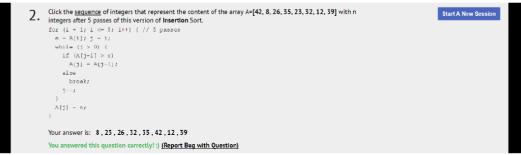
# Sorting



First n+1 numbers will be sorted. Insertion sort is **STABLE**, so the rest of the elements is unaffected.

3.	Suppose the a of the original [101, 104, 106 From the resu	integers is being sorted in ascending order using <b>Quicksort</b> .  Lgorithm has just finished the first pass of partitioning and pivot swapping thus changing the content array into the following array: 1, 109, 103, 106, 108, 115, 113, 107]  Iting array above, determine how many integers could have been the pivot? s == pivot are partitioned to the right.	Start A New Session
	1		
	You answered	this question correctly!:) (Report Bug with Question)	

Find number such that every number to the left is lesser than (or equal to) itself and every number to the right is greater than (or equal to) itself



Find number such that every number to the left is lesser than (or equal to) itself and every number to the right is greater than (or equal to) itself

Brute force count swaps

Selection sort is UNSTABLE. Keep track of where elements are swapped to.

```
11. Click the sequence of integers that represent the content of the array A=[38, 4, 6, 5, 29, 34, 32, 26] with n integers after 7 passes of this version of Bubble Sort.

for (j = 0; j < 7; ji+) // 7 passes
for (i = 0; i < n-j-1; i++)

if (A[i] > A[i+1])

ewap (A[i], A[i+1];

Your answer is: 4,5,6,26,29,32,34,38

You answered this question correctly!:) (Report Bug with Question)
```

Biggest 7 elements are bubbled right. So, this sequence of 8 integers will be fully sorted lol

Become the algorithm

```
How many pass(es) is/are required to sort an array of n=5 integers: [35, 26, 37, 9, 7] using this version of Bubble

Sort?

j = 0;

do (

manapped = falme; j++;

for (i = 0; i < n-j; i++) // each completion of this for-loop is counted as one pass

if [A[i] > A[i+1]);

swspped = true;

}
while (amapped);

You answered this question wrongly.

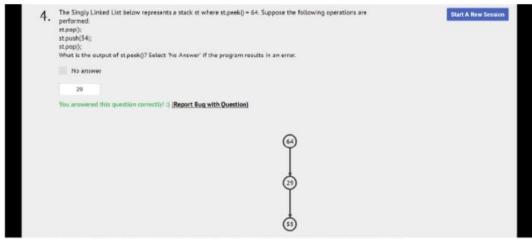
The correct answer is: 5
```

Become the algorithm. This is the "smart" version which breaks if no new swaps are made in the pass-through.

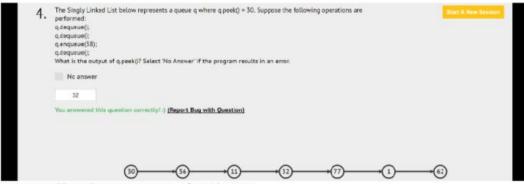
# **Linked list**



Just be careful. Tail pointer is usually not updated. Final\_element.next will lead to NULL (still valid), but attempting to NULL.next or dereference NULL will lead to compilation error -> NO ANSWER



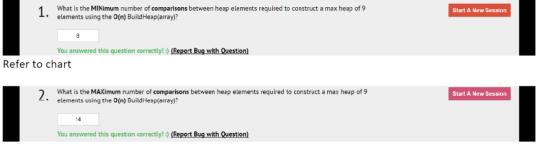
Add and remove stuff from the TOP (64 is the element at the top in this example)



Enqueue adds at the RIGHT, removes from the LEFT

## **Binary Heap**

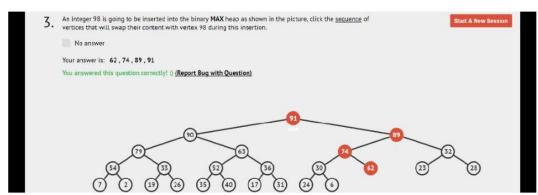
Valid Max Heap: Make sure parent bigger than child



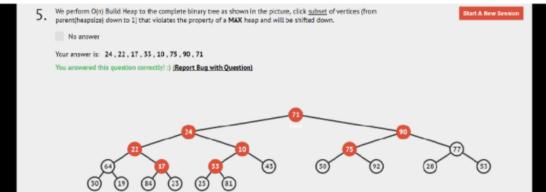
Refer to chart



#### Refer to chart



Add to bottom-most and left-most available spot. Bubble it up.

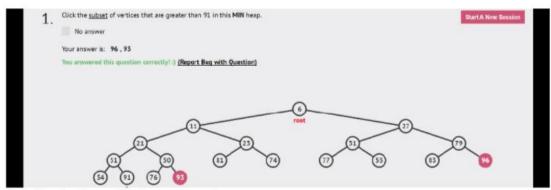


Vertices that are smaller than any of their descendants will be shifted down.



Smallest 17 - 13 = 4 elements are remaining

#### Matthew's VA Cheat Sheet



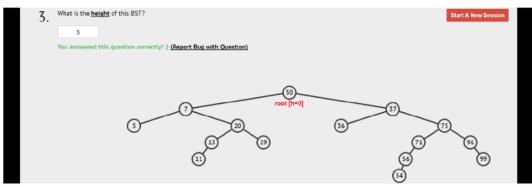
Find all vertex bigger than the query number

# **Binary Search Tree**

BST: Left child < current < Right child.

Successor: Most direct way to the next element bigger than the node of interest Predecessor: Most direct way to the next element smaller than the node of interest

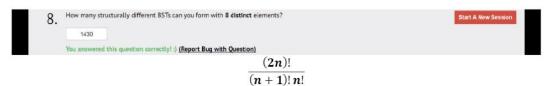
Preorder: Print first, then check children Inorder: Check left, then print, then check right Postorder: Check children first, then print.



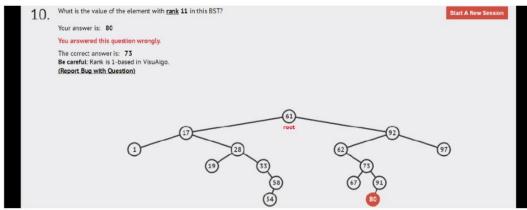
Largest number of edges from root to leaf



If target is bigger, next number should be bigger; if target smaller, next number should be smaller.

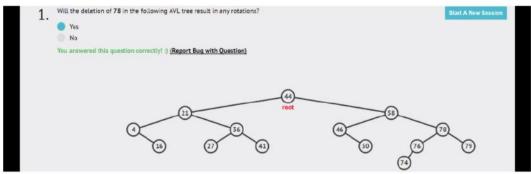


Matthew's VA Cheat Sheet

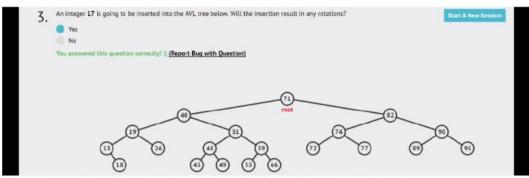


indexed-1. Leftmost is 1. Flatten the tree and count until rank.

# **AVL Tree**

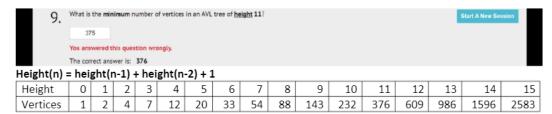


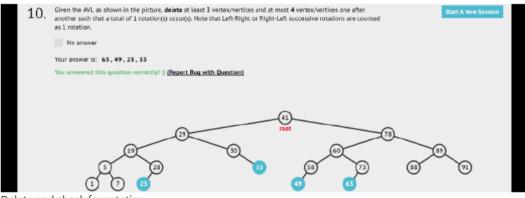
Search upwards starting from the replaced node all the way to root and see if anyone complains



Insert into correct position and check for rotations. Max 1 "fix" required for insertion.

#### Matthew's VA Cheat Sheet





Delete and check for rotation.

## **Graph DS**

Directed acyclic, just make sure small number -> bigger number.

Tree: V vertices, V-1 edges, acyclic.

Complete graph: An edge between any pair of vertices.

Connected graph: From any vertex, can get to all other vertex



Need to enumerate neighbours, BUT NOT ENOUGH MEMORY (we need V^2 entries for AM)



Adjacency Matrix enumerate not as fast but still accepted, so it depends on memory.

#### Matthew's VA Cheat Sheet

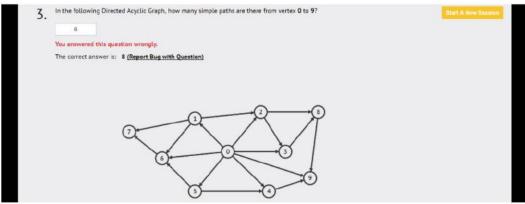
7.	Which best graph D5(es) should you use to store a simple undirected graph with 200 vertices, 19900 edges, and the existence of <b>edge(u, v)</b> is frequently asked?  Suppose your computer only has enough memory to store 40000 entries.  Adjacency Matrix	Start A New Session
	Adjacency List	
	Edge List	
	You answered this question correctly!:) (Report Bug with Question)	
Ack about	existence of edge > AM/Good for Floyd Warshall's 4 line wonder :D)	_

Ask about existence of edge -> AM (Good for Floyd-Warshall's 4 line wonder :D)

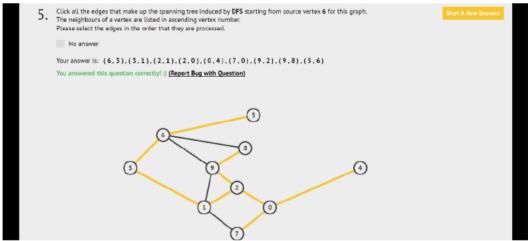


Edges in an edge list can be easily sorted by weight

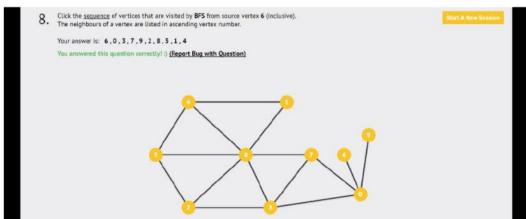
## **Graph traversal**



Write down all the paths

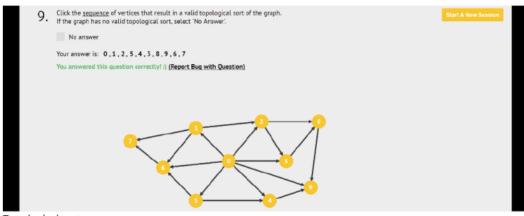


DFS: Go deep, recurse back and go deep for neighbour. Take note of order of neighbours.

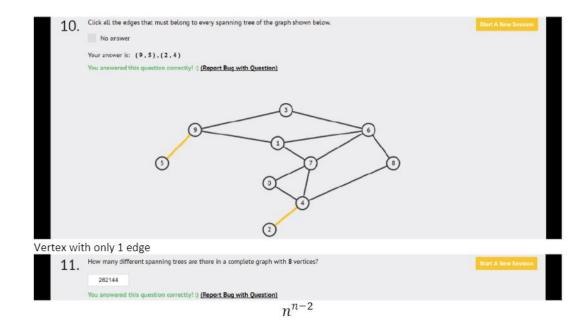


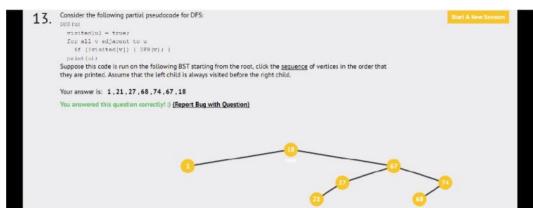
BFS: Visit all neighbours, then all neighbours of neighbours, and so on. Keep track of the order of exploration of neighbours

## Matthew's VA Cheat Sheet



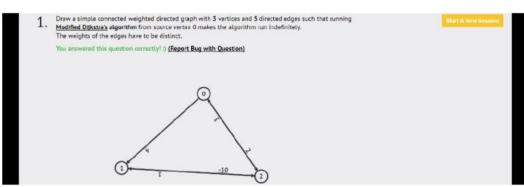
Topological sort



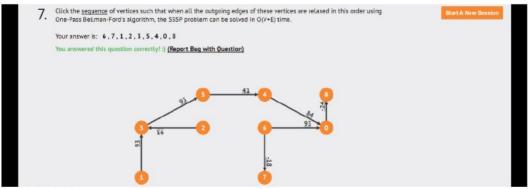


Print comes after exploring all neighbours -> Post-order traversal

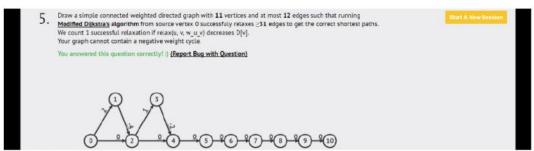
## **SSSP**



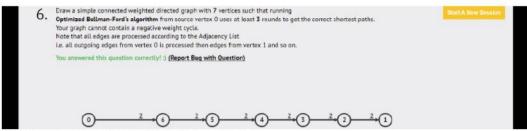
Create a negative cycle. MAKE SURE THE WEIGHTS ARE DISTINCT



Topological sort



Dijkstra Killer. THERE ARE (E-V+1) TRIANGLES.



Bellman-Ford Killer. 0 -> (N-1) -> (N-2) -> ... -> 2 -> 1

# SSSP- Path Weight Criteria

	Bellman-Ford	Original Dijkstra	Modified Dijkstra
Terminate	Always	Always	Does not terminate when there's negative weight cycle
Wrong when there's	Negative eight cycle	Negative edges and Negative weight cycle	Negative weight cycle- it doesn't even terminate in the first place (3)

#### Note:

If the graph contains negative edges but does not contain any negative cycles, Bellman Ford and modified Dijkstra will give the correct answer all the time.

Original Dijkstra *could* still give the correct answer (run Bellman Ford and original Dijkstra and compare the answers)

## <u>Tables</u>

## Max Heap max swaps

Elements	6	7	8	9	10	11	12	13	14
Swaps	4	4	7	7	8	8	10	10	11

## Max Heap max comparisons

Elements	6	7	8	9	10	11	12	13	14
Compare	7	8	11	14	15	16	18	20	21

## Max Heap Min Comparisons: N-1

Elements	6	7	8	9	10	11	12	13	14
Compare	5	6	7	8	9	10	11	12	13

## Binary Search Tree: how many permutations

Elements	1	2	3	4	5	6	7	8	9
Catalan #	1	2	5	14	42	132	429	1430	4862

AVL Tree: Minimum # of vertices (Height(n) = height(n-1) + height(n-2) + 1)

Height	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Vertices	1	2	4	7	12	20	33	54	88	143	232	376	609	986	1596	2583

# **Graph Structures**

Neighbours Frequently Enumerated AND HAS ENOUGH MEMORY: AL + AM (We need V^2 memory for AM)

Neighbours Frequently Enumerated AND NOT ENOUGH MEMORY: AL ONLY (memory <  $V^2$ )

Existence of edge(u,v): AM ONLY

Edges need to be sorted: EDGE LIST ONLY

# Number of spanning trees in complete graph n^(n-2)

Vertic	es	3	4	5	6	7	8	9	10	11
Ans		3	16	625	1296	16807	262144	4782969	100000000	2357947691

## Path Weight Criteria

	Bellman-Ford	Original Dijkstra	Modified Dijkstra
Terminate	Always	Always	Does not terminate when there's negative weight cycle
Wrong when there's	Negative weight cycle	Negative weight (might still be correct if the negative edge doesn't mess up any of the nodes)	Negative weight cycle- it doesn't even terminate in the first place 😟

# Modulo Tables for Hashing

%11

.	0	1	2	3	4	5	6	7	8	9	10
	0	1	2	3	4	5	6	7	8	9	10
	11	12	13	14	15	16	17	18	19	20	21
	22	23	24	25	26	27	28	29	30	31	32
	33	34	35	36	37	38	39	40	41	42	43
	44	45	46	47	48	49	50	51	52	53	54
	55	56	57	58	59	60	61	62	63	64	65
	66	67	68	69	70	71	72	73	74	75	76
	77	78	79	80	81	82	83	84	85	86	87
	88	89	90	91	92	93	94	95	96	97	98
	99	100	101	102	103	104	105	106	107	108	109

%13

3	0	1	2	3	4	5	6	7	8	9	10	11	12
	0	1	2	3	4	5	6	7	8	9	10	11	12
	13	14	15	16	17	18	19	20	21	22	23	24	25
	26	27	28	29	30	31	32	33	34	35	36	37	38
	39	40	41	42	43	44	45	46	47	48	49	50	51
	52	53	54	55	56	57	58	59	60	61	62	63	64
	65	66	67	68	69	70	71	72	73	74	75	76	77
	78	79	80	81	82	83	84	85	86	87	88	89	90
	91	92	93	94	95	96	97	98	99	100	101	102	103
	104	105	106	107	108	109	110	111	112	113	114	115	116

%12

٠.	_			_		_	_	_				
۷	0	1	2	3	4	5	6	7	8	9	10	11
	0	1	2	3	4	5	6	7	8	9	10	11
	12	13	14	15	16	17	18	19	20	21	22	23
	24	25	26	27	28	29	30	31	32	33	34	35
	36	37	38	39	40	41	42	43	44	45	46	47
	48	49	50	51	52	53	54	55	56	57	58	59
	60	61	62	63	64	65	66	67	68	69	70	71
	72	73	74	75	76	77	78	79	80	81	82	83
	84	85	86	87	88	89	90	91	92	93	94	95
	96	97	98	99	100	101	102	103	104	105	106	107
	108	109	110	111	112	113	114	115	116	117	118	119

%14

<u>ا</u> إ	0	1	2	3	4	5	6	7	8	9	10	11	12	13
	0	1	2	3	4	5	6	7	8	9	10	11	12	13
	14	15	16	17	18	19	20	21	22	23	24	25	26	27
	28	29	30	31	32	33	34	35	36	37	38	39	40	41
	42	43	44	45	46	47	48	49	50	51	52	53	54	55
	56	57	58	59	60	61	62	63	64	65	66	67	68	69
	70	71	72	73	74	75	76	77	78	79	80	81	82	83
	84	85	86	87	88	89	90	91	92	93	94	95	96	97
	98	99	100	101	102	103	104	105	106	107	108	109	110	111

# Kattis the Cat for good luck

