

Proof. Suppose to the contrary, some $\|g\|_2 > M$, $g \in \partial f(x)$

Look at $y = x + g$.

$$\begin{aligned} f(y) &\geq f(x) + g^T(y-x) \\ &= f(x) + \|g\|_2^2 \\ &> f(x) + \|g\|_2 \cdot M \end{aligned}$$

$$\Rightarrow f(x+g) - f(x) > M \|g\|_2 \cdot \cancel{\text{ }} \Rightarrow \text{Bound subgrad. } \square$$

Theorem For any convex f that is M -Lipschitz, then for α_k s.t.

$$\sum_{i=0}^{\infty} \alpha_i = \infty, \quad \sum_{i=0}^{\infty} \alpha_i^2 \leq C$$

we have $\lim_{k \rightarrow \infty} \min_{i \leq k} \{f(x_i) - f(x^*)\} \rightarrow 0$.

(some subsequence converges in obj gap.)
or some $f(x_k) = f(x^*)$

By our Lemma,

$$\text{Proof. } 0 \leq \|x_{k+1} - x^*\|_2^2 \leq \|x_k - x^*\|_2^2 - 2\alpha_k(f(x_k) - f(x^*)) + \alpha_k^2 \|g_k\|_2^2$$

$$\text{(induction)} \leq \|x_0 - x^*\|_2^2 - 2 \sum_{i=0}^k \alpha_i (f(x_i) - f(x^*)) + \sum_{i=0}^k \alpha_i^2 \|g_i\|_2^2$$

$$\Rightarrow \frac{\sum_{i=0}^k \alpha_i (f(x_i) - f(x^*))}{\sum_{i=0}^k \alpha_i} \leq \frac{\|x_0 - x^*\|_2^2 + \sum_{i=0}^k \alpha_i^2 \|g_i\|_2^2}{2 \sum_{i=0}^k \alpha_i}$$

Want $\rightarrow 0$.

$\min_{i \leq k} \{f(x_i) - f(x^*)\} \leq \text{Weighted Average of gaps.}$

$$\min_{i \leq K} \{f(x_i) - f(x^*)\} \leq \frac{\|x_0 - x^*\|^2 + M^2 \sum_{i=0}^K \alpha_i^2}{2 \sum_{i=0}^K \alpha_i} \quad (\text{by 2nd Lemma})$$

$$\rightarrow 0 \quad \text{whenever} \quad \sum_{i=0}^{\infty} \alpha_i = \infty$$

$$\sum_{i=0}^{\infty} \alpha_i^2 \leq C$$

$$(\alpha_k = \frac{1}{k}, \quad C = \frac{\pi^2}{6})$$

Finite-Time Guarantees.

Constant $\alpha_k = \alpha$.

$$\Rightarrow \min_{i \leq K} \{f(x_i) - f(x^*)\} \leq \frac{\|x_0 - x^*\|^2 + M^2 \alpha^2 \cdot K}{2 \alpha K}$$

$$= \frac{\|x_0 - x^*\|^2}{2 \alpha K} + \frac{M^2 \alpha}{2}.$$

Want LHS $\leq \epsilon$, need $\frac{\|x_0 - x^*\|^2}{2 \alpha K} + \frac{M^2 \alpha}{2} = \epsilon$

aim $\frac{\epsilon}{2} + \frac{\epsilon}{2}$

$$\frac{M^2 \|x_0 - x^*\|^2}{2 \epsilon \cdot K} \quad \text{Need } \boxed{\alpha = \frac{\epsilon}{M^2}}$$

For comparison, GD: $f(x_k) - f(x^*) \leq \frac{\epsilon}{K}$

$$\Rightarrow K \geq \frac{\epsilon}{\epsilon}.$$

Accel: $\Rightarrow K = \sqrt{C/\epsilon}.$

$$\Leftrightarrow \epsilon \leq \frac{M \|x_0 - x^*\|}{\sqrt{K}}.$$

Sadly this is as good as possible.

"Theorem" = there exists a bad function $f: \mathbb{R}^d \rightarrow \mathbb{R}$ and subgradient oracle $g(x) \in \partial f(x)$ s.t.

using $k+1 = d$, after k steps

$$f(x_k) - f(x^*) \geq \frac{M \|x_0 - x^*\|}{2(1 + \sqrt{k+1})}.$$

Proof. Nesterov's Convex OPT book.

$$f(x) = \max\{x_1, \dots, x_d\} + \frac{M}{2} \|x\|_2^2.$$

\Rightarrow The subgradient method is optimal
(but still slow).

Extra Directions

Speedup under Strong Convexity: $O(1/k)$

Nonconvex Guarantees [Davis, 2018] $O(1/k^{1/4})$
 \uparrow adviser

Overview So Far

1. Necessary + Sufficient Optimality Cond
2. First-Order Methods

Smooth OPT

(linesearchs,

GD,

Accel)

Equiv Characterizations for Smooth, Convex,
Strongly, Convex.

Proximal First-Order Methods

(Mirrors smooth setting)

Nonsmooth OPT more generally

Subgradients

Subgrad Method is optimal.

Midterm

Stochastic OPT

3. Second-Order Methods

Newton

"Quasi-Newton"