# Presentation-5

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#### For a **single qubit**, the oracle outputs the following-

1. 
$$U_f |0\rangle = \frac{1}{\sqrt{2}}(|1\rangle + |0\rangle)$$

2. 
$$U_f |1\rangle = |0\rangle$$

and its matrix is 
$$U_{f1}=\begin{bmatrix} \frac{1}{\sqrt{2}} & 1 \\ \frac{1}{\sqrt{2}} & 0 \end{bmatrix}$$

Similarly, for two qubits,  $U_f$  gives -

1. 
$$U_f |00\rangle = \frac{1}{2}(|11\rangle + |00\rangle + |01\rangle + |10\rangle)$$

2. 
$$U_f |10\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |01\rangle)$$

3. 
$$U_f |01\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |10\rangle)$$

4. 
$$U_f |11\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

For single & two-qubit systems [Continued]

Taking the tensor product of  $U_{f1}$  gives -

$$U_{f2}^{'}=U_{f1}\otimes U_{f1}=egin{bmatrix} rac{1}{2} & rac{1}{\sqrt{2}} & rac{1}{\sqrt{2}} & 1 \ rac{1}{2} & 0 & rac{1}{\sqrt{2}} & 0 \ rac{1}{2} & rac{1}{\sqrt{2}} & 0 & 0 \ rac{1}{2} & 0 & 0 & 0 \end{bmatrix}$$

#### For three qubits, the oracle outputs-

1. 
$$U_f |000\rangle = \frac{1}{2\sqrt{2}}(|000\rangle + |001\rangle + |010\rangle + |011\rangle + |100\rangle + |101\rangle + |110\rangle + |111\rangle)$$

2. 
$$U_f |001\rangle = \frac{1}{2}(|000\rangle + |010\rangle + |100\rangle + |110\rangle)$$

3. 
$$U_f |010\rangle = \frac{1}{2}(|000\rangle + |001\rangle + |100\rangle + |101\rangle)$$

4. 
$$U_f |011\rangle = \frac{1}{2}(|000\rangle + |011\rangle + |100\rangle + |111\rangle)$$

5. 
$$U_f |100\rangle = \frac{1}{2}(|000\rangle + |001\rangle + |010\rangle + |011\rangle)$$

6. 
$$U_f |101\rangle = \frac{1}{2}(|000\rangle + |010\rangle + |101\rangle + |111\rangle)$$

7. 
$$U_f |110\rangle = \frac{1}{2}(|000\rangle + |001\rangle + |110\rangle + |111\rangle)$$

8. 
$$U_f |111\rangle = \frac{1}{2} (|000\rangle + |011\rangle + |101\rangle + |110\rangle)$$

For single, two, three-qubits systems [Continued]

And its matrix is - 
$$U_{f3}=$$
 
$$\begin{bmatrix} \frac{1}{2\sqrt{2}} & \frac{1}{2} \\ \frac{1}{2\sqrt{2}} & 0 & \frac{1}{2} & 0 & \frac{1}{2} & 0 & \frac{1}{2} & 0 \\ \frac{1}{2\sqrt{2}} & \frac{1}{2} & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 \\ \frac{1}{2\sqrt{2}} & \frac{1}{2} & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & \frac{1}{2} \\ \frac{1}{2\sqrt{2}} & \frac{1}{2} & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\ \frac{1}{2\sqrt{2}} & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & \frac{1}{2} \\ \frac{1}{2\sqrt{2}} & \frac{1}{2} & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2\sqrt{2}} & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix}$$



For single, two, three-qubits systems [Continued]

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\begin{array}{l} |0\rangle \longrightarrow 0,1 \text{ and } |1\rangle \longrightarrow 0 \\ |00\rangle \longrightarrow 0,1,2,3, \ |01\rangle \longrightarrow 0,2, \ |10\rangle \longrightarrow 0,1 \text{ and } |11\rangle \longrightarrow 0,3 \\ |000\rangle \longrightarrow 0,1..7, \ |001\rangle \longrightarrow 0,2,4,6, \ |010\rangle \longrightarrow 0,1,4,5, \\ |011\rangle \longrightarrow 0,3,4,7, \ |100\rangle \longrightarrow 0,1,2,3, \ |101\rangle \longrightarrow 0,2,5,7, \\ |110\rangle \longrightarrow 0,1,6,7 \text{ and } |111\rangle \longrightarrow 0,3,5,6 \\ |0000\rangle \longrightarrow 0,1..15, \ |0001\rangle \longrightarrow 0,2,..14, \\ |0010\rangle \longrightarrow 0,1,4,5,8,9,12,13, \ |0011\rangle \longrightarrow 0,3,4,7,8,11,12,15, \\ |0100\rangle \longrightarrow 0,1,2,3,8,9,10,11, \ |0100\rangle \longrightarrow 0,, \\ |1101\rangle \longrightarrow 0,2,5,7,9,11,12,14, \dots \end{array}
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For single, two, three-qubits systems [Continued]

The oracle's solutions corresponds to -

- 1. If a bit in the input string is zero, then corresponding to its index, it can be 0 or 1.
- 2. If a bit is 1 and the input has odd number of 1's, then 2n non-zero bits of the strings can be 1, the remaining one non-zero bit string must be zero.
- 3. If a bit is 1 and the input has even number of 1's, then the corresponding non-zero bit strings must be 1.

The input bit string has to be reversed to follow Qiskit's ordering. The occurrence of the first non-zero string bit of the reversed input string must be stored in a variable [called *flagbit*]. Its oracle is implemented by first copying the qubits of input register to that of the second register using CNOT gates, i.e.  $g1_a \longrightarrow g2_a$ . Then, whenever, we have a non-zero string bit occurring in the reversed string, we perform a *CNOT* operation with control as input register qubit whose index corresponds to flagbit and target as second register qubit whose index corresponds to the index of the occurrence of a non-zero bit string in the reversed input string.

When input is "011", we get the following circuit -

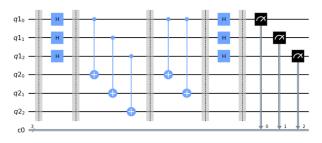


Figure: Circuit for Simon's Algorithm