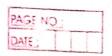
			Jessica Bragas	720	Charles and Charles and Charles
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		Cloud Computing Ass	ignment z		
D Expla	in in brief	Implementation lev	en cua promonente del cuan inconse perior promonente institución de perior.		the table of surgicines in the sure of specific times and
Ans			CO OF VIEGOTI	20.40 Davidence and American	
purposition per comment across North District (Section (Section Consultation	APPLICATIO	N LEVEL			Silve dall'an Tamoni più candida della
		JVM/ . NET CLR / F	anot		
F. 5.	LIBPADY	1050 100			
		JSER LEVEL API) LEV			
	100	WINE / WABI / LXRUN /	Visual Main Win	/ VCUDA	
	OPERATING	SYSTEM LEVEL			
244.	1	Jain / Virtual Engronn	nent / Ensim's VF	S FUM	
		1 5/11/2 5/12			
	HARDWARE	ABSTRACTION LAYER	(HAL) LEVEL	penney &	
	v	mware / Virtual PC/D	enale / Xen/ L4/	Plux 86	
	lu	er mock linux / copper	ative linux.	1.0	
		1800 July 1 - 1873 18 18	maria indi	okans silt	
	INSTRUCTION	SET ARCHITECTURE	(ISA) LEVEL	2) 4 3/4 •	
	B.	ochs / Chisas - DEMU	BIRD / DYNAMO	7 5 7	
1 / 1			4.3 4.4.0		
1 1 2 1 2	Fig	Implementation levels	of virtualization	10	
	ng california	Miles adl Branco de		V#1/1-2-30	
1) Instructi	on Set Axch	itecture level (ISA)	r <del>å</del> jag (s		
° ISA v	rtualization	can work through	13A emulatio	n. This is w	<u>led</u>
nux at	many lego	my water that w	ene muitten f	on different	
config u	nation of he	undware,		Samuel (A	
These u	odes suin on	any virtual mas	hine wing	the ISA : wit	h this
		at originally her			
					and a supply

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to 2	un is now rapable of surning on x86 machine.
	basic emulation, an interpreter is needed, which interprets the
	is code and then converts it into hardware format that
	be read. This then allows processing.
2) Ha	udware abstraction level (HAL) -
	e the hardware.
	is makes use of a hypernison which is used for functioning the
	adware using the process of intualization.
	allows the intualization of each of the handware component
wh	uch could be the input -output device, the memory, the
- Pr	Diceson etc.
• Mu	Itiple were will not be able to use the same hardware and
als	ob use multiple instalization instances at the same time.
3) Op	enating system level -
• At	the level of the operating system, the virtualization model
ŭ	capable of creating a layer that is abstraction between
the	e operating system and the application.
· Thu	is is an isolated contained that is on the operating system
an	d the physical server, which makes use of the software and
ha	ordurare. Each of these then functions in the form of a server.
	zen there are several users and no one wants to share the
	raware, then this is where the virtualization level is used.
	ry user will get this virtual environment using a virtual
hax	durance resource that is dedicated.
4) Libs	kary rever-
	e operating system is cumbersome and this is when the
	plication makes use of the API that is from the absorbes
	-a user level.
	#####################################



. These Api's are documented well and this why the library
irrhualization level is preformed in these scenarios
· API hosts make it possible as it controls the link of communication
From the application to the system.
5) Application level-
. The application well virtualization is used when there is a
dervice to irrtualize only one application and is the last of
the implementation levels of virtualization in cloud computing.
This is generally used when you run virtual machine that
use high-level languages.
· The application will sit above virtualization layer, which is
turn sits on the application program.
. It lets the high level language purgrams compiled to be used
in the application level of the virtual mathine run, seamlessly.
2) List and explain vmm design requirements.
Ans: 4 There are 3 requirements of vmm-
i) First, a vmm should provide an environment for programs which
is essentially identical to the original machine.
2) Second, programs sun in this environment should show at
2) Second, programs sun in this environment should show at worst, only minor decreases in speed.
그는 이 살았다. 그는 이 그는 이 그는 이 이 그리고 들었다. 그런 그런 그런 그는 것은 아무리를 하는 것을 것을 하는 것을 하는 것을 하는 것을 하는 것을
worst, only minor decreases in speed.
2) Third, a vmm should be in compute to that which it want on
2) Third, a vmm should be in compute to that which it sums on  The original machine directly?
worst, only minor decreases in speed.  2) Third, a vmm should be in compute to that which it was an  The original machine directly:  The later qualification is required obecause of the interviencing
worst, only minor decreases in speed.  3) Third, a vmm should be in compute to that which it sums an  The original machine directly:  The later qualification is required obecause of the interviencing  Level of software and the effect of any other vms concernently  existing on the same hardware.
worst, only minor decreases in speed.  3) Third, a vmm should be in compute to that which it sums on  The original machine directly:  The later qualification is required obscause of the interviening  Level of software and the effect of any other vmc concernently
worst, only minor decreases in speed.  3) Third, a vmm should be in compute to that which it runs on  The original machine directly:  • The later qualification is required obecause of the interviencing  Level of software and the effect of any other vmc concerently  existing on the same hardware.  • B. vmm should domonstrate efficiency in using the vmc.  • Complete control of these resources by a vm includes:
worst, only minor decreases in speed.  3) Third, a vmm should be in compute to that which it sums on  The original machine directly:  The later qualification is required obscause of the interviencing  Level of software and the effect of any other vms concernently  existing on the same hardware.  A vmm should domonstrate efficiency in using the vms

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	bresowies abready allocated.
The state of the s	all processors satisfy there requirements for a vmm.
31.	
	Explain in brief virtual uniters and resource management.
	physical cluster is ionsidered as a group of physical
	nachine called as servers which are intermonnected with each
6,	they with the help of physical betweek such as a LAN
公の日本語	Three withial design issues of virtual chusters
<u> </u>	We migration of vms
	when there are mixed nodes of host and guest systems
	in a cluster the street of nost and guest systems
	the whole this a
	the whole thing on the physical machine in the situation
	at tacture et one vm another vm en a different node
	replace its rate until both the ums are executing with
ь)	memory and fill migrations -in and the territory
	there is very high initial insting of ownership
e in he	factore like space, power wordifiching as well as woling
	equipment. Hence to share this cost, the leaving on
4.96	shaving access of a comment
	shaving access of a common cluster is a good option when
1 STATE	Dunamic daily over time remarks and the same
	distinct approximent of virtual distance
	The tourses
	117112 - 1119 - COD-COLL
	multiple umas
	when more vme ining and joined
	teads to problems with overland in ineffective configuration
, 100 ere	leads to problems with overloading on under utilization.
	TOTAL AND THE PROPERTY OF THE