

14. State-Space Search

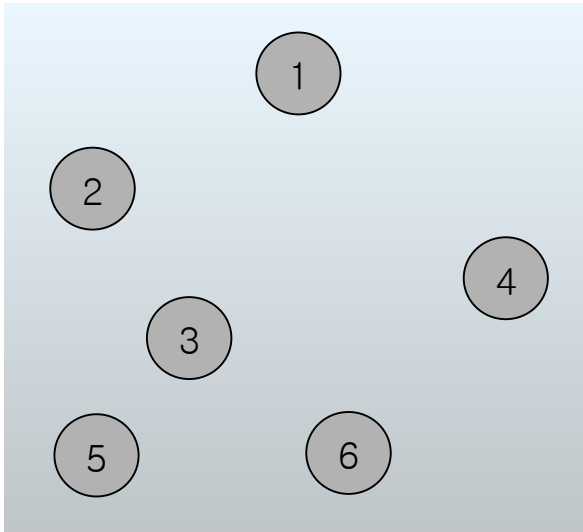
Goals

- Understand state-space tree.
- Learn backtracking.
- Learn branch-and-bound.
- Learn A* algorithm.

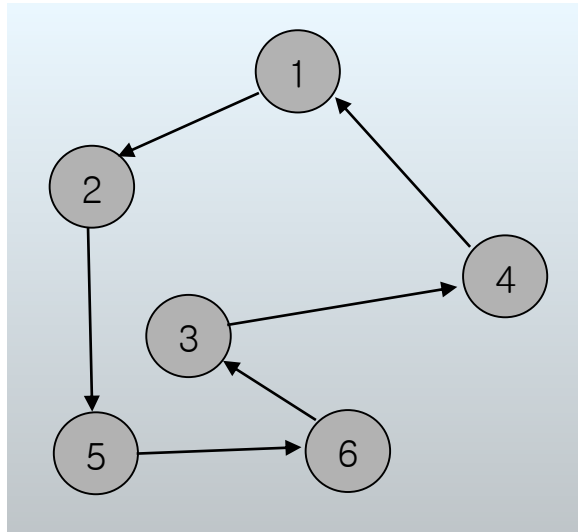
State-Space Tree

- State space: set of states that are generated in problem solving process
- State space tree: tree which represents states of problem solving process in a systematic way
 - Solution space: set of candidate solutions for the problem (leaves in a state space tree). Some nodes in the solution space are answers to the problem
 - Partial solution: internal node
- Search techniques for state space
 - Backtracking
 - Branch-and-bound
 - A* algorithm

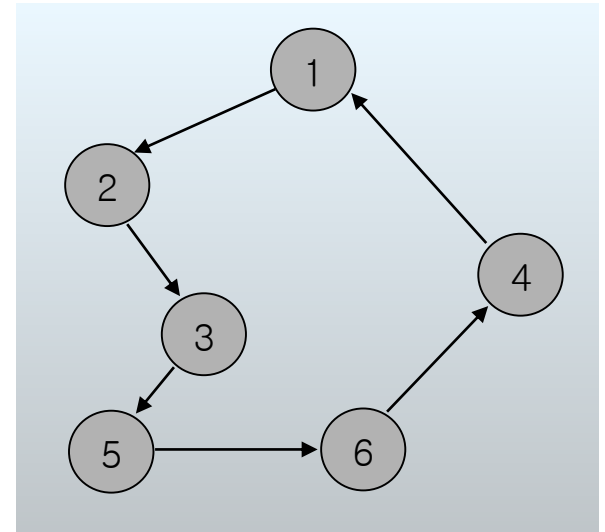
TSP



(a) instance of TSP



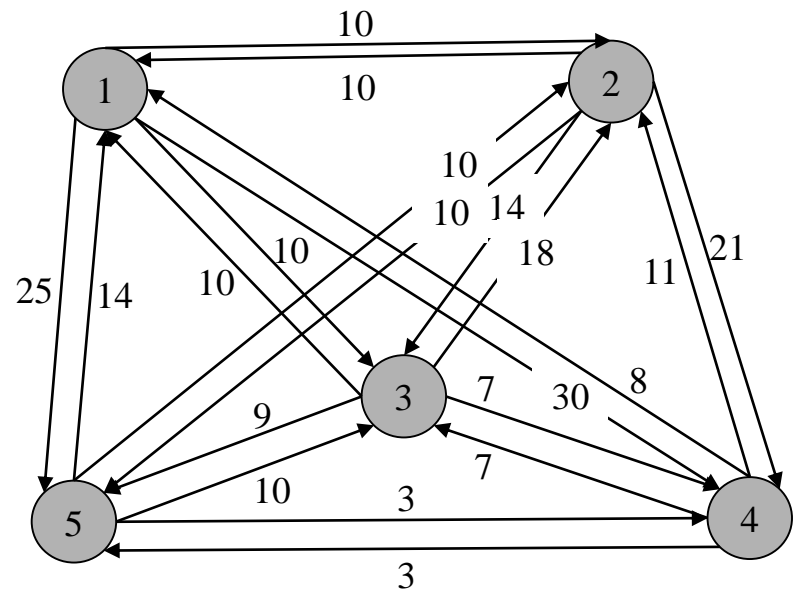
(b) a solution



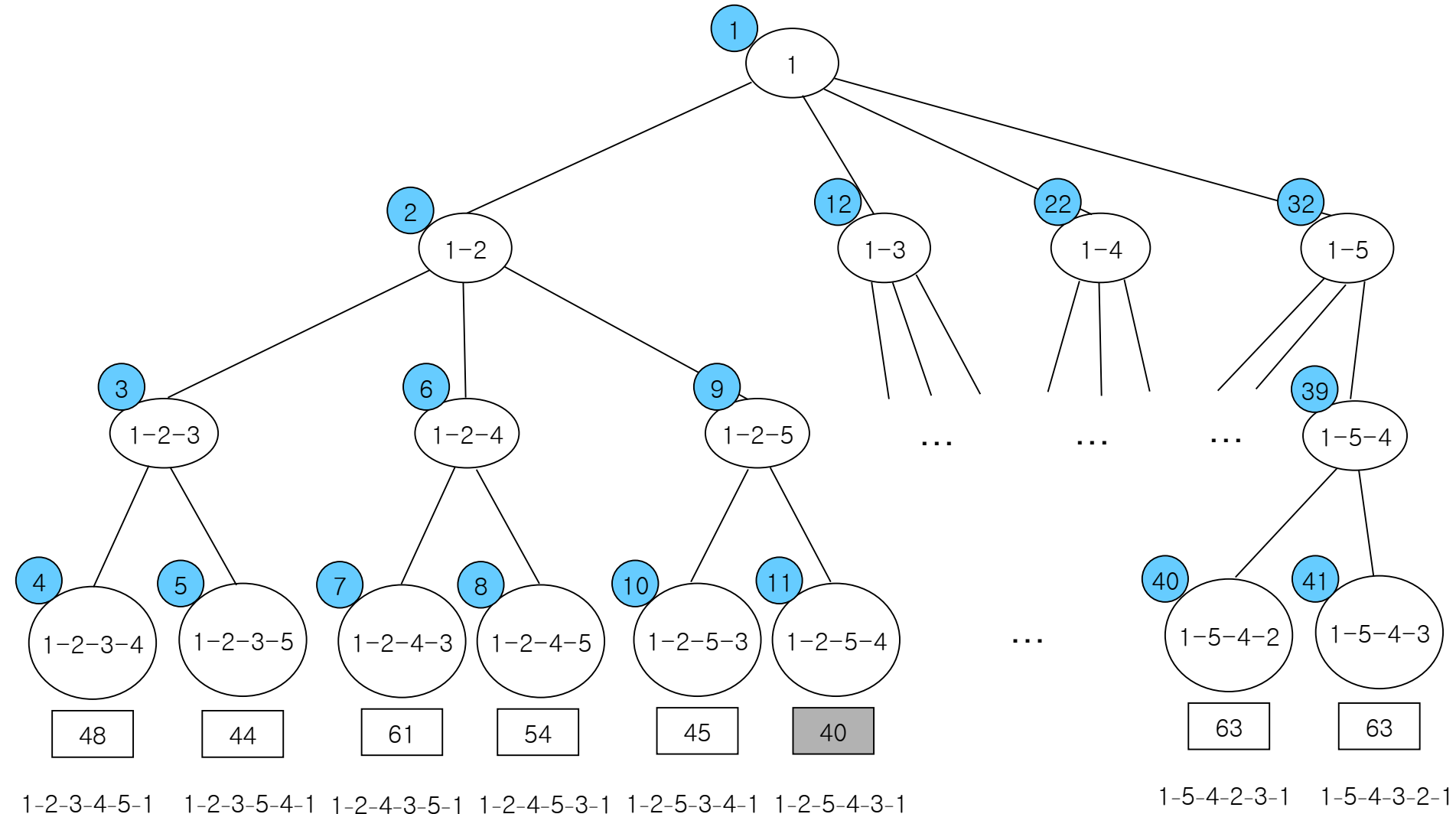
(c) optimal solution

TSP and Adjacency Matrix

	1	2	3	4	5
1	0	10	10	30	25
2	10	0	14	21	10
3	10	18	0	7	9
4	8	11	7	0	3
5	14	10	10	3	0



Lexicographic order search of state space



Backtracking

- Current state $(x_1, x_2, \dots, x_{i-1})$: a path from root to a node in a state space tree
- $T(x_1, \dots, x_{i-1})$: set of all possible values for x_i
- $f(x_1, \dots, x_i) = \text{false}$ only if the path cannot be extended to reach an answer node

```
Backtrack( $n$ ) {  
     $i \leftarrow 1$ ;  
    while ( $i > 0$ ) {  
        if there is an untried  $x_i$  such that  
         $x_i \in T(x_1, \dots, x_{i-1})$  and  $f(x_1, \dots, x_i)$  then  
            if  $(x_1, \dots, x_i)$  is an answer node then print  $(x_1, \dots, x_i)$ ;  
            else  $i \leftarrow i+1$ ;  
        else  $i \leftarrow i-1$ ;  
    }  
}
```

Backtracking

- Recursive version of backtracking
- Initial call: BacktrackR(1)
- Current state $(x_1, x_2, \dots, x_{i-1})$

```
BacktrackR(i) {  
    for each  $x_i$  such that  $x_i \in T(x_1, \dots, x_{i-1})$  and  $f(x_1, \dots, x_i)$  {  
        if  $(x_1, \dots, x_i)$  is an answer node then print  $(x_1, \dots, x_i)$ ;  
        else BacktrackR(i+1);  
    }  
}
```


8-Queens Problem

- Place 8 queens on an 8×8 chessboard so that no two can attack (i.e., no two are on the same row, column, or diagonal)

			Q				
					Q		
							Q
	Q						
						Q	
Q							
		Q					
				Q			

8-Queens Problem

- Queen i is placed on row i .
- All answers are represented by (x_1, \dots, x_8) , where x_i is the column on which queen i is placed.
- x_i increases from 1 to $n+1$.
- Current state $(x_1, x_2, \dots, x_{i-1})$

```
Queens( $n$ ) {  
     $i \leftarrow 1$ ;  
    while ( $i > 0$ ) {  
         $x_i \leftarrow$  next column on which queen  $i$  can be placed  
        if  $x_i \leq n$  then  
            if  $i = n$  then print  $(x_1, \dots, x_i)$ ;  
            else  $i \leftarrow i+1$ ;  
        else  $i \leftarrow i-1$ ;  
    }  
}
```

8-Queens Problem

- Current state: $(x_1 = 2, x_2 = 4, x_3 = 1)$
- x_4 can be 3, 7, 8.
- Next current state: $(x_1 = 2, x_2 = 4, x_3 = 1, x_4 = 3)$
- x_5 can be 5, 8.

	1						
			2				
3							
		X				X	X

	1						
			2				
3							
		4					
				X			X

Branch-and-Bound

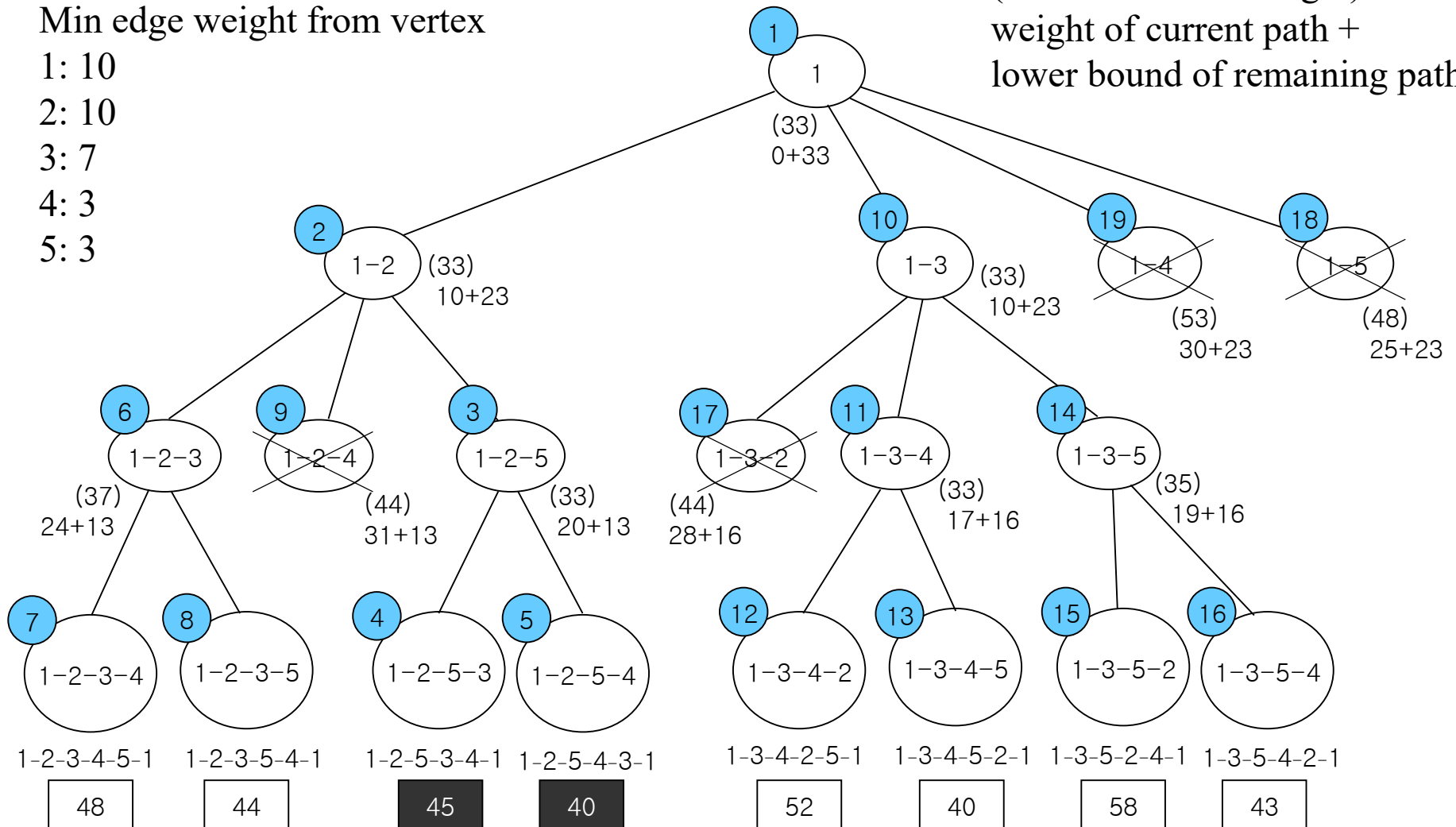
- Combination of ‘branch’ and ‘bound’
 - A cost function exists
 - A minimum cost answer node is to be found
- Comparison with backtracking
 - common
 - Require a method to list states
 - different
 - Backtracking – backtrack when there is no further way to go
 - Branch-and-bound – don’t branch if it is guaranteed that there is no optimal solution in that branch

State-Space Tree of Branch-and-Bound for TSP

Min edge weight from vertex

- 1: 10
- 2: 10
- 3: 7
- 4: 3
- 5: 3

(lower bound of weight)
weight of current path +
lower bound of remaining path



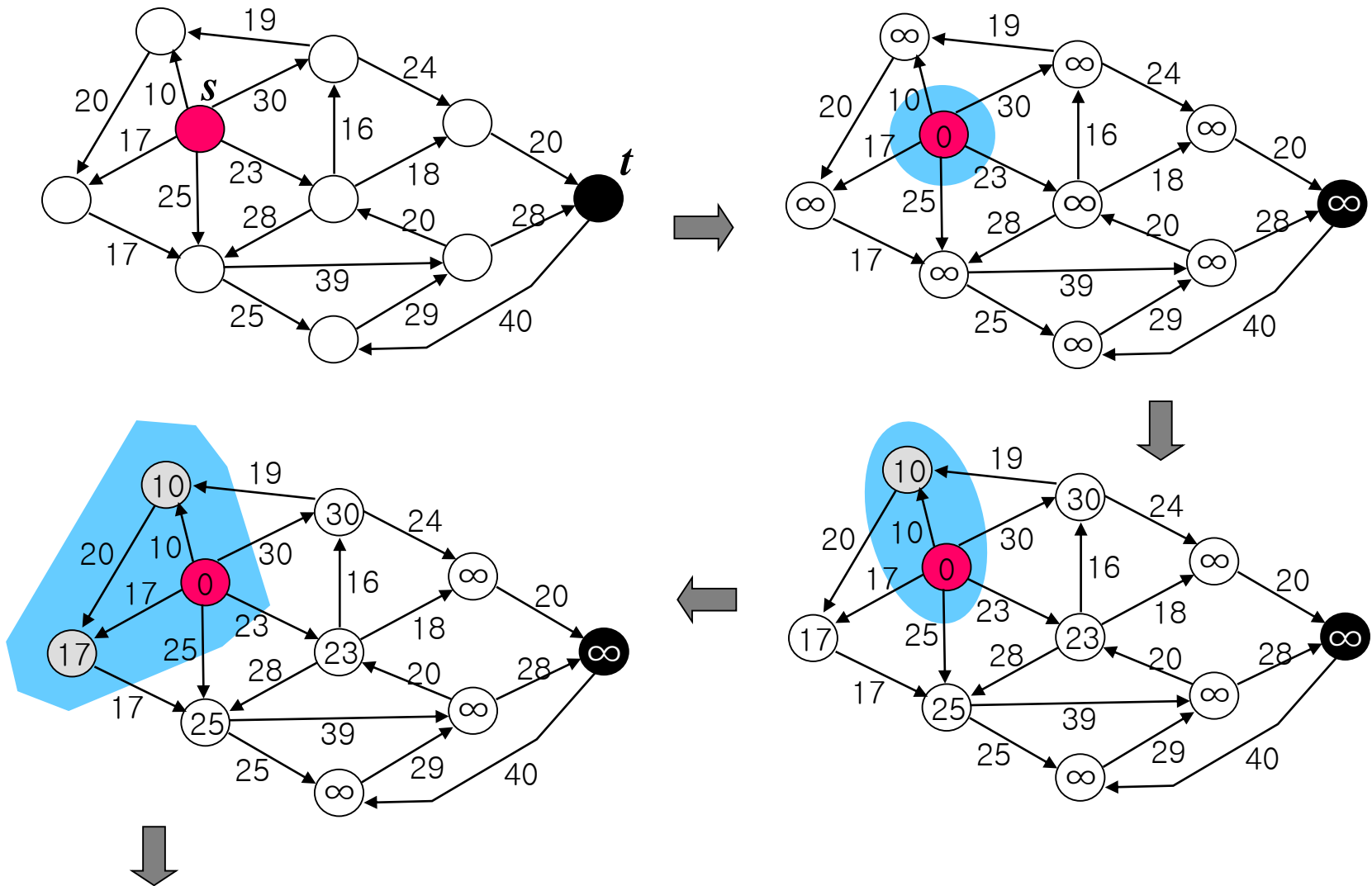
A* Algorithm

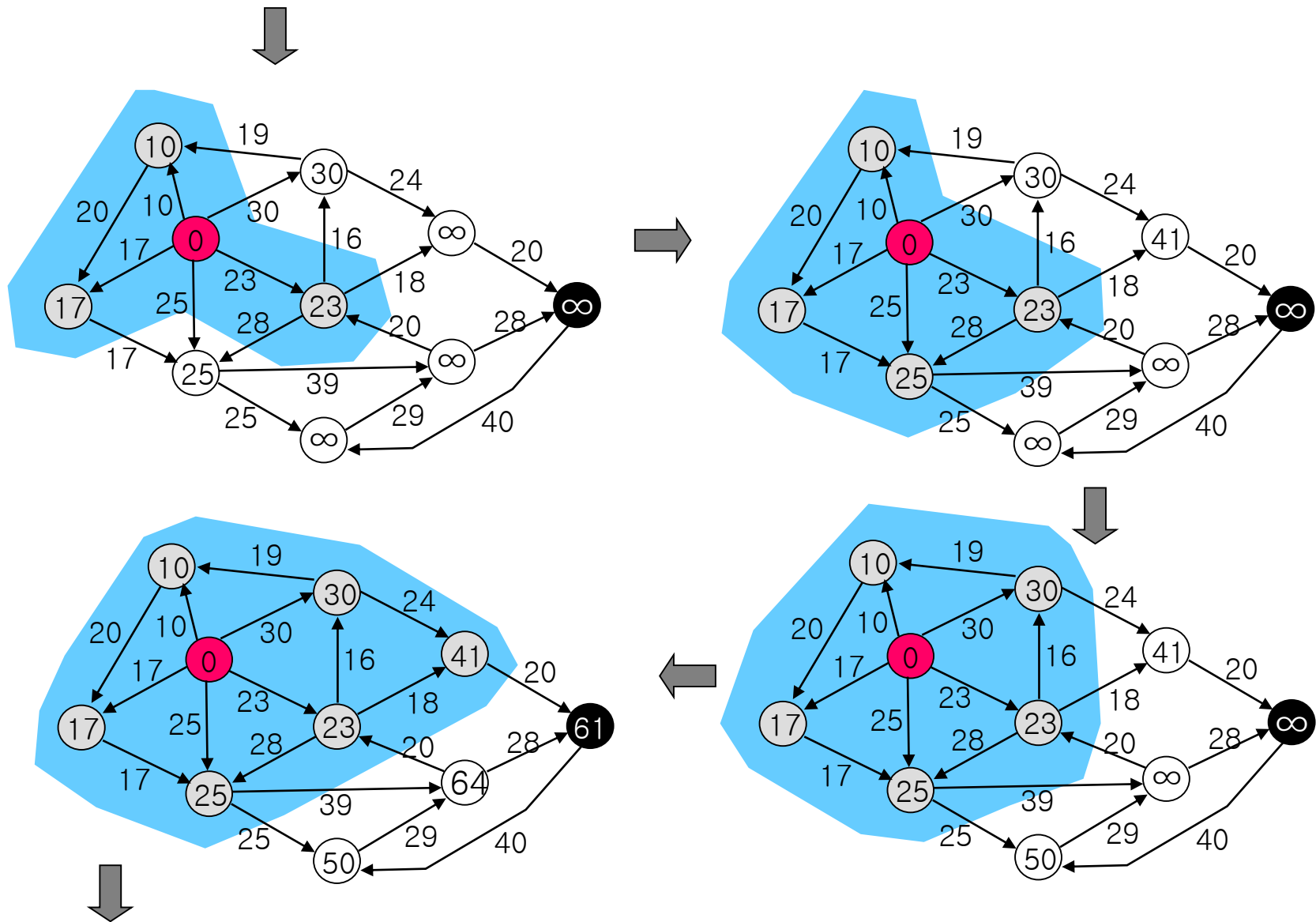
- Find the shortest path from a source to a destination
- Can be applied to NP-hard and P problems
- cf. Dijkstra algorithm
 - Single source
 - All destinations

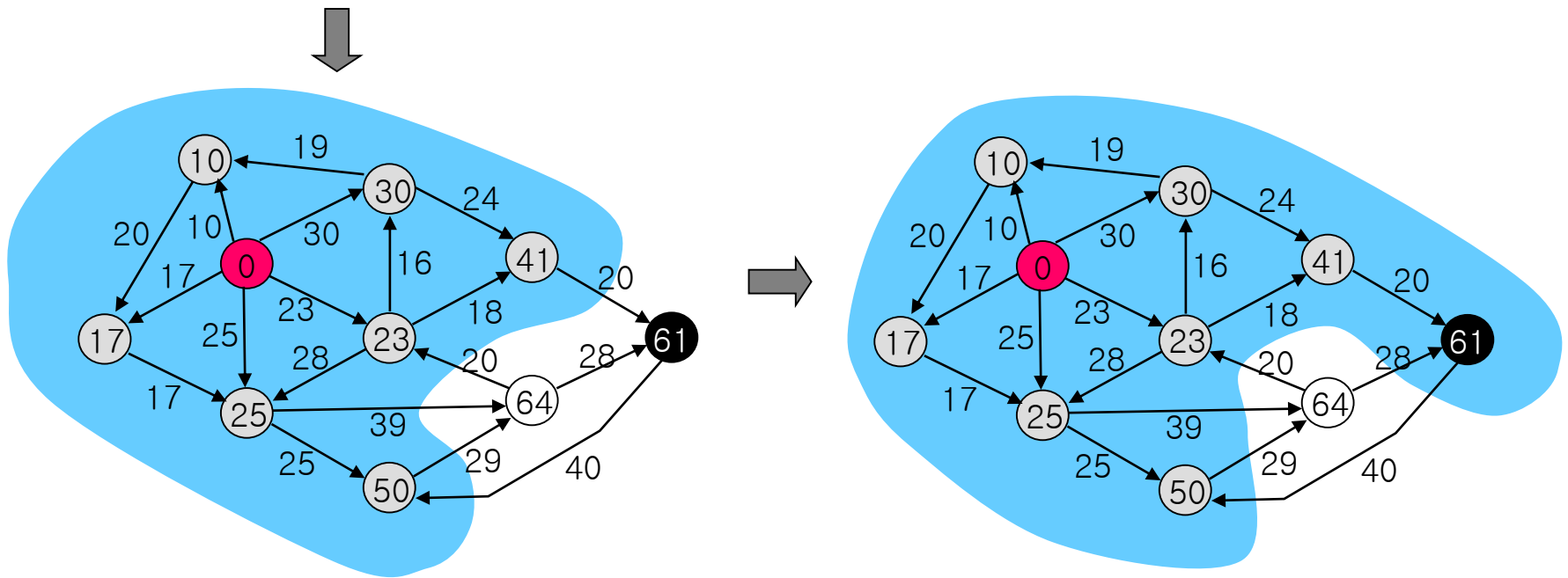
A* Algorithm

- Best-first search
 - Each vertex x has $g(x)$: cost (shortest path weight) from source to x
 - Each vertex x has $h(x)$: estimate of cost from x to destination. Properties of $h(x)$:
 - $h(x)$ must be less than or equal to actual cost from x to destination
 - For all x, y , $h(x) \leq w(x,y) + h(y)$
 - A* always selects a vertex x that minimizes $g(x) + h(x)$

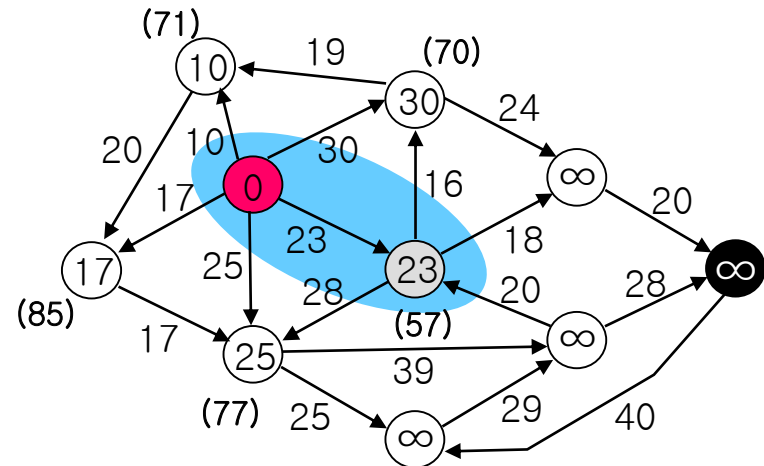
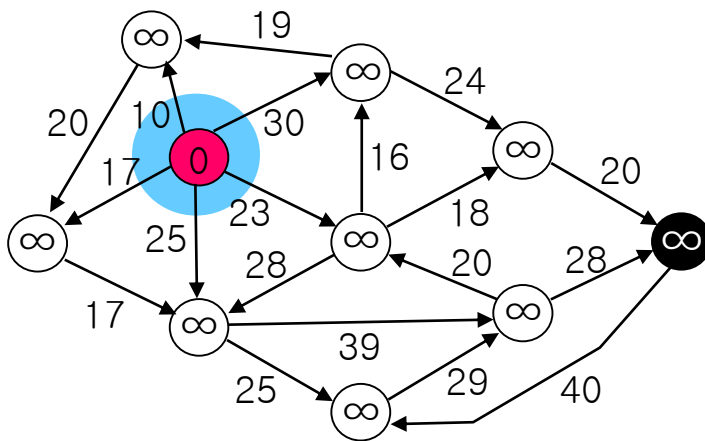
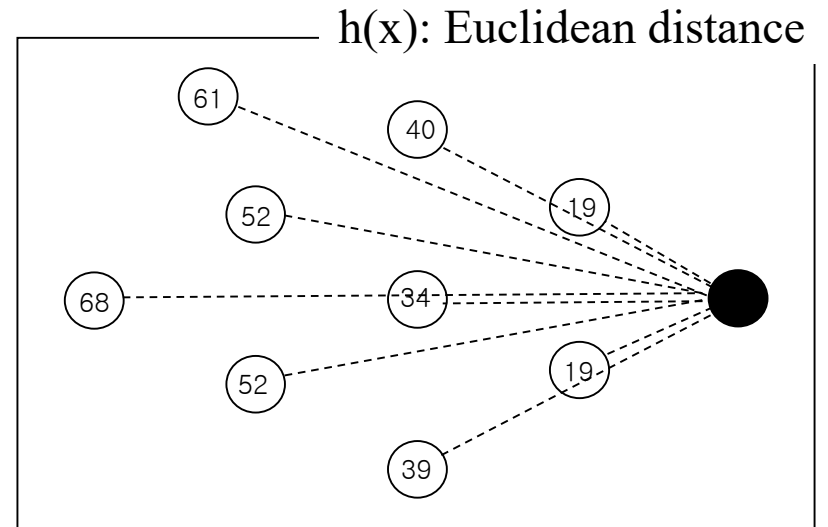
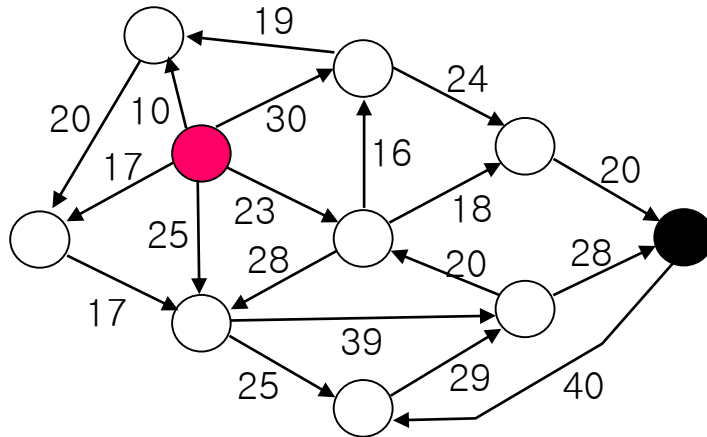
Dijkstra Algorithm

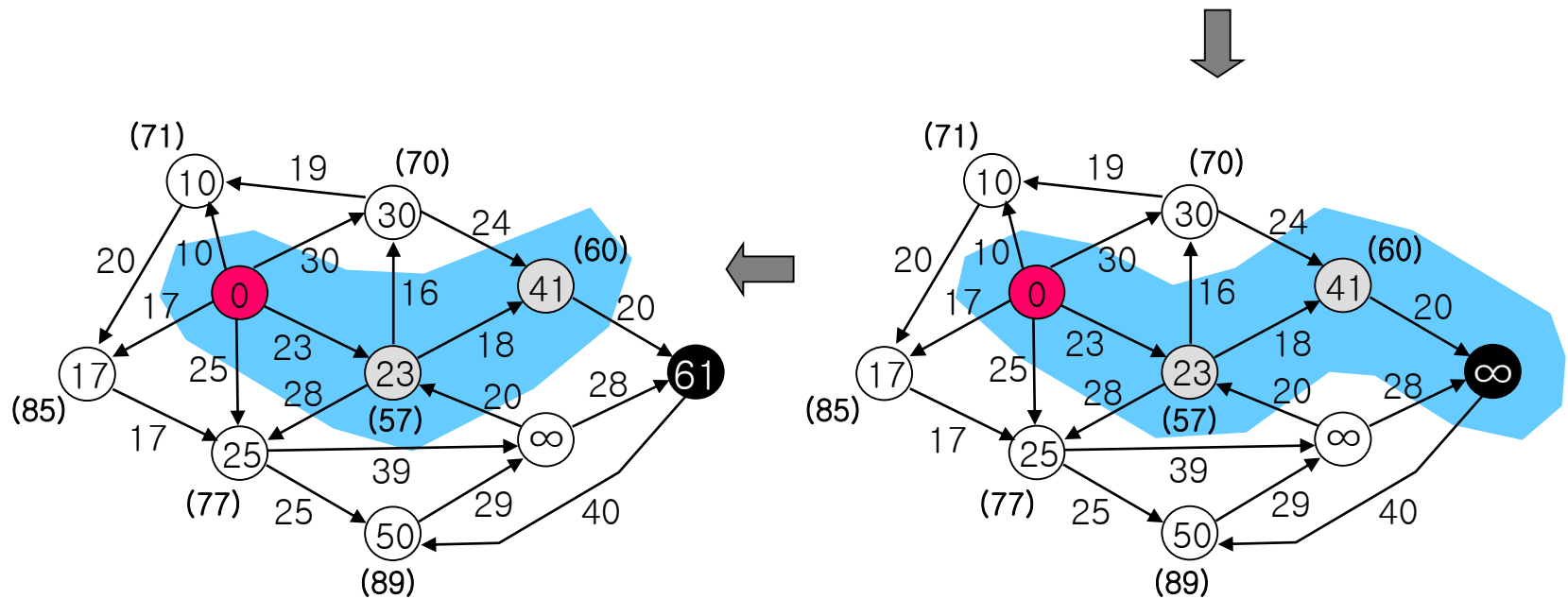






A* Algorithm





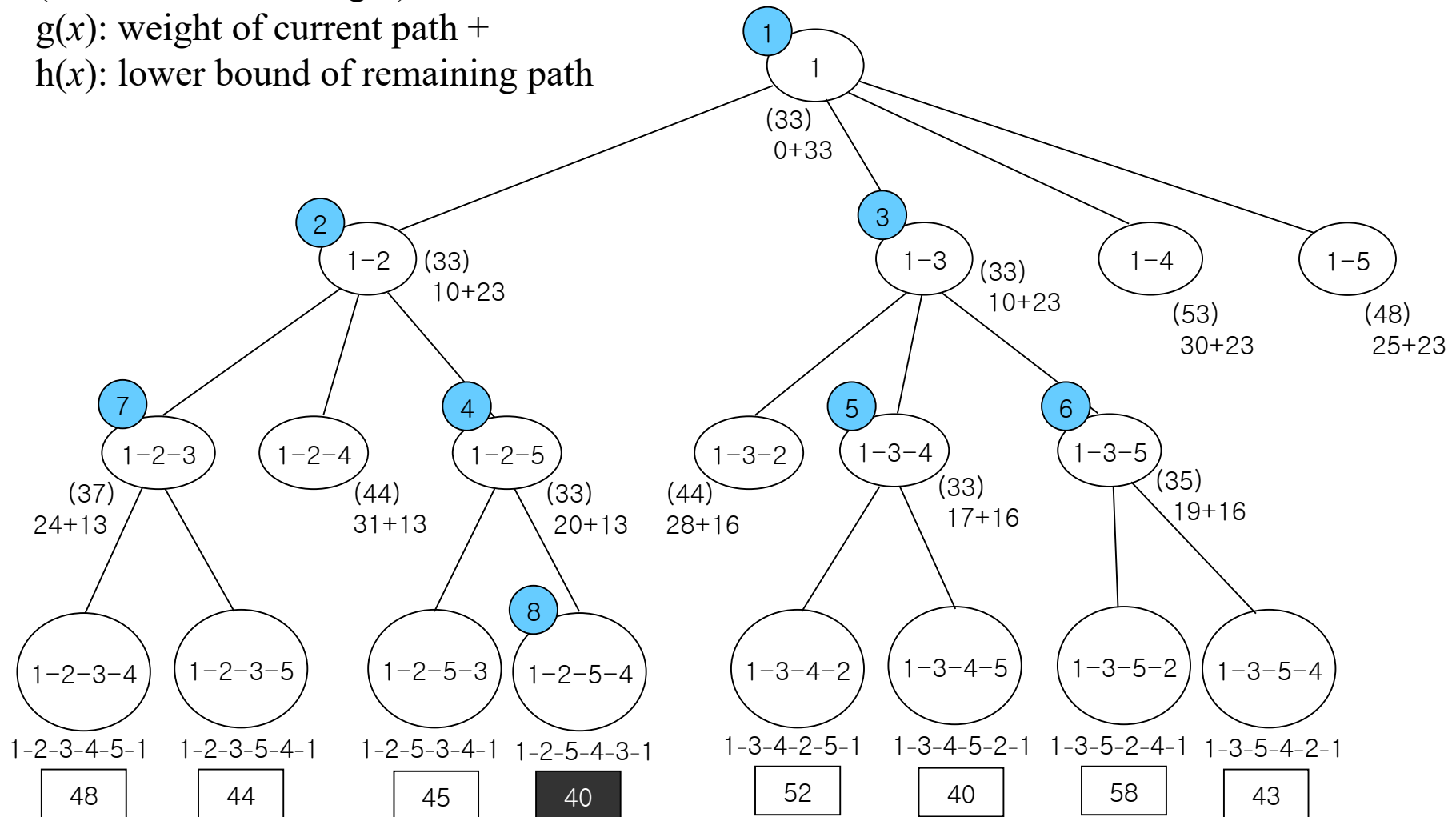
✓ A* is faster than Dijkstra by using $h(x)$

State-Space Tree of A* Algorithm for TSP

(lower bound of weight)

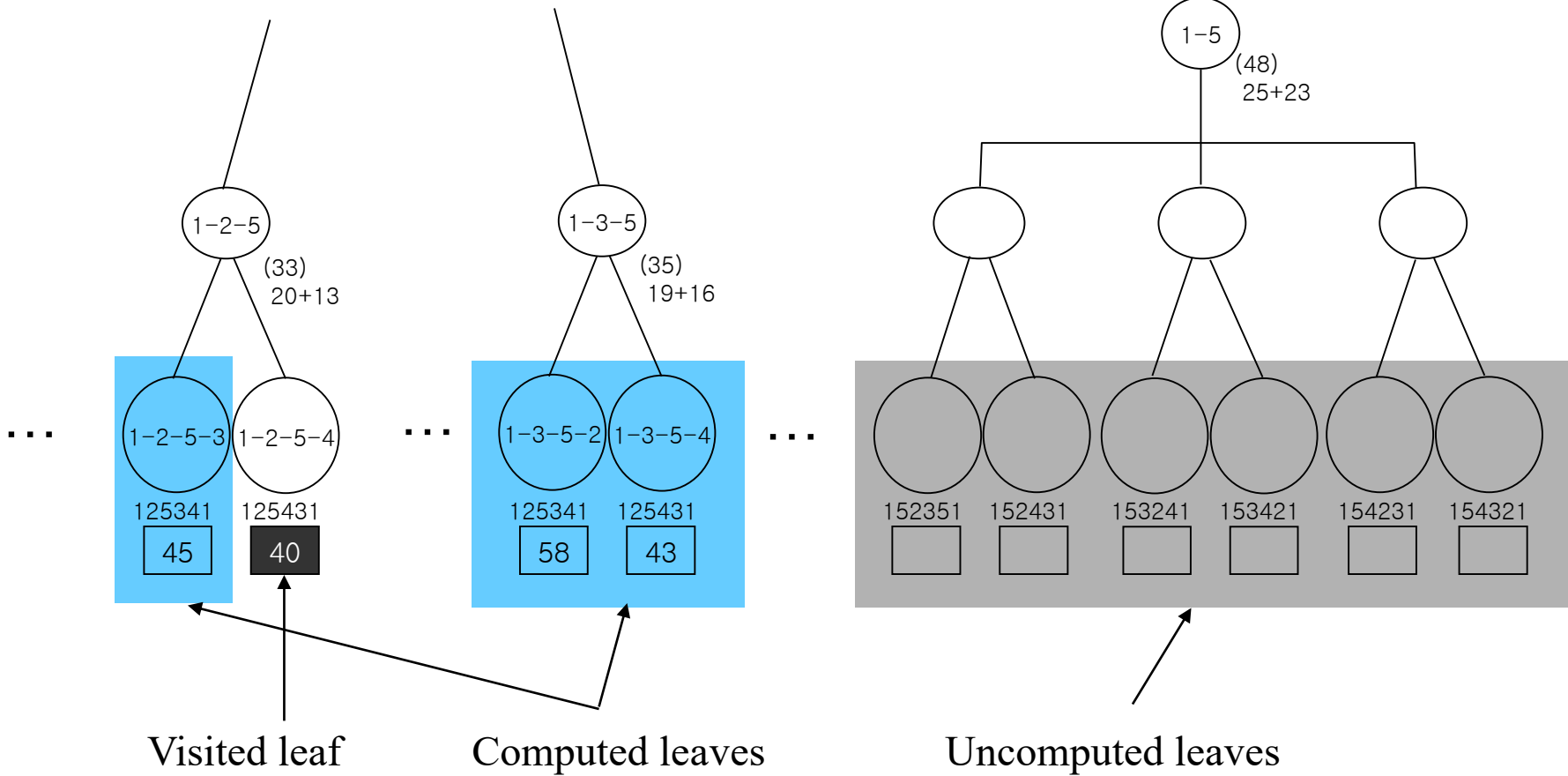
$g(x)$: weight of current path +

$h(x)$: lower bound of remaining path



A* Algorithm Terminates When It Visits First Leaf

leaves and leaves cannot be smaller than 40





Thank you
