### Question 1 (a).

Since ABC  $\cap$  CD = C and C is a superkey of R 2 (due to the FD C  $\rightarrow$  D), the decomposition is lossless.

## Question 1 (b).

Since AC  $\cap$  ABD = A and A is a superkey of R 1 (due to the FD A  $\rightarrow$  BCD), the decomposition is lossless.

### Question 1 (c).

A decomposition is lossless join if there exists a sequence of binary lossless-join decompositions that generates that decomposition. It turns out that there is one such decomposition. First, decompose R into R3(A,C,D) and R4(A,B,C,E). This is lossless because R3  $\cap$  R4 = AC, and AC is a superkey of R3. Next, decompose R4 into R1(A,B,C) and R2(A,B,E). This is also lossless because R1  $\cap$  R2 = AB, and AB is a superkey of R1. Therefore, {R1(A,B,C), R2(A,B,E), R3 (A,C,D)} is a lossless-join decomposition.

#### Question 2 (a).

R is not in BCNF because A  $\rightarrow$  E violates BCNF: A is not a superkey of R as {A}+ = ABDE.

#### Question 2 (b).

The decomposition is lossless-join because  $R1 \cap R2 = \{E\}$  and  $E \rightarrow BDE$ .

### Question 2 (c).

Note that {A}+=ABDE. Therefore, on R2, we have {A}+=AE. This indicates that R2 is not in BCNF.

### Question 2 (d).

- 1. {E}+ = {EDB}, which indicates that R is not in BCNF. Accordingly, we decompose R into R1(E, B, D) and R2(E, A, C).
- 2. The only non-trivial FD on R1 is  $E \rightarrow DB$ , and E is a key of R1. Thus, R1 is in BCNF.
- 3. {A}+={AEDB}, and hence, we have {A}+={AE} on R2. This indicates that R2 is not in BCNF. Accordingly, we decompose R2 into R3(A, E) and R4(A, C).
- 4. R3 and R4 are in BCNF since each of them contains only two attributes.
- 5. Therefore, { R1(E, D, B), R3(A, E), R4(A, C) } is a BCNF decomposition of R.

# Question 3.

- 1. {B}+=BCDE, which indicates that R is not in BCNF. Accordingly, we decompose R into R1(B, C, D, E) and R2(B, A).
- 2. Since R2 has only two attributes, it is in BCNF.
- 3. For R1, let's examine the closure of each attribute subset:
  - {B}+={BCDE}, which indicates that B is a key of R1.
  - {C}+={BCDE}, which indicates that C is a key of R1.
  - {D}+={D}
  - {E}+={E}
  - {DE}+={DE}
  - All other attribute subsets are superkeys since they are supersets of either B or C.

Therefore, R1 is in BCNF.

4. As a consequence, {R1(B, C, D, E), R2(A, B)} is a BCNF decomposition of R.