

1、初始化

```
private static final int MAXIMUM_CAPACITY = 1 << 30; //数组最大的容量，01000000
00000000 00000000 00000000
private static final int DEFAULT_CAPACITY = 16; //初始化容量
static final int MAX_ARRAY_SIZE = Integer.MAX_VALUE - 8; //
private static final int DEFAULT_CONCURRENCY_LEVEL = 16; //默认的并发级别，默认是16个
tsegment所以每个segment都可以被加上一把锁
private static final float LOAD_FACTOR = 0.75f; //默认的加载因子
static final int TREEIFY_THRESHOLD = 8; //树化的阈值
static final int UNTREEIFY_THRESHOLD = 6; //树退化的阈值
static final int MIN_TREEIFY_CAPACITY = 64; //树化的最小segment长度
private static final int MIN_TRANSFER_STRIDE = 16; //扩容线程每次最少要迁移16个hash桶
private static final int RESIZE_STAMP_BITS = 16; //
private static final int MAX_RESIZERS = (1 << (32 - RESIZE_STAMP_BITS)) - 1; //最
多有多多少个线程可以同时Map进行扩容
private static final int RESIZE_STAMP_SHIFT = 32 - RESIZE_STAMP_BITS; //

static final int MOVED      = -1; // hash for forwarding nodes
static final int TREEBIN    = -2; // hash for roots of trees
static final int RESERVED   = -3; // hash for transient reservations
static final int HASH_BITS  = 0x7fffffff; // usable bits of normal node hash

transient volatile Node<K,V>[] table; //table数组
private transient volatile Node<K,V>[] nextTable; //resize的时候，记录下一个将使用的
table;
private transient volatile long baseCount; //

private transient volatile int sizeCtl; //控制数组大小
private transient volatile int transferIndex; //扩容的时候新数组下标位置
private transient volatile int cellsBusy; //标识当前cell数组是否在初始化或扩容中的CAS标
志位
private transient volatile CounterCell[] counterCells; //记录每一个segment的结点数量
counterCells数组，总数值的分值分别存在每个cell中

// views
private transient KeySetView<K,V> keySet;
private transient ValuesView<K,V> values;
private transient EntrySetView<K,V> entrySet;
```

1.1 sizeCtl详解

- 未初始化: sizeCtl=0: 表示没有指定初始容量。sizeCtl>0: 表示初始容量。
- 初始化中: sizeCtl=-1, 标记作用, 告知其他线程, 正在初始化
- 正常状态: sizeCtl=0.75n, 扩容阈值
- 扩容中: sizeCtl < 0 : 表示有其他线程正在执行扩容
- sizeCtl = (resizeStamp(n) << RESIZE_STAMP_SHIFT)+2 表示此时只有一个线程在执行扩容;

2、ConcurrentHashMap中的Node结点

- 这个和HashMap中的Node一致;

```

static class Node<K,V> implements Map.Entry<K,V> {
    final int hash;
    final K key;
    volatile V val;
    volatile Node<K,V> next;

    Node(int hash, K key, V val, Node<K,V> next) {
        this.hash = hash;
        this.key = key;
        this.val = val;
        this.next = next;
    }

    public final K getKey()      { return key; }
    public final V getValue()    { return val; }
    public final int hashCode()  { return key.hashCode() ^ val.hashCode(); }

    public final String toString(){ return key + "=" + val; }
    public final V setValue(V value) {
        throw new UnsupportedOperationException();
    }

    public final boolean equals(Object o) {
        Object k, v, u; Map.Entry<?,?> e;
        return ((o instanceof Map.Entry) &&
            (k = (e = (Map.Entry<?,?>)o).getKey()) != null &&
            (v = e.getValue()) != null &&
            (k == key || k.equals(key)) &&
            (v == (u = val) || v.equals(u)))));
    }

    /**
     * virtualized support for map.get(); overridden in subclasses.
     */
    Node<K,V> find(int h, Object k) {
        Node<K,V> e = this;
        if (k != null) {
            do {
                K ek;
                if (e.hash == h &&
                    ((ek = e.key) == k || (ek != null && k.equals(ek))))
                    return e;
            } while ((e = e.next) != null);
        }
        return null;
    }
}

```

3、ForwardingNode结点

```

//A node inserted at head of bins during transfer operations.
static final class ForwardingNode<K,V> extends Node<K,V> {
    final Node<K,V>[] nextTable;
    ForwardingNode(Node<K,V>[] tab) {

```

```

        super(MOVED, null, null, null);
        this.nextTable = tab;
    }

    Node<K,V> find(int h, Object k) {
        // loop to avoid arbitrarily deep recursion on forwarding nodes
        outer: for (Node<K,V>[] tab = nextTable;;) {
            Node<K,V> e; int n;
            if (k == null || tab == null || (n = tab.length) == 0 ||
                (e = tabAt(tab, (n - 1) & h)) == null)
                return null;
            for (;;) {
                int eh; K ek;
                if ((eh = e.hash) == h &&
                    ((ek = e.key) == k || (ek != null && k.equals(ek))))
                    return e;
                if (eh < 0) {
                    if (e instanceof ForwardingNode) {
                        tab = ((ForwardingNode<K,V>)e).nextTable;
                        continue outer;
                    }
                    else
                        return e.find(h, k);
                }
                if ((e = e.next) == null)
                    return null;
            }
        }
    }
}

```

3、扰动函数

- 在HashMap中spread函数，最后会与HASH_BITS做与运算，这是因为MAXIMUM_CAPACITY = 01000000 00000000 00000000 00000000，保证HASH值

HASH_BITS = 0X7FFFFFFF

01111111 11111111 11111111 11111111

hash1: 00000000 00110101 11001100 01100010

hash1>>>16: 00000000 00000000 00000000 00110101

做异或运算

 00000000 00110101 11001100 01010111

与HASH_BITS做与运算:

00000000 00110101 11001100 01010111

01111111 11111111 11111111 11111111

做&运算

00000000 00110101 11001100 01010111

```
//ConcurrentHashMap扰动函数，
static final int spread(int h) {
    return (h ^ (h >>> 16)) & HASH_BITS; //因为MAXIMUM_CAPACITY最高位必须为0，所以计算完后还需要与HASH_BITS做与操作，将hash最高位置0，而最高位是为了帮助扩容的，因为sizeCtl < 0时最高位为1，此时正在扩容；
}

//HashMap中的扰动函数，
static final int hash(Object key) {
    int h;
    return (key == null) ? 0 : (h = key.hashCode()) ^ (h >>> 16);
}
```

4、构造函数

```
public ConcurrentHashMap(int initialCapacity) {
    if (initialCapacity < 0)
        throw new IllegalArgumentException();
    int cap = ((initialCapacity >= (MAXIMUM_CAPACITY >>> 1)) ?
        MAXIMUM_CAPACITY :
        tableSizeFor(initialCapacity + (initialCapacity >>> 1) + 1)); //计算大于1.5*initialCapacity的2的幂次方
    this.sizeCtl = cap;
}

public ConcurrentHashMap(int initialCapacity,
    float loadFactor, int concurrencyLevel) {
    if (!(loadFactor > 0.0f) || initialCapacity < 0 || concurrencyLevel <= 0)
        throw new IllegalArgumentException();
    if (initialCapacity < concurrencyLevel) // Use at least as many bins
        initialCapacity = concurrencyLevel; // as estimated threads
    long size = (long)(1.0 + (long)initialCapacity / loadFactor); //
    int cap = (size >= (long)MAXIMUM_CAPACITY) ?
        MAXIMUM_CAPACITY : tableSizeFor((int)size);
    this.sizeCtl = cap;
}
```

5、get方法

```
public V get(Object key) {
    Node<K,V>[] tab; Node<K,V> e, p; int n, eh; K ek;
    int h = spread(key.hashCode()); //计算Hash值
    if ((tab = table) != null && (n = tab.length) > 0 &&
        (e = tabAt(tab, (n - 1) & h)) != null) { //找到key对应的segment下标
        if ((eh = e.hash) == h) { //
            if ((ek = e.key) == key || (ek != null && key.equals(ek))) //地址相等或者值相等
                return e.val;
        }
    }
}
```

```

        else if (eh < 0) //hash值最高为1,
            return (p = e.find(h, key)) != null ? p.val : null;
        while ((e = e.next) != null) { //在链表中查找
            if (e.hash == h &&
                ((ek = e.key) == key || (ek != null && key.equals(ek))))
                return e.val;
        }
    }
    return null;
}

```

6、putval函数（向map中添加元素）

向concurrenthashmap中添加元素时，需要保证添加的key和value均不为null;

```

final V putVal(K key, V value, boolean onlyIfAbsent) {
    if (key == null || value == null) throw new NullPointerException();
    //key、value均不能为null
    int hash = spread(key.hashCode()); //首先计算key的hash值
    int binCount = 0;
    for (Node<K,V>[] tab = table;;) { //遍历table
        Node<K,V> f; int n, i, fh; //fh记录结点f的hash
        if (tab == null || (n = tab.length) == 0) //如果table为空，说明map还没有
            申请内存空间
            tab = initTable(); //初始化table;
        else if ((f = tabAt(tab, i = (n - 1) & hash)) == null) { //如果下标位
            置为空
            //如果此位置没有元素，利用cas往下标位置加，如果为空才加，否则加不了；
            if (casTabAt(tab, i, null,
                new Node<K,V>(hash, key, value, null))) //cas操作尝试
                将插入结点到此位置
                break; // no lock when adding to empty bin
        }
        else if ((fh = f.hash) == MOVED) //如果当前map正在扩容，那么先协助扩容再更
            新值：
            tab = helpTransfer(tab, f);
        else {
            V oldVal = null;
            synchronized (f) { //对桶的头节点加锁，某个线程在插入过程中，不允许其他线程
                对桶中元素进行修改
                if (tabAt(tab, i) == f) { //这里判断是因为为了防止添加元素后变成了
                    一颗树
                    if (fh >= 0) { //是普通的链表结构
                        binCount = 1;
                        for (Node<K,V> e = f;; ++binCount) { //循环遍历链表
                            K ek;
                            //如果key相等则直接替换value，然后break;
                            if (e.hash == hash &&
                                ((ek = e.key) == key ||
                                 (ek != null && key.equals(ek)))) {
                                oldVal = e.val;
                                if (!onlyIfAbsent)
                                    e.val = value;
                                break;
                            }
                        }
                    }
                }
            }
        }
    }
}

```

```

//不相等则继续查找下一个元素
Node<K,V> pred = e;
//查找到尾结点发现还是不相等，则插入新结点到末尾结点之后
if ((e = e.next) == null) {
    pred.next = new Node<K,V>(hash, key,
                               value, null);
    break;
}
}
}
else if (f instanceof TreeBin) { //如果fh < 0则不是链表结构，为树结构，
    Node<K,V> p;
    binCount = 2;
    if ((p = ((TreeBin<K,V>)f).putTreeval(hash, key,
                                             value)) != null) {
        oldval = p.val;
        if (!onlyIfAbsent)
            p.val = value;
    }
}
}
}
if (binCount != 0) {
    if (binCount >= TREEIFY_THRESHOLD) //如果是链表，判断链表长度是否达到了树化的阈值，如果达到了就需要树化
        treeifyBin(tab, i);
    if (oldval != null)
        return oldval;
    break;
}
}
}
addCount(1L, binCount); //
return null;
}

```

6.1 initTable申请内存空间

- initTable是在多线程情况下，假设同时有多个线程尝试initTable；
- 首先每个线程都会判断table是否为null，如果为空就会尝试初始化table；
- 如果当前sizeCtl < 0，那么说明table正在被其他线程创建，此线程就通过Thread.yield()函数挂起，让出cpu的执行权；
- 如果说sizeCtl >= 0，那么当前table还没有被初始化，当前线程尝试CAS修改sizeCtl的内存值，如果修改失败就重新进入循环；
- 如果当前线程修改sizeCtl内存值成功，说明当前线程获取到了初始化table的权限，尝试为table申请内存空间；
- 在申请前还要再次判断tab是否为空，
- 判断sizeCtl是否大于0，不大于0就使用默认的初始化容量；
- 申请内存空间；
- 将sizeCtl赋值为0.75初始化容量大小；

```

//初始化表的过程
private final Node<K,V>[] initTable() {
    Node<K,V>[] tab; int sc;

```

```

        while ((tab = table) == null || tab.length == 0) { //如果表为空
            if ((sc = sizeCtl) < 0) //判断sizeCtl大小，我们根据sizeCtl来申请具体大小的
内存空间
                Thread.yield(); // lost initialization race; just spin
            else if (U.compareAndSwapInt(this, SIZECTL, sc, -1)) { //SIZECTL表示
sizeCtl的偏移量//通过CAS
                try {
                    if ((tab = table) == null || tab.length == 0) { //第二次还需要
判断是为了避免重复初始化;
                        int n = (sc > 0) ? sc : DEFAULT_CAPACITY;
                        @SuppressWarnings("unchecked")
                        Node<K,V>[] nt = (Node<K,V>[])new Node<?,?>[n];
                        table = tab = nt;
                        sc = n - (n >>> 2); //0.75n
                    }
                } finally {
                    sizeCtl = sc; //记录扩容阈值
                }
                break;
            }
        }
        return tab;
    }
}

```

6.2 CAS过程

o: 传入的对象;

offset: 偏移量

expected: 期待值;

x: 如果o对象offset处值为expected，那么将其修改为x;

```
compareAndSwapInt(Object o, long offset, int expected, int x);
```

6.3 treeifyBin函数树化过程

```

private final void treeifyBin(Node<K,V>[] tab, int index) {
    Node<K,V> b; int n, sc;
    if (tab != null) {
        if ((n = tab.length) < MIN_TREEIFY_CAPACITY) //如果table最大长度小于64,
则不会树化，而是选择扩容
            tryPresize(n << 1); //扩容到原来的两倍

        //这里才会进行树化操作，因为是多线程情况下，所以一旦某个线程在树化，就会对Index位
        置的桶加锁，其他线程无法树化
        else if ((b = tabAt(tab, index)) != null && b.hash >= 0) {
            synchronized (b) {
                if (tabAt(tab, index) == b) {
                    TreeNode<K,V> hd = null, tl = null;
                    //循环将链表结构转换为树结构，这里不是真正的树化过程，只是通过
                    TreeNode的next指针将树连接起来
                    for (Node<K,V> e = b; e != null; e = e.next) {
                        TreeNode<K,V> p =

```

```

        new TreeNode<K,V>(e.hash, e.key, e.val,
                           null, null);
        if ((p.prev = t1) == null)
            hd = p;
        else
            t1.next = p;
        t1 = p;
    }
    //调用TreeBin实现真正的树化过程，并且将树的根结点存放到index下标
    的桶中
    setTabAt(tab, index, new TreeBin<K,V>(hd));
}
}
}
}
}
}
}

```

//扩容过程,这里传入的size是原数组大小扩容为2倍的值

```

private final void tryPresize(int size) {
    //首先判断扩容后数组大小是否超过了MAXIMUM_CAPACITY，大于的话调整为
    MAXIMUM_CAPACITY
    //否则，调用tableSizeFor调整为2的幂次方
    int c = (size >= (MAXIMUM_CAPACITY >>> 1)) ? MAXIMUM_CAPACITY :
        tableSizeFor(size + (size >>> 1) + 1);
    int sc; //记录sizeCtl
    while ((sc = sizeCtl) >= 0) { //自旋尝试扩容
        Node<K,V>[] tab = table; int n;
        //如果数组为null，或者数组长度为0，说明数组还没有初始化呢，那么进行初始化操作
        if (tab == null || (n = tab.length) == 0) {
            n = (sc > c) ? sc : c; //n为sizeCtl和c较大值
            //尝试CAS将sizeCtl的值替换为-1，也就是一个标记为，如果为-1说明有线程正在扩
            容，CAS成功，说明当前线程获得了扩容权限
            if (U.compareAndSwapInt(this, SIZECTL, sc, -1)) {
                try {
                    //双重判断，当前线程进入后再判断一次table是否为空，为空才会申请内存
                    if (table == tab) {
                        @SuppressWarnings("unchecked")
                        Node<K,V>[] nt = (Node<K,V>[])new Node<?,?>[n];
                        table = nt;
                        sc = n - (n >>> 2);
                    }
                } finally {
                    //sizeCtl = 0.75*capacity;
                    sizeCtl = sc;
                }
            }
        }
    }
    //数组不为空，扩容值小于sizeCtl或者大于最大容量，直接Break;
    else if (c <= sc || n >= MAXIMUM_CAPACITY)
        break;
    //table不为空的扩容过程
    else if (tab == table) {
        int rs = resizeStamp(n);
        if (sc < 0) {
            Node<K,V>[] nt;

            //sc右移动16位后和rs判读是否值相等
            if ((sc >>> RESIZE_STAMP_SHIFT) != rs || sc == rs + 1 ||

```



```

        sc == rs + MAX_RESIZERS || (nt = nextTable) == null ||
        transferIndex <= 0)
            break;
        if (U.compareAndSwapInt(this, SIZECTL, sc, sc + 1))
            transfer(tab, nt);
    }
    else if (U.compareAndSwapInt(this, SIZECTL, sc,
                                (rs << RESIZE_STAMP_SHIFT) + 2))
        transfer(tab, null);
    }
}

```

```

TreeBin(TreeNode<K,V> b) {
    super(TREEBIN, null, null, null);
    this.first = b; //first取b的头节点
    TreeNode<K,V> r = null; //r记录根结点
    //循环遍历树结点,
    for (TreeNode<K,V> x = b, next; x != null; x = next) {
        next = (TreeNode<K,V>)x.next;
        x.left = x.right = null;
        if (r == null) { //如果开始r为空, 那么r = x;
            x.parent = null;
            x.red = false;
            r = x;
        }
        //
        else {
            K k = x.key;
            int h = x.hash;
            Class<?> kc = null;
            for (TreeNode<K,V> p = r;;) { //p指向已经树化的那部分结点
                int dir, ph;
                K pk = p.key;
                if ((ph = p.hash) > h) //判断ph的hash值和h的hash值, 大于dir =
-1:
                    dir = -1;
                else if (ph < h) //小于 , dir = 1;
                    dir = 1;
                //如果ph == h也就是hash值相等, 那么调用tieBreakOrder判断需要插
入左子树还是右子树
                else if ((kc == null &&
                    (kc = comparableClassFor(k)) == null) ||
                    (dir = compareComparables(kc, k, pk)) == 0)
                    dir = tieBreakOrder(k, pk);
                TreeNode<K,V> xp = p;
                //如果dir <= 0, 并且p.left为空, 那么就将结点插入到左子树。
                // 如果dir > 0, 并且p.right为空, 那么就将结点插入到右子树。
                if ((p = (dir <= 0) ? p.left : p.right) == null) {
                    x.parent = xp;
                    if (dir <= 0)
                        xp.left = x;
                    else
                        xp.right = x;
                    //插值完成后, 可能会破坏红黑树的特性, 这里需要平衡一次
                    r = balanceInsertion(r, x);
                }
            }
        }
    }
}

```

```

        break;
    }
}
}
}
this.root = r;
assert checkInvariants(root);
}

/**
 * 用这个方法比较两个对象，返回值要么大于0，要么小于0，不会为0
 * 也就是说这一步一定能确定要插入的节点要么是树的左节点，要么是右节点，不然就无继续满足二叉树结构了
 * 先比较两个对象的类名，类名是字符串对象，就按字符串的比较规则
 * 如果两个对象是同一个类型，那么调用本地方法为两个对象生成hashCode值，再进行比较，hashCode相等的话返回-1
 */
static int tieBreakOrder(Object a, Object b) {
    int d;
    if (a == null || b == null ||
        (d = a.getClass().getName().
            compareTo(b.getClass().getName())) == 0)
        d = (System.identityHashCode(a) <= System.identityHashCode(b) ?
            -1 : 1);
    return d;
}

```

//平衡过程和HashMap的一样的，因为能够进入这个函数的时候，当前线程已经对桶加了锁，所以这里面不需要加锁考虑并发问题

```

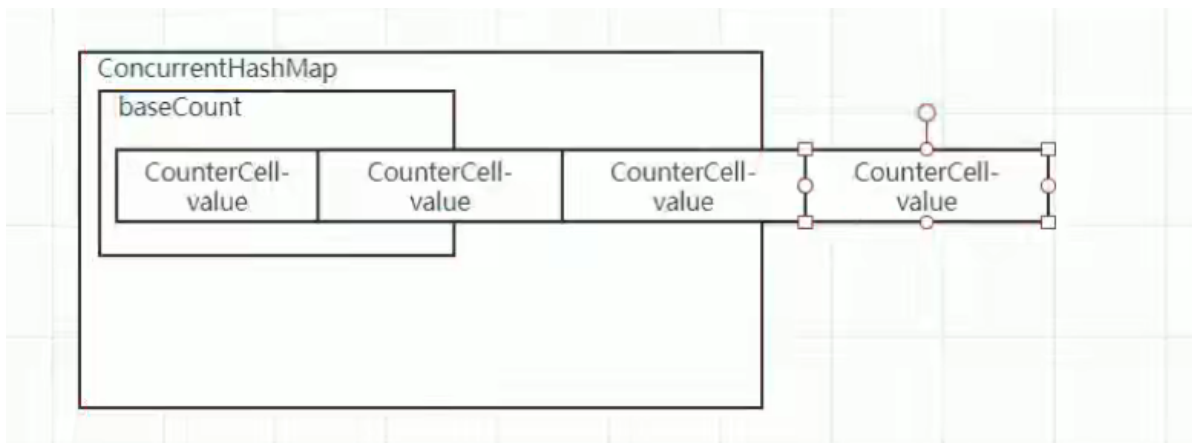
static <K,V> TreeNode<K,V> balanceInsertion(TreeNode<K,V> root,
                                           TreeNode<K,V> x) {
    x.red = true;
    for (TreeNode<K,V> xp, xpp, xpp1, xppr;;) {
        if ((xp = x.parent) == null) {
            x.red = false;
            return x;
        }
        else if (!xp.red || (xpp = xp.parent) == null)
            return root;
        if (xp == (xpp1 = xpp.left)) {
            if ((xppr = xpp.right) != null && xppr.red) {
                xppr.red = false;
                xp.red = false;
                xpp.red = true;
                x = xpp;
            }
        }
        else {
            if (x == xp.right) {
                root = rotateLeft(root, x = xp);
                xpp = (xp = x.parent) == null ? null : xp.parent;
            }
            if (xp != null) {
                xp.red = false;
                if (xpp != null) {
                    xpp.red = true;
                    root = rotateRight(root, xpp);
                }
            }
        }
    }
}

```

```

    }
}
else {
    if (xpp1 != null && xpp1.red) {
        xpp1.red = false;
        xp.red = false;
        xpp.red = true;
        x = xpp;
    }
    else {
        if (x == xp.left) {
            root = rotateRight(root, x = xp);
            xpp = (xp = x.parent) == null ? null : xp.parent;
        }
        if (xp != null) {
            xp.red = false;
            if (xpp != null) {
                xpp.red = true;
                root = rotateLeft(root, xpp);
            }
        }
    }
}
}
}
```

7.1 addCount判断是否需要扩容



- 首先尝试对baseCount进行加操作，如果没有加成功则找对应的桶下标的CounterCell做加一操作；
- 这里之所以需要使用counterCells是因为多个线程尝试对baseCount增加压力过大，使用counterCells数组分担压力，最后计算具体的count值，只需要将baseCount的值和counterCells里面的值求和就行；

```
private final void addCount(long x, int check) {
    CounterCell[] as; long b, s; //as中记录的是每个桶的结点数量

    //尝试把数量加到baseCount上，如果失败了再加到分段的CounterCell上
    if ((as = counterCells) != null ||
        !U.compareAndSwapLong(this, BASECOUNT, b = baseCount, s = b + x)) {
        CounterCell a; long v; int m;
        boolean uncontended = true;
        //如果当前as为空，或者长度为0，或者当前线程所在的段为null，或者当前线程的段上加数量失败
        if (as == null || (m = as.length - 1) < 0 ||
            (a = as[ThreadLocalRandom.getProbe() & m]) == null ||
            !(uncontended =
                U.compareAndSwapLong(a, CELLVALUE, v = a.value, v + x))) {
            //强制增加数量，不同线程对应的不同段都更新失败说明冲突已经发生，对counterCells进行扩容，减少多线程hash到同一段的概率；
            fullAddCount(x, uncontended);
            return;
        }
        //check大于1，为添加元素，小于等于1则不是
        if (check <= 1)
            return;
        //计算数组元素个数
        s = sumCount();
    }
    if (check >= 0) {
        Node<K,V>[] tab, nt; int n, sc;
        while (s >= (long)(sc = sizeCtl) && (tab = table) != null &&
            (n = tab.length) < MAXIMUM_CAPACITY) {

            //resizeStamp会将n向移动15为得到值给rs
            int rs = resizeStamp(n);
            if (sc < 0) { //小于0说明有其他线程正在执行扩容操作，线程会协助扩容
                if ((sc >>> RESIZE_STAMP_SHIFT) != rs || sc == rs + 1 ||
                    sc == rs + MAX_RESIZERS || (nt = nextTable) == null ||
                    transferIndex <= 0)
                    break;
                if (U.compareAndSwapInt(this, SIZECTL, sc, sc + 1))

```

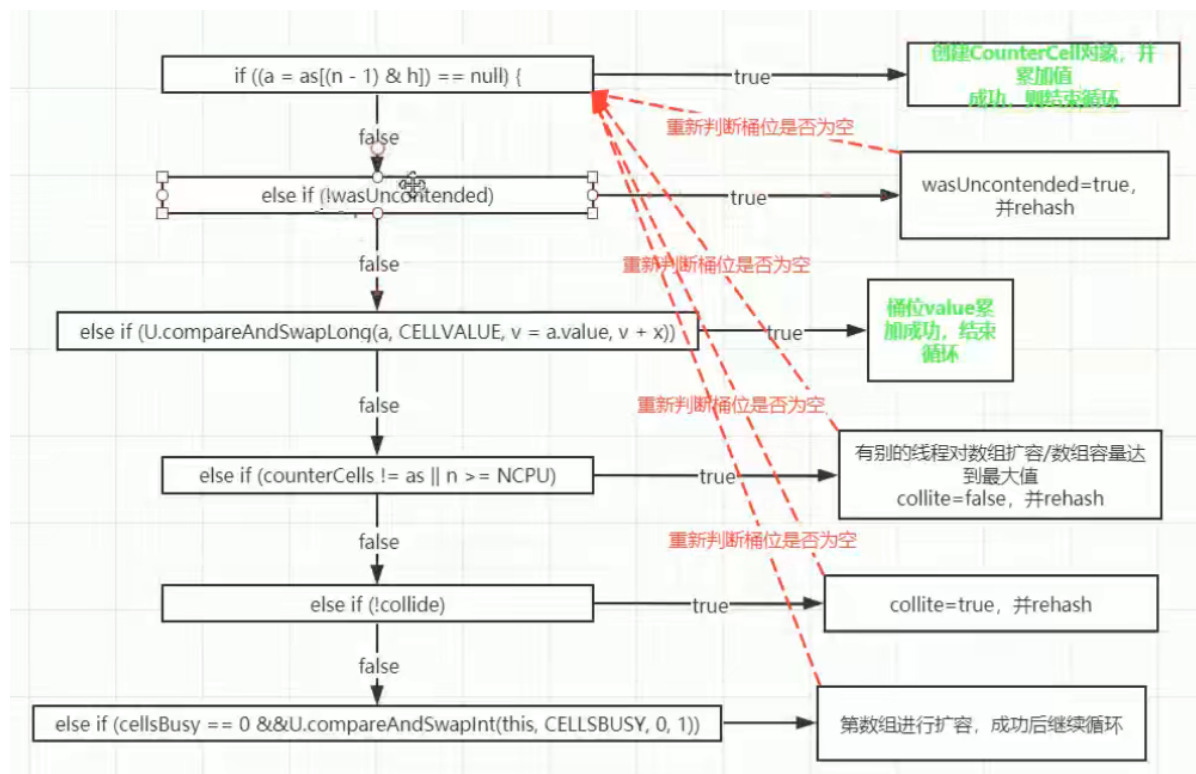
```

        transfer(tab, nt); //协助扩容
    }
    //这里rs向左移动16位, 加上之前移动的15位那么就是移动了31位, 最高位将会变成
    1,:
    //所以这里是尝试将sizeCtl修改为负值, 当sizeCtl为负值;
    else if (U.compareAndSwapInt(this, SIZECTL, sc,
        (rs << RESIZE_STAMP_SHIFT) + 2))
        //如果修改成功, 则执行扩容过程
        transfer(tab, null); //进行扩容
    s = sumCount();
    }
}
}

```

7.2 fullAddCount函数 (countercells为空, 或者当线程无法对countercells加值时)

强制增加数量, 不同线程对应的不同段都更新失败说明冲突已经发生, 对counterCells进行扩容, 减少多线程hash到同一段的概率;



```

private final void fullAddCount(long x, boolean wasUncontended) {
    int h;
    //计算数组角标h
    if ((h = ThreadLocalRandom.getProbe()) == 0) {
        ThreadLocalRandom.localInit(); // force initialization
        h = ThreadLocalRandom.getProbe();
        wasUncontended = true;
    }
    boolean collide = false; // True if last slot nonempty
    for (;;) {
        CounterCell[] as; CounterCell a; int n; long v;

```

```

//如果counterCells不为空，并且as的长度大于0
if ((as = counterCells) != null && (n = as.length) > 0) {
    //如果当前线程对应的counterCell位置中槽位为空，在此位置添加一个
    CounterCell元素；
    //再并发环境下，可能存在多线程同时尝试对counterCells槽中位置添加
    counterCell所以cellBusy是一个标记位
    //如果cellsBusy为0，才可以添加counterCell；
    if ((a = as[(n - 1) & h]) == null) {
        //如果当前CounterCell没有再扩容，使用Double Check；
        if (cellsBusy == 0) { // Try to attach new Cell
            CounterCell r = new CounterCell(x); // Optimistic create
            //
            if (cellsBusy == 0 &&
                U.compareAndSwapInt(this, CELLSBUSY, 0, 1)) {
                boolean created = false;
                //尝试在j位置添加counterCell；
                try { // Recheck under lock
                    CounterCell[] rs; int m, j;
                    if ((rs = counterCells) != null &&
                        (m = rs.length) > 0 &&
                        rs[j = (m - 1) & h] == null) {
                        rs[j] = r;
                        created = true;
                    }
                } finally {
                    //添加成功后将cellsBusy置为0
                    cellsBusy = 0;
                }
                if (created)
                    break;
                continue; // Slot is now non-empty
            }
        }
        collide = false;
    }
    //如果对应counterCells槽位置不为空
    else if (!wasUncontended) // CAS already known to fail
        wasUncontended = true; // Continue after rehash
    //尝试CAS修改counterCells对应槽中的value值；
    else if (U.compareAndSwapLong(a, CELLVALUE, v = a.value, v + x))
        break;

    //如果数组的长度大于cpu个数，没必要进行扩容了
    else if (counterCells != as || n >= NCPU)
        collide = false; // At max size or stale
    //为扩容做准备
    else if (!collide)
        collide = true;
    //如果这些线程尝试向counterCells中添加元素失败了，说明这个counterCells太小了，需要进行扩容操作。CAS + 自旋尝试对counterCells扩容，扩容为原来的两倍，并且将原来的counterCells数组元素拷贝到新数组中；
    else if (cellsBusy == 0 &&
        U.compareAndSwapInt(this, CELLSBUSY, 0, 1)) {
        try {
            if (counterCells == as) { // Expand table unless stale
                CounterCell[] rs = new CounterCell[n << 1];
                for (int i = 0; i < n; ++i)
                    rs[i] = as[i];
            }
        }
    }
}

```

```

        counterCells = rs;
    }
} finally {
    cellsBusy = 0;
}
collide = false;
continue;           // Retry with expanded table
}
h = ThreadLocalRandom.advanceProbe(h);
}
//如果CounterCells没有被初始化，没有其他线程对其进行初始化，那么当前线程尝试CAS
修改cellsBusy为1
//默认初始化为2
else if (cellsBusy == 0 && counterCells == as &&
    U.compareAndSwapInt(this, CELLSBUSY, 0, 1)) {
    boolean init = false;
    try {
        // Initialize table
        if (counterCells == as) {
            CounterCell[] rs = new CounterCell[2];
            rs[h & 1] = new CounterCell(x);
            counterCells = rs;
            init = true;
        }
    } finally {
        cellsBusy = 0;
    }
    if (init)
        break;
}
//如果数组正在被创建的过程中，那么当选线程优先往basecount里面加值；
else if (U.compareAndSwapLong(this, BASECOUNT, v = baseCount, v +
x))
    break;           // Fall back on using base
}
}
}

```

7.3 helpTransfer协助扩容

```

final Node<K,V>[] helpTransfer(Node<K,V>[] tab, Node<K,V> f) {
    Node<K,V>[] nextTab; int sc;
    if (tab != null && (f instanceof ForwardingNode) &&
        (nextTab = ((ForwardingNode<K,V>)f).nextTable) != null) {
        int rs = resizeStamp(tab.length); //n左移动15位
        //如果有线程正在执行扩容
        while (nextTab == nextTable && table == tab &&
            (sc = sizeCtl) < 0) {
            //判断是否已经有足够的线程进行扩容，也就是slot不够
            if ((sc >>> RESIZE_STAMP_SHIFT) != rs || sc == rs + 1 ||
                sc == rs + MAX_RESIZERS || transferIndex <= 0)
                break;
            //Cas尝试协助扩容；
            if (U.compareAndSwapInt(this, SIZECTL, sc, sc + 1)) {
                transfer(tab, nextTab);
                break;
            }
        }
    }
}

```

```

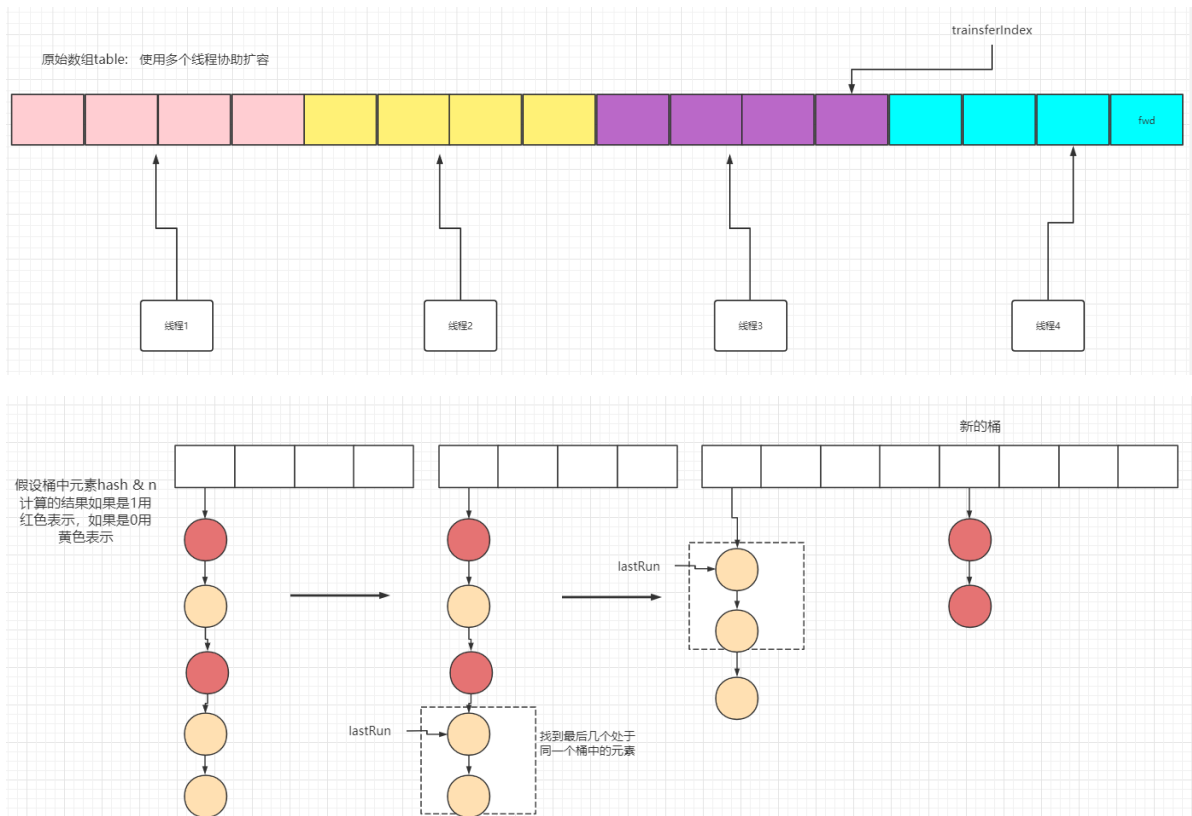
    }
    return nextTab;
}
return table;
}

```

7.4 transfer迁移元素

每个线程负责一部分的元素进行迁移，最小的任务单元的16;

如果table某个slot位置的元素已经被迁移完成，则将其位置设置为ForwardingNode结点; 结点内部hash值为MOVED



```

private final void transfer(Node<K,V>[] tab, Node<K,V>[] nextTab) {
    int n = tab.length, stride;
    //最小的任务单元是16个
    if ((stride = (NCPUs > 1) ? (n >>> 3) / NCPUs : n) < MIN_TRANSFER_STRIDE)
        stride = MIN_TRANSFER_STRIDE; // subdivide range
    //如果是正在扩容的线程而不是协助扩容的线程
    if (nextTab == null) { // initiating
        try {
            @SuppressWarnings("unchecked")
            //新建一个数组，容量为原来数组的两倍
            Node<K,V>[] nt = (Node<K,V>[])new Node<?,?>[n << 1];
            nextTab = nt;
        } catch (Throwable ex) { // try to cope with OOME
            sizeCtl = Integer.MAX_VALUE;
            return;
        }
        //nextTable记录新数组
        nextTable = nextTab;
        //记录旧数组的size
        transferIndex = n;
    }
}

```



```

int nextn = nextTab.length;
ForwardingNode<K,V> fwd = new ForwardingNode<K,V>(nextTab);
boolean advance = true;
boolean finishing = false; // to ensure sweep before committing nextTab
for (int i = 0, bound = 0;;) {
    Node<K,V> f; int fh;

    //计算当前线程需要负责哪些部分
    while (advance) {
        int nextIndex, nextBound;
        if (--i >= bound || finishing)
            advance = false;
        else if ((nextIndex = transferIndex) <= 0) {
            i = -1;
            advance = false;
        }
        //判断集合长度是否大于步长, 大于nextBound = nextIndex - stride否则等于
        0, 每次都会修改transferindex
        else if (U.compareAndSwapInt
            (this, TRANSFERINDEX, nextIndex,
             nextBound = (nextIndex > stride ?
                          nextIndex - stride : 0))) {
            bound = nextBound;
            i = nextIndex - 1;
            advance = false;
        }
    }
    //判断是否扩容完成了
    if (i < 0 || i >= n || i + n >= nextn) {
        int sc;
        //如果扩容完成了, 那么将新的Table赋值给原来的table;
        if (finishing) {
            nextTable = null;
            table = nextTab;
            sizeCtl = (n << 1) - (n >>> 1); //sizeCtl = 0.75* n; n<<1会
            改变n的值, n>>>1不会改变n的值
            return;
        }
        //如果扩容动作做完了, 那么才会将finish标记为true
        if (U.compareAndSwapInt(this, SIZECTL, sc = sizeCtl, sc - 1)) {
            if ((sc - 2) != resizeStamp(n) << RESIZE_STAMP_SHIFT)
                return;
            finishing = advance = true;
            i = n; // recheck before commit
        }
    }
    //这一部分内容是具体的迁移过程
    else if ((f = tabAt(tab, i)) == null) //如果i的位置元素为空, 那么不需要进行
    迁移
        advance = castTabAt(tab, i, null, fwd); //直接往当前位置设置fwd;
    else if ((fh = f.hash) == MOVED)
        advance = true; // already processed
    else {
        synchronized (f) { //对迁移的位置加锁, 这是为了防止多线程环境下一个线程正在
        迁移元素, 而其他线程对此位置进行元素添加或者删除操作

            //首先判断还是之前那个slot吗, 是的话继续迁移
            if (tabAt(tab, i) == f) {

```

```

Node<K,V> ln, hn;
// 第一个元素的hash值大于等于0
// 说明该桶中元素是以链表形式存储的
// 这里与HashMap迁移算法基本类似
// 唯一不同的是多了一步寻找lastRun
// 这里的lastRun是提取出链表后面不用处理再特殊处理的子链表
// 比如所有元素的hash值与桶大小n与操作后的值分别为 0 0 4 4 0 0 0
// 则最后后面三个0对应的元素肯定还是在同一个桶中
// 这时lastRun对应的就是倒数第三个节点
// 至于为啥要这样处理，我也没太搞明白
if (fh >= 0) { //如果hash值大于0
    //等于0的放到低位链表(low)中，不等于0的放到高位链表(high)中
    int runBit = fh & n; //记录在新的桶中迁移是迁移到元素下标位
置还是原始下标位置加原始数组大小

    Node<K,V> lastRun = f; //记录桶中最近一个迁移元素

    //找到最后几个必定在同一个桶的元素
    for (Node<K,V> p = f.next; p != null; p = p.next) {
        int b = p.hash & n; //记录下一个元素的下标位置
        if (b != runBit) { //如果下一个下标和前一个下标位置不
            一致

                runBit = b; //当前迁移下标
                lastRun = p; //记录当前迁移元素
            }
        }
        if (runBit == 0) { //runBit == 0说明在新数组中迁移到位置
            不变

                ln = lastRun;
                hn = null;
            }
        else { //runBit == 1说明在新数组中迁移到位置改变了
                hn = lastRun;
                ln = null;
            }
        }

        //遍历链表，把hash&n == 0的放在Low链表中、等于1的放入高链表
        中;

        for (Node<K,V> p = f; p != lastRun; p = p.next) {
            int ph = p.hash; K pk = p.key; V pv = p.val;
            if ((ph & n) == 0)
                ln = new Node<K,V>(ph, pk, pv, ln);
            else
                hn = new Node<K,V>(ph, pk, pv, hn);
        }
        setTabAt(nextTab, i, ln); //将ln存放到低位
        setTabAt(nextTab, i + n, hn); //将hn存放到低位
        setTabAt(tab, i, fwd); //迁移完成的原来数组i位置设置为fwd;
        advance = true;
    }

    //如果当前槽是树结构
    // lo 记录低位头结点，loTail记录低位末尾节点，lc记录低位结点数
    // hi 记录高位头结点，hiTail记录高位末尾节点，hc记录高位结点数
    else if (f instanceof TreeBin) {
        TreeBin<K,V> t = (TreeBin<K,V>)f;
        TreeNode<K,V> lo = null, loTail = null;
        TreeNode<K,V> hi = null, hiTail = null;
        int lc = 0, hc = 0;
        for (Node<K,V> e = t.first; e != null; e = e.next) {

```

```

        int h = e.hash;
        TreeNode<K,V> p = new TreeNode<K,V>
            (h, e.key, e.val, null, null);
        if ((h & n) == 0) {
            if ((p.prev = loTail) == null)
                lo = p;
            else
                loTail.next = p;
            loTail = p;
            ++lc;
        }
        else {
            if ((p.prev = hiTail) == null)
                hi = p;
            else
                hiTail.next = p;
            hiTail = p;
            ++hc;
        }
    }

    //判断低位是否需要树退化，不需要退化就重新建立低位的红黑树结构
    ln = (lc <= UNTREEIFY_THRESHOLD) ? untreeify(lo) :
        (hc != 0) ? new TreeBin<K,V>(lo) : t;
    //判断高位是否需要树退化，不需要退化就建立高位的红黑树结构
    hn = (hc <= UNTREEIFY_THRESHOLD) ? untreeify(hi) :
        (lc != 0) ? new TreeBin<K,V>(hi) : t;

    setTabAt(nextTab, i, ln); //设置低位slot结点为root
    setTabAt(nextTab, i + n, hn); //设置高位slot结点
    setTabAt(tab, i, fwd); //设置原来数组位置为fwd;
    advance = true;
}

}

}

}

}

```