

# Cryptanalysis of Vigenère

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## Find key length t

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Here use Kasiski Examination to find the GCD of distance for all patterns of length 3

```
def find_repeated_patterns(text, length):
    pattern_positions = defaultdict(list)

    for i in range(len(text) - length + 1):
        pattern = text[i:i + length]
        pattern_positions[pattern].append(i)

    repeated_patterns = {pattern: positions for pattern, positions in pattern_positions.items()
if len(positions) > 1}

    return repeated_patterns

def find_gcd_of_distances(positions):
    distances = [positions[i+1] - positions[i] for i in range(len(positions) - 1)]
    if distances:
        return reduce(gcd, distances)
    return None

def kasiski_examination(text, length=2):
    repeated_patterns = find_repeated_patterns(text, length)
    gcd_results = {}

    for pattern, positions in repeated_patterns.items():
        gcd_value = find_gcd_of_distances(positions)
        if gcd_value:
            gcd_results[pattern] = gcd_value

    return gcd_results

# Cipher Text
cipher_text =
"FRSGKIPMJLZNZJTAJZAQBJXLZKODZFGSCPERSUSILEBNVTNQZXKLHJZGCMFJDGNSVEYOBNZILGFVZTIXSNSQRWOQGYWCEJL

# Running Kasiski Examination for patterns of length 3
kasiski_results_length_2 = kasiski_examination(cipher_text, length=3)
```

Plot the Top 10 Most Frequent GCD Values in Kasiski Examination Results

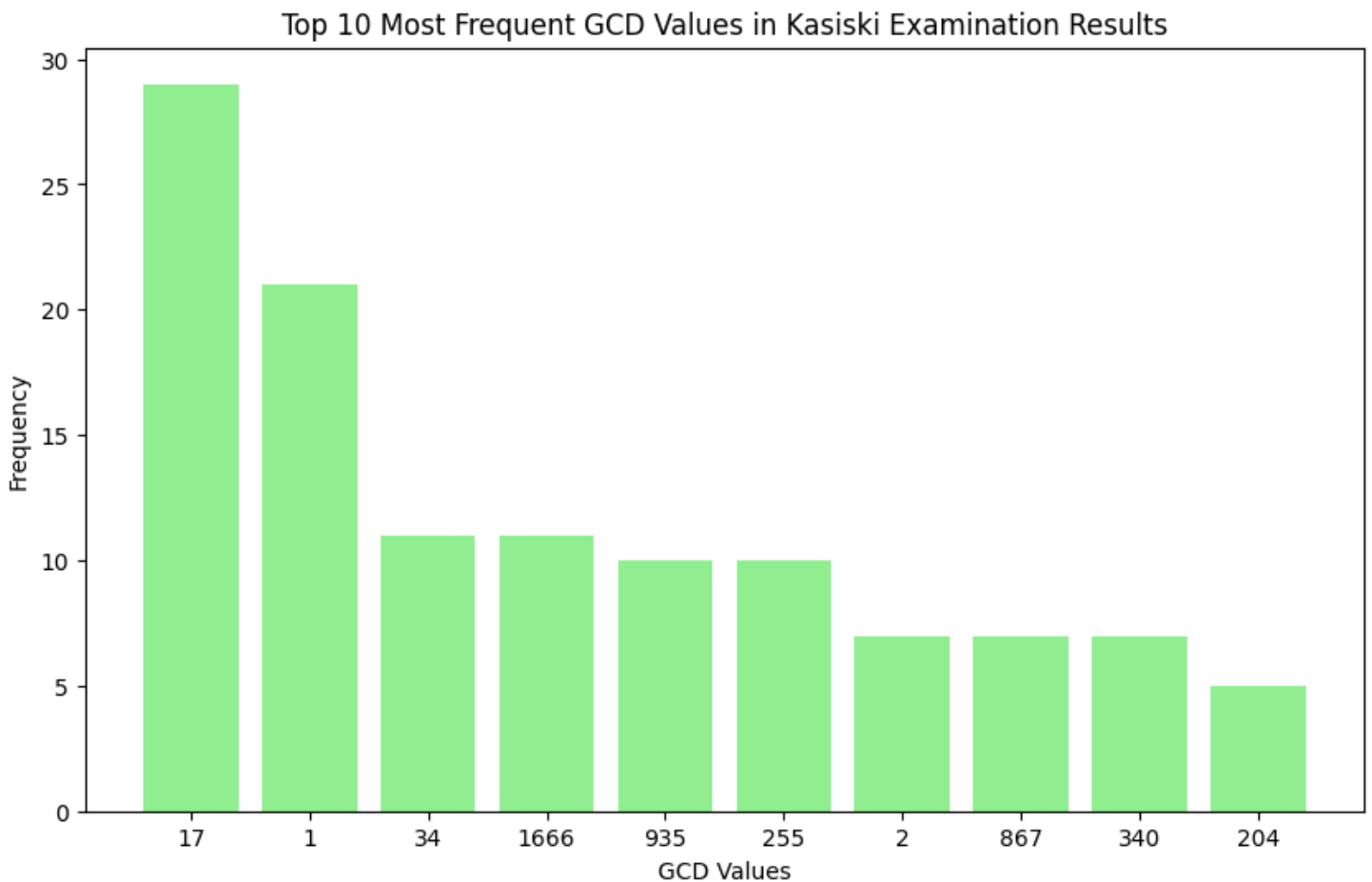
```
gcd_values = list(kasiski_results_length_4.values())
gcd_frequency = defaultdict(int)
for gcd_value in gcd_values:
    gcd_frequency[gcd_value] += 1

sorted_gcd_frequency = sorted(gcd_frequency.items(), key=lambda x: x[1], reverse=True)

top_10_gcd = sorted_gcd_frequency[:10]

gcds_top_10 = [item[0] for item in top_10_gcd]
frequencies_top_10 = [item[1] for item in top_10_gcd]

plt.figure(figsize=(10, 6))
plt.bar(range(len(gcds_top_10)), frequencies_top_10, color='lightgreen')
plt.xlabel('GCD Values')
plt.ylabel('Frequency')
plt.title('Top 10 Most Frequent GCD Values in Kasiski Examination Results')
plt.xticks(range(len(gcds_top_10)), gcds_top_10)
plt.show()
```



Here we find the the most frequent GCD of distance for all patterns of length 3 is 17 and there are also many multiple of 17 were very frequent. So we make the decision to choose the key length  $t = 17$  here.

## Verify key length t

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We can also verify the correctness of  $t = 17$  using the method index of coincidence

Here print the sum  $S = \sum_{i=0}^{25} q_i^2$  for all sequence  $c_a c_{a+t} c_{a+2t} \dots$ ,  $a \in [1, 17]$

```
# Define the alphabet (A = 0, ..., Z = 25)
alphabet = {chr(i + ord('A')): i for i in range(26)}

def calculate_S_values(ciphertext, period=17):
    n = len(ciphertext)
    S_values = []

    for i in range(period):
        sequence = [ciphertext[j] for j in range(i, n, period) if ciphertext[j] in alphabet]

        frequency_count = Counter(sequence)

        total_count = sum(frequency_count.values())
        q_i = {alphabet[char]: count / total_count for char, count in frequency_count.items()}

        S_i = sum((q_i.get(i, 0) ** 2) for i in range(26))

        S_values.append(S_i)

    return S_values

S_values = calculate_S_values(cipher_text, period=17)
S_values
```

```
[0.07843076058799663,
 0.07711994007677664,
 0.06987754071455624,
 0.06943567731347051,
 0.07183436434793586,
 0.0706350208307032,
 0.07669486175987882,
 0.0765686150738543,
 0.07063502083070318,
 0.06470142658755208,
 0.07820982199217269,
 0.07088751420275215,
 0.06665825022093169,
 0.06482767327357657,
 0.0807978790556748,
 0.06949880065648276,
 0.07852543870723391]
```

We can see sums were all  $\approx 0.065$  when  $t = 17$

Also verify sum  $\approx 0.038$  when  $t \neq 17$  (not shown here since it's a lot...)

So the key length  $t = 17$  should be correct

## compute key k

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Use the method that calculating  $I_j = \sum_{i=0}^{25} p_i \cdot q_{i+j}$ , and find k by determining  $I_k \approx 0.065$  and  $I_j < 0.0455, j \neq k$  for all sequence  $c_a c_{a+t} c_{a+2t} \dots, a \in [1, 17]$

```

# English letter frequencies, p_j for j=0 to 25 corresponding to A to Z
english_frequencies = np.array([
    0.08167, 0.01492, 0.02782, 0.04253, 0.12702, # A - E
    0.02228, 0.02015, 0.06094, 0.06966, 0.00153, # F - J
    0.00772, 0.04025, 0.02406, 0.06749, 0.07507, # K - O
    0.01929, 0.00095, 0.05987, 0.06327, 0.09056, # P - T
    0.02758, 0.00978, 0.02360, 0.00150, 0.01974, # U - Y
    0.00074 # Z
])

# Alphabet to index dictionary
alphabet = {chr(i): i - ord('A') for i in range(ord('A'), ord('Z') + 1)}
index_to_char = {i: chr(i + ord('A')) for i in range(26)}

# Function to calculate frequency of each letter in a sequence
def calculate_frequencies(sequence):
    total_count = len(sequence)
    frequencies = np.zeros(26)

    counter = Counter([letter for letter in sequence if letter in alphabet])

    for letter, count in counter.items():
        frequencies[alphabet[letter]] = count / total_count

    return frequencies

sequences = []
for i in range(17):
    sequence = cipher_text[i::17]
    sequences.append(sequence)

k_i = []

for i in range(17):
    sequence = sequences[i]
    q_i_j = calculate_frequencies(sequence)

    max_S_i = -float('inf')
    best_k = 0

    for k in range(26):
        S_i = np.sum(english_frequencies * np.roll(q_i_j, -k))

        if S_i >= 0.06:
            best_k = k
            break

    k_i.append(best_k)

key_letters = ''.join([index_to_char[k] for k in k_i])

print(f"The key is: {key_letters}")

```

The key is: QAOBKGLTCDVRRRTZVL

## Decryption

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Decrypt the cipher text using the key we found

```
# Function to decrypt Vigenère cipher
def vigenere_decrypt(ciphertext, key):
    decrypted_text = []
    for i in range(len(ciphertext)):
        cipher_char = ord(ciphertext[i]) - ord('A')
        key_char = ord(key[i % len(key)]) - ord('A')
        plain_char = (cipher_char - key_char + 26) % 26
        decrypted_text.append(chr(plain_char + ord('A')))
    return ''.join(decrypted_text)

key = "QAOBKGLTCDVRRRTZVL"

# Decrypt the ciphertext
plain_text = vigenere_decrypt(cipher_text, key)
plain_text
```

The plaintext is:

PREFACE THIS IS A BOOK ON INFORMATION THEORETICALLY SECURE MULTIPARTY COMPUTATION MPC AND SECRET SHARING AND ABOUT THE INTIMATE AND FASCINATING RELATIONSHIP BETWEEN THE TWO NOTIONS WE DECIDED TO WRITE THE BOOK BECAUSE WE FELT THAT A COMPREHENSIVE TREATMENT OF UNCONDITIONALLY SECURE TECHNIQUES FOR MPC WAS MISSING IN THE LITERATURE IN PARTICULAR BECAUSE SOME OF THE FIRST GENERAL PROTOCOLS WERE FOUND BEFORE APPROPRIATE DEFINITIONS OF SECURITY HAD CRYSTALLIZED PROOFS OF THOSE BASIC SOLUTIONS HAVE BEEN MISSING SO FAR WE PRESENT THE BASIC FEASIBILITY RESULTS FOR UNCONDITIONALLY SECURE MPC FROM THE LATE EIGHTIES GENERALIZATION TO ARBITRARY ACCESS STRUCTURES USING LINEAR SECRET SHARING AND A SELECTION OF MORE RECENT TECHNIQUES FOR EFFICIENCY IMPROVEMENTS WE ALSO PRESENT OUR OWN VARIANT OF THE UC FRAMEWORK IN ORDER TO BE ABLE TO GIVE COMPLETE AND MODULAR PROOFS FOR THE PROTOCOLS WE PRESENT WE ALSO PRESENT A GENERAL TREATMENT OF THE THEORY OF SECRET SHARING AND IN PARTICULAR SECRET SHARING SCHEMES WITH ADDITIONAL ALGEBRAIC PROPERTIES ALSO AS SUBJECT WHERE A TREATMENT ON TEXT BOOK LEVEL SEEMS TO BE MISSING THE LITERATURE ONE OF THE THINGS WE FOCUS ON IS ASYMPTOTIC RESULTS FOR MULTIPLICATIVE SECRET SHARING WHICH HAS VARIOUS INTERESTING APPLICATIONS THAT WE PRESENT IN THE MPC PART OUR AMBITION HAS BEEN TO CREATE A BOOK THAT WILL BE OF INTEREST TO BOTH COMPUTER SCIENTISTS AND MATHEMATICIANS AND CAN BE USED FOR TEACHING AT SEVERAL DIFFERENT LEVELS WHERE DIFFERENT PARTS OF THE BOOK WILL BE USED WE HAVE THEREFORE TRIED TO MAKE BOTH MAIN PARTS BE SELF-CONTAINED EVEN IF THIS IMPLIES SOME OVERLAP BETWEEN THE PARTS THIS MEANS THAT THERE ARE SEVERAL DIFFERENT WAYS TO READ THE BOOK AND WE GIVE A FEW SUGGESTIONS FOR THIS BELOW IN PARTICULAR THE CONCEPT OF SECRET SHARING OF COURSE APPEARS PROMINENTLY IN BOTH PARTS IN THE FIRST MPC PART HOWEVER IT IS INTRODUCED ONLY AS A TOOL ON AN EDITORIAL BASIS IN THE SECOND PART WHERE IT IS INTRODUCED THEN BUT NOW AS A GENERAL CON

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