

# **SYS2(OS) Lab 5: Threads & Synchronisation**

**Dr Poonam Yadav and Sven Signer**

# POSIX Threads

---

Last week's lab session looked at creating new processes with the `fork()` system call, and sending simple signals to processes.

---

This practical looks at implementing multiple threads within a process using the POSIX pthreads library.

---

Unlike separate processes created with `fork()`, threads run in the same address space.

---

Unexpected behaviour can occur if access to shared memory is not carefully coordinated.

---

# Thread Management

```
#include <pthread.h>
```

```
int pthread_create(pthread_t *thread, const pthread_attr_t *attr,  
                  void *(*start_routine) (void *), void *arg);
```

- Creates a new thread in the current process.
- Function pointer `start_routine` is invoked in the new thread with arguments `arg`.
- Thread ID is stored into `thread` argument.
- Function returns 0 on successful thread creation.

```
#include <pthread.h>
```

```
int pthread_join(pthread_t thread, void **retval);
```

- Makes the current thread wait for another thread (specified by `thread` argument) to complete.
- If `retval` is not null, the return value of the joined thread is stored into it.
- Function returns 0 on success.

# Thread Creation Example

```
#include <pthread.h>
```

```
void *some_fn(void *arg) { ... }
```

← Function containing  
thread implementation

```
int main() {
```

```
    ...
```

```
    pthread_t tid;
```

```
    pthread_create(&tid, NULL, some_fn, arg);
```

Create a new thread  
that executes the specified  
function.

```
    ...
```

```
    pthread_join(tid, NULL);
```

← Wait for thread completion.

```
    ...
```

# Thread Synchronisation

- Threads operate in the same address space, meaning global memory can be accessed from different threads simultaneously.
- Without thread synchronisation, it is easy to introduce race conditions, where program behaviour depends on the random scheduling behaviour.
- Mutexes allow the program to eliminate race conditions by ensuring that only one thread at a time will execute in a critical region.

# Synchronisation with Mutexes

```
#include <pthread.h>
```

```
int pthread_mutex_init(pthread_mutex_t *mutex, const pthread_mutexattr_t *attr);  
int pthread_mutex_destroy(pthread_mutex_t *mutex);  
int pthread_mutex_lock(pthread_mutex_t *mutex);  
int pthread_mutex_trylock(pthread_mutex_t *mutex);  
int pthread_mutex_unlock(pthread_mutex_t *mutex);
```


- Once initialised, mutexes can be locked/unlocked to ensure only a single thread executes in a critical region at any given time.
- As always, see the man pages for specifics on the arguments.

# Mutex Example Program

```
#include <pthread.h>
pthread_mutex_t mutex;
```


```
void *thread_fn(void *param) {
    ...
    pthread_mutex_lock(&mutex);
    // <Critical Section>
    pthread_mutex_unlock(&mutex);
    ...
}
```

Thread(s) acquire a mutex before a critical section to guarantee mutual exclusion.



```
int main(int argc, char *argv[]) {
    pthread_mutex_init(&mutex, NULL);
    // <Launch Threads>
    pthread_mutex_destroy(&mutex);
}
```

Mutex must be initialised before it can be used, for example in the main function before launching worker threads.



# Condition Variables

- While mutexes implement synchronisation by controlling thread access to data, condition variables allow threads to synchronise based upon the actual value of data.
- A program might need to acquire a lock (mutex) on some shared data, but might require this to be in a certain state before it can proceed.
- Condition variables allow program to efficiently guarantee mutual exclusion while testing some predicate.



# Condition Variables

```
#include <pthread.h>
```

```
int pthread_cond_wait(pthread_cond_t *cond, pthread_mutex_t *mutex);
```

- Releases the mutex and waits for the condition variable to be signalled, then reacquires the mutex prior to continuing.
- Spurious wakeups are explicitly permitted, so typically included inside a while loop that re-checks some predicate.

```
#include <pthread.h>
```

```
int pthread_cond_signal(pthread_cond_t *cond);
```

- Signal a thread waiting on the condition variable to unblock.
- Has no effect if no thread is waiting on the condition.

# Condition Variable Example

```
#include <pthread.h>
```

```
pthread_cond_t cond;
```

```
pthread_mutex_t mutex;
```

```
void *thread1(void *param);
```

```
void *thread2(void *param);
```

```
int main() {
```

```
    pthread_mutex_init(&mutex, NULL);
```

```
    pthread_cond_init(&cond, NULL);
```

```
    // <Launch & Wait for Threads>
```

```
    pthread_mutex_destroy(&mutex);
```

```
    pthread_cond_destory(&cond);
```


```
}
```

← Like mutexes, condition variables need to be initialised before they can be used.

# Condition Variable Example

```
void *thread1(void *param) {  
    ...  
    pthread_mutex_lock(&mutex);  
    while (!predicate) {  
        pthread_cond_wait(&cond, &mutex);  
    }  
  
    // <Critical Section>  
  
    pthread_mutex_unlock(&mutex);  
    ...  
}
```


Thread locks mutex, then waits for predicate to hold before executing the critical section. Condition variable handles releasing and reacquiring the mutex.



# Condition Variable Example

```
void *thread2(void *param) {  
    ...  
    pthread_mutex_lock(&mutex);  
  
    // <Critical Section>  
  
    if (predicate) {  
        pthread_cond_signal(&cond);  
    }  
    pthread_mutex_unlock(&mutex);  
    ...  
}
```

Other thread(s) check the predicate after completing their own critical section, signalling any blocked threads to continue if the predicate now holds.



# Assignment



This lab session looks at thread creation and race conditions that can be caused by simultaneous access to global variables.

An example program `simple_threads.c` is provided on the VLE to get you started.