

# Nachos Assignment

## Project 3 - Memory Management

Qi Zheng, Lin.

Advisor: Farn, Wang.

# Outline

- Memory Management
- Report
- Policy

# Memory Management

About the virtual memory manager, swap, page fault handler and etc.

# Goal of Memory Management

- Both the test cases require lots of memory.
  - `/test/matmult.c`
  - `/test/sort.c`
  - Larger than the physical memory
- Goal:
  - Run this two programs concurrently, and get the correct result.
  - Remember to specify the scheduling algorithm you used in the report.
  - **DO NOT** pass by simply modify the memory size in “machine”.

# Hints

- File system – swap space.
- Maintain three tables to record the information.
  - PageTable
  - FrameTable
  - SwapTable
- Think how to “load” the page correctly and deal with the PageFaultException
- Design your own “Virtual Memory Manager”
- Some of you have figure out some point in project 1. Remember to write them down in report

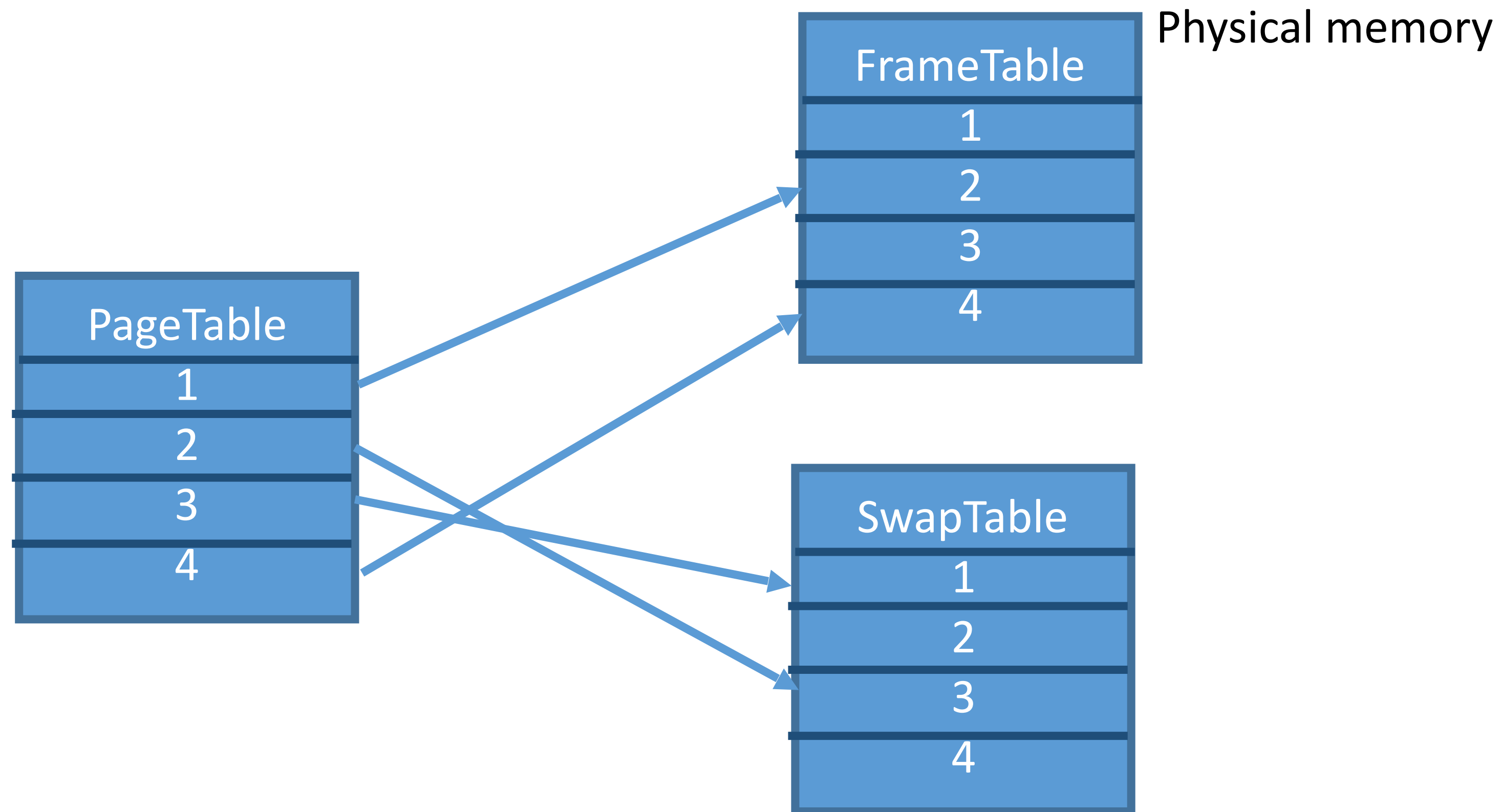
# Swap Space

- swap = `new SynchDisk` in your kernel.
- Use the disk as the swap space.
- Access the virtual memory in the disk by...
  - `kernel->swap->WriteSector`
  - `kernel->swap->ReadSector`
- See “`synchdisk.cc`” and other files in `/filesys` for details!
- Read the header in those files first.

# Maintain Three Tables

- **PageTable**
  - One page table per process.
  - Decide your virtual page number.
- **FrameTable**
  - Record every physical page's information.
  - Each frame represent one physical page.
- **SwapTable**
  - Record every sector's information in swap.
  - The number of entries in SwapTable is the same as the swap sectors.
  - Each entry represent one frame in the disk.

# Maintain Three Tables Cont.





# PageTable

- For each entry in PageTable
- TranslationEntry {
  - unsigned int virtualPage; //virtual memory page
  - unsigned int physicalPage; //if in physical memory
  - bool valid; //if in physical memory
  - bool use; //been used(read or write)
  - bool dirty; //been modified

# FrameTable

- For each entry in FrameTable
- FrameInfoEntry {
  - Bool valid; //if being used
  - Bool lock;
  - AddrSpace \*addrSpace; //which process is using this page
  - Unsigned int vpn; //which virtual page of the process is stored in  
//this page
- };

# SwapTable

- For each entry in SwapTable
  - FrameInfoEntry { //exactly same as FrameTable
    - Bool valid; //if being used
    - Bool lock;
    - AddrSpace \*addrSpace; //which process is using this page
    - Unsigned int vpn; //which virtual page of the process is stored in  
//this page
- };

# Addrspac::Load

- Note: Load 1 page a time.
- First ask for 1 page, if the FrameTable is full, select 1 frame and put it into SwapTable.
  - Design your own **page replacement method** to get the frame.
  - Specify the method in your report.
- Mapping virtual address to physical address.
- Invoke “executable->ReadAt(&(kernel->machine->mainMemory[physical address]), sizeToLoadNow, inFileAddr)”

# Address Mapping

- Map the Virtual Address to Physical Address.
- Like the way you used in **project 1**.
- Physical Address =  
$$\text{pageTable}[(\text{virtual address} / \text{PageSize})].\text{physicalPage} * \text{PageSize} + (\text{virtual address} \% \text{PageSize})$$

# Maybe Create a New Class “Memory Manager”

- Class MemoryManager {

- Int TransAddr(AddrSpace \*space, int virtAddr);

- //return phyAddr (translated from virtAddr)

- Bool AcquirePage(AddrSpace \*space, int vpn);

- //ask a page (frame) for vpn

- Bool ReleasePage(AddrSpace \*space, int vpn);

- //free a page

- Void PageFaultHandler();

- //will be called when manager want to swap a page from SwapTable

- //to FrameTable and the FrameTable is full.

- };

# Page-Fault Handler

- Put the pages in SwapTable into FrameTable.
- When all physical memory frames are occupied, design a page replacement method to get a frame.
  - Remember to specify it in your report!
- You can do this work in your own way.

# Some File May Be Useful

- For the disk usage details, see:
  - `/filesys/synchdisk.h`
  - `/filesys/synchdisk.cc`
- Other files in `/filesys` (Optional).
- For the swap space initialization:
  - `/userprog/userkernel.h`
  - `/userprog/userkernel.cc`
- For the table maintaining, see:
  - `/machine/machine.h` and `/machine/machine.cc`
  - `/machine/translate.h` and `/machine/translate.cc`



# Some File May Be Useful Cont

- For the table maintaining, see:
  - /machine/machine.h and /machine/machine.cc
  - /machine/translate.h and /machine/translate.cc
- For the loading of pages.
  - `userprog/addrspace.h`
  - `userprog/addrspace.cc`
- Always **see the header and comments** first.
- Based on your implementation, there might be different files that your need to see and modify.

# Report & Policy

Report contents, grading policy

# Report

- Report
  - Motivation and the problem analysis
  - What's your plan to deal with the problem (high-level)
  - You can including some important code segments and comments
  - Experiment result and some discussion
  - Tell me **why and how** it happened
  - Remember there are two parts in project3
- Please saved as [Student ID]\_report.pdf
  - E.g. r04921119\_report.pdf

# Code and Report

- Upload to CEIBA
  - Do not mail me the homework please
- Source code and report BOTH
  - create a folder and follow the structure below  
/r04921119\_Nachos3
    - |\_\_\_ /nachos-4.0
    - |\_\_\_ r04921119\_report.pdf
- `tar zcvf r04921119_Nachos2.tar.gz ./r04921119_Nachos2`

# Policy

- Nachos source code: (40%)
- Report: (60%)
  - Important !!! Tell as detail as you can
  - **why you choose, what you do, and why it works**
- Bonus
  - The replacement method you design.
  - Extra observation or modification on nachos.
  - Tell me how to switch and show in report