kernel hacker 修炼之道——李万鹏

男儿立志出乡关, 学不成名死不还。 埋骨何须桑梓地, 人生无处不青山。 ——西乡隆盛诗

Linux驱动修炼之道-SPI驱动框架源码分析(下)

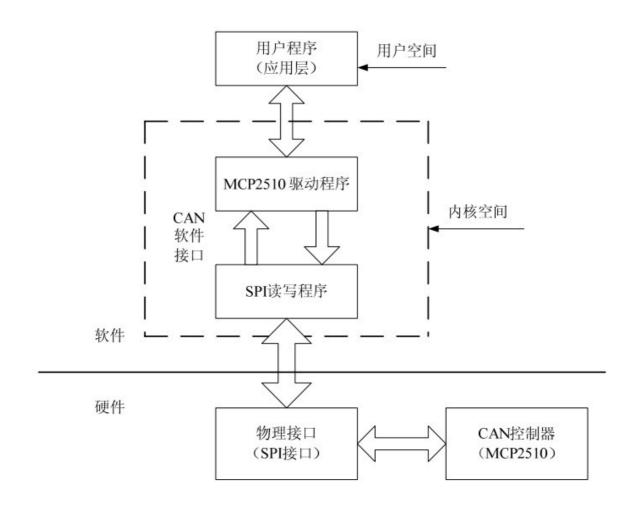
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这篇文档主要介绍spi数据传输过程。

当应用层要向设备传输数据的时候,会通过ioctl向设备驱动发送传输数据的命令。如图,向SPI从设备发送读写命令,实际的读写操作还是调用了主机控制器驱动的数据传输函数。transfer函数用于spi的IO传输。但是,transfer函数一般不会执行真正的传输操作,而是把要传输的内容放到一个队列里,然后调用一种类似底半部的机制进行真正的传输。这是因为,spi总线一般会连多个spi设备,而spi设备间的访问可能会并发。如果直接在transfer函数中实现传输,那么会产生竞态,spi设备互相间会干扰。所以,真正的spi传输与具体的spi控制器的实现有关,spi的框架代码中没有涉及。像spi设备的片选,根据具体设备进行时钟调整等等都在实现传输的代码中被调用。spi的传输命令都是通过结构spi_message定义,设备程序调用transfer函数将spi_message交给spi总线驱动,总线驱动再将message传到底半部排队,实现串行化传输。



在spidev.c中实现了file_operations:

```
static struct file_operations spidev_fops = {
.owner = THIS_MODULE,
.write = spidev write,
.read = spidev_read,
.unlocked ioctl = spidev ioctl,
.open = spidev_open,
.release = spidev release,
};这里看spidev_ioctl的实现:
static long
spidev_ioctl(struct file *filp, unsigned int cmd, unsigned long
arg)
{
int err = 0;
int retval = 0;
struct spidev data *spidev;
struct spi device *spi;
u32 tmp;
unsigned n ioc;
struct spi_ioc_transfer *ioc;
```

```
/*查看这个命令的幻数字段是否为'k'*/
if ( IOC TYPE(cmd) != SPI IOC MAGIC)
return -ENOTTY;
/*如果方向是用户空间从内核读,即内核向用户空间写,则检查用户空间的
地址是否有效*/
if (IOC DIR(cmd) & IOC READ)
err = !access ok(VERIFY WRITE,
(void user *) arg, IOC SIZE(cmd));
/*如果方向是用户空间向内核写,即内核读用户空间,则检查用户空间的地
址是否有效*/
if (err == 0 && IOC DIR(cmd) & IOC WRITE)
err = !access ok (VERIFY READ,
(void user *) arg, IOC SIZE(cmd));
if (err)
return -EFAULT;
/* guard against device removal before, or while,
* we issue this ioctl.
*/
spidev = filp->private_data;
spin_lock_irq(&spidev->spi lock);
spi = spi_dev_get(spidev->spi);
spin unlock irg(&spidev->spi lock);
if (spi == NULL)
return -ESHUTDOWN;
mutex lock(&spidev->buf lock);
switch (cmd) {
/* read requests */
case SPI IOC RD MODE:
/*因为已经进行了地址是否有效的检查,所以这里使用
put user, get user, copy from user可以节省几个时钟周期呢*/
retval = put user(spi->mode & SPI MODE MASK,
( u8 user *) arg);
break:
case SPI IOC RD LSB FIRST:
retval = put user((spi->mode & SPI LSB FIRST) ? 1 : 0,
( u8 user *) arg);
break:
case SPI IOC RD BITS PER WORD:
```

```
retval = __put_user(spi->bits_per_word, (__u8 __user *)arg);
break:
case SPI IOC RD MAX SPEED HZ:
retval = put user(spi->max speed hz, ( u32 user *) arg);
break;
/*设置SPI模式*/
case SPI_IOC_WR_MODE:
retval = __get_user(tmp, (u8 __user *)arg);
if (retval == 0) {
/*先将之前的模式保存起来,一旦设置失败进行回复*/
u8 save = spi->mode;
if (tmp & ~SPI MODE MASK) {
retval = -EINVAL;
break;
tmp |= spi->mode & ~SPI_MODE_MASK;
spi-mode = (u8) tmp:
retval = spi setup(spi);
if (retval < 0)
spi->mode = save;
else
dev dbg(&spi->dev, "spi mode %02x\n", tmp);
}
break:
case SPI_IOC_WR_LSB_FIRST:
retval = __get_user(tmp, (__u8 __user *)arg);
if (retval == 0) {
u8 save = spi->mode;
if (tmp)
spi->mode |= SPI LSB FIRST;
else
spi->mode &= ~SPI LSB FIRST;
retval = spi setup(spi);
if (retval < 0)
spi->mode = save;
dev_dbg(&spi->dev, "%csb first\n",
tmp ? '1' : 'm');
}
break;
```

```
case SPI_IOC_WR_BITS_PER_WORD:
retval = __get_user(tmp, (__u8 __user *)arg);
if (retval == 0) {
u8 save = spi->bits per word;
spi->bits_per_word = tmp;
retval = spi setup(spi);
if (retval < 0)
spi->bits per word = save;
dev dbg(&spi->dev, "%d bits per word\n", tmp);
}
break;
case SPI_IOC_WR_MAX_SPEED_HZ:
retval = __get_user(tmp, (__u32 __user *)arg);
if (retval == 0) {
u32 save = spi->max speed hz;
spi->max speed hz = tmp;
retval = spi setup(spi);
if (retval < 0)
spi->max_speed_hz = save;
else
dev_dbg(&spi->dev, "%d Hz (max)\n", tmp);
break;
default:
/* segmented and/or full-duplex I/O request */
if (_IOC_NR(cmd) != _IOC_NR(SPI_IOC_MESSAGE(0))
| IOC_DIR(cmd) != IOC_WRITE) {
retval = -ENOTTY;
break;
/*得到用户空间数据的大小*/
tmp = IOC SIZE(cmd);
/*如果这些数据不能分成spi ioc transfer的整数倍,则不能进行传
输, spi io transfer是对spi transfer的映射*/
if ((tmp % sizeof(struct spi ioc transfer)) != 0) {
retval = -EINVAL;
break;
/*计算出能分多少个spi_ioc transfer*/
n ioc = tmp / sizeof(struct spi ioc transfer);
```

```
if (n ioc == 0)
break:
/*在内核中分配装载这些数据的内存空间*/
ioc = kmalloc(tmp, GFP KERNEL);
if (!ioc) {
retva1 = -ENOMEM:
break;
/*把用户空间的数据拷贝过来*/
if (copy from user(ioc, (void user *) arg, tmp)) {
kfree(ioc):
retval = -EFAULT;
break;
}
/*进行数据传输*/
retval = spidev message(spidev, ioc, n ioc);
kfree(ioc);
break:
mutex unlock(&spidev->buf lock);
spi_dev_put(spi);
return retval:
}
下面跟踪spidev_message看看: static int spidev_message(struct
spidev data *spidev,
struct spi_ioc_transfer *u_xfers, unsigned n_xfers)
struct spi_message msg;
struct spi transfer *k xfers;
struct spi transfer *k tmp;
struct spi ioc transfer *u tmp;
unsigned n, total;
u8 *buf;
int status = -EFAULT;
/*初始化spi message的tranfers链表头*/
spi message init(&msg);
/*分配n个spi transfer的内存空间,一个spi message由多个数据段
spi message组成*/
k xfers = kcalloc(n xfers, sizeof(*k tmp), GFP KERNEL);
if (k xfers == NULL)
```

```
return -ENOMEM;
buf = spidev->buffer;
total = 0;
/*这个for循环的主要任务是将所有的spi transfer组装成一个
spi message*/
for (n = n xfers, k tmp = k xfers, u tmp = u xfers;
n:
n--, k tmp++, u tmp++) {
/*u tmp是从用户空间传下来的spi ioc message的大小, spi ioc message是
对spi message的映射*/
k \text{ tmp-} > 1 \text{en} = u \text{ tmp-} > 1 \text{en}:
/*统计要传输数据的总量*/
total += k tmp \rightarrow len;
if (total > bufsiz) {
status = -EMSGSIZE;
goto done;
/*spi transfer是一个读写的buffer对,如果是要接收则把buffer给接收的
rx buf*/
if (u tmp->rx buf) {
k \text{ tmp-} > rx \text{ buf = buf};
if (!access ok(VERIFY WRITE, (u8 user *)
(uintptr_t) u_tmp->rx_buf,
u tmp->len))
goto done;
/*如果要传输,这个buffer给tx buf使用,从用户空间拷过来要传输的数据
*/
if (u tmp \rightarrow tx buf) {
k \text{ tmp->}tx \text{ buf = buf;}
if (copy from user(buf, (const u8 user *)
(uintptr t) u tmp->tx buf,
u tmp->len))
goto done;
}
/*指向下一段内存*/
buf += k tmp \rightarrow len;
/*最后一个transfer传输完毕是否会影响片选*/
k tmp->cs change = !!u tmp->cs change;
/*每字长的字节数*/
k tmp->bits per word = u tmp->bits per word;
/*一段数据传输完需要一定的时间等待*/
k tmp->delay usecs = u tmp->delay usecs;
```

```
/*初始化传输速度*/
k_{tmp} > speed_hz = u_tmp > speed_hz:
/*将spi_transfer通过它的transfer_list字段挂到spi_message的transfer
队列上*/
spi message add tail(k tmp, &msg);
/*调用底层的传输函数*/
status = spidev_sync(spidev, &msg);
if (status < 0)
goto done;
/* copy any rx data out of bounce buffer */
buf = spidev->buffer;
/*把传输数据拷贝到用户空间打印出来,可以查看是否传输成功*/
for (n = n\_xfers, u\_tmp = u\_xfers; n; n--, u\_tmp++) {
if (u tmp->rx buf) {
if (_copy_to_user((u8 __user *)
(uintptr t) u tmp->rx buf, buf,
u tmp \rightarrow len) {
status = -EFAULT:
goto done;
buf += u_tmp->len;
status = total;
done:
kfree(k xfers);
return status:
看spidev sync的实现:
static ssize t
spidev sync(struct spidev data *spidev, struct spi message
*message)
{
/*声明并初始化一个完成量*/
DECLARE COMPLETION ONSTACK (done);
int status;
/*指定spi message使用的唤醒完成量函数*/
message->complete = spidev complete;
message->context = &done;
```

```
spin lock irg(&spidev->spi lock);
if (spidev->spi == NULL)
status = -ESHUTDOWN;
else
/*调用spi核心中的函数进行数据传输*/
status = spi async(spidev->spi, message);
spin_unlock_irq(&spidev->spi_lock);
if (status == 0) {
/*等待完成量被唤醒*/
wait_for_completion(&done);
status = message->status;
if (status == 0)
status = message->actual_length;
return status;
spi async在spi.h中定义的:
static inline int
spi_async(struct spi_device *spi, struct spi_message *message)
message->spi = spi;
return spi->master->transfer(spi, message);
这里的master->transfer是在spi bitbang start中进行赋值的:
bitbang->master->transfer= spi_bitbang_transfer;
看spi bitbang transfer的实现: int spi bitbang transfer(struct
spi_device *spi, struct spi_message *m)
{
struct spi bitbang *bitbang;
unsigned long flags;
int status = 0;
m-actual length = 0;
m->status = -EINPROGRESS;
/*在spi alloc master函数中调用spi master set devdata把struct
s3c24xx spi结构存放起来,而struct spi_bitbang正是struct s3c24xx_spi
结构所包含的第一个结构*/
bitbang = spi master get devdata(spi->master);
spin lock irqsave (&bitbang->lock, flags);
```

```
if (!spi->max_speed_hz)
status = -ENETDOWN:
else {
/*把message加入到bitbang的等待队列中*/
list add tail(&m->queue, &bitbang->queue);
/*把bitbang-work加入bitbang->workqueue中,调度运行*/
queue work(bitbang->workqueue, &bitbang->work);
spin unlock irgrestore(&bitbang->lock, flags);
return status;
EXPORT SYMBOL GPL(spi bitbang transfer);
分析工作队列的处理函数:
                                    daemon0
  SPI 控制器 0
                    cpu workq
  spi bitbang 的
                    ueue struct
                                     work
                                               work
                                                        (work
  workqueue
                                    daemonl
  SPI 控制器 1
                    cpu_workq
  spi bitbang 的
                    ueue struct
                                    work
                                               work
                                                         work
  workqueue
static void bitbang work(struct work struct *work)
struct spi bitbang *bitbang =
container of (work, struct spi bitbang, work);
unsigned long flags;
spin lock irqsave (&bitbang->lock, flags);
/*设置成忙状态*/
bitbang->busy = 1;
/*对bitqueue中的每一个spi message进行处理*/
while (!list empty(&bitbang->queue)) {
struct spi message *m;
struct spi device *spi;
unsigned nsecs;
```

```
struct spi transfer *t = NULL;
unsigned tmp:
unsigned cs change;
int status;
int (*setup transfer) (struct spi device *,
struct spi_transfer *);
m = container_of(bitbang->queue.next, struct spi_message,
queue):
/*从队列中驱动这个spi message*/
list del init(&m->queue);
spin unlock irgrestore (&bitbang->lock, flags);
nsecs = 100;
spi = m->spi;
tmp = 0;
cs change = 1;
status = 0;
setup transfer = NULL:
/*对spi message的transfers上的每个spi transfer进行处理*/
list_for_each_entry (t, &m->transfers, transfer_list) {
0 0 0 0 0 0 0 0 0 0 0 0 0 0
if (t->len) {
if (!m->is dma mapped)
t\rightarrow rx dma = t\rightarrow tx dma = 0;
/*调用bitbang->txrx bufs进行数据的传输, bitbang->txrx_bufs =
s3c24xx spi txrx;这个在s3c24xx spi probe中进行赋值的*/
status = bitbang->txrx bufs(spi, t);
0 0 0 0 0 0 0 0 0 0 0 0 0 0
m->status = status;
/*传输完成,唤醒刚才的那个完成变量*/
m->complete(m->context);
/* restore speed and wordsize */
if (setup transfer)
setup_transfer(spi, NULL);
if (!(status == 0 && cs change)) {
ndelay(nsecs);
bitbang->chipselect(spi, BITBANG CS INACTIVE);
ndelay(nsecs);
```

```
spin lock irqsave (&bitbang->lock, flags);
bitbang \rightarrow busy = 0;
spin unlock irgrestore (&bitbang->lock, flags);
这个工作队列的处理函数中调用了spi controller driver中的传输函数:
static int s3c24xx spi txrx(struct spi device *spi, struct
spi_transfer *t)
struct s3c24xx_spi *hw = to_hw(spi);
dev_dbg(&spi->dev, "txrx: tx %p, rx %p, len %d\n",
t->tx buf, t->rx buf, t->len);
hw->tx = t->tx buf; //发送指针
hw->rx = t->rx buf; //接收指针
hw->len = t->len; //需要发送/接收的数目
hw->count = 0; //存放实际spi传输的数据数目
/*初始化了完成量*/
init completion(&hw->done);
/*
*只需发送第一个字节(如果发送为空,则发送0xff),中断中就会自动发送完
其他字节(并接受数据)
*直到所有数据发送完毕且所有数据接收完毕才返回
*/
writeb(hw txbyte(hw, 0), hw->regs + S3C2410 SPTDAT);
/*等待完成量被唤醒*/
wait for completion(&hw->done);
return hw->count;
static inline unsigned int hw_txbyte(struct s3c24xx_spi *hw, int
count)
return hw->tx ? hw->tx[count] : 0xff;
//如果还有数据没接收完且要发送的数据经已发送完毕,发送空数据0xFF
下面来分析中断函数:
static irgreturn t s3c24xx spi irg(int irg, void *dev)
struct s3c24xx spi *hw = dev;
/*读取spi的状态寄存器*/
unsigned int spsta = readb(hw->regs + S3C2410 SPSTA);
```

```
unsigned int count = hw->count;
/*检测冲突*/
if (spsta & S3C2410 SPSTA DCOL) {
dev dbg(hw->dev, "data-collision\n");
/*唤醒完成量*/
complete(&hw->done);
goto irq done;
/*设备忙*/
if (!(spsta & S3C2410_SPSTA READY)) {
dev dbg(hw->dev, "spi not ready for tx?\n");
/*唤醒完成量*/
complete(&hw->done);
goto irq done:
hw->count++;
/*接收数据*/
if (hw->rx)
hw\rightarrow rx[count] = readb(hw\rightarrow regs + S3C2410 SPRDAT);
count++:
/*如果count小于需要发送或接收数据的数目,发送其他数据*/
if (count < hw->len)
writeb(hw txbyte(hw, count), hw->regs + S3C2410_SPTDAT);
else
/*唤醒完成量,通知s3c24xx spi txrx函数*/
complete(&hw->done);
irq done:
return IRQ HANDLED;
}
至此spi数据传输过程完成,如果不想为自己的SPI设备写驱动,那么可以用
Linux自带的spidev. c提供的驱动程序,只要在登记时,把设备名设置成
spidev就可以了。spidev. c会在device目录下自动为每一个匹配的SPI设备创
建设备节点,节点名"spi%d"。之后,用户程序可以通过字符型设备的通用接
口控制SPI设备。需要注意的是,spidev创建的设备在设备模型中属于虚拟设
备,他的class是spidev class,他的父设备是在boardinfo中定义的spi设
备。
```