

Crafting a Compiler

5.5

Original Grammar:

1. DeclList → DeclList ; Decl
2. | Decl
3. Decl → IdList : Type
4. IdList → IdList , id
5. | id
6. Type → ScalarType
7. | array (ScalarTypeList) of Type
8. ScalarType → id
9. | Bound .. Bound
10. Bound → Sign intconstant
11. | id
12. Sign → +
13. | -
14. |
15. ScalarTypeList → ScalarTypeList , ScalarType
16. | ScalarType

Answer:

Left recursion is the issue: there are no common prefixes
Left recursion is an issue in the bolded steps

1. DeclList → Decl DeclList'
2. DeclList' → ; DeclList
3. |
4. Decl → IdList : Type
5. IdList → id IdList'
6. IdList' → , IdList
7. |
8. Type → ScalarType
9. | array (ScalarTypeList) of Type
10. ScalarType → id
11. | Bound .. Bound
12. Bound → Sign intconstant
13. | id
14. Sign → +
15. | -
16. |
16. ScalarTypeList → ScalarType ScalarTypeList'
17. ScalarTypeList' → , ScalarTypeList
18. |

ScalarTypeList , ScalarType
16. | ScalarType

6.51? ←Optional?

The bottom-up parsing techniques given in this chapter are more powerful than top-down techniques given in Chapter 5.

Using the alphabet { a, b }, devise a language that is not LL(k) for any k but is LR(k) for some k. What property of LR(k) parsing allows such a grammar to be constructed?

$A \rightarrow a + B$
| $a - b$
 $B \rightarrow a + b$

This is not LL(k) because it is left recursive and has common prefixes.

It can be LR(k) because it is not right recursive, which means it has a unique prefix when reading right to left.

Dragon

4.5.3

Give bottom-up parses for the following input strings and grammars:

a)

Ex 4.5.1:

$S \rightarrow 1. 0 S 1 \mid 2. 0 1$

The input 000111 according to the grammar of Ex 4.5.1

$S \Rightarrow 1. 0 S 1 \Rightarrow 1. 0 0 S 1 1 \Rightarrow 2. 0 0 0 1 1 1$

b)

Ex 4.5.2:

$S \rightarrow 1. S S + \mid 2. S S * \mid 3. a$

The input $aaa*a++$ according to the grammar of Ex 4.5.2

Shift Reduce:

$a \Rightarrow S \Rightarrow a S \Rightarrow S S \Rightarrow a S S \Rightarrow S S S \Rightarrow S S S * \Rightarrow S S \Rightarrow S S a \Rightarrow S S S \Rightarrow S S S + \Rightarrow S S \Rightarrow S S + \Rightarrow S$

4.6.5

Show that the following grammar:

$S \rightarrow A a A b \mid B b B a$

$A \rightarrow \epsilon$

$B \rightarrow \epsilon$

Is LL(1) but not SLR(1) ['Don't worry about SLR' – Alan Labouseur]

Parse Table:

	a	b
S	1	2
A		
B		