

Week 5

Start at 8: 05, Scan the QR code

P1: Head Guards

Double inclusion

File "grandparent.h"

```
struct foo {  
    int member;  
};
```

File "parent.h"

```
#include "grandparent.h"
```

File "child.c"

```
#include "grandparent.h"  
#include "parent.h"
```

Result

```
struct foo {  
    int member;  
};  
struct foo {  
    int member;  
};
```

[compilation error](#), since the structure type `foo` will thus be defined twice

#include guards

File "grandparent.h"

```
#ifndef GRANDPARENT_H  
#define GRANDPARENT_H  
  
struct foo {  
    int member;  
};  
  
#endif /* GRANDPARENT_H */
```

File "parent.h"

```
#include "grandparent.h"
```

File "child.c"

```
#include "grandparent.h"  
#include "parent.h"
```

```
#ifndef GRANDPARENT_H // => no exits, start if statement  
#define GRANDPARENT_H // => define
```

```

struct foo {
    int member;
};

#endif /* GRANDPARENT_H */ // => end if

#ifndef GRANDPARENT_H      // => already exist? skip the next until
#endif
#define GRANDPARENT_H

struct foo {
    int member;
};

#endif /* GRANDPARENT_H */

```

Result

```

struct foo {
    int member;
};

```

`#ifndef` test returns false, the preprocessor skips down to the `#endif`

P2: Stack and Heap

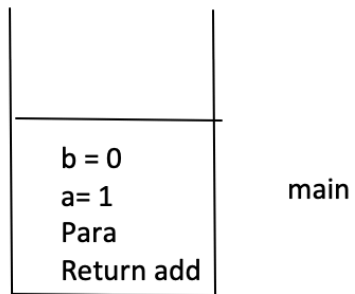
Stack

To keep track of the current memory place, there is a special processor register called **Stack Pointer**. Every time you need to save something — like a variable or the return address from a function — it pushes and moves the stack pointer up. Every time you exit from a function, it pops everything from the stack pointer until

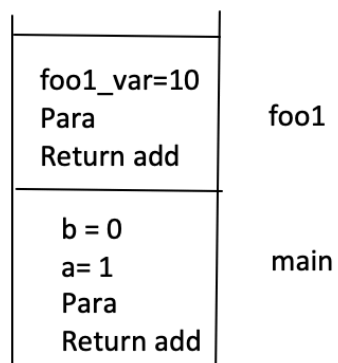
the saved return address from the function.

```
void fool(char c){  
    int fool_var = 10;  
    return;  
}  
  
int main(){  
    int a = 1;  
    int b = 0;  
    fool('x');  
}
```

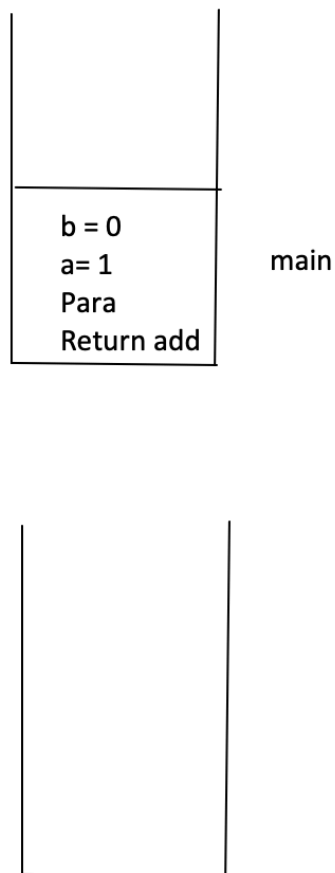
The memory for stack



The memory for stack



The memory for stack



A bad example:

```
#include <stdio.h>

int* createArray(int size){
    int arr[size];
    return arr;
}

int main(){
    int s = 10;
    int* arr = createArray(s);
```

```

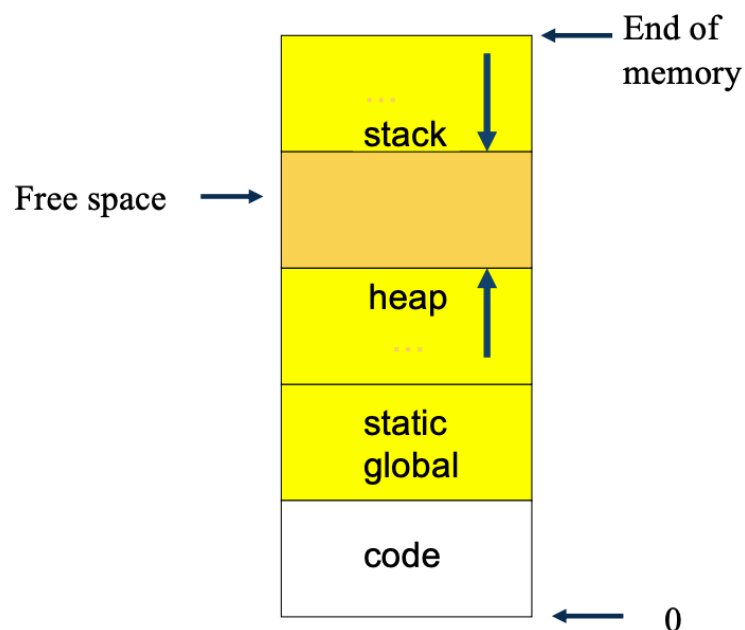
for (int i = 0; i < s; i++){
    arr[i] = i;
}
return 0;
}

```

Heap

Exist independently of functions and scopes

Memory Layout



```

#include <stdlib.h>

/*
  Returns a pointer to the allocated memory, if
  successful, or a NULL pointer if unsuccessful

  int* arr = (int*) malloc(sizeof(int) * 10);
*/
void* malloc(size_t size);

```

```

/*
  It has two arguments:
  - num specifies the number of "blocks" of contiguous
  memory
  - size specifies the size of each block

  • The allocated memory is cleared (set to '0').
  // assume you are create an array with num elements and each of them
  has size
*/
void* calloc(size_t num, size_t size);

/*
  This takes previously-allocated memory and
  attempts to resize it. ==> will free the original memory

  This may require a new block of memory to be
  found, so it returns a new void pointer to memory.

  Contents are preserved

*/
void* realloc(void *ptr, size_t size);

void free(void *ptr);

```

Change the bad example

```

#include <stdio.h>
#include <stdlib.h>

int* createArray(int size){
    int* arr = (int*) malloc(sizeof(int) * 10);

```



```
    return arr;
}

int main(){
    int s = 10;
    int * arr = createArray(s);

    for (int i = 0; i < s; i++){
        arr[i] = i;
    }

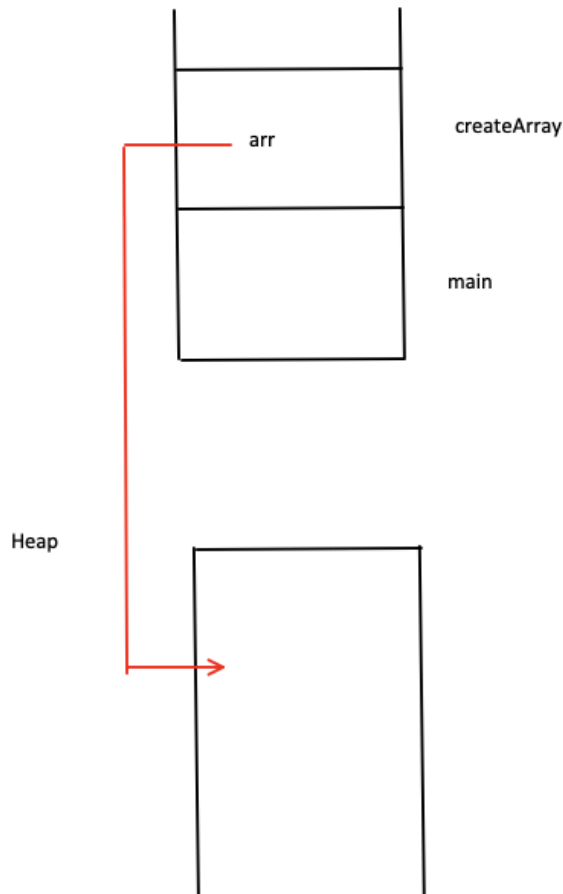
    free(arr);

    arr = NULL;

    return 0;
}
```

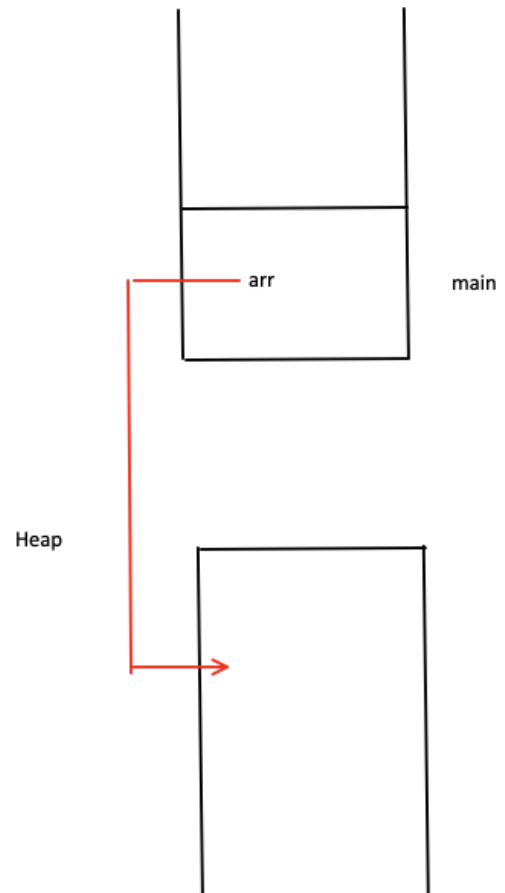
Before createArray return

The memory for stack



After createArray return

The memory for stack



Once malloc remember to free once ! Memory Leak/....

If we have two pointers points to the same memory, only free once.

2. The order for free

3. Compare with Java

Java use keyword `new` for dynamic memory, but we do not need to worry about `free` since In *Java*, process of deallocating memory is handled automatically by the *garbage collector*

==> Tool

Do Question 5

P3: Linked List

Linked list is combined by nodes and we link nodes with pointers

```
#include <stdio.h>
#include <stdlib.h>

typedef struct node node;

struct node{
    node* next;
    int v;
}

node* list_init(int value){
    node* n = malloc(sizeof(struct node));
    // check the return value for malloc
    n->v = value;
    n->next = NULL;
    return n;
}
```

```
}
```

```
void list_add(node* h, int value){  
    if (h != NULL){  
        node* c = h;  
        while (c->next != NULL){  
            c = c -> next;  
        }  
  
        c -> next = list_init(value);  
    }  
}
```

```
void list_delete(node** h, node* n){  
    if (h != NULL){  
        if (*h != NULL){  
            node* prev = NULL;  
            node* cur = *h;  
            // find the node to be deleted and find the prev  
            while (cur != n && cur != NULL){  
                prev = cur;  
                cur = cur -> next;  
            }  
  
            if (cur != NULL){  
                // Two cases  
                // 1. The node to be deleted is the head  
                if (prev == NULL){  
                    /*  
  
                                node0 ----> node1 ----> node2 ----> node3  
prev          cur  
                                deleted  
  
                    */  
                    node* new_head = (*h) -> next;  
                    free(*h);  
                    // Here is the reason why we need node** head  
                    *h = new_head;  
                }  
            }  
        }  
    }  
}
```

```

        }else{

            /*

                node0 ----> node1 ----> node2 ----> node3
                    prev          cur
                        deleted

            */
            // 2. The node to be deleted is not the head
            prev->next = cur->next;
            free(cur);
        }
    }
}

node* list_next(node* n){
    node* r = NULL;
    if (n){
        r = n -> next;
    }
    return r;
}

void list_free(node* h){
    node* t = NULL;
    while (h){
        // record the next firstly
        t = h -> next;
        // free the cur
        free(h);
        h = t;
    }
}

```

```
// // #####
// void changeTheValue(int* a){
//     *a = 2;
// }

// int main(){
//     int a = 1;
//     changeTheValue(&a);
// }
// // #####
// void changeTheHead(node** head){
//     *head = new_head;
// }
// int main(){
//     node* head = ...;
// }
// // #####
```

Tutorial/Week4/Q2, Q3, Q5, Q7

Tool

Check memory leak using gcc flag

```
$ gcc -fsanitize=address -Wall -Werror -std=gnull
```

```
=====
==37==ERROR: LeakSanitizer: detected memory leaks

Direct leak of 40 byte(s) in 1 object(s) allocated from:
#0 0x7f22dc1d3279 in __interceptor_malloc /build/gcc/src/gcc/libsanitizer/asan/asan_malloc_linux.cpp:145
#1 0x556c6f1da17a in main (/home/a.out+0x117a)
#2 0x7f22dbf77b24 in __libc_start_main (/usr/lib/libc.so.6+0x27b24)
SUMMARY: AddressSanitizer: 40 byte(s) leaked in 1 allocation(s).
```

Other error like: heap buffer overflow

GDB

A debugger for C (and C++)

It allows you to do things like run the program up to a certain point then stop and print out the values of certain variables at that point, or step through the program one line at a time and print out the values of each variable after executing each line.

1. Installed ?

```
$ gdb --version
```

2. Compile your program with `-g` or `-g -O0`

3. Some common operations

```
# start up gdb in the directory of executable file
$ cd path/to/a.out
$ gdb

# specify the debug file
(gdb) file [executable_file]

# print the file
(gdb) list

# set break point
(gdb) break [Line]

# run the program
(gdb) run

# print variable
(gdb) print [variable_name]
```

```
# continue to run until the next break point
(gdb) continue

# continue to run next line
(gdb) next

# quit
(gdb) quit
```

Cheat Sheet: <https://darkdust.net/files/GDB%20Cheat%20Sheet.pdf>

Valgrind

`valgrind` supports lots of tools, like `memcheck`, `addrcheck`, `cachegrind`...

If we do not specify which one, `memcheck` is the default.

`memcheck` detects problems with memory management in programs. It checks all read/write operations to memory and intercepts all malloc/new/free/delete calls. So the memcheck tool is able to detect the following problems:

- use uninitialized memory
- Read/write memory that has been freed
- read/write memory out of bounds
- Read/write inappropriate memory stack space
- memory leak
- Using malloc/new/new[] and free/delete/delete[] do not match.
- Overlap of src and dst

1. Installed ?

```
$ valgrind --version
```


2. Compile program with debug option `-g`

3. Use `valgrind`

```
$ valgrind --leak-check=full --show-leak-kinds=all ./a.out
```

- `--leak-check=full`: "each individual leak will be shown in detail"
- `--show-leak-kinds=all`: Show all of "definite, indirect, possible, reachable" leak kinds in the "full" report
 - *"indirectly lost" means **your program is leaking memory in a pointer-based structure**. (E.g. if the root node of a binary tree is "definitely lost", all the children will be "indirectly lost".)*

4. Example

```
#include <stdlib.h>

int main()
{
    char *x = (char*)malloc(20); // line 5

    return 0;
}
```

```
==42== HEAP SUMMARY:
==42==    in use at exit: 20 bytes in 1 blocks
==42==   total heap usage: 1 allocs, 0 frees, 20 bytes allocated
==42==
==42== 20 bytes in 1 blocks are definitely lost in loss record 1 of 1
==42==    at 0x483E899: malloc (in /usr/lib/valgrind/vgpreload_memcheck-amd64-linux.so)
==42==    by 0x10914A: main (memory_leak.c:5)
==42==
==42== LEAK SUMMARY:
==42==    definitely lost: 20 bytes in 1 blocks
==42==    indirectly lost: 0 bytes in 0 blocks
==42==    possibly lost: 0 bytes in 0 blocks
==42==    still reachable: 0 bytes in 0 blocks
==42==    suppressed: 0 bytes in 0 blocks
==42==
==42== For lists of detected and suppressed errors, rerun with: -s
==42== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 0 from 0)
```

