

# Tutorial 01 - Basic C++, Basic OOP, Analysis

CS2040C Semester 1 2020/2021

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# Question 1

(Basic) List ADT

## List Array ADT

Imagine *list* as a *chain* of beads

You can add beads to the front and back

(and in between)

We will elaborate more about “**ADT**” in Tut 02.

For now, understand the syntax of C++ classes.

## List Array ADT - Methods

`get(i)`

Gets the  $i$ -th element from the front (0-indexed)

`search(v)`

Return the **first** index which contains  $v$

## List Array ADT - Methods

`insert(i, v)`

Insert element *v* at index *i*.

`remove(i)`

Remove the element at index *i*.

## List Array ADT - Methods

`printList()`

Prints the list from front to back.

`sortList()`

Sort the list, in default ascending order.

## Question 1a, 1b, 1c

```
class ListArray {  
    private:  
        int N;  
        int A[10];           // question 1a  
    public:  
        ListArray() : N(0) {} // question 1b  
        int get(int i) {  
            return A[i];      // question 1c  
        }  
}
```

## Question 1a

```
class ListArray {  
    private:  
        int N;  
        int A[10];  
    public:  
        ListArray() : N(0) {}  
        int get(int i) {  
            return A[i];  
        }  
}
```

Anything wrong with this line?



## Question 1a

```
class ListArray {  
    private:  
        int N;  
        int A[10];  
    public:  
        ListArray() : N(0) {}  
        int get(int i) {  
            return A[i];  
        }  
}
```

Anything wrong with this line?

Limited to **10** items!

## Question 1b

```
class ListArray {  
    private:  
        int N;  
        int A[10];  
    public:  
        ListArray() : N(0) {}  
        int get(int i) {  
            return A[i];  
        }  
}
```

What does this line mean?

## Question 1b

```
class ListArray {  
    private:  
        int N;  
        int A[10];  
    public:  
        ListArray() : N(0) {}  
        int get(int i) {  
            return A[i];  
        }  
}
```

What does this line mean?

Pass 0 to constructor of  
integer N. Effectively:

```
ListArray() {  
    N = 0;  
}
```

## Question 1c

```
class ListArray {  
    private:  
        int N;  
        int A[10];  
    public:  
        ListArray() : N(0) {}  
        int get(int i) {  
            return A[i];  
        }  
}
```

Anything potential issues with this line?

## Question 1c

```
class ListArray {  
    private:  
        int N;  
        int A[10];  
    public:  
        ListArray() : N(0) {}  
        int get(int i) {  
            return A[i];  
        }  
}
```

Anything potential issues with this line?

No “safeguard”!

What if `i > N - 1`?

What if `i` is negative?

## Question 1d, 1e, 1f

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N)) // question 1d  
        return;  
    for (int j = i; j <= N-1; j++) // question 1e  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}  
  
void remove(int i) {  
    for (int j = i; j < N-1; j++) // question 1f  
        A[j] = A[j+1];  
    N--;  
}
```

## Question 1d

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N))  
        return;  
    for (int j = i; j <= N-1; j++)  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}
```

What does this line mean?

## Question 1d

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N))  
        return;  
    for (int j = i; j <= N-1; j++)  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}
```

What does this line mean?

If ...  
N is 10 or  
i is negative or  
i > N

Stop inserting.



## Question 1d

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N))  
        return;  
    for (int j = i; j <= N-1; j++)  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}
```

What does this line mean?

If ...  
N is 10 or  
i is negative or  
i > N

Stop inserting.

Is there anything wrong with this?

How many possible values of i are accepted?

## Question 1d

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N))  
        return;  
    for (int j = i; j <= N-1; j++)  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}
```

What does this line mean?

If ...  
N is 10 or  
i is negative or  
i > N

Stop inserting.

Is there anything wrong with this?

How many possible values of i are accepted?

N+1 possible insertion points

## Question 1e

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N))  
        return;  
    for (int j = i; j <= N-1; j++)  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}
```

Any potential issues with this line?

## Question 1e

```
void insert(int i, int v) {  
    if ((N == 10) || (i < 0) || (i > N))  
        return;  
    for (int j = i; j <= N-1; j++)  
        A[j+1] = A[j];  
    A[i] = v;  
    N++;  
}
```

Any potential issues with this line?

Values are being overwritten in the wrong order!

```
A[i+1] = A[i];  
A[i+2] = A[i+1];  
A[i+3] = A[i+2];
```

... etc

## Question 1f

```
void remove(int i) {  
    for (int j = i; j < N-1; j++)  
        A[j] = A[j+1];  
    N--;  
}
```

Any potential issues with this line?

## Question 1f

```
void remove(int i) {  
    if (i < 0 || i >= N)  
        return;  
    for (int j = i; j < N-1; j++)  
        A[j] = A[j+1];  
    N--;  
}
```

Any potential issues with this line?

No. This is correct.

```
A[i]    = A[i+1];  
A[i+1]  = A[i+2];  
A[i+2]  = A[i+3];  
... etc;
```

Can also add safeguard `i` in `[0..N-1]`

## Question 1g

```
void sortList() { // sort array A, question 1g  
    // ...  
}
```

Implement this routine using any sorting algorithm that you know!

## Question 1g

```
void sortList() { // sort array A, question 1g
    sort(A, A+N);
}
```

Implement this routine using any sorting algorithm that you know!

Approach: Use std library



## Question 1h

```
int main() {  
    ListArray* LA = new ListArray();  
    LA->insert(0, 5);  
    LA->insert(0, 1);  
    // ...  
}
```

Can we just write `ListArray LA;` in this line?

## Question 1h

```
int main() {  
    ListArray* LA = new ListArray();  
    LA->insert(0, 5); // (*LA).insert(0,5)  
    LA->insert(0, 1);  
    // ...  
}
```

Can we just write `ListArray LA;` in this line?

Yes, but we need to modify the way we call methods:

```
ListArray LA;  
LA.insert(0, 5);  
cout << LA.get(1) << endl;
```

## Question 1h

What's the difference between the two methods?

```
ListArray LA;  
LA.insert(0, 5);  
cout << LA.get(1) << endl;
```

```
ListArray *LA = new ListArray();  
LA->insert(0, 5);  
cout << LA->get(1) << endl;
```

## Question 1g

```
void sortList() { // sort array A, question 1g
    for (int i = 0; i < N; i++) {
        for (int j = 1; j < N; j++) {
            if (A[j-1] > A[j])
                swap(A[j-1], A[j]);
        }
    }
}
```

Implement this routine using any sorting algorithm that you know!

Approach: manual implementation

# Question 3

Analysis/ Order of Growth

# Complexity analysis

- Is a rough estimate of how execution time will grow with size of input. I.E. *Order of growth*
- *Time complexity* is commonly used as a metric for comparing the performance of different algorithms on the same task
- When order of growth is applied to measure **memory consumption** of algorithms, we call that *space complexity*

# Two rules

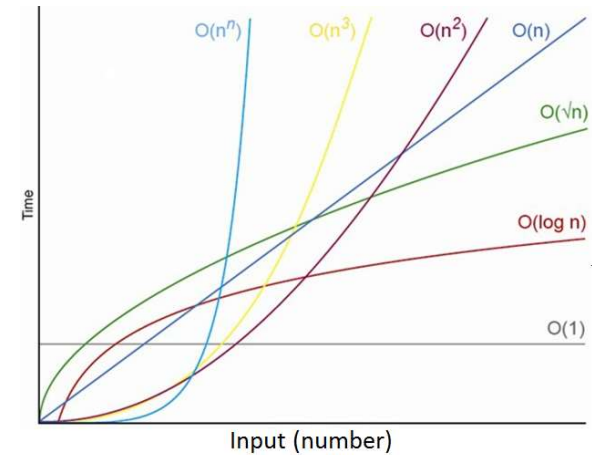
## Two rules:

1. Constants can be ignored

$$O(3n) = O(n)$$

2. Smaller complexity terms can be ignored

$$O(n + \log(n) + n^2) = O(n^2)$$



## Real life application?

Single thread computer:  $\sim 10^8$  operations/sec

Sorting  **$N = 10^6$**  numbers

$O(n^2)$   $\rightarrow$   $\sim 2.5$  hours

$O(n \times \log(n))$   $\rightarrow$   $\sim 0.2$  seconds

Kattis runs submissions at around  $4 \times 10^8$  operations/second.



## Question 3

What is the bound to the following functions?

$$a. F(n) = \log(2^n) + \sqrt{n} + 100\,000\,000$$

$$b. F(n) = n + \frac{1}{2}n + \frac{1}{3}n + \frac{1}{4}n + \dots + 1$$

$$c. G(n) = n + \frac{1}{2}n + \frac{1}{4}n + \frac{1}{8}n + \dots + 1$$

### Question 3.a)

$$F(n) = \log(2^n) + \sqrt{n} + 100\,000\,000$$

$$O(F(n)) = O(\log(2^n) + \sqrt{n} + 100\,000\,000)$$

$$= O(n * \mathbf{\log 2} + \sqrt{n})$$

$$= O(n + \sqrt{n})$$

$$= \mathbf{O(n)}$$

### Question 3.b)

$$F(n) = n + \frac{1}{2}n + \frac{1}{3}n + \frac{1}{4}n + \dots + 1$$

$$O(F(n)) = O(n + \frac{1}{2}n + \frac{1}{3}n + \frac{1}{4}n + \dots + 1)$$

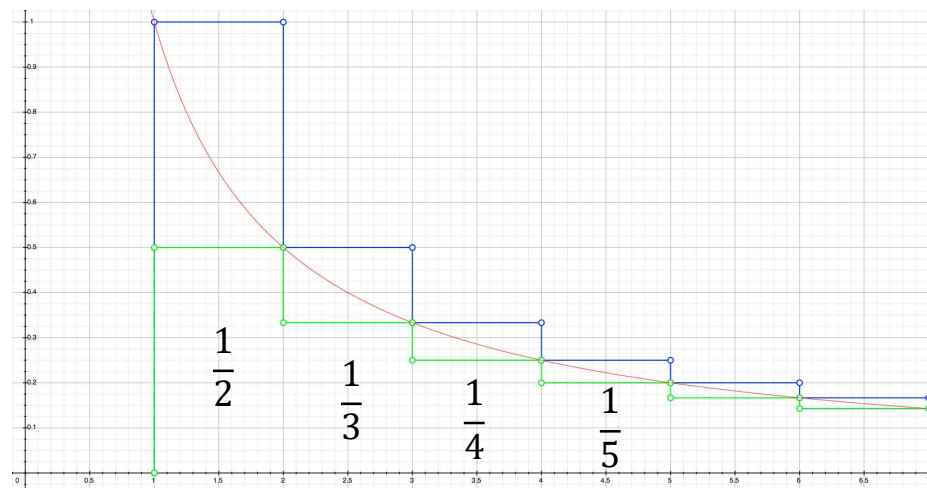
$$= O(n (\underbrace{1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}}_{\text{How to deal with this?}}))$$

How to deal with this?

# Harmonic Series

$$\sum_{n=1}^{\infty} \frac{1}{n} = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \frac{1}{5} + \dots$$

Using integration:



$$\log(N+1) < \sum_{n=1}^N \frac{1}{n} = 1 + \sum_{n=1}^N \frac{1}{n+1} < 1 + \int_1^N \frac{dx}{x+1} = 1 + \log(N+1) = O(\log N)$$

You can try prove the lower bound as well.

### Question 3.b)

*From Harmonic Series*

$$O(1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n}) = O(\log n)$$

$$\begin{aligned} O(F(n)) &= O(n (1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n})) \\ &= \mathbf{O(n \log n)} \end{aligned}$$

### Question 3.c)

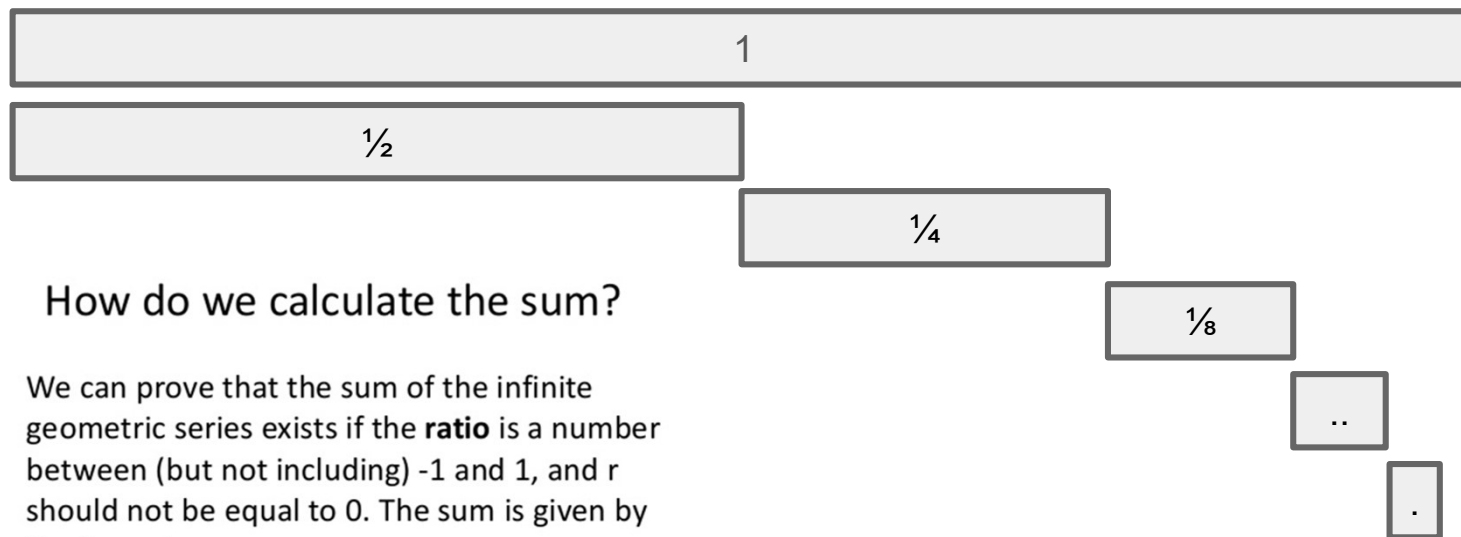
$$G(n) = n + \frac{1}{2}n + \frac{1}{4}n + \frac{1}{8}n + \dots + 1$$

$$O(G(n)) = O(n + \frac{1}{2}n + \frac{1}{4}n + \frac{1}{8}n + \dots + 1)$$

$$= O(n (\underbrace{1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + 1/n}_{\text{How to deal with this?}}))$$

How to deal with this?

# Convergent Geometric Series



How do we calculate the sum?

We can prove that the sum of the infinite geometric series exists if the **ratio** is a number between (but not including) -1 and 1, and  $r$  should not be equal to 0. The sum is given by the formula:

$$\sum_{k=0}^{\infty} ar^k = a + ar + ar^2 + ar^3 + \dots = \frac{a}{1-r}$$

### Question 3.c)

*From Convergent Geometric Series*

$$O(1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{n}) = O(1 / (1 - \frac{1}{2}))$$

$$= O(2)$$

$$= O(1)$$

$$O(G(n)) = O(n (1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots + \frac{1}{n}))$$



# Additional questions

Example (AY17/18 S1 Midterm Paper)

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = N; i >= 1; i--) {
    for (int j = 1; j <= N/i; j++) {
        counter++;
    }
}
cout << counter << endl;
```

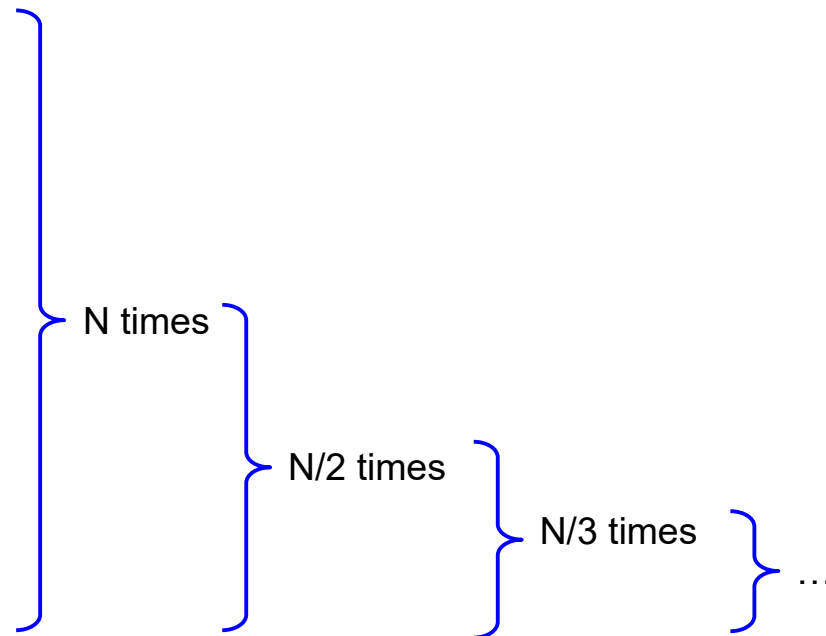
What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = N; i >= 1; i--) {
    for (int j = 1; j <= N/i; j++) {
        counter++;
    }
}
cout << counter << endl;
```

***$O(N \log N)$***

# Analysis

- $i == N, j == 1$
- $i == N-1, j == 1$
- $i == N-2, j == 1$
- ...
- $i == N/2, j == 1, j == 2$
- ...
- $i == N/3, j == 1, j == 2, j == 3$
- ...
- $i == 1, j == 1 \text{ to } j == N$



$$N + N/2 + N/3 + N/4 + N/5 + N/6 + \dots + 1$$

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = 0; i < N; i++) {
    for (int j = 0; j < i; j++) {
        counter += j;
    }
}
cout << counter << endl;
```

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = 0; i < N; i++) {
    for (int j = 0; j < i; j++) {
        counter += j;
    }
}
cout << counter << endl;
```

**$O(N^2)$**

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = 0; i < N; i++) {
    for (int j = 0; j < N; j++) {
        counter++;
        i++;
    }
}
cout << counter << endl;
```

What's the time complexity?

```
int N, counter = 0;
cin >> N;
for (int i = 0; i < N; i++) {
    for (int j = 0; j < N; j++) {
        counter++;
        i++;
    }
}
cout << counter << endl;
```

$O(N)$



# Questions

`/basicprogramming1`