRoll(2);`

{1, 1}	{3, 1}	{5, 1}
{1, 2}	{3, 2}	{5, 2}
{1, 3}	{3, 3}	{5, 3}
{1, 4}	{3, 4}	{5, 4}
{1, 5}	{3, 5}	{5, 5}
{1, 6}	{3, 6}	{5, 6}
{2, 1}	{4, 1}	{6, 1}
{2, 2}	{4, 2}	{6, 2}
{2, 3}	{4, 3}	{6, 3}
{2, 4}	{4, 4}	{6, 4}
{2, 5}	{4, 5}	{6, 5}
{2, 6}	{4, 6}	{6, 6}

`diceRoll(3);`



{1, 1, 1}
{1, 1, 2}
{1, 1, 3}
{1, 1, 4}
{1, 1, 5}
{1, 1, 6}
{1, 2, 1}
{1, 2, 2}
...
{6, 6, 4}
{6, 6, 5}
{6, 6, 6}

- How is this problem recursive (self-similar)?

The Backtracking Checklist

- ❑ **Find what choice(s) we have at each step.** What different options are there for the next step?

For each valid choice:

- ❑ **Make it and explore recursively.** Pass the information for a choice to the next recursive call(s).
- ❑ **Undo it after exploring.** Restore everything to the way it was before making this choice.
- ❑ **Find our base case(s).** What should we do when we are out of decisions?

The Backtracking Checklist

- ❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each val

What die value should I choose next?

- ❑ Make it for a choice to the next recursive call(s).
- ❑ Undo it after exploring. Restore everything to the way it was before making this choice.
- ❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

❑ **Make it and explore recursively.** Pass the information for a choice to the next recursive call(s).

❑ Undo it after exploring. Restore everything to the way it was before making this choice.

❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

- ❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

- ❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).

- ❑ Undo it : We need to communicate it was before to the way

- ❑ Find our base case when we

are out of decisions?

**We need to communicate
the dice chosen so far to
the next recursive call.**

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).

❑ Undo it after exploring. Restore everything to the way it was before making this choice.

❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

- Find what choice(s) we have at each step. What different options do we have? For each value:
 - Make it a choice. Remove the die we added to our roll so far. information for a choice to the next recursive call(s).
 - Undo it after exploring. Restore everything to the way it was before making this choice.
- Find our base case(s). What should we do when we are out of decisions?

We need to be able to remove the die we added to our roll so far.

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).

❑ Undo it after exploring. Restore everything to the way it was before making this choice.

❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

- ❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

- ❑ Make it and explore recursively. Pass the information for a choice.
- ❑ Undo it
- ❑ Find our base case(s). What should we do when we are out of decisions?

When we have no dice left to choose, print them out.

Examining the problem

- We want to generate all possible sequences of values.

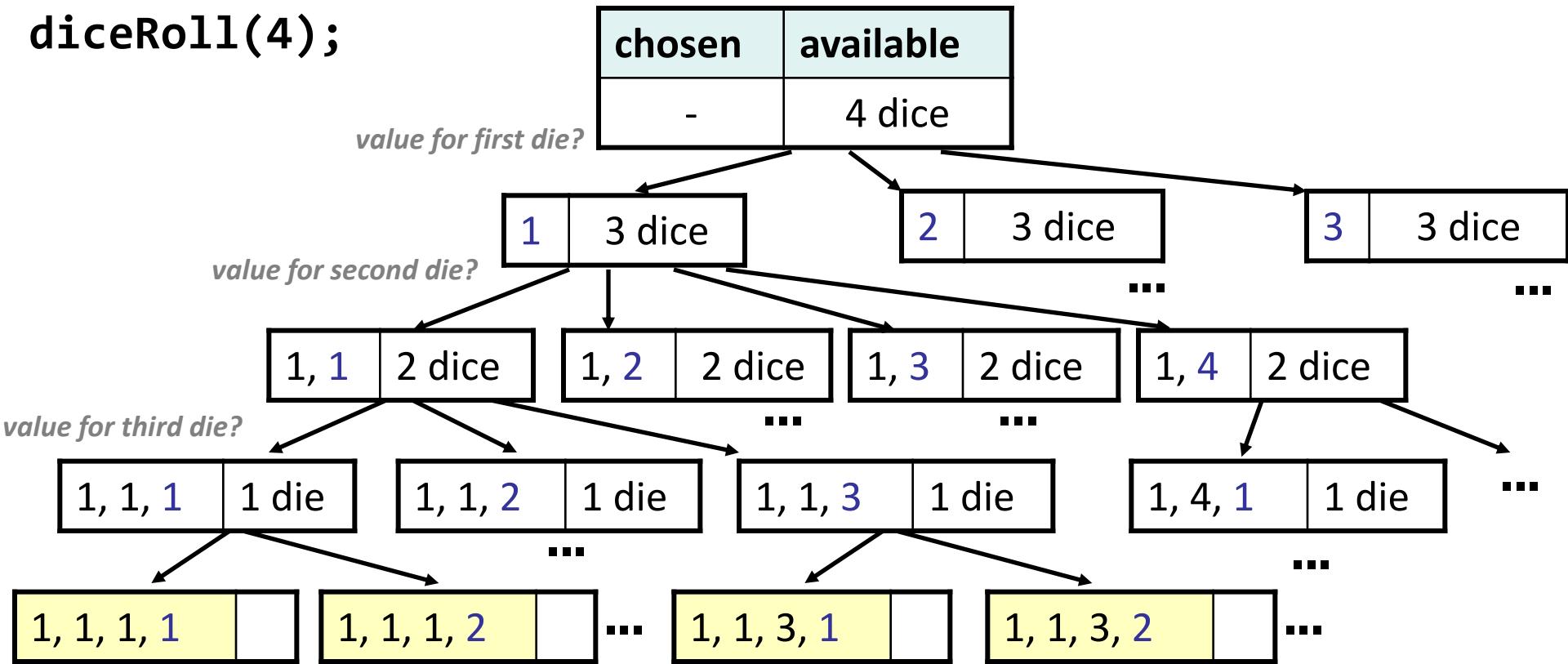
```
for (each possible first die value):  
    for (each possible second die value):  
        for (each possible third die value):  
            ...  
            print!
```



- This is called a **depth-first search**
- You may think that **for** loops are the way to approach this problem.
 - But how many loops are needed?
- How can we completely explore such a large search space?

A decision tree

`diceRoll(4);`



diceRolls pseudocode

```
function diceRolls(dice):  
    if dice == 0:  
        nothing to do.  
    else:  
        // handle all roll values for a single die; let recursion do the rest.  
        for each die value i in range [1..6]:  
            choose that the current die will have value i.  
            diceRolls(dice-1).      // explore the remaining dice.  
            un-choose the value i.
```

- How do we keep track of our choices?

diceRolls solution

```
// Prints all possible outcomes of rolling the given
// number of six-sided dice in {#, #, #} format.
void diceRolls(int dice) {
    Vector<int> chosen;
    diceRollHelper(dice, chosen);
}

// private recursive helper to implement diceRolls logic
void diceRollHelper(int dice, Vector<int>& chosen) {
    if (dice == 0) {
        cout << chosen << endl;                                // base case
    } else {
        for (int i = 1; i <= 6; i++) {
            chosen.add(i);                                      // choose
            diceRollHelper(dice - 1, chosen);                    // explore
            chosen.remove(chosen.size() - 1);                  // un-choose
        }
    }
}
```

Exercise: Dice roll sum

- Write a function **diceSum** similar to **diceRoll**, but it also accepts a desired sum and prints only combinations that add up to exactly that sum.

```
diceSum(2, 7);
```

```
{1, 6}  
{2, 5}  
{3, 4}  
{4, 3}  
{5, 2}  
{6, 1}
```



```
diceSum(3, 7);
```

```
{1, 1, 5}  
{1, 2, 4}  
{1, 3, 3}  
{1, 4, 2}  
{1, 5, 1}  
{2, 1, 4}  
{2, 2, 3}  
{2, 3, 2}  
{2, 4, 1}  
{3, 1, 3}  
{3, 2, 2}  
{3, 3, 1}  
{4, 1, 2}  
{4, 2, 1}  
{5, 1, 1}
```

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

- ❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).
- ❑ Undo it after exploring. Restore everything to the way it was before making this choice.
- ❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

- ❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

- ❑ Make it and explore recursively. Pass the information for a choice. When we have no dice left to choose, print them out *only if they equal our sum.*
- ❑ Undo it
- ❑ Find our base case(s). What should we do when we are out of decisions?

Minimal modification

```
void diceSum(int dice, int desiredSum) {
    Vector<int> chosen;
    diceSumHelper(dice, desiredSum, chosen);
}

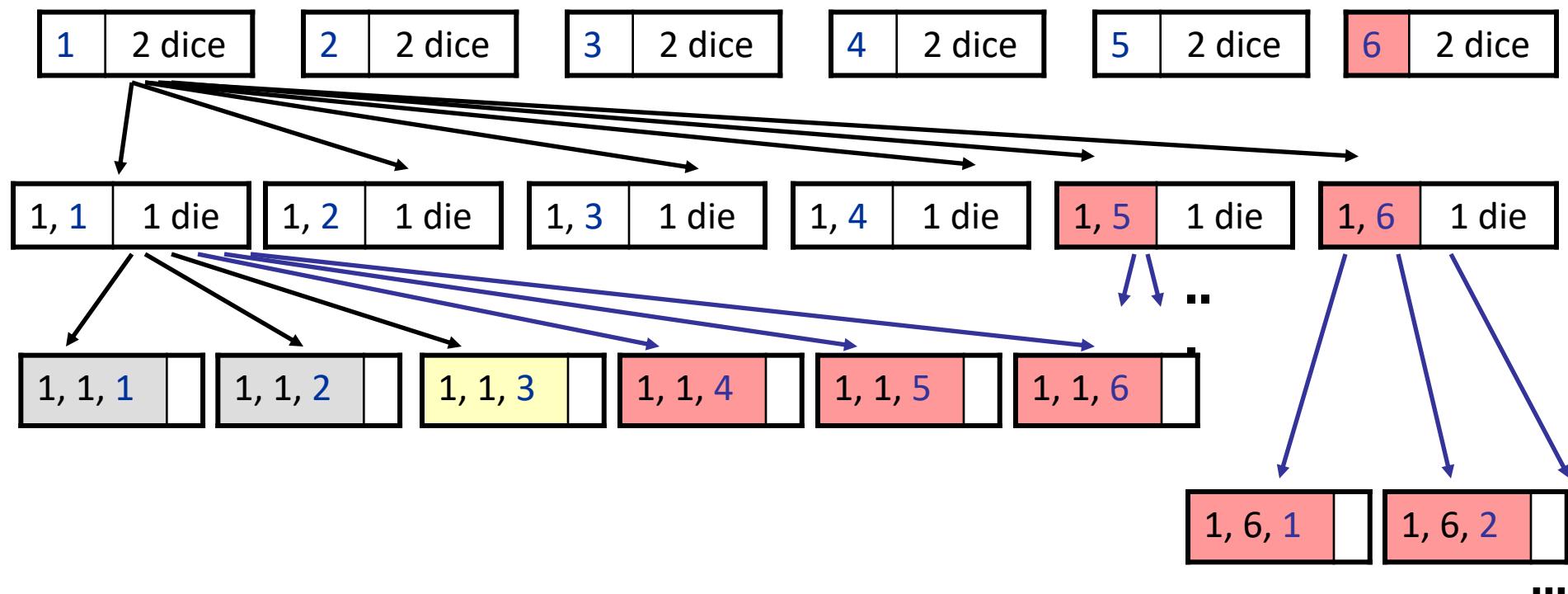
void diceSumHelper(int dice, int desiredSum, Vector<int>& chosen) {
    if (dice == 0) {
        if (sumAll(chosen) == desiredSum) {
            cout << chosen << endl; // base case
        }
    } else {
        for (int i = 1; i <= 6; i++) {
            chosen.add(i); // choose
            diceSumHelper(dice - 1, desiredSum, chosen); // explore
            chosen.remove(chosen.size() - 1); // un-choose
        }
    }
}

int sumAll(const Vector<int>& v) {
    int sum = 0;
    for (int k : v) { sum += k; }
    return sum;
}
```

Wasteful decision tree

`diceSum(3, 5);`

chosen	available	desired sum
-	3 dice	5



Optimizations

- We need not visit every branch of the decision tree.
 - Some branches are clearly not going to lead to success.
 - We can preemptively stop, or **prune**, these branches.
- Inefficiencies in our dice sum algorithm:
 - Sometimes the current sum is already **too high**.
 - (Even rolling 1 for all remaining dice would exceed the desired sum.)
 - Sometimes the current sum is already **too low**.
 - (Even rolling 6 for all remaining dice would fall short of the desired sum.)
 - The code must **re-compute** the sum many times.
 - ($1+1+1 = \dots, 1+1+2 = \dots, 1+1+3 = \dots, 1+1+4 = \dots, \dots$)

diceSum solution

```
void diceSum(int dice, int desiredSum) {  
    Vector<int> chosen;  
    diceSumHelper(dice, 0, desiredSum, chosen);  
}  
  
void diceSumHelper(int dice, int sum, int desiredSum, Vector<int>& chosen) {  
    if (dice == 0) {  
        if (sum == desiredSum) {  
            cout << chosen << endl; // base case  
        }  
    } else if (sum + 1*dice <= desiredSum  
              && sum + 6*dice >= desiredSum) {  
        for (int i = 1; i <= 6; i++) {  
            chosen.add(i); // choose  
            diceSumHelper(dice - 1, sum + i, desiredSum, chosen); // explore  
            chosen.remove(chosen.size() - 1); // un-choose  
        }  
    }  
}
```

For you to think about...

- How would you modify `diceSum` so that it prints only unique combinations of dice, ignoring order?
 - (e.g. don't print both `{1, 1, 5}` and `{1, 5, 1}`)

`diceSum2(2, 7);`

`{1, 6}`
`{2, 5}`
`{3, 4}`
`{4, 3}`
`{5, 2}`
`{6, 1}`



`diceSum2(3, 7);`

`{1, 1, 5}`
`{1, 2, 4}`
`{1, 3, 3}`
`{1, 4, 2}`
`{1, 5, 1}`
`{2, 1, 4}`
`{2, 2, 3}`
`{2, 3, 2}`
`{2, 4, 1}`
`{3, 1, 3}`
`{3, 2, 2}`
`{3, 3, 1}`
`{4, 1, 2}`
`{4, 2, 1}`
`{5, 1, 1}`

diceSum solution

```
void diceSum(int dice, int desiredSum) {  
    Vector<int> chosen;  
    diceSumHelper(dice, 0, desiredSum, chosen, 1);  
}  
  
void diceSumHelper(int dice, int sum, int desiredSum, Vector<int>& chosen,  
    int minDieValue) {  
    if (dice == 0) {  
        if (sum == desiredSum) {  
            cout << chosen << endl; // base case  
        }  
    } else if (sum + 1*dice <= desiredSum  
        && sum + 6*dice >= desiredSum) {  
        for (int i = minDieValue; i <= 6; i++) {  
            chosen.add(i); // choose  
            diceSumHelper(dice - 1, sum + i, desiredSum, chosen, i); // explore  
            chosen.remove(chosen.size() - 1); // un-choose  
        }  
    }  
}
```

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Announcements

- Pair Programming on HW4

- Pair programming means that two people work *together* on an assignment, completely.
- Pair programmers must never be working on the assignment independently, and should both be looking at the same screen, with one of the students typing (they should take turns).
- Students may ask conceptual questions in the LaIR and on Piazza independently, but if you are in a pair you must get help on the code together.
- If one student has taken the course before, there can be no overlapping code from that student's prior work unless they worked on it with the exact same people then as now.

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Exercise: subsets

printSubVectors

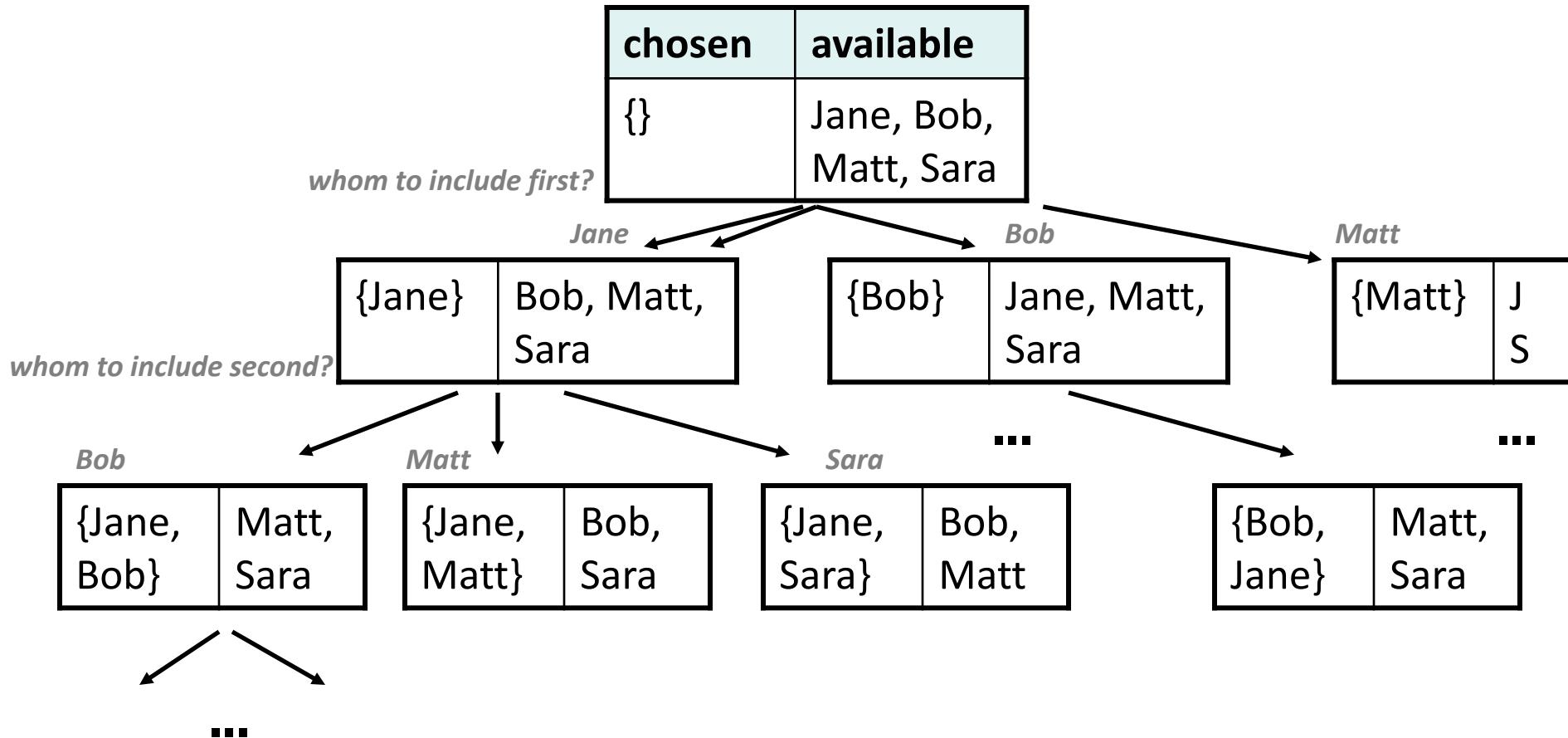
- Write a function **subsets** that finds every possible sub-list of a given vector. A sub-list of a vector V contains ≥ 0 of V 's elements.
 - Example: if V is {Jane, Bob, Matt, Sara}, then the call of **subsets(V)**; prints:

{Jane, Bob, Matt, Sara}
{Jane, Bob, Matt}
{Jane, Bob, Sara}
{Jane, Bob}
{Jane, Matt, Sara}
{Jane, Matt}
{Jane, Sara}
{Jane}

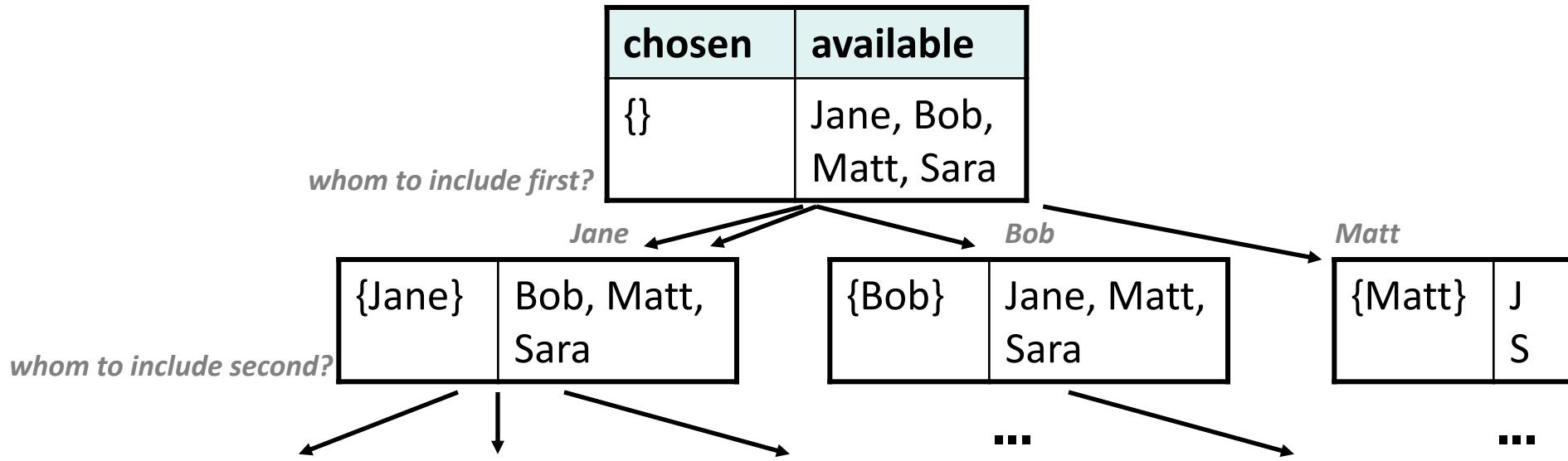
{Bob, Matt, Sara}
{Bob, Matt}
{Bob, Sara}
{Bob}
{Matt, Sara}
{Matt}
{Sara}
{}

- You can print the subsets out in any order, one per line.
 - *What are the "choices" in this problem? (choose, explore)*

Decision tree?



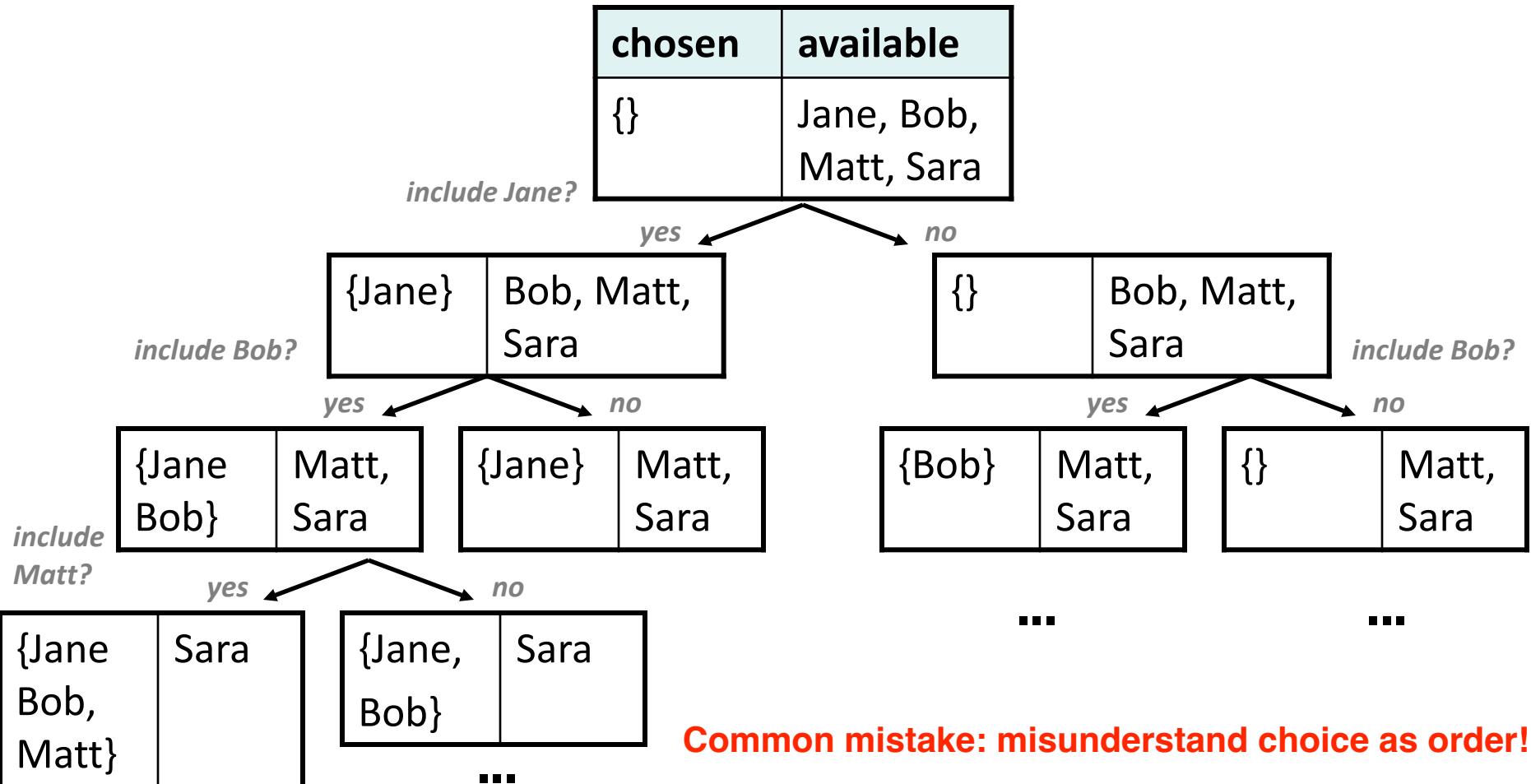
Wrong decision tree



Q: Why isn't this the right decision tree for this problem?

- A. It does not actually end up finding every possible subset.
- B. It does find all subset, but it finds them in the wrong order.
- C. It does find all subset, but it finds them multiple times.
- D. None of the above

Better decision tree



- Each decision is: "Include Jane or not?" ... "Include Bob or not?" ...
 - The **order** of people chosen does not matter; only the **membership**.

The Backtracking Checklist

- ❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each val

- Should I include the next person in the subset?**
- ❑ Make it for a choice to the next recursive call(s).
- ❑ Undo it after exploring. Restore everything to the way it was before making this choice.
- ❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

❑ **Make it and explore recursively.** Pass the information for a choice to the next recursive call(s).

❑ Undo it after exploring. Restore everything to the way it was before making this choice.

❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).

❑ Undo it : **Find all subsets with this choice made.** ; to the way it was before

❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).

❑ Undo it after exploring. Restore everything to the way it was before making this choice.

❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

- Find what choice(s) we have at each step. What different options are there? For each value:
 - Make sure the chosen set is the same as it was before.
 - Make it for a choice to the next recursive call(s).
- Undo it after exploring. Restore everything to the way it was before making this choice.
- Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

- ❑ Make it and explore recursively. Pass the information for a choice to the next recursive call(s).
- ❑ Undo it after exploring. Restore everything to the way it was before making this choice.
- ❑ Find our base case(s). What should we do when we are out of decisions?

The Backtracking Checklist

- ❑ Find what choice(s) we have at each step. What different options are there for the next step?

For each valid choice:

- ❑ Make it and explore recursively. Pass the information for a choice.
- ❑ Undo it
- ❑ Find our base case(s). What should we do when we are out of decisions?

When we have no choices left, print out the subset.

sublists solution

```
void subSets(const Set<string>& masterSet) {
    Set<string> chosen;
    listSubsetsRec(masterSet, chosen);
}

void listSubsetsRec(const Set<string>& masterSet, const Set<string>& used) {
    if (masterSet.isEmpty()) {
        cout << used << endl;
    } else {
        string element = masterSet.first();

        listSubsetsRec(masterSet - element, used);           // Without
        listSubsetsRec(masterSet - element, used + element); // With
    }
}
```

sublists solution

```
void subSets(Set<string>& masterSet) {
    Set<string> chosen;
    listSubsetsRec(masterSet, chosen);
}

void listSubsetsRec(Set<string>& masterSet, Set<string>& used) {
    if (masterSet.isEmpty()) {
        cout << used << endl;
    } else {
        string element = masterSet.first();

        masterSet.remove(element);
        listSubsetsRec(masterSet, used);           // Without

        used.add(element);
        listSubsetsRec(masterSet, used); // With
        masterSet.add(element);
        used.remove(element);
    }
}
```

Recap

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Next time: more backtracking

Overflow

Exercise: Permute Vector

- Write a function **permute** that accepts a Vector of strings as a parameter and outputs all possible rearrangements of the strings in that vector. The arrangements may be output in any order.
 - Example: if v contains {"a", "b", "c", "d"}, your function outputs these permutations:

{a, b, c, d}	{b, a, c, d}	{c, a, b, d}	{d, a, b, c}
{a, b, d, c}	{b, a, d, c}	{c, a, d, b}	{d, a, c, b}
{a, c, b, d}	{b, c, a, d}	{c, b, a, d}	{d, b, a, c}
{a, c, d, b}	{b, c, d, a}	{c, b, d, a}	{d, b, c, a}
{a, d, b, c}	{b, d, a, c}	{c, d, a, b}	{d, c, a, b}
{a, d, c, b}	{b, d, c, a}	{c, d, b, a}	{d, c, b, a}

Permute solution

```
// Outputs all permutations of the given vector.
void permute(Vector<string>& v) {
    Vector<string> chosen;
    permuteHelper(v, chosen);
}

void permuteHelper(Vector<string>& v, Vector<string>& chosen) {
    if (v.isEmpty()) {
        cout << chosen << endl;      // base case
    } else {
        for (int i = 0; i < v.size(); i++) {
            string s = v[i];
            v.remove(i);
            chosen.add(s);           // choose
            permuteHelper(v, chosen); // explore
            chosen.remove(chosen.size() - 1); // un-choose
            v.insert(i, s);
        }
    }
}
```