

International Institute of Information Technology - Hyderabad

lazy three

Mohammed Faisal, Harshvardhan Rana, Tanay Gad

DLOCAL A.cpp"

Contest (1)

template.h

```
8c379d, 55 lines
// #ifdef LOCAL
// #include "include/include.h"
// #else
// #include <bits/stdc++.h>
// #include <ext/pb ds/assoc container.hpp>
// #endif
// #pragma GCC target("bmi,bmi2,lzcnt,popcnt")
// #pragma GCC optimize("02,unroll-loops")
// #pragma GCC target("avx2")
//#pragma GCC optimize("02")
//#pragma GCC optimize("Ofast")
//#pragma GCC target("avx,avx2,fma")
using namespace std;
using namespace __gnu_pbds;
typedef tree<int, null type, less<int>, rb tree tag,
tree order statistics node update> o set;
// order_of_key (val): returns the no. of values less than
↓ val
// find_by_order (k): returns the kth largest element.(0-
based)
template <typename T>
using minHeap = priority queue<T, vector<T>, greater<T>>;
template <typename T>
using maxHeap = priority queue<T>;
#define int long long
#define all(s) s.begin(), s.end()
#define sz(s) (int)s.size()
#define testcases \
  cin >> tt;
 for (i = 1; i <= tt; i++)
#define fast
  ios_base::sync_with_stdio(0); \
  cin.tie(NULL);
```

```
cout.tie(NULL)
#define deb(a) cerr << #a << " = " << (a) << endl;
#define deb1(a)
  cerr << #a << " = [ ";</pre>
  for (auto it = a.begin(); it != a.end(); it++) cerr << *</pre>

it << " "; \
</pre>
 cerr << "]" << endl;</pre>
typedef vector<int> vi;
typedef vector<vector<int>> vvi;
typedef pair<int, int> pii;
typedef vector<pair<int, int>> vpii;
typedef vector<bool> vb;
const int INF = LONG LONG MAX;
const int M = 1e9 + 7;
void solve(int tt) {
int32 t main() {
 fast:
  int tt = 1;
 int i = 1;
 testcases
      solve(i);
flags.txt
                                                 425523, 4 lines
## Put this in end of ~/.bashrc
alias s="g++ -Wall -Wextra -pedantic -std=c++17 -02 -
Wshadow -Wformat=2 -Wfloat-equal -Wconversion -Wlogical-
bwcast-align -D GLIBCXX_DEBUG -D_GLIBCXX_DEBUG_PEDANTIC -
D FORTIFY SOURCE=2 -fsanitize=address -fsanitize=
bundefined -fno-sanitize-recover -fstack-protector -
```

```
alias f="g++ -std=c++17 -O2 -DLOCAL A.cpp"

alias d="g++ -std=c++17 -g -Wall -Wextra -pedantic -

\[ \Perc \width{W} \text{Shadow} - \width{W} \text{Inon-unused-parameter} - \width{W} \text{onversion} - \width{W} \text{logical-op} -

\[ \perc \width{W} \text{Shift-overflow=2} - \width{W} \text{duplicated-cond} - \width{W} \text{cast-qual} - \width{W} \text{cast-qual} - \width{W} \text{cast-qual} - \width{W} \text{cast-qual} - \width{D} \text{LIBCXX_DEBUG_PEDANTIC} -

\[ \perc \D_FORTIFY_SOURCE=2 - \text{fsanitize=address} - \text{fsanitize=} \]

\[ \perc \text{undefined} - \text{fno-sanitize-recover} - \text{fstack-protector} -

\[ \perc \DLOCAL - \text{fdiagnostics-color=always} \]
```

Mathematics (2)

Data structures (3)

SegmentTree.h

7660bf, 28 lines

}

void rebuild(int i) {

```
template <typename T, typename F>
struct SegTree {
 int n;
 vector<T> t;
 const T id;
 F f:
  SegTree (const vector<T> &a, T id, F f) : n(sz(a)), t(2 *
  \downarrow n), id(id), f(f) {
   for (int i = 0; i < n; i++) t[n + i] = a[i];
    for (int i = n - 1; i >= 1; i--)
      t[i] = f(t[2 * i], t[2 * i + 1]);
 T query(int 1, int r) {
    T resl(id), resr(id);
    for (1 += n, r += n; 1 <= r; 1 >>= 1, r >>= 1) {
     if (1 == r) {
        res1 = f(res1, t[1]);
        break:
      if (1 & 1) resl = f(resl, t[1++]);
      if (!(r & 1)) resr = f(t[r--], resr);
```

```
return f(resl, resr);
 }
 void update(int v, T value) {
    for (t[v += n] = value; v >>= 1;)
      t[v] = f(t[2 * v], t[2 * v + 1]);
 }
};
LazySegTree.h
                                                   9376a1, 90 lines
template <typename T, typename U>
struct seg tree lazy {
 int S, H;
  T zero;
  vector<T> value;
 U noop;
  vector<bool> dirty;
  vector<U> prop;
  seg_tree_lazy(int _S, T _zero = T(), U _noop = U()) {
    zero = _zero, noop = _noop;
    for (S = 1, H = 1; S < _S;) S *= 2, H++;
    value.resize(2 * S, zero);
    dirty.resize(2 * S, false);
    prop.resize(2 * S, noop);
  void set leaves(vector<T> &leaves) {
    copy(leaves.begin(), leaves.end(), value.begin() + S);
    for (int i = S - 1; i > 0; i--)
      value[i] = value[2 * i] + value[2 * i + 1];
  void apply(int i, U &update) {
    value[i] = update(value[i]);
    if (i < S) {
      prop[i] = prop[i] + update;
      dirty[i] = true;
```

```
for (int 1 = i / 2; 1; 1 /= 2) {
    T combined = value [2 * 1] + value [2 * 1 + 1];
    value[1] = prop[1] (combined);
 }
}
void propagate(int i) {
  for (int h = H; h > 0; h--) {
   int 1 = i \gg h;
   if (dirty[1]) {
      apply(2 * 1, prop[1]);
      apply(2 * 1 + 1, prop[1]);
      prop[1] = noop;
      dirty[1] = false;
    }
  }
void upd(int i, int j, U update) {
  i += S, i += S;
  propagate(i), propagate(j);
  for (int 1 = i, r = j; 1 \le r; 1 \ne 2, r \ne 2) {
  if ((1 & 1) == 1) apply(1++, update);
   if ((r \& 1) == 0) apply(r--, update);
  rebuild(i), rebuild(j);
T query(int i, int j) {
  i += S, j += S;
  propagate(i), propagate(j);
  T res left = zero, res right = zero;
  for (; i <= j; i /= 2, j /= 2) {
   if ((i & 1) == 1) res_left = res_left + value[i++];
   if ((j & 1) == 0) res right = value[j--] + res right
    ٠,
  }
```

```
return res left + res right;
 }
};
struct node {
 int sum, width;
 node operator+(const node &n) {
   // Change 1
   return {sum + n.sum, width + n.width};
 }
};
struct update {
 bool type; // 0 for add, 1 for reset
 int value;
 node operator()(const node &n) { // apply update on n
   // Change 2
   if (type)
      return {n.width * value, n.width};
      return {n.sum + n.width * value, n.width};
  update operator+(const update &u) { // u is the recent
  \update, *this is the older update
   // Change 3
   if (u.type) return u;
    return {type, value + u.value};
 }
};
```

RMQ.h

d15d6d, 16 lines

```
jmp[k][j] = min(jmp[k-1][j], jmp[k-1][j+pw])
   }
  }
 T query(int a, int b) {
    assert (a \leq b); // tie(a, b) = minimax(a, b)
    int dep = 31 - builtin clz(b-a+1);
    return min(jmp[dep][a], jmp[dep][b - (1 << dep) + 1]);
 }
};
```

FenwickTree.h 6856de, 41 lines

```
template <typename T>
struct Fenwick {
 vector<T> bit;
 vector<T>& original;
 Fenwick(vector<T>& _arr) : bit(_arr.size(), OLL),

original(_arr) {
   int n = sz(arr);
   for (int i = 0; i < n; i++) {
    bit[i] = bit[i] + _arr[i];
     if ((i | (i + 1)) < n) bit [(i | (i + 1))] = bit [(i | const int mxn = 1000;
      (i + 1))] + bit[i];
  }
  // returns smallest index i, st. sum[0..i] >= x, returns
  \rightarrow -1 if no such i exists
 // returns n if x >= sum of array
  // ASSUMES NON NEGATIVE ENTRIES IN TREE
 int lower bound(int x) {
   if (x < 0) return -1;
   if (x == 0) return 0;
    int pos = 0;
    for (int pw = 1LL << 20; pw; pw >>= 1)
     if (pw + pos \le sz(bit)) and bit[pos + pw - 1] < x)
        pos += pw, x -= bit[pos - 1];
    return pos;
```

```
T query(int r) {
   assert(r < sz(bit));</pre>
   int ret = 0;
   for (r++; r > 0; r &= r - 1) ret += bit[r - 1];
   return ret:
 T query(int 1, int r) {
   T ret = query(r);
   if (1 != 0) ret -= query(1 - 1);
   return ret;
 void update(int i, int x) {
   int n = bit.size();
   T diff = x - original[i];
   original[i] = x;
   for (; i < n; i = i | i + 1) bit[i] += diff;
};
```

FenwickTree2d.h

deb44b, 17 lines

```
int grid[mxn + 1][mxn + 1];
int bit[mxn + 1][mxn + 1];
void update(int row, int col, int d) {
 grid[row][col] += d;
 for (int i = row; i \le mxn; i += (i \& -i))
   for (int j = col; j <= mxn; j += (j & -j))
     bit[i][j] += d;
int sum(int row, int col) {
 // calculates sum from [1,1] till [row,col]
 int res = 0;
 for (int i = row; i > 0; i -= (i & -i))
   for (int j = col; j > 0; j -= (j \& -j))
      res += bit[i][j];
  return res:
```

```
}
```

DSU.h

524177, 23 lines

```
struct DSU {
  int n;
 vector<int> parent;
  vector<int> size;
  DSU(int n): n(n), parent(n), size(n, 1) { iota(parent
  begin(), parent.end(), 0); }
  int find set(int x) {
    if (parent[x] == x) return x;
    return parent[x] = find set(parent[x]);
  }
  int getSize(int x) { return size[find set(x)]; } //
  returns size of component of x
  void union sets(int x, int y) {
    x = find set(x);
   y = find_set(y);
    if (x == v) return;
    if (size[x] > size[y]) {
     parent[y] = x;
      size[x] += size[y];
   } else {
      parent[x] = y;
      size[v] += size[x];
    }
  }
};
```

Mos.h

1cb2c0, 34 lines

```
int mblock = 1 / BLOCK, oblock = o.1 / BLOCK;
    return (mblock < oblock) or</pre>
             (mblock == oblock and mblock % 2 == 0 and r < o
             \rightarrow.r) or
             (mblock == oblock and mblock % 2 == 1 and r > o
             };
// Solve
void solve() {
  vector<Query> queries;
  queries.reserve(q);
  for (int i = 0; i < q; i++) {
    int 1, r; cin >> 1 >> r;
    1--, r--;
    queries.emplace back(l, r, i);
  sort (all (queries));
  int ans = 0;
  auto add = [&](int v) {};
  auto rem = [&](int v) {};
  vector<int> out (q); // Change out type if necessary
  int cur_1 = 0, cur_r = -1;
  for (auto &[l, r, id] : queries) {
    while (\operatorname{cur} 1 > 1) add(--\operatorname{cur} 1);
    while (\operatorname{cur} 1 < 1) \operatorname{rem}(\operatorname{cur} 1 + +);
    while (cur r < r) add(++cur r);</pre>
    while (cur r > r) rem(cur r--);
    out[id] = ans;
```

Graph (4)

4.1 Network flow MinCostMaxFlow.h

9a3690, 177 lines

IIIT-H

```
template <const int MAX N, typename flow t, typename
Gost t, flow t FLOW INF, cost t COST INF, const int
\RightarrowSCALE = 16>
struct CostScalingMCMF {
#define sz(a) a.size()
#define zero stl(v, sz) fill(v.begin(), v.begin() + (sz),
Ļ0)
 struct Edge {
    int v:
    flow t c;
    cost t d;
    int r:
    Edge() = default;
    Edge(int v, flow t c, cost t d, int r) : v(v), c(c), d
    \rightarrow (d), \mathbf{r}(\mathbf{r}) {}
  };
  vector<Edge> g[MAX N];
  cost t negativeSelfLoop;
  array<cost_t, MAX_N> pi, excess;
  array<int, MAX_N> level, ptr;
  CostScalingMCMF() { negativeSelfLoop = 0; }
  void clear() {
    negativeSelfLoop = 0;
    for (int i = 0; i < MAX_N; i++) g[i].clear();</pre>
  void addEdge(int s, int e, flow t cap, cost t cost) {
   if (s == e) {
      if (cost < 0) negativeSelfLoop += cap * cost;</pre>
      return:
    }
    g[s].push back(Edge(e, cap, cost, sz(g[e])));
    q[e].push back(Edge(s, 0, -cost, sz(q[s]) - 1));
  flow t getMaxFlow(int V, int S, int T) {
    auto BFS = [&]() {
      zero stl(level, V);
      queue<int> q;
```

```
q.push(S);
    level[S] = 1;
    for (q.push(S); !q.empty(); q.pop()) {
      int v = q.front();
      for (const auto &e : g[v])
        if (!level[e.v] && e.c) q.push(e.v), level[e.v]
        \rightarrow = level[v] + 1;
    return level[T];
  };
  function<flow t(int, flow t)> DFS = [&](int v, flow t
  ∮f1) {
    if (v == T || fl == 0) return fl;
    for (int &i = ptr[v]; i < (int)g[v].size(); i++) {</pre>
      Edge &e = q[v][i];
      if (level[e.v] != level[v] + 1 || !e.c) continue;
      flow t delta = DFS(e.v, min(fl, e.c));
      if (delta) {
        e.c -= delta;
        g[e.v][e.r].c += delta;
        return delta;
    return flow_t(0);
  flow t maxFlow = 0, tmp = 0;
  while (BFS()) {
    zero stl(ptr, V);
    while ((tmp = DFS(S, FLOW INF))) maxFlow += tmp;
  return maxFlow;
pair<flow t, cost t> maxflow(int N, int S, int T) {
  flow t maxFlow = 0;
  cost t eps = 0, minCost = 0;
  stack<int, vector<int>> stk;
  auto c_pi = [&](int v, const Edge &edge) { return edge
```

```
auto push = [&](int v, Edge &edge, flow_t delta, bool
∫flaq) {
 delta = min(delta, edge.c);
 edge.c -= delta;
 g[edge.v][edge.r].c += delta;
 excess[v] -= delta;
  excess[edge.v] += delta;
 if (flag && 0 < excess[edge.v] && excess[edge.v] <=</pre>

delta) stk.push(edge.v);
};
auto relabel = [&](int v, cost t delta) { pi[v] -=

delta + eps; };
auto lookAhead = [&](int v) {
 if (excess[v]) return false;
 cost t delta = COST INF;
 for (auto &e : q[v]) {
   if (e.c <= 0) continue;</pre>
   cost_t cp = c_pi(v, e);
   if (cp < 0)
     return false:
   else
      delta = min(delta, cp);
  relabel(v, delta);
  return true;
};
auto discharge = [&](int v) {
  cost t delta = COST INF;
 for (int i = 0; i < sz(g[v]); i++) {
   Edge &e = q[v][i];
   if (e.c <= 0) continue;</pre>
   cost t cp = c pi(v, e);
   if (cp < 0) {
     if (lookAhead(e.v)) {
        i--;
        continue:
```

```
push(v, e, excess[v], true);
         if (excess[v] == 0) return;
       } else
         delta = min(delta, cp);
     relabel(v, delta);
     stk.push(v);
  zero stl(pi, N);
   zero stl(excess, N);
  for (int i = 0; i < N; i++)
     for (auto &e : q[i]) minCost += e.c * e.d, e.d *=
     \rightarrowMAX N + 1, eps = max(eps, e.d);
  maxFlow = getMaxFlow(N, S, T);
  while (eps > 1) {
     eps /= SCALE;
     if (eps < 1) eps = 1;
     stk = stack<int, vector<int>>();
     for (int \mathbf{v} = 0; \mathbf{v} < \mathbf{N}; \mathbf{v} + +)
       for (auto &e : q[v])
         if (c_pi(v, e) < 0 && e.c > 0) push(v, e, e.c,
         \false):
     for (int \mathbf{v} = 0; \mathbf{v} < \mathbf{N}; \mathbf{v}^{++})
       if (excess[v] > 0) stk.push(v);
     while (stk.size()) {
       int top = stk.top();
       stk.pop();
       discharge(top);
  for (int v = 0; v < N; v++)
    for (auto &e : q[v]) e.d /= MAX N + 1, minCost -= e.
     \c * e.d;
  minCost = minCost / 2 + negativeSelfLoop;
  return {maxFlow, minCost};
}
```

```
} ;
void solve() {
  CostScalingMCMF<102, int, int, 100, 100> flow;
 int n, m;
  cin >> n >> m;
  int start = 0;
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < m; j++) {
     int inp;
      cin >> inp;
     if (inp) {
        flow.addEdge(i + 1, n + 1 + j, 1, 0);
        start++;
     } else
        flow.addEdge(i + 1, n + 1 + j, 1, 1);
    }
  }
  int counta = 0, countb = 0;
  for (int i = 0; i < n; i++) {
    int inp;
    cin >> inp;
    counta += inp;
    flow.addEdge(0, i + 1, inp, 0);
  for (int i = 0; i < m; i++) {
    int inp;
    cin >> inp;
    countb += inp;
    flow.addEdge(n + i + 1, n + m + 1, inp, 0);
  if (counta != countb) {
    cout << -1 << endl;</pre>
    return;
  pii t = flow.maxflow(102, 0, n + m + 1);
  if (t.first != counta) {
```

```
cout << -1 << endl;</pre>
    return:
 }
  cout << t.second + start + t.second - counta << endl;</pre>
Dinic.h
             Complexity O(VE \log U) where U = \max |\text{cap}|.
Description:
O(\min(E^{1/2}, V^{2/3})E) if U = 1; O(\sqrt{V}E) for bipartite matching.
template <class T = int>
class Dinic {
 public:
 struct Edge {
    Edge(int a, T b) {
     to = a;
      cap = b;
    int to;
    T cap;
 };
  Dinic(int n) {
    edges.resize(n);
    this->n = n;
 }
  T maxFlow(int src, int sink) {
    T ans = 0;
    while (bfs(src, sink)) {
      T flow;
      pt = vector<int>(n, 0);
      while ((flow = dfs(src, sink))) {
        ans += flow;
    return ans:
```

IIIT-H

h[src] = 0;

```
void addEdge(int from, int to, T cap = 1) {
                                                               queue<int> q;
   edges[from].push_back(list.size());
                                                               q.push(src);
  list.push_back(Edge(to, cap));
                                                               while (!q.empty()) {
  edges[to].push_back(list.size());
                                                                 int on = q.front();
  list.push_back(Edge(from, 0));
                                                                 q.pop();
}
                                                                 for (auto a : edges[on]) {
                                                                   if (list[a].cap == 0) {
                                                                     continue;
private:
int n;
vector<vector<int>> edges;
                                                                   int to = list[a].to;
vector<Edge> list;
                                                                   if (h[to] > h[on] + 1) {
vector<int> h, pt;
                                                                     h[to] = h[on] + 1;
T dfs(int on, int sink, T flow = 1e9) {
                                                                     q.push(to);
  if (flow == 0) {
    return 0;
  if (on == sink) {
                                                               return h[sink] < n;</pre>
    return flow;
                                                             }
                                                           };
   for (; pt[on] < sz(edges[on]); pt[on]++) {</pre>
                                                           void solve() {
    int cur = edges[on][pt[on]];
                                                             int n, m;
    if (h[on] + 1 != h[list[cur].to]) {
                                                             cin >> n >> m;
       continue;
                                                             vi a(n);
                                                             for (int i = 0; i < n; i++) {
     T got = dfs(list[cur].to, sink, min(flow, list[cur].
                                                               cin >> a[i];
     □cap));
    if (got) {
                                                             Dinic<int> flow(n + 2);
                                                             map<int, map<int, int>> factors;
      list[cur].cap -= got;
      list[cur ^ 1].cap += got;
                                                             for (int i = 0; i < n; i++) {
                                                               for (int j = 2; j * j <= a[i]; j++) {
       return got;
    }
                                                                 while (a[i] % j == 0) {
                                                                   factors[j][i + 1]++;
  return 0;
                                                                   a[i] /= j;
bool bfs(int src, int sink) {
  h = vector<int>(n, n);
                                                               if (a[i] > 1) {
```

```
factors[a[i]][i + 1]++;
  }
}
for (int i = 0; i < m; i++) {
 int u, v;
  cin >> u >> v;
 if (u % 2 == 0) {
    swap(u, v);
  flow.addEdge(u, v, 100);
}
int ans = 0;
for (auto t : factors) {
  Dinic<int> tempflow = flow;
  for (auto t1 : t.second) {
   if (t1.first % 2 == 0) {
      tempflow.addEdge(t1.first, n + 1, t1.second);
   } else {
      tempflow.addEdge(0, t1.first, t1.second);
   }
  ans += tempflow.maxFlow(0, n + 1);
cout << ans << endl;</pre>
```

4.2 DFS algorithms

SCC.h

cc0b7e, 72 lines

```
struct SCC {
 int n;
 vvi &adjLists, transposeLists;
 vi scc, leader;
 int sccCount = 0;
 vi sccSize:
 SCC(vvi& adjLists) : n(sz(_adjLists)), adjLists(
 └-1) {
```

```
for (int u = 0; u < n; u++) {
 for (int v : adjLists[u]) transposeLists[v].
  bpush back(u);
vb visited(n);
stack<int> topoSort;
function<void(int) > topoDFS = [&](int from) {
 visited[from] = true;
 for (auto to : adjLists[from]) {
   if (visited[to]) continue;
   topoDFS(to);
 topoSort.push(from);
};
for (int i = 0; i < n; i++)
  if (not visited[i]) topoDFS(i);
visited.assign(n, false);
int sccPtr = 0;
sccSize.assign(n, 0);
function<void(int) > sccDFS = [&](int from) {
  scc[from] = sccPtr;
  sccSize[sccPtr]++;
 visited[from] = true;
  for (auto to : transposeLists[from]) {
    if (visited[to]) continue;
    sccDFS(to);
};
while (not empty(topoSort)) {
  int i = topoSort.top();
  topoSort.pop();
  if (visited[i]) continue;
  sccDFS(i);
  leader[sccPtr] = i;
  sccPtr++;
```

```
int timer;
    sccCount = sccPtr;
                                                             void dfs(int v, int p = -1) {
  int size(int index) { // Returns size of scc of index
                                                              visited[v] = true;
    return sccSize[scc[index]];
                                                              tin[v] = low[v] = timer++;
  }
                                                               for (int to : adj[v]) {
  const int& operator[](int index) {
                                                                 if (to == p) continue;
    return scc[index];
                                                                 if (visited[to]) {
                                                                   low[v] = min(low[v], tin[to]);
  vi indexInCycle;
                                                                 } else {
  void sccEnumeration() {
                                                                   dfs(to, v);
    indexInCycle.assign(n, 0);
                                                                   low[v] = min(low[v], low[to]);
    vb visited(n);
                                                                   if (low[to] > tin[v])
                                                                     IS_BRIDGE(v, to);
    int index = 0;
    function<void(int, int) > sccDFS = [&](int from, int sc
    →) {
      indexInCycle[from] = index++;
      visited[from] = true;
      for (auto to : adjLists[from]) {
                                                             void find_bridges() {
        if (scc[to] != sc) continue;
                                                               timer = 0;
        if (visited[to]) continue;
                                                               visited.assign(n, false);
        sccDFS(to, sc);
                                                               tin.assign(n, -1);
      }
                                                               low.assign(n, -1);
                                                               for (int i = 0; i < n; ++i) {
    };
    for (int i = 0; i < sccCount; i++) {</pre>
                                                                 if (!visited[i])
      index = 0;
                                                                   dfs(i);
      sccDFS(leader[i], i);
                                                               }
    }
                                                             // ARTICULATION POINTS:
};
                                                             int n;
                                                             vector<vector<int>> adj;
bridges.h
                                                   1bcaee, 69 lines | vector<bool> visited;
int n:
                           // number of nodes
                                                             vector<int> tin, low;
vector<vector<int>>> adj; // adjacency list of graph
                                                             int timer;
                                                             void dfs(int v, int p = -1) {
vector<bool> visited;
                                                               visited[v] = true;
vector<int> tin, low;
```

RMO<int> rmg;

└ ((dfs(C, 0, -1), ret)) {}

tin[v] = low[v] = timer++;

```
int children = 0;
  for (int to : adj[v]) {
   if (to == p) continue;
   if (visited[to]) {
      low[v] = min(low[v], tin[to]);
   } else {
      dfs(to, v);
      low[v] = min(low[v], low[to]);
      if (low[to] >= tin[v] && p != -1)
        IS CUTPOINT (v);
      ++children;
    }
  }
  if (p == -1 \&\& children > 1)
    IS CUTPOINT (v);
void find cutpoints() {
 timer = 0;
 visited.assign(n, false);
 tin.assign(n, -1);
 low.assign(n, -1);
 for (int i = 0; i < n; ++i) {
   if (!visited[i])
      dfs(i);
  }
4.3 Trees
LCA.h
"../data-structures/RMQ.h"
                                                   993da6, 24 lines
struct LCA {
 int T = 0;
 vi st, path, ret;
 vi en, d;
```

```
void dfs(vvi& adj, int v, int par) {
    st[v] = T++;
    for (auto to : adj[v])
     if (to != par) {
        path.pb(v), ret.pb(st[v]);
        d[to] = d[v] + 1;
        dfs(adj, to, v);
    en[v] = T - 1;
 }
 bool anc(int p, int c) { return st[p] <= st[c] and en[p]
 \Rightarrow >= en[c]; }
 int lca(int a, int b) {
   if (a == b) return a;
   tie(a, b) = minmax(st[a], st[b]);
   return path[rmq.query(a, b - 1)];
 int dist(int a, int b) { return d[a] + d[b] - 2 * d[lca(

¬a, b)]; }

};
```

HLD.h **Description:** Root must be 0.

c2f6cd, 38 lines

```
struct HLD {
                                                                    int n, timer = 0;
                                                                    vi top, tin, p, sub;
                                                                    HLD(vvi \& adj) : n(sz(adj)), top(n), tin(n), p(n, -1),
                                                                    \rightarrow sub (n, 1) {
                                                                      vi ord(n + 1);
                                                                      for (int i = 0, t = 0, v = ord[i]; i < n; v = ord[++i]
                                                                       └])
                                                                        for (auto &to : adj[v])
                                                                           if (to != p[v]) p[to] = v, ord[++t] = to;
                                                                       for (int i = n - 1, v = ord[i]; i > 0; v = ord[--i])
                                                                       \rightarrow sub[p[v]] += sub[v];
LCA(vector < vi > \& C) : st(sz(C)), en(sz(C)), d(sz(C)), rmq
                                                                      for (int \mathbf{v} = 0; \mathbf{v} < \mathbf{n}; \mathbf{v} + +)
                                                                        if (sz(adj[v])) iter_swap(begin(adj[v]), max_element
```

```
(all(adj[v]), [&](int a, int b) { return make pair
      (a != p[v], sub[a]) < make_pair(b != p[v], sub[b])
      →; }));
    function<void(int)> dfs = [&](int v) {
      tin[v] = timer++;
      for (auto &to : adj[v])
        if (to != p[v]) {
          top[to] = (to == adj[v][0] ? top[v] : to);
          dfs(to);
        }
    } ;
    dfs(0);
  int lca(int u, int v) {
    return process(u, v, [](...) {});
  template <class B>
  int process(int a, int b, B op, bool ignore lca = false)
  └ {
    for (int v;; op(tin[v], tin[b]), b = p[v]) {
     if (tin[a] > tin[b]) swap(a, b);
      if ((v = top[b]) == top[a]) break;
    if (int 1 = tin[a] + ignore_lca, r = tin[b]; 1 <= r)</pre>
    \rightarrow op(1, r);
    return a;
  template <class B>
  void subtree(int v, B op, bool ignore lca = false) {
    if (sub[v] > 1 or !ignore_lca) op(tin[v] + ignore_lca,
    \downarrow \, tin[v] + sub[v] - 1);
  }
};
```

KthAnc.h

28228d, 55 lines

```
struct LCA {
  int n;
```

```
vvi& adjLists;
int lg;
vvi up;
vi depth;
LCA(vvi& _adjLists, int root = 0) : n(sz(_adjLists)),

¬adjLists( adjLists) {
  lq = 1;
  int pw = 1;
  while (pw <= n) pw <<= 1, lg++;</pre>
  // 1q = 20
  up = vvi(n, vi(lg));
  depth.assign(n, -1);
  function<void(int, int)> parentDFS = [&](int from, int

  parent) {
    depth[from] = depth[parent] + 1;
    up[from][0] = parent;
    for (auto to : adjLists[from]) {
      if (to == parent) continue;
      parentDFS(to, from);
   }
  };
  parentDFS(root, root);
  for (int j = 1; j < lq; j++) {
    for (int i = 0; i < n; i++) {
      up[i][j] = up[up[i][j-1]][j-1];
int kthAnc(int v, int k) {
  int ret = v;
  int pw = 0;
  while (k) {
    if (k & 1) ret = up[ret][pw];
    k >>= 1;
    pw++;
  return ret:
```

```
}
  int lca(int u, int v) {
    if (depth[u] > depth[v]) swap(u, v);
    v = kthAnc(v, depth[v] - depth[u]);
    if (u == v) return v;
    while (up[u][0] != up[v][0]) {
      int i = 0;
      for (; i < lq - 1; i++) {
        if (up[u][i + 1] == up[v][i + 1]) break;
      u = up[u][i], v = up[v][i];
    return up[u][0];
  };
  int dist(int u, int v) {
    return depth[u] + depth[v] - 2 * depth[lca(u, v)];
  }
};
```

4.4 Math

4.4.1 Number of Spanning Trees

Create an $N \times N$ matrix mat, and for each edge $a \to b \in G$, do mat [a] [b] --, mat [b] [b] ++ (and mat [b] [a] --, mat [a] [a] ++ if G is undirected). Remove the ith row and column and take the determinant; this yields the number of directed spanning trees rooted at i (if G is undirected, remove any row/column).

4.4.2 Erdős–Gallai theorem

A simple graph with node degrees $d_1 \ge \cdots \ge d_n$ exists iff $d_1 + \cdots + d_n$ is even and for every $k = 1 \dots n$,

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k).$$

Number theory (5)

5.1 Modular arithmetic

ModularArithmetic.h

79388d, 20 lines

```
int add(int x, int y, int m = M) {
 int ret = (x + y) % m;
 if (ret < 0) ret += m;
 return ret:
int mult(int x, int y, int m = M) {
 int ret = (x * y) % m;
 if (ret < 0) ret += m;
 return ret;
int pw(int a, int b, int m = M) {
 int ret = 1;
 int p = a;
 while (b) {
    if (b & 1) ret = mult(ret, p, m);
   b >>= 1;
    p = mult(p, p, m);
 }
 return ret:
```

gcdextended.h

Description: Finds two integers x and y, such that $ax + by = \gcd(a, b)$. If a and b are coprime, then x is the inverse of $a \pmod{b}$.

```
int euclid(int a, int b, int &x, int &y) {
  if (!b) return x = 1, y = 0, a;
  int d = euclid(b, a % b, y, x);
  return y -= a / b * x, d;
}
```

5.2 Primality spf.h

```
#define SIEVE_TILL (int) 1e6
vector<int> primes;
vector<int> spf;
void sieve() {
  spf = vector<int>(SIEVE_TILL + 1, 0);
  for (int i = 2; i <= SIEVE TILL; i++) {</pre>
    if (spf[i] == 0) primes.push_back(i), spf[i] = i;
    for (int j = 0; j < sz(primes) and i * primes[j] <=</pre>
    →SIEVE_TILL; j++) {
      spf[i * primes[j]] = primes[j];
      if (spf[i] == primes[j]) break;
  }
bool isPrime(int n) {
  if (n <= 1) return false;</pre>
  return spf[n] == n;
}
```

Miller Rabin.h

Description: Deterministic for numbers up to $7 \cdot 10^{18}$

Time: 7 times the complexity of $a^b \mod c$.

859df8, 13 lines

```
using ull = uint64_t;
bool isPrime(ull n) {
  if (n < 2 | | n % 6 % 4 != 1) return (n | 1) == 3;
 ull A[] = \{2, 325, 9375, 28178, 450775, 9780504,
  →1795265022}.
      s = builtin ctzll(n - 1), d = n >> s;
  for (ull a : A) {
    ull p = pw(a % n, d, n), i = s;
    while (p != 1 && p != n - 1 && a % n && i--)
      p = mult(p, p, n);
    if (p != n - 1 && i != s) return 0;
  return 1:
```

8647ab, 17 lines **5.2.1 Primes**

Primitive roots exist modulo any p^a , except p=2, a>2, and there are $\phi(\phi(p^a)).$

5.2.2 Estimates

 $\left|\sum_{d|n} d = O(n \log \log n)\right|$.

5.2.3 Mobius Function

Mobius Inversion:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d)g(n/d)$$

Other useful formulas/forms:

$$g(n) = \sum_{n|d} f(d) \Leftrightarrow f(n) = \sum_{n|d} \mu(d/n)g(d)$$
$$g(n) = \sum_{1 \le m \le n} f(\lfloor \frac{n}{m} \rfloor) \Leftrightarrow f(n) = \sum_{1 \le m \le n} \mu(m)g(\lfloor \frac{n}{m} \rfloor)$$

Strings (6)

Manacher.h

Description: p[0][i] = half length of longest even palindrome around positions.[i, p[1][i] = longest odd (half rounded down).

e7ad79, 13 lines

```
array<vi, 2> manacher(const string& s) {
  int n = sz(s);
  array < vi, 2 > p = \{vi(n+1), vi(n)\};
  rep(z, 0, 2) for (int i=0, l=0, r=0; i < n; i++) {
    int t = r-i+!z;
    if (i<r) p[z][i] = min(t, p[z][1+t]);</pre>
    int L = i-p[z][i], R = i+p[z][i]-!z;
    while (L > = 1 \&\& R + 1 < n \&\& s[L-1] == s[R+1])
      p[z][i]++, L--, R++;
    if (R>r) l=L, r=R;
  return p;
```

Trie.h

```
struct trieobject {
  trieobject() {
    children[0] = NULL;
    children[1] = NULL;
    numelems = 0;
  };
  struct trieobject* children[2];
  int numelems;
};
struct trie {
  trieobject base;
  trie() {
    trieobject base;
  void add(int x) {
    int pow2 = (111 << 3111);
    trieobject* temp = &base;
    while (pow2 > 0) {
      if (temp->children[1 && (x & pow2)] == NULL) {
        temp->children[1 && (x & pow2)] = new trieobject;
      temp->children[1 && (x & pow2)]->numelems++;
      temp = temp->children[1 && (x & pow2)];
      pow2 /= 2;
  // ADD FUNCTION BELOW
};
```

Combinatorial (7)

7.1 Permutations

6d8465, 29 lines **7.1.1** Cycles

Let $q_S(n)$ be the number of n-permutations whose cycle lengths all belong to the set S. Then $\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$

7.1.2 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D_n = (n-1)(D_{n-1} + D_{n-2}) = nD_{n-1} + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

7.1.3 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of | X up to symmetry equals $\frac{1}{|G|} \sum_{g \in G} |X^g|$, where X^g are the elements fixed by q(q.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$ to get $g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k) \phi(\frac{n}{k}).$

7.2 Partitions and subsets

7.2.1 Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write $|n = n_k p^k + ... + n_1 p + n_0$ and $m = m_k p^k + ... + m_1 p + m_0$. Then $\binom{n}{m} \equiv \prod_{i=0}^{k} \binom{n_i}{m_i} \pmod{p}$.

7.3 General purpose numbers

7.3.1 Stirling numbers (1)

|# permutations on n items with k cycles.

$$c_{n,k} = c_{n-1,k-1} + (n-1)c_{n-1,k}, \ c_{0,0} = 1$$
$$\sum_{k=0}^{n} c_{n,k} x^k = x(x+1) \dots (x+n-1)$$

7.3.2 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$. $E_{n,k} = (n-k)E_{n-1,k-1} + (k+1)E_{n-1,k}$ $E_{n,0} = E_{n,n-1} = 1$ $E_{n,k} = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$

7.3.3 Stirling numbers (2)

Partitions of n distinct elements into exactly k groups.

$$S_{n,k} = S_{n-1,k-1} + kS_{n-1,k} S_{n,1} = S_{n,n} = 1$$

$$S_{n,k} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

7.3.4 Bell numbers

Ways to partition of n distinct elements. $B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$

7.3.5 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$
 $C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \ C_{n+1} = \sum C_i C_{n-i} \cdot \text{sub-diagonal}$
monotone paths in an $n \times n$ grid. \cdot strings with n pairs of parenthesis. \cdot binary trees with with $n+1$ leaves (0 or 2 children). \cdot ordered trees with $n+1$ vertices. \cdot ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines. \cdot permutations of $[n]$ with no 3-term increasing subseq.

Various (8)

8.0.1 Bit hacks

• x&-x is the least bit in x. • for (int x=m; x; x=(x-1)&m) loops over all subset masks of m. • c=x&-x, r=x+c; (((r^x) >> 2)/c) | r is the next number after x with the same number of bits set. • rep(b, 0, K) rep(i, 0, (1<<K)) if (i & 1 << b) D[i] += D[i^(1 << b)]; computes all sums of subsets.

Numerical (9)

9.1 Fourier transforms

FastFourierTransform.h

Description: fft(a) computes $\hat{f}(k) = \sum_x a[x] \exp(2\pi i \cdot kx/N)$ for all k. N must be a power of 2. conv (a, b) = c, $c[x] = \sum a[i]b[x-i]$. For convolution of complex numbers or more than two vectors: FFT, multiply pointwise, divide by n, reverse(start+1, end), FFT back. Rounding is safe if $(\sum a_i^2 + \sum b_i^2) \log_2 N < 9 \cdot 10^{14}$ (in practice 10^{16} ; higher for random inputs).

00ced6, 35 lines

```
typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C>& a) {
  int n = sz(a), L = 31 - builtin clz(n);
  static vector<complex<long double>> R(2, 1);
  static vector<C> rt(2, 1); // (^ 10% faster if double)
  for (static int k = 2; k < n; k *= 2) {
    R.resize(n); rt.resize(n);
    auto x = polar(1.0L, acos(-1.0L) / k);
    rep(i,k,2*k) rt[i] = R[i] = i&1 ? R[i/2] * x : R[i/2];
  vi rev(n);
  rep(i,0,n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
  rep(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
  for (int k = 1; k < n; k *= 2)
    for (int i = 0; i < n; i += 2 * k) rep(j, 0, k) {
       \mathbf{C} \mathbf{z} = \mathbf{rt} [\mathbf{j} + \mathbf{k}] * \mathbf{a} [\mathbf{i} + \mathbf{j} + \mathbf{k}]; // (25\% \text{ faster if hand-}
       \rolled)
       \mathbf{a}[\mathbf{i} + \mathbf{j} + \mathbf{k}] = \mathbf{a}[\mathbf{i} + \mathbf{j}] - \mathbf{z};
       a[i + j] += z;
vd conv(const vd& a, const vd& b) {
  if (a.empty() || b.empty()) return {};
  vd res(sz(a) + sz(b) - 1);
  int L = 32 - __builtin_clz(sz(res)), n = 1 << L;</pre>
```

```
vector<C> in(n), out(n);
copy(all(a), begin(in));
rep(i,0,sz(b)) in[i].imag(b[i]);
fft(in);
for (C& x : in) x *= x;
rep(i,0,n) out[i] = in[-i & (n - 1)] - conj(in[i]);
fft(out);
rep(i,0,sz(res)) res[i] = imag(out[i]) / (4 * n);
return res;
}
```

NTT.h

Description: Can be used for convolutions modulo specific nice primes of the form $2^ab + 1$, where the convolution result has size at most 2^a . (125000001 << 3) + 1 = 1e9 + 7, therefore do not use this for M = 1e9 + 7. For $p < 2^30$ there is also e.g. (5 << 25, 3), (7 << 26, 3), For other primes/integers, use two different primes and combine with CRT. (479 << 21, 3) and (483 << 21, 5). The last two are > 10^9 . Inputs must be in [0, mod).

```
f9d990, 46 lines
// Requires mod func
const int M = 998244353;
const int root = 3;
// (119 << 23) + 1, root = 3; // for M = 998244353
void ntt(int* x, int* temp, int* roots, int N, int skip)
 if (N == 1) return;
 int n2 = N / 2;
 ntt(x, temp, roots, n2, skip * 2);
 ntt(x + skip, temp, roots, n2, skip * 2);
 for (int i = 0; i < N; i++) temp[i] = x[i * skip];</pre>
 for (int i = 0; i < n2; i++) {
    int s = temp[2 * i], t = temp[2 * i + 1] * roots[skip]
    →* i];
    x[skip * i] = (s + t) % M;
    x[skip * (i + n2)] = (s - t) % M;
  }
```

```
void ntt(vi& x, bool inv = false) {
 int e = pw(root, (M - 1) / sz(x));
 if (inv) e = pw(e, M - 2);
  vi roots(sz(x), 1), temp = roots;
  for (int i = 1; i < sz(x); i++) roots[i] = roots[i - 1]
  →* e % M;
  ntt(&x[0], &temp[0], &roots[0], sz(x), 1);
// Usage: just pass the two coefficients list to get a*b (
→modulo M)
vi conv(vi a, vi b) {
  int s = sz(a) + sz(b) - 1;
 if (s <= 0) return {};</pre>
  int L = s > 1 ? 32 - builtin clzl1(s - 1) : 0, n = 1
  →<< L;
  if (\mathbf{s} \le 200) { // (factor 10 optimization for |\mathbf{a}|, |\mathbf{b}|
  └= 10)
    vi c(s);
    for (int i = 0; i < sz(a); i++)
      for (int i = 0; i < sz(b); i++)
        c[i + j] = (c[i + j] + a[i] * b[j]) % M;
    return c:
  a.resize(n);
  ntt(a);
  b.resize(n);
  ntt(b);
  vi c(n);
  int d = pw(n, M - 2);
  for (int i = 0; i < n; i++) c[i] = a[i] * b[i] % M * d %
  ↓ M;
  ntt(c, true);
  c.resize(s);
  return c;
```

Geometry (10)

10.1 Geometric primitives

Point.h

Description: Avoid T = int47ec0a, 31 lines

```
template <class T>
int sgn(T x) { return (x > 0) - (x < 0); }
template <class T>
struct Point {
  typedef Point P;
  T x, y;
  explicit Point (\mathbf{T} \mathbf{x} = 0, \mathbf{T} \mathbf{y} = 0) : \mathbf{x}(\mathbf{x}), \mathbf{y}(\mathbf{y}) {}
  bool operator \langle (P p) \rangle const \{ return tie(x, y) < tie(p.x, y) \}

¬p.y); }

 bool operator == (P p) const { return tie(x, y) == tie(p.x vector < P > convexHull(vector < P > pts) {
  \rightarrow, p.y); }
  P operator+(P p) const { return P(x + p.x, y + p.y); }
  P operator-(P p) const { return P(x - p.x, y - p.y); }
  P operator*(T d) const { return P(x * d, y * d); }
  P 	ext{ operator}/(T 	ext{ d}) 	ext{ const } \{ 	ext{ return } P(x / d, y / d); \}
  T dot(P p) const { return x * p.x + y * p.y; }
  T cross(P p) const { return x * p.y - y * p.x; }
  T cross(P a, P b) const { return (a - *this).cross(b - *

this); }

  T dist2() const { return x * x + y * y; }
  double dist() const { return sqrt((double)dist2()); }
  // angle to x-axis in interval [-pi, pi]
  double angle() const { return atan2(y, x); }
  P unit() const { return *this / dist(); } // makes dist
  4()=1
  P perp() const { return P(-y, x); }
                                                   // rotates
  4+90 degrees
  P normal() const { return perp().unit(); }
  // returns point rotated 'a' radians ccw around the
  4origin
  P rotate(double a) const {
    return P(x * cos(a) - y * sin(a), x * sin(a) + y * cos
```

```
\rightarrow (a));
friend ostream& operator<<(ostream& os, P p) {</pre>
  return os << "(" << p.x << "," << p.y << ")";
```

ConvexHull.h

Description: Returns a vector of the points of the convex hull in counterclockwise order. Points on the edge of the hull between two other points are not considered part of the hull.

Time: $\mathcal{O}(n \log n)$

310954, 14 lines

```
// Needs point
typedef Point<11> P;
 if (sz(pts) <= 1) return pts;</pre>
 sort (all (pts));
 vector<P> h(sz(pts)+1);
 int s = 0, t = 0;
 for (int it = 2; it--; s = --t, reverse(all(pts)))
    for (P p : pts) {
      while (t >= s + 2 \&\& h[t-2].cross(h[t-1], p) <= 0) t
      └--:
      h[t++] = p;
 return {h.begin(), h.begin() + t - (t == 2 && h[0] == h
 →[1])};
```

ClosestPair.h

Description: Finds the closest pair of points.

fb15fd, 18 lines

```
// Requires point
typedef Point<int> P;
pair<P, P> closest(vector<P> v) {
  assert (sz(v) > 1);
 set <P> S;
```

IIIT-H

```
20
```

```
sort(all(v), [](P a, P b) { return a.y < b.y; });
pair < int, pair < P, P >> ret { LLONG_MAX, {P(), P()} };
int j = 0;
for (P p : v) {
   P d{1 + (int) sqrtl(ret.first), 0};
   while (v[j].y <= p.y - d.x) S.erase(v[j++]);
   auto lo = S.lower_bound(p - d), hi = S.upper_bound(p + d);
   for (; lo != hi; ++lo)
      ret = min(ret, {(*lo - p).dist2(), {*lo, p}});
   S.insert(p);
}
return ret.second;
}</pre>
```