

# Introducción a la Criptografía y a la Seguridad de la Información

#### Sesión 5

Number Theory Background

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## **Session 5**

#### Number Theory Background

- $\triangleright$  Congruent Modulo n
- ightharpoonup Equivalent Class Modulo n
- $\triangleright$  Integer Modulo n ( $\mathbb{Z}_n$ )
- > Factorization
- ⊳ GCD
- ⊳ Relatively Prime
- ightharpoonup Multiplicative Group of  $\mathbb{Z}_n$

- ⊳ PowerMod

- ightharpoonup Primitive Elements in  $\mathbb{Z}_p^*$

### **Division Theorem**

 $\forall a, b \in \mathbb{Z}, \exists \text{ unique } q, r \in \mathbb{Z} : a = qb + r, 0 < r < |b|.$ 

- q = |a/b| is called **quotient** of the division.
- $r = a \mod b$  and is called **remainder** (or **residue**).

**Example:** a = 36, b = 16

$$a = qb + r$$
$$36 = 2 \cdot 16 + 4$$

$$q = 2, r = 4$$

a|b (read a divides b), if  $\exists c \in \mathbb{Z} : b = a \cdot c$ .

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# Divisor, GCD, Prime, Composite

**Divisor:** c is a common divisor of a, b if  $c|a \wedge c|b$ .

**Greatest Common Divisor (GCD):**  $d = \gcd(a, b)$  if d is a common divisor of a and b, and  $\forall c, c | a \land c | b \land c | d$ , Note that d > 1.

The integer p > 1 is a **prime** if its only divisors are 1 and p.

An integer a > 1 that is not a prime is called a **composite number** (or a **composite**).

Integer 1 (one) is neither prime nor composite but a *unit*.

Integer 2 (two) is a prime (the only even one).

# Congruent Modulo n

$$a \equiv b \pmod{n}$$
 iff  $\begin{cases} n | (a - b) \\ a \mod{n} = b \mod{n} \end{cases}$ 

#### Proof:

$$a \mod n = b \mod n$$
  
 $a - kn = b - k'n$   
 $a - b = k''n \Rightarrow n|(a - b)$ 

 $\begin{array}{c} a \bmod n = \\ a - n \lfloor a/n \rfloor = \\ a - nk \end{array}$ 

**Example**:  $24 \equiv 9 \pmod{5}$ 

$$a = b \mod n \Rightarrow a \equiv b \pmod n$$
  
 $a \equiv b \pmod n \Rightarrow a = b \mod n$ 

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# **Equivalence Class Modulo** n

$$[r]_n = \{r + kn : k \in \mathbb{Z}\}$$

#### Example:

 $a \in [b]_n$  is equivalent to writing  $a \equiv b \pmod{n}$ .

# Integers Modulo n $(\mathbb{Z}_n)$

$$\mathbb{Z}_n = \{ [r]_n : 0 \le r \le n-1 \} = \{0, 1, 2, \dots, n-1 \}$$

#### Example:

$$\begin{split} \mathbb{Z}_3 &= \{0,1,2\} \\ \mathbb{Z}_7 &= \{0,1,2,3,4,5,6\} \\ \mathbb{Z}_{14} &= \{0,1,2,3,4,5,6,7,8,9,10,11,12,13\} \\ \mathbb{Z}_{18} &= \{0,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17\} \end{split}$$

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# **Multiplicative Inverse**

$$x \in \mathbb{Z}_n \text{ s.t. } ax \equiv 1 \pmod{n}$$

x is denoted by  $a^{-1}$ 

**Example:**  $\mathbb{Z}_4 = \{0, 1, 2, 3\}, \ 3x \equiv 1 \ (\text{mod } 4), \ \boxed{x=3}$ 

Fact 1:  $a \in \mathbb{Z}_n$ ; a is invertible iff gcd(a, n) = 1

$$(\Leftarrow)$$
  $ax + ny = 1$   $n(-y) = ax - 1$   $\rightarrow$   $n|(ax - 1)$   $\rightarrow$   $ax \equiv 1 \pmod{n}$ .

**Exercise:** in  $\mathbb{Z}_9$  which integers are invertible and what are their inverses.

## **Factorization**

 $n \ge 2$  has a *unique* factorization as a product of distinct prime powers.

$$n = p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k}, \ p_i = \text{prime}, \ e_i \in \mathbb{Z}^+ \ 1 \le i \le k$$

Example: 24

$$24 = 2^3 3^1$$

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**GCD** 

$$a = p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k}$$
$$b = p_1^{f_1} p_2^{f_2} \cdots p_k^{f_k}$$

$$gcd(a,b) = p_1^{\min(e_1,f_1)} p_2^{\min(e_2,f_2)} \cdots p_k^{\min(e_k,f_k)}$$

**Example:** Compute gcd(210, 126)

$$gcd(210, 126) = 2^1 3^1 5^0 7^1 = 2 \cdot 3 \cdot 7 = 42$$

# **Relatively Prime**

Two integers a, b are called **relatively prime** if gcd(a, b) = 1

#### **Example:**

- 234 and 67 are relatively prime because gcd(234, 67) = 1
- 321 and 34 are relatively prime because gcd(321, 34) = 1
- 762 and 105 are NOT relatively prime because gcd(762, 105) = 3

Exercise: Are 123 and 45 relatively prime?

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# Multiplicative Group of $\mathbb{Z}_n$

$$\mathbb{Z}_n^* = \{ a \in \mathbb{Z}_n : \gcd(a, n) = 1 \}$$

 $\phi(n)=|\mathbb{Z}_n^*|=$  number of integers [0,n-1] which are relatively prime to n

- a)  $\phi(p) = p 1$  if p is prime
- **b)**  $\phi(nm) = \phi(n)\phi(m)$  if gcd(n,m) = 1
- c)  $\phi(n) = n(1 \frac{1}{p_1})(1 \frac{1}{p_2})\cdots(1 \frac{1}{p_k})$  if  $n = p_1^{e_1}p_2^{e_2}\cdots p_k^{e_k}$

**Example:** Find  $\phi(21)$ 

$$\mathbb{Z}_{21}^* = \{1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, 20\}, \ \phi(21) = \phi(3)\phi(7) = 12$$

## **Euler's Theorem**

if 
$$a \in \mathbb{Z}_n^*$$
,  $a^{\phi(n)} \equiv 1 \pmod n$ 

#### Proof:

$$g^{\phi(n)} \equiv 1 \pmod{n}$$

$$a \in \mathbb{Z}_n^* \Rightarrow \exists x : a \equiv g^x \pmod{n}$$

$$a^{\phi(n)} \equiv (g^x)^{\phi(n)} \pmod{n} \equiv (g^{\phi(n)})^x \equiv 1 \pmod{n}$$

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## Fermat's Little Theorem

if 
$$gcd(a, p) = 1$$
,  $a^{p-1} \equiv 1 \pmod{p}$ 

p is prime.

#### Proof:

Using Euler's Theorem

# **EEA** - Extended Euclidean Algorithm

#### Pseudo-code:

```
1 procedure EEA(a,b) { q \leftarrow \lfloor a/b \rfloor }
2 begin
3 if b=0 then return (a,1,0)
4 (d',x',y') \leftarrow EEA(b,a \mod b)
5 (d,x,y) \leftarrow (d',y',x'-qy')
6 return (d,x,y)
7 end
```

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# **EEA** - Extended Euclidean Algorithm

Example

Compute EEA(372, 321)

a	b	q	d	x	y	
372	321	1	3	-44	51	$> -44 \times 372 + 51 \times 321 = 3$
321	51	6	3	7	-44	$ > 7 \times 321 + -44 \times 51 = 3 $
51	15	3	3	-2	7	
15	6	2	3	1	-2	$ > 1 \times 15 + -2 \times 6 = 3 $
6	3	2	3	0	1	$ > 0 \times 6 + 1 \times 3 = 3 $
3	0	_	3	1	0	$\triangleright 1 \times 3 + 0 \times 0 = 3$

# **PowerMod - Modular Exponentiation**

#### Pseudo-code:

```
procedure PowerMod(a,b,n) { \langle b_k,b_{k-1},\ldots,b_0\rangle_2\leftarrow b,\ z\leftarrow 1}
begin
for i\leftarrow k downto 0 do
if b_i=1 then z\leftarrow (z^2\times a) mod n
else z\leftarrow z^2 mod n
od
return z
end
```

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# PowerMod - Modular Exponentiation

Example

Compute PowerMod(5,18,17)

$$a = 5 
b = 18_{10} = \langle 10010 \rangle_2 
n = 17$$

$$\frac{i \quad 4 \quad 3 \quad 2 \quad 1 \quad 0}{b_i \quad 1 \quad 0 \quad 0 \quad 1 \quad 0} 
z \quad 5 \quad 8 \quad 13 \quad 12 \quad 8$$

Then  $5^{18} \mod 17 = 8$ 

Exercise: Compute PowerMod(7,452,31)

#### **CRT** - Chinese Remainder Theorem

The following problem was posed by Sunzi [Sun Tsu] (4th century AD) in the book Sunzi Suanjing:

"There are certain things whose number is unknown. Repeatedly divided by 3, the remainder is 2; by 5 the remainder is 3; and by 7 the remainder is 2. What will be the number?"

CRT was commonly known as General Sun counting the soldiers or General Han counting the soldiers.

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Oystein Ore mentions another puzzle with a dramatic element from Brahma-Sphuta-Siddhanta (Brahma's Correct System) by Brahmagupta (born 598 AD):

"An old woman goes to market and a horse steps on her basket and crashes the eggs. The rider offers to pay for the damages and asks her how many eggs she had brought. She does not remember the exact number, but when she had taken them out two at a time, there was one egg left. The same happened when she picked them out three, four, five, and six at a time, but when she took them seven at a time they came out even. What is the smallest number of eggs she could have had?"

Problems of this kind are all examples of CRT

## **CRT** - Definition

Let  $n_1, n_2, \ldots, n_k$  be pairwise relatively prime integers. If  $a_1, a_2, \ldots, a_k$  are any integers, then the system of simultaneous congruences

$$x \equiv a_i \pmod{n_i} \quad \forall i \in \{1 \dots k\}$$

has a unique solution modulo  $N = n_1 n_2 \dots n_k$ 

$$x = \sum_{i=1}^{k} N_i y_i a_i \bmod N$$

where  $N_i = N/n_i$  and  $y_i = N_i^{-1} \mod n_i$ 

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# **Example**

$$k=2, n_1=5, n_2=3$$

$$N=n_1n_2=5\times 3=15$$

$$\pi(x) = (x \mod 5, x \mod 3) : \mathbb{Z}_{15} \to \mathbb{Z}_5 \times \mathbb{Z}_3$$

$$\begin{array}{l} N_1 = N/n_1 = 15/5 = 3 \\ N_2 = N/n_2 = 15/3 = 5 \\ \\ y_1 = N_1^{-1} \mod n_1 = 3^{-1} \mod 5 = 2 \\ y_2 = N_2^{-1} \mod n_2 = 5^{-1} \mod 3 = 2 \\ \\ x = \pi^{-1}(a_1, a_2) = (N_1 y_1 a_1 + N_2 y_2 a_2) \mod N \\ \qquad = (3 \times 2 a_1 + 5 \times 2 a_2) \mod 15 \\ \qquad = (6 a_1 + 10 a_2) \mod 15 \\ \\ \text{for } a_1 = 1 \text{ and } a_2 = 2 \text{ we get} \\ \\ x \equiv 1 \pmod 5 \\ x \equiv 2 \pmod 3 \\ \\ x = \pi^{-1}(1, 2) = (6 \times 1 + 10 \times 2) \mod 15 \\ \qquad = (6 + 20) \mod 15 \\ \qquad = 26 \mod 15 \\ \qquad = 11 \\ \\ 11 \mod 5 = 1 \\ \\ 11 \mod 3 = 2 \\ \end{array}$$

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# The Order of an Integer

Let  $a \in \mathbb{Z}_n^*$ .

$$\operatorname{ord}(a) = \min(t : a^t \equiv 1 \pmod{n})$$

Exercise: Find the order of 5 with the following moduli

# Primitive elements in $\mathbb{Z}_p^*$

for p=prime,  $\alpha$  is called a **primitive element modulo** p if

$$\mathbb{Z}_p^* = \{\alpha^i \bmod p : i \in \mathbb{Z}_p^*\}$$

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# **Example**

For p=13 find all the possible primitive elements modulo 13.

t	1	2	3	4	5	6	7	8	9	10	11	12
$2^t \mod 13$	2	4	8	3	6	12	11	9	5	10	7	1
$3^t \mod 13$	3	9	1	3	9	1	3	9	1	3	9	1
$4^t \mod 13$	4	3	12	9	10	1	4	3	12	9	10	1
$5^t \bmod 13$	5	12	8	1	5	12	8	1	5	12	8	1
$6^t \bmod 13$	6	10	8	9	2	12	7	3	5	4	11	1
$7^t \bmod 13$	7	10	5	9	11	12	6	3	8	4	2	1
$8^t \bmod 13$	8	12	5	1	8	12	5	1	8	12	5	1
$9^t \bmod 13$	9	3	1	9	3	1	9	3	1	9	3	1
$10^t \bmod 13$	10	9	12	3	4	1	10	9	12	3	4	1
$11^t \bmod 13$	11	4	5	3	7	12	2	9	6	8	10	1
$12^t \mod 13$	12	1	12	1	12	1	12	1	12	1	12	1

a	1	2	3	4	5	6	7	8	9	10	11	12
ord(a)	1	12	3	6	4	12	12	4	3	6	12	2

Hence, the primitive elements modulo 13 are 2, 6, 7 and 11.