

1. Consider an undirected graph  $G = (V, E)$  where  $V = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$  and  $E = \{\{1, 2\}, \{2, 3\}, \{4, 5\}, \{6, 7\}, \{6, 8\}, \{7, 8\}, \{8, 9\}\}$ .
  - (i) Draw  $G$ .
  - (ii) Draw a subgraph  $(V', E')$  that is not a tree but satisfies  $|E'| = |V'| - 1$ .
  - (iii) Draw a connected subgraph  $(V'', E'')$  that is not a tree, but has an edge  $\{x, y\}$  such that  $(V'', E'' \setminus \{\{x, y\}\})$  is not connected.
  - (iv) Let  $T$  be a rooted tree such that  $T$ 's edges are in  $E$  and  $T$ 's root is 7. Draw all possible  $T$ s.
  - (v) Draw another graph  $G'' = (\{1, 2, 3, 4, 5, 6, 7, 8, 9\}, E'')$  such that  $E'' \neq E$  but  $G''$  is isomorphic to  $G$ .
  - (vi) Determine the number of graphs  $G'' = (\{1, 2, 3, 4, 5, 6, 7, 8, 9\}, E'')$  such that  $E'' \neq E$  but  $G''$  is isomorphic to  $G$ .
- 2.\*
  - (i) Draw all trees with  $n$  nodes for  $n = 1, 2, 3, 4$ . What is the general formula for the number of trees with  $n$  nodes?
  - (ii) Determine the number of nonisomorphic trees with  $n$  nodes, for  $n = 1, 2, 3, 4$ .  
What is the relationship between (i) and (ii)?
3. For two rooted trees to be isomorphic under a permutation  $\pi$ , we require that the image  $\pi(u)$  of a root  $u$  must also be a root. Determine the number of nonisomorphic rooted trees with  $n$  nodes, for  $n = 1, 2, 3, 4$ .
- 4.\* Let  $G = (V, E)$  be an undirected graph. Prove that if  $G$  is connected, then  $|E| \geq |V| - 1$ . Is the converse true?
5. Let  $G = (V, E)$  be an undirected graph. Prove that if  $G$  is acyclic, then  $|E| \leq |V| - 1$ . Is the converse true?
- 6.\* Prove that a loopless undirected graph is a tree if and only if there is exactly one path between every pair of nodes.
7. Recall from Tutorial 9 (Problem 8) that a complete graph has all loops and all edges. How many spanning trees are there for the complete graph with 4 nodes?
- 8.\* Consider a rooted tree  $T$  in which every parent has at most  $b$  children ( $b \in \mathbb{Z}^+$ ; a binary tree has  $b = 2$ ).  
State and prove a result relating the number of leaves and number of parents in  $T$ , if each parent has exactly  $b$  children.

9. For  $n \geq 2$ , a directed graph  $(\{v_1, v_2, \dots, v_n\}, \{(v_1, v_2), \dots, (v_{n-1}, v_n), (v_n, v_1)\})$  is called a **cycle**. A directed graph  $G = (V, D)$  is **cyclic** if it contains a loop or a cycle as a subgraph; otherwise, it is **acyclic**.
- Prove that if  $G$  is acyclic, then  $D$  is antisymmetric.
  - Prove or disprove the converse of (i).
  - Prove that if  $D$  is a partial order, then  $G$  does not contain any cycles.  
[We can hence arrange the edges in the graph for a partial order, so they all point in one direction.]
  - \* Prove that, for any  $n \geq 2$ , there is a directed acyclic graph with  $n$  nodes and  $\frac{1}{2}n(n-1)$  edges.  
[Contrast this with the  $|E| = |V| - 1$  characterization for undirected acyclic graphs.]
  - \* Prove that any directed graph with  $n$  nodes and more than  $\frac{1}{2}n(n-1)$  edges must be cyclic.

10. Let  $\mathcal{D}_3$  be the set of all directed graphs whose nodes are  $a, b, c$ . Suppose  $G = (\{a, b, c\}, D) \in \mathcal{D}_3$ . Determine the number of possible  $G$ 's such that:

- |                            |                            |
|----------------------------|----------------------------|
| (i)* $G$ has a loop;       | (ii) $G$ is acyclic.;      |
| (iii)* $D$ is reflexive;   | (iv) $D$ is symmetric;     |
| (v)* $D$ is antisymmetric; | (vi) $D$ is a total order. |

- 11.\* Consider a loopless undirected graph  $G = (V, E)$  and  $V' \subseteq V$ . We say  $V'$  **covers** an edge  $\{x, y\} \in E$  iff  $x \in V'$  or  $y \in V'$ , and  $V'$  is a **vertex cover** iff  $V'$  covers every edge. For example,  $\{00, 02, 11, 12, 21, 22\}$  is a vertex cover for the graph in (a).

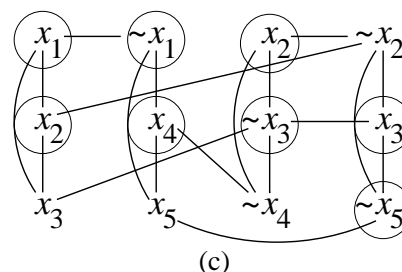
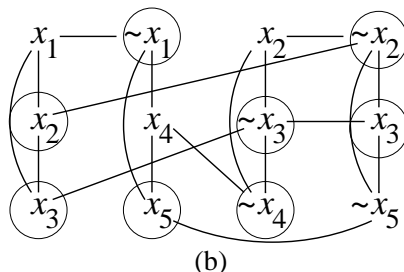
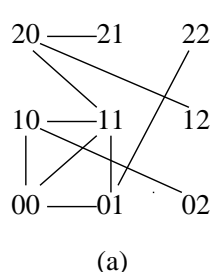
- (i) For the graph in (a), find a vertex cover of smallest possible size.

A Boolean expression  $\alpha$  is **3CNF** iff  $\alpha$  is a conjunction of clauses with exactly 3 literals per clause. A 3CNF expression  $\alpha$  with  $k$  clauses can be transformed into a loopless undirected graph  $G_\alpha$ , so that  $\alpha$  is satisfiable (Tutorial 9) iff  $G_\alpha$  has a vertex cover of size at most  $2k$ . The following illustrates this transformation: Suppose, for  $k = 4$ ,

$$\alpha = (x_1 \vee x_2 \vee x_3) \wedge (\sim x_1 \vee x_4 \vee x_5) \wedge (x_2 \vee \sim x_3 \vee \sim x_4) \wedge (\sim x_2 \vee x_3 \vee \sim x_5).$$

Then the graph in (b) and in (c) (without the circles) is  $G_\alpha$ .

- Consider the truth assignment  $f(x_1) = f(x_2) = f(x_4) = T$  and  $f(x_3) = f(x_5) = F$ . Verify that  $f$  satisfies  $\alpha$ . Let  $C_f$  be the set of vertices in (b) that are circled. Verify that  $C_f$  is a vertex cover of size 8.
- Consider the set  $C$  of vertices in (c) that are circled. Verify that  $C$  is a vertex cover of size 8. Define a truth function  $f_C$  by  $f_C(x_1) = f_C(x_3) = f_C(x_5) = T$  and  $f_C(x_2) = f_C(x_4) = F$ . Verify that  $f_C$  satisfies  $\alpha$ .
- How is  $G_\alpha$  derived from  $\alpha$ ? How is  $C_f$  derived from  $f$ ? How is  $f_C$  derived from  $C$ ?



[This problem suggests that finding a vertex cover is “harder” than finding a satisfying truth assignment. Actually, they are “equally hard”, since both are NP-Complete.]