

[SWE2015-41] Introduction to Data Structures (자료구조개론)

#### **Variants of Linked Lists**

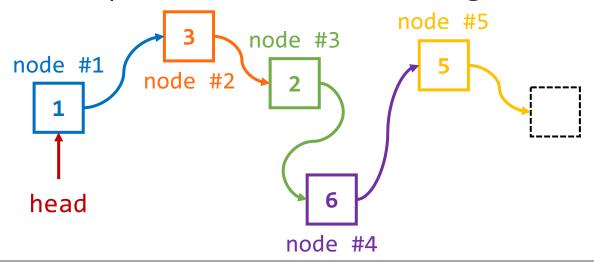
**Department of Computer Science and Engineering** 

Instructor: Hankook Lee (이한국)

# (Recap) Linked Lists



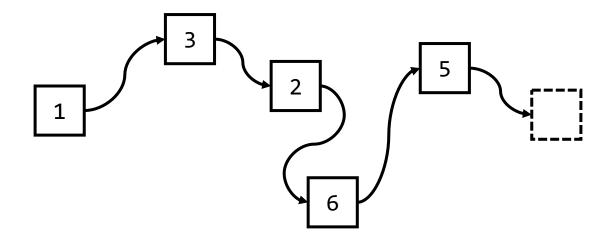
- A linked list is a collection of sequentially-connected elements
  - The elements are not required to be stored in contiguous memory



```
typedef struct _Node {
    int value;
    struct _Node *next;
} Node;
typedef struct _LinkedList {
    Node *head;
} LinkedList;
```

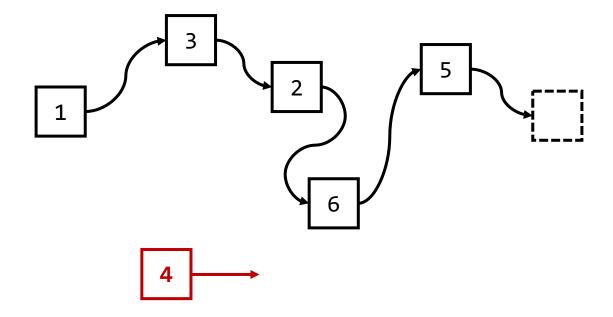


- How to insert an item at the middle of the linked list?
  - Example: Insert "4" at the 2<sup>nd</sup> position





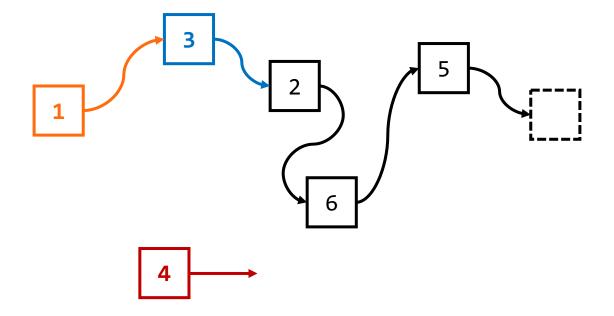
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Create new node



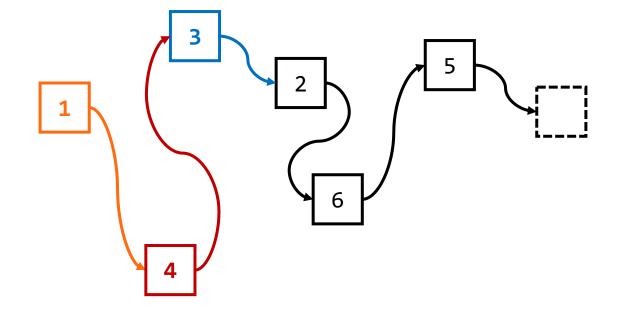
- How to insert an item at the middle of the linked list?
  - Example: Insert "4" at the 2<sup>nd</sup> position



- 1. Create new node
- 2. Find the 1st node and its next node



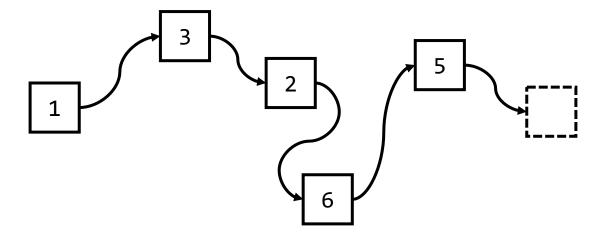
- How to insert an item at the middle of the linked list?
  - Example: Insert "4" at the 2<sup>nd</sup> position



- 1. Create new node
- 2. Find the 1st node and its next node
- 3. Connect them: node → node → node

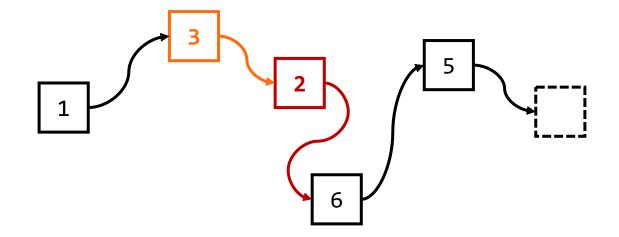


- How to delete an item from the linked list?
  - Example: delete "2" at the 3<sup>rd</sup> position





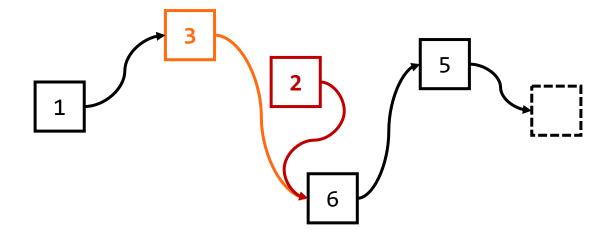
- How to delete an item from the linked list?
  - Example: delete "2" at the 3<sup>rd</sup> position



1. Find the 2<sup>nd</sup> node and its next node



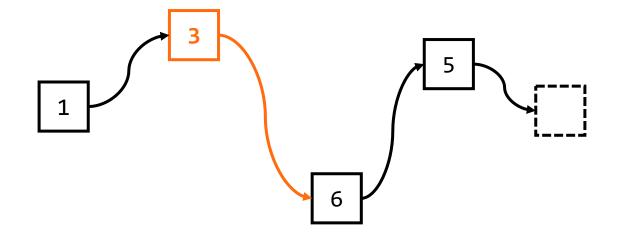
- How to delete an item from the linked list?
  - Example: delete "2" at the 3<sup>rd</sup> position



- 1. Find the 2<sup>nd</sup> node and its next node
- 2. Skip the connection between node and node



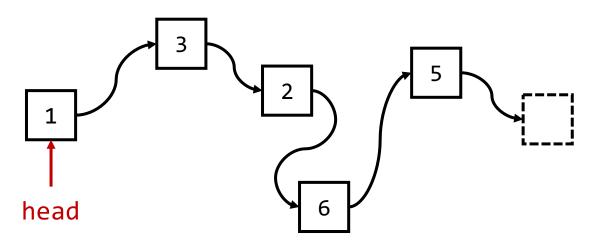
- How to delete an item from the linked list?
  - Example: delete "2" at the 3<sup>rd</sup> position



- 1. Find the 2<sup>nd</sup> node and its next node
- 2. Skip the connection between node and node
- 3. Free the node's memory



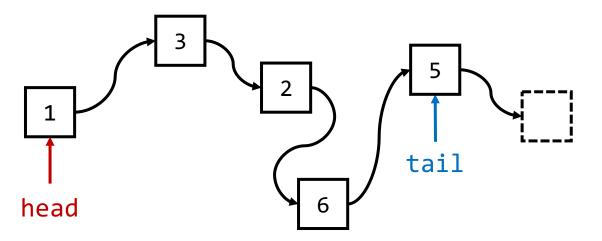
- Singly Linked Lists
  - The simplest linked list structure
  - One-directional → cannot go backward
  - O(n) for insertion at the end





- Singly Linked Lists with the tail pointer
  - The simplest linked list structure
  - One-directional → cannot go backward
  - *O*(1) for insertion at the end

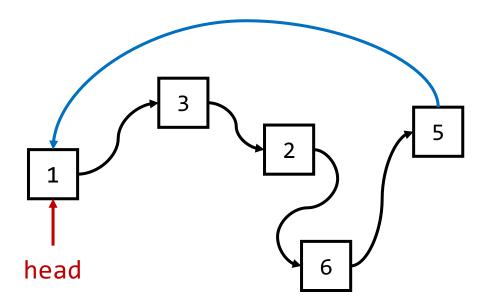
```
typedef struct _LinkedList {
   Node *head, *tail;
} LinkedList;
```





- Circular Linked Lists
  - Every node has non-NULL next pointer → makes handling corner cases easier

```
typedef struct _Node { int value; struct _Node *next; } Node;
typedef struct _LinkedList { Node *head; } LinkedList;
```

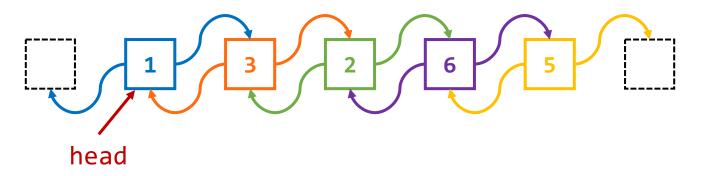




- Doubly Linked Lists
  - Every node has prev and next pointers to point previous and next nodes

```
typedef struct _Node {
   int value;
   struct _Node *prev, *next;
} Node;
```

Two-directional → can move forward and backward



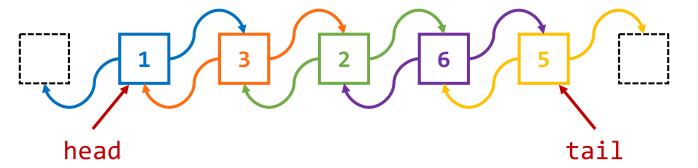


Let's implement a doubly linked list with the tail pointer

```
typedef struct _Node {
   int value;
   struct _Node *prev, *next;
} Node;

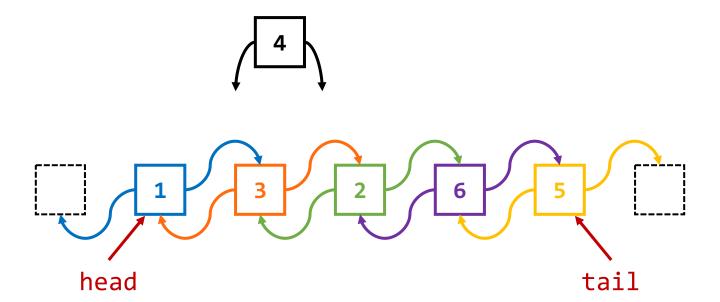
typedef struct _LinkedList {
   Node *head, *tail;
} LinkedList;

void insert(LinkedList *list, int item, int index);
void delete(LinkedList *list, int index);
```



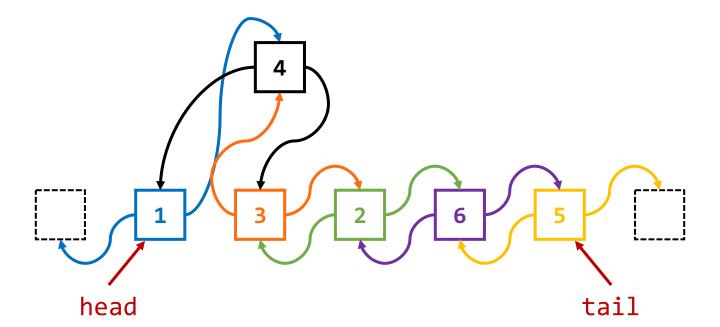


- Let's implement a doubly linked list with the tail pointer
  - How to insert an element 4 between two nodes 1 and 3?



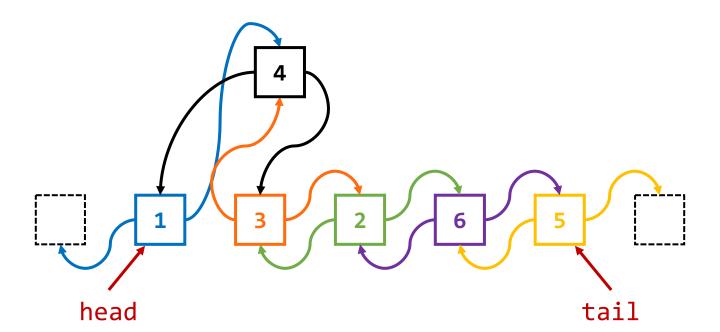


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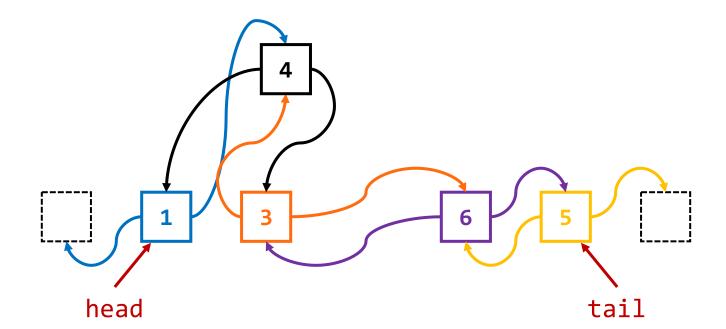


- Let's implement a doubly linked list with the tail pointer
  - How to insert an element 4 between two nodes 1 and 3?
  - How to delete the element 2?



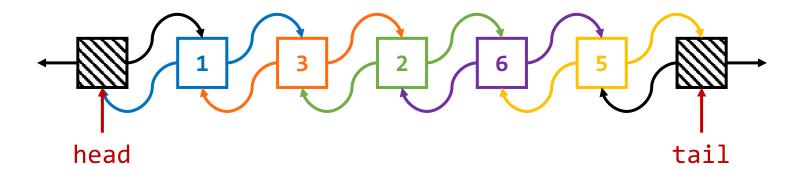


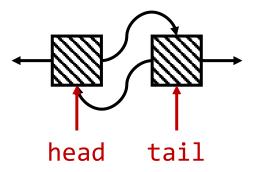
- Let's implement a doubly linked list with the tail pointer
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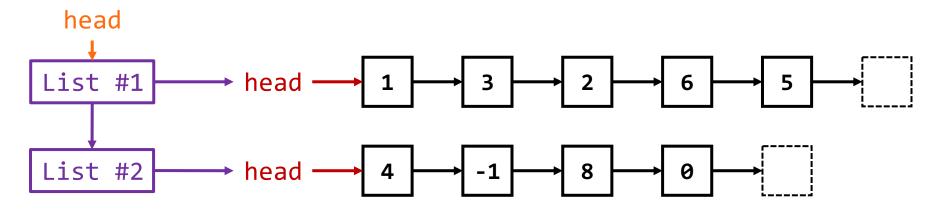
- A technique for easy implementation: Dummy nodes
  - This guarantees that head and tail are always non-NULL even if the list is empty
  - This guarantees that intermediate nodes have non-NULL prev and next pointers







Two-dimensional Linked Lists (like 2D arrays)

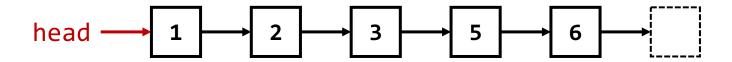


```
typedef struct _Node { int value; struct _Node *next; } Node;
typedef struct _LinkedList { // This represents a "list" node
   Node *head;
   struct _LinkedList *next; // This connects two linked lists
} LinkedList;

typedef struct _LinkedList2D {
   LinkedList *head;
} LinkedList2D;
```

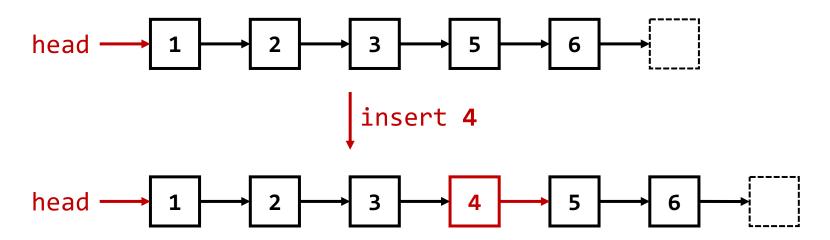


- Sorted Linked Lists
  - A linked list whose elements should be always sorted



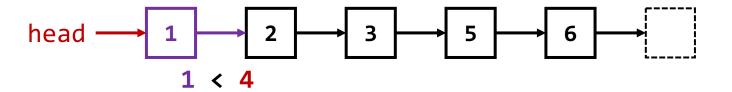


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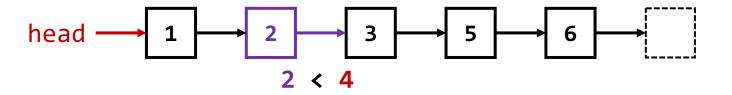


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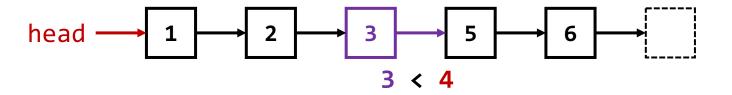


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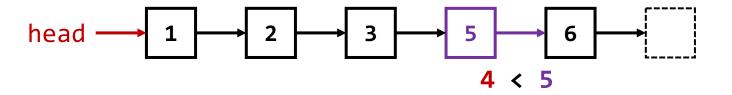


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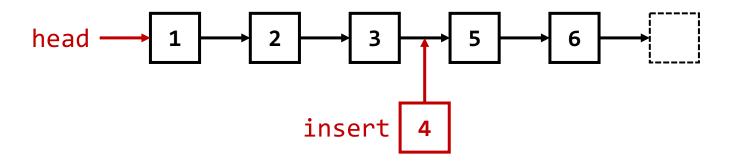


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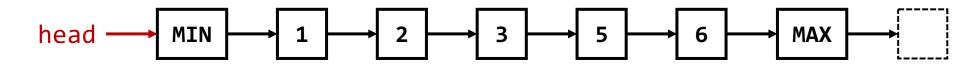


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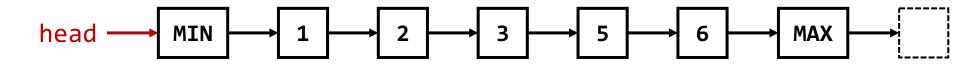
For easy implementation, one can consider dummy nodes, MIN and MAX

```
typedef struct _Node { int value; struct _Node *next; } Node;
typedef struct _LinkedList { Node *head; } LinkedList;

void search(LinkedList *list, int value);
void insert(LinkedList *list, int value);
void delete(LinkedList *list, int value);
```



- Sorted Linked Lists
  - A linked list whose elements should be always sorted



For easy implementation, one can consider dummy nodes, MIN and MAX

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typedef struct _Node { int value; struct _Node *next; } Node;
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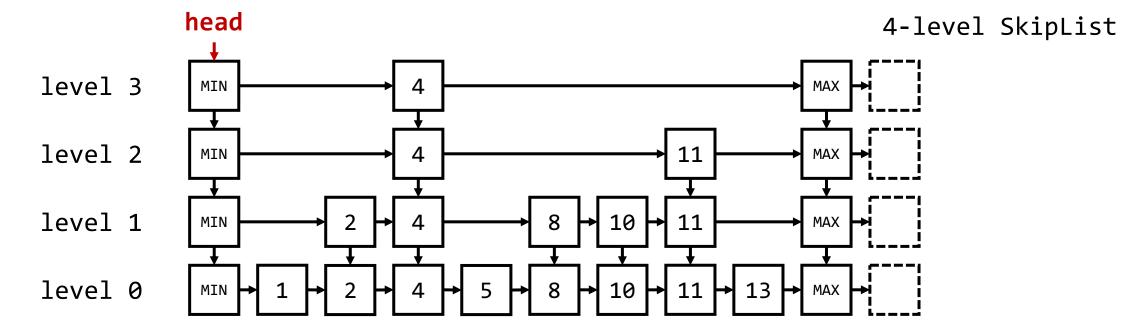
• The search, insertion, deletion operations require O(n) time complexity



- A Skip List is an advanced variant of ordered/sorted linked lists
  - This can be implemented by a two-dimensional linked list

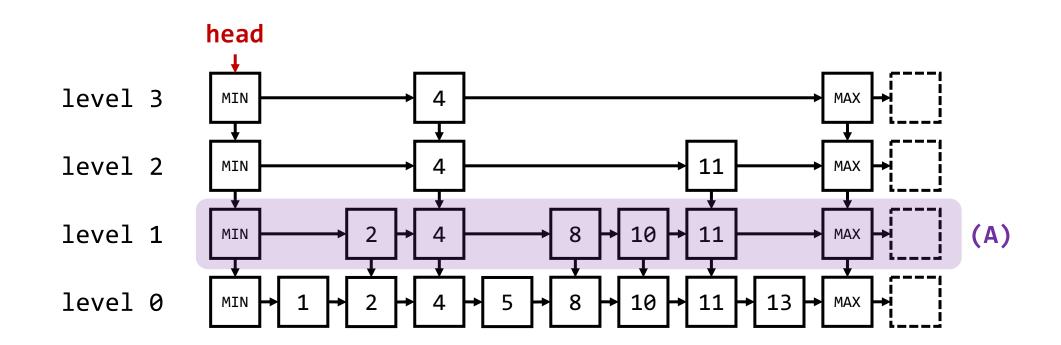
```
typedef struct _Node {
   int value;
   struct _Node *next, *lower;
} Node;
Node ? next
lower
```

```
typedef struct _SkipList {
    Node *head;
    int num_levels;
} SkipList;
```



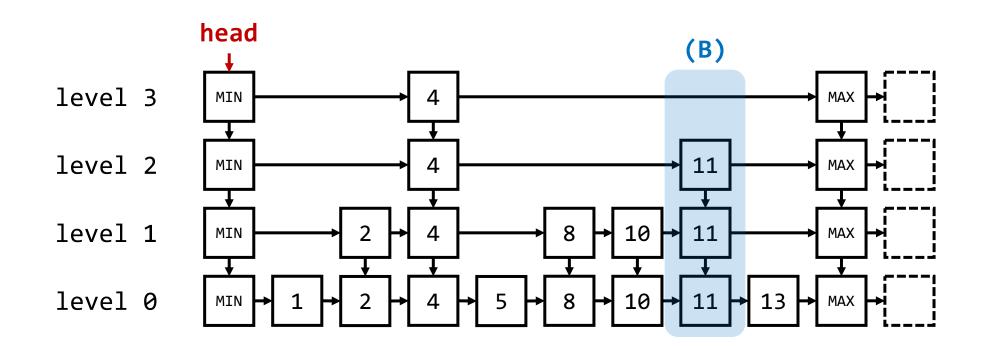


- A Skip List is an advanced variant of ordered/sorted linked lists
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  - Two properties:
    - (A) Elements at each level (i.e., row) are sorted



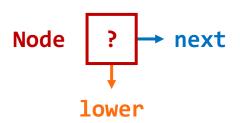


- A Skip List is an advanced variant of ordered/sorted linked lists
  - This can be implemented by a two-dimensional linked list
  - Two properties:
    - (A) Elements at each level (i.e., row) are sorted
    - (B) Elements at each column exist consecutively

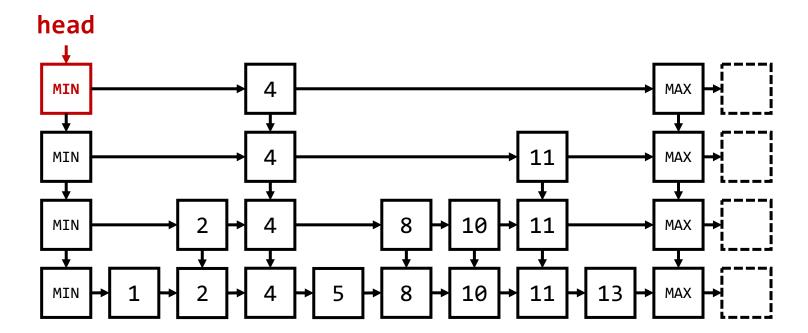




- The search procedure in Skip Lists
  - Starting from head,

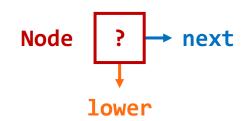


target: 10

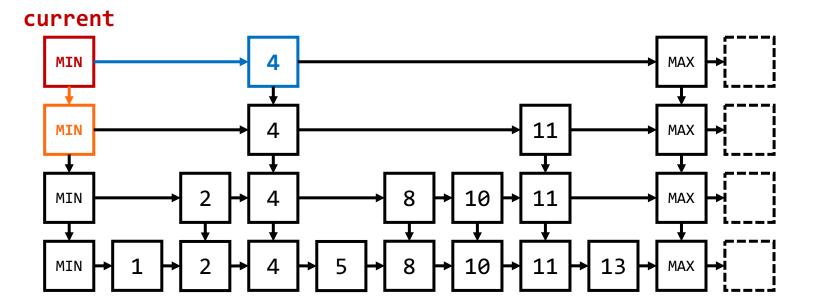




- The search procedure in Skip Lists
  - Starting from head,
  - 2. Compare the target value with the next node

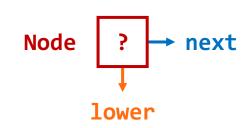


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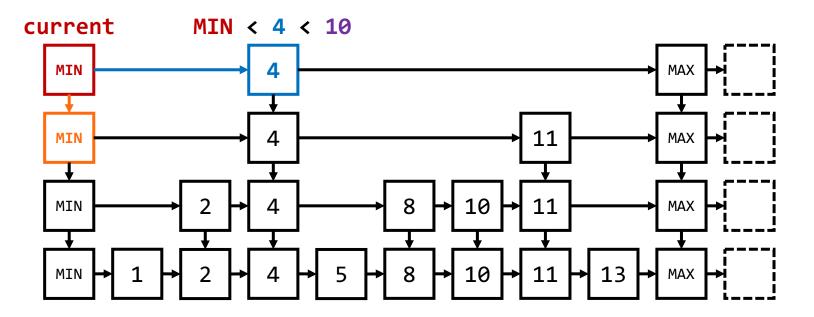




- The search procedure in Skip Lists
  - Starting from head,
  - 2. Compare the **target** value with the **next** node
    - 1. If current < next <= target, then move to the next node

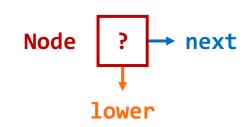


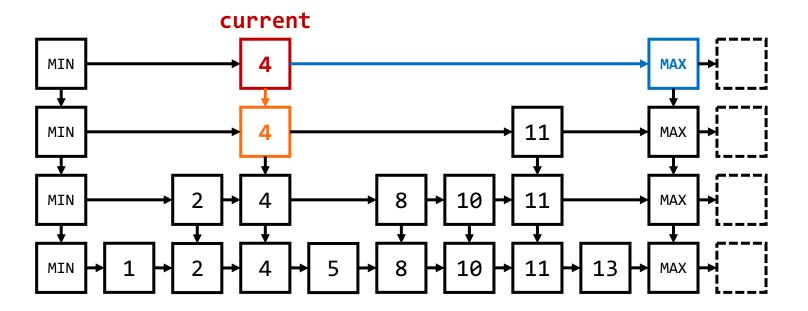
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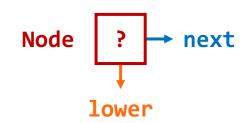
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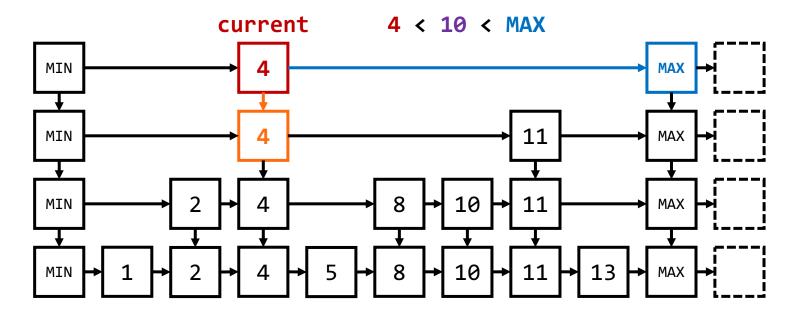






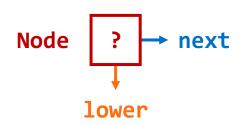
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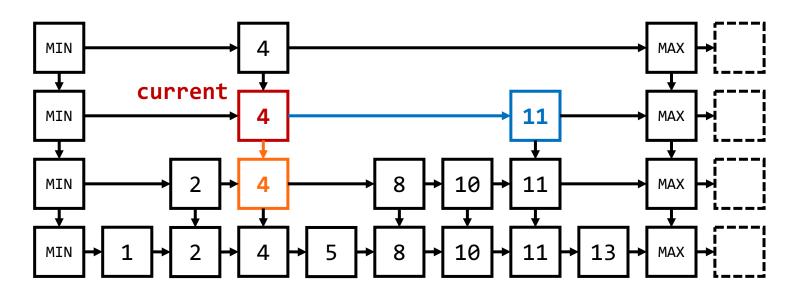






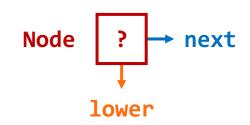
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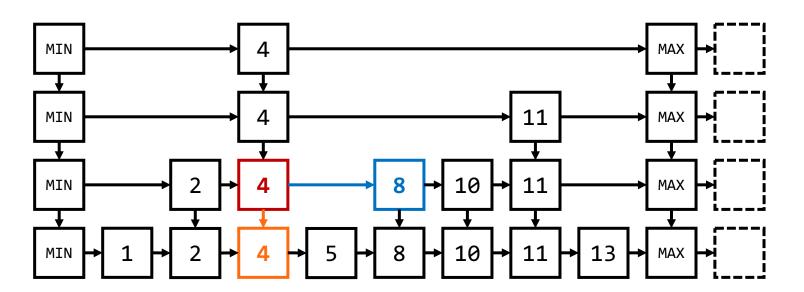






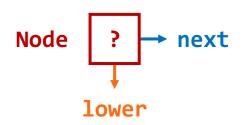
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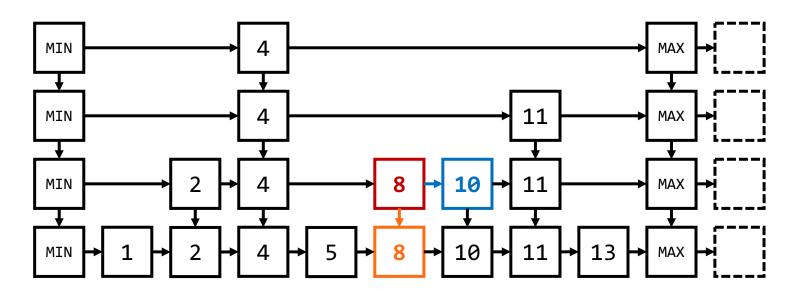






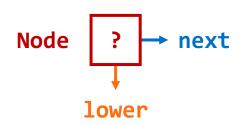
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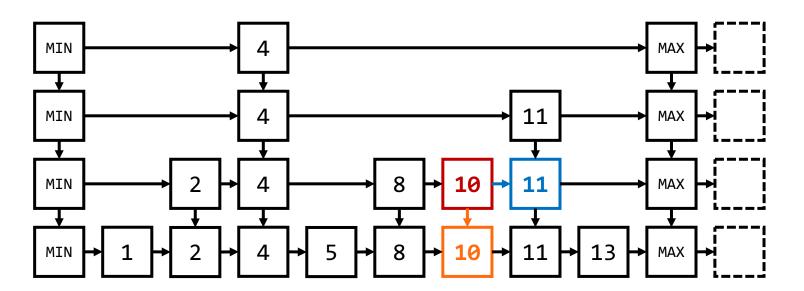






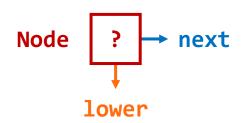
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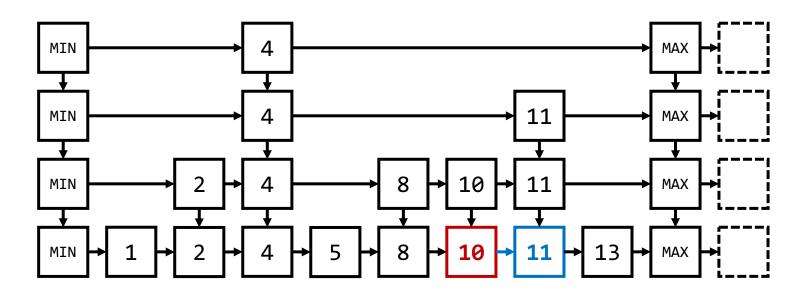






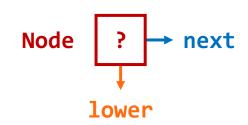
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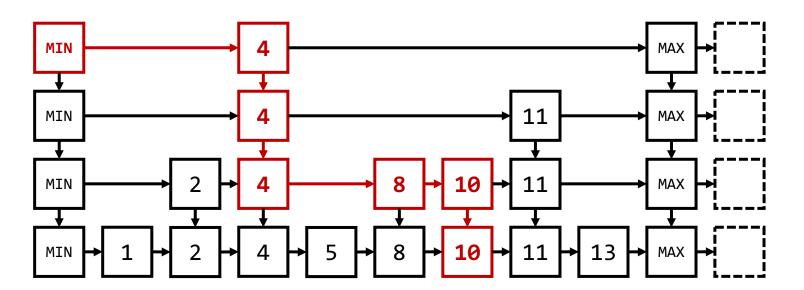






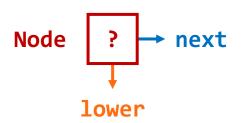
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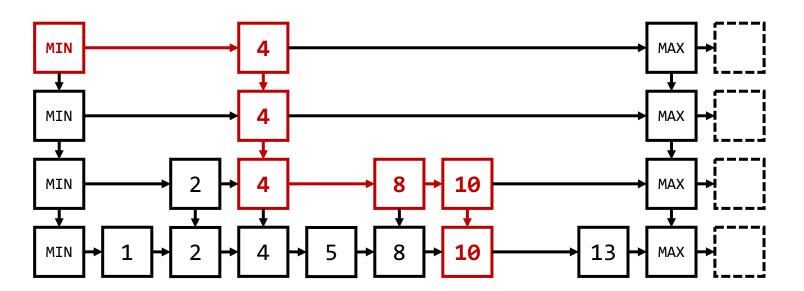






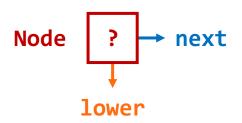
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  - 1. For all levels, find rightmost **nodes < target**

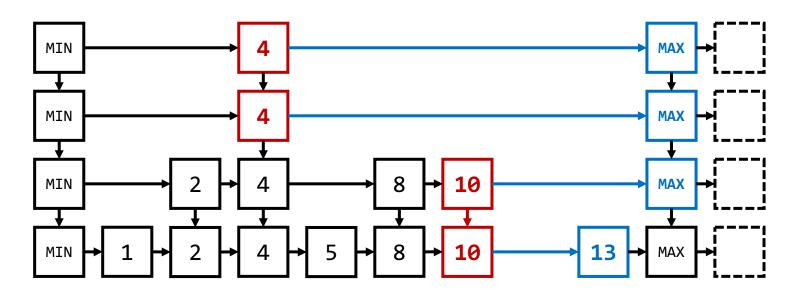






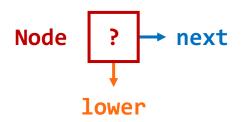
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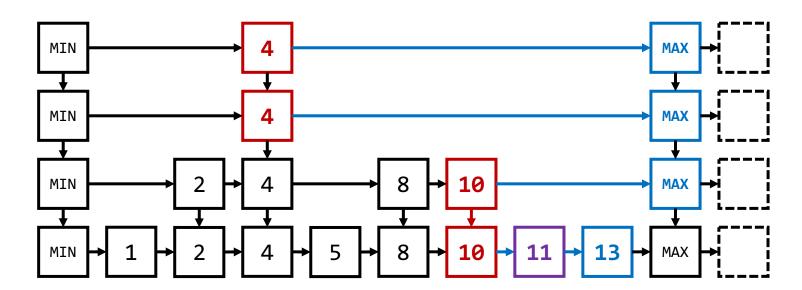






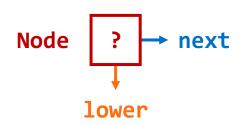
- The insertion procedure in Skip Lists
  - 1. For all levels, find rightmost **nodes < target**
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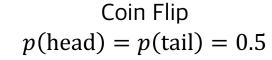




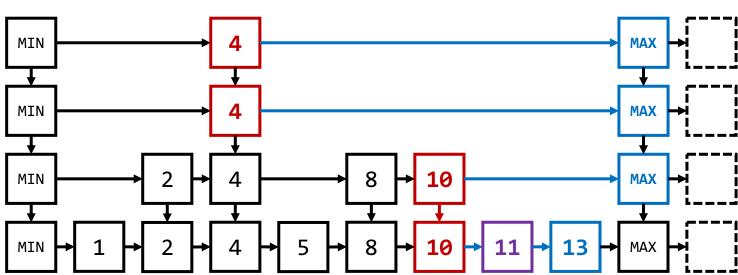


- The insertion procedure in Skip Lists
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    - 1. If the flipped coin shows head, continue to insert



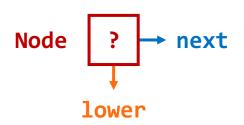


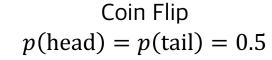




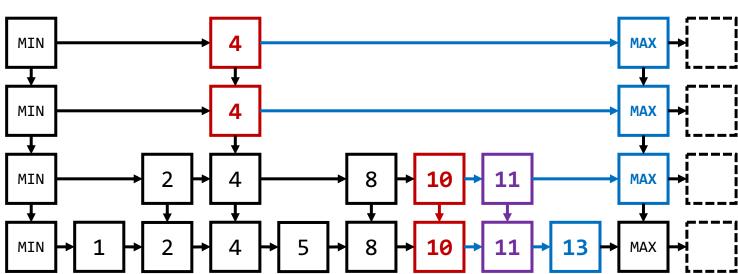


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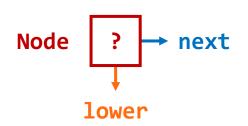


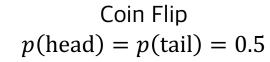




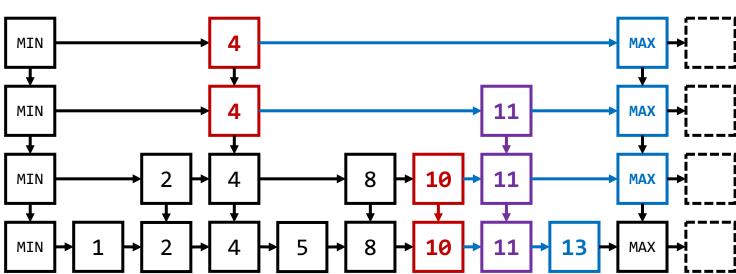


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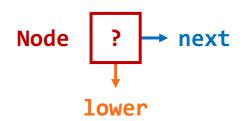


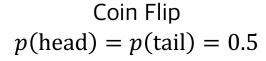




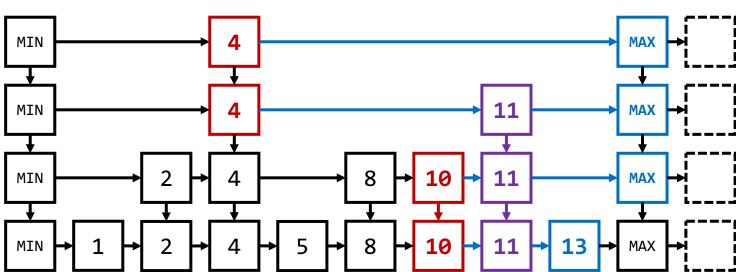


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  - 2. Insert **nodes** from the bottom to the top
    - 1. If the flipped coin shows head, continue to insert
    - 2. If the flipped coin shows tail, stop insertion



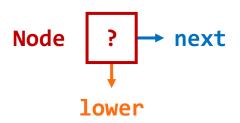


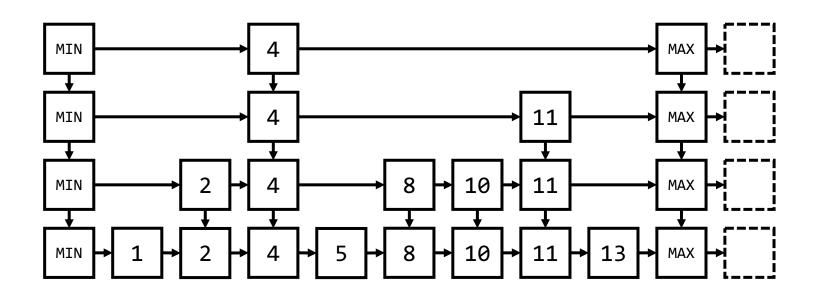






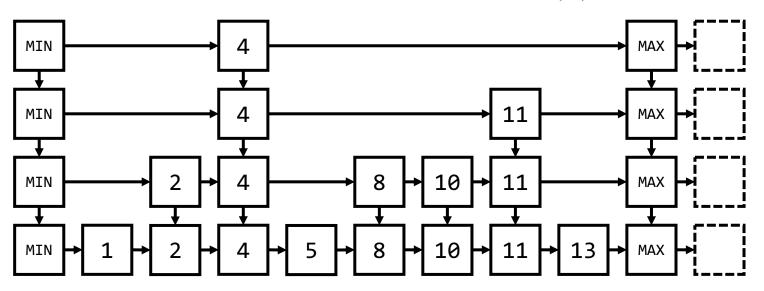
- The insertion procedure in Skip Lists
  - 1. For all levels, find rightmost **nodes < target**
  - 2. Insert **nodes** from the bottom to the top
    - 1. If the flipped coin shows head, continue to insert
    - 2. If the flipped coin shows tail, stop insertion







- Skip Lists A probabilistic data structure for an ordered sequence
  - This allows  $O(\log n)$  average time complexity for search
  - This allows  $O(\log n)$  average time complexity for insertion/deletion while maintaining the order
  - # of nodes is roughly N + N/2 + N/4 + ... = 2N
  - Therefore, the space complexity of the skip list is O(n)



# **Time & Space Complexities**



#### Time Complexity

Operation	(standard) Array	(standard) Linked List	Skip List (average case)	Skip List (worst case)
Insertion	O(n)	0(1)	$O(\log n)$	O(n)
Deletion	O(n)	0(1)	$O(\log n)$	O(n)
Search by Index	0(1)	O(n)	$O(\log n)$	O(n)
Search by Value	O(n)	O(n)	$O(\log n)$	O(n)
Reversion	O(n)	O(n)	-	-

#### Space Complexity

(standard)	(standard)	Skip List	Skip List
Array	Linked List	(average case)	(worst case)
O(n)	O(n)	O(n)	$O(n \log n)$

• Note. The array operations can be improved to  $O(\log n)$  with advanced structures

# **Any Questions?**

