

Lemma: Any algorithm satisfying assumption 2

$$x_{i+1} \in \text{Lin} \{ \nabla f_k(x_0), \dots, \nabla f_k(x_{i-1}) \}$$

has  $\text{Lin} \{ \nabla f_k(x_0), \dots, \nabla f_k(x_i) \} = \mathbb{R}^{i+1} \times \{0\}^{d-i-1}$   
for  $i \leq k$ .

Proof. Trivially for  $i=0$ ,  $\nabla f_k(x_0) = A_k x_0 - e_1$   
 $= -e_1$ .

Inductively, for  $i \geq 0$ ,  $\nabla f_k(x_i) = A_k \cdot x_{i-1} - e_1$

$$\in A_k \cdot \text{Lin} \{ \nabla f(x_0), \dots, \nabla f(x_{i-1}) \}$$

$$= A_k \cdot \mathbb{R}^i \times \{0\}^{d-i}$$

$$\text{(since } A_k \text{ is tri-diagonal)} = \mathbb{R}^{i+1} \times \{0\}^{d-i-1} \quad \square$$

Check  $f_k$  is convex: for  $s \in \mathbb{R}^d$

$$s^T \nabla^2 f_k(x) s = \frac{1}{4} s^T A_k s = \frac{1}{4} \left[ (s^{(1)})^2 + \sum_{i=1}^{k-1} (s^{(i)} - s^{(i+1)})^2 + (s^{(k)})^2 \right] \geq 0. \quad \checkmark$$

Check  $f_k$  is  $L$ -smooth: for  $s \in \mathbb{R}^d$

$$\begin{aligned} s^T \nabla^2 f_k(x) s &= \dots = \frac{1}{4} \left[ (s^{(1)})^2 + \sum (s^{(i+1)} - s^{(i)})^2 + (s^{(k)})^2 \right] \\ &\stackrel{(a-b)^2 \leq 2a^2 + 2b^2}{\leq} \frac{1}{4} \left[ (s^{(1)})^2 + \sum [2(s^{(i)})^2 + 2(s^{(i+1)})^2] + (s^{(k)})^2 \right] \\ &\leq \frac{1}{4} \sum 4(s^{(i)})^2 = L \|s\|_2^2. \quad \checkmark \end{aligned}$$

How poorly are solutions with only the first  $i$  coordinates nonzero?

Claim:  $\nabla f_k(x) = 0 = \frac{L}{4} (A_k x - e_1)$  has ~~unique~~ sol. ← only in first  $k$  variables

$$\bar{x}_k^{(1)} = \begin{cases} 1 - \frac{1}{k+1} & \text{for } i = 1, \dots, k \\ 0 & \text{otherwise.} \end{cases}$$

$\Rightarrow$  Optimal obj of  $f_k(\cdot)$  comes from plugging in

$$\begin{aligned} f_k^* &:= f_k(\bar{x}_k) \\ &= \frac{L}{4} \left( \frac{1}{2} \bar{x}_k^T A_k \bar{x}_k - e_1^T \bar{x}_k \right) \\ &= \frac{L}{4} \left( \frac{1}{2} \left( \left(1 - \frac{1}{k+1}\right)^2 + (k-1) \left(\frac{1}{k+1}\right)^2 + \left(1 - \frac{k}{k+1}\right)^2 - \left(1 - \frac{1}{k+1}\right) \right) \right) \\ &= \frac{L}{8} \left( -1 + \frac{1}{k+1} \right) \end{aligned}$$

$$\text{Likewise, } \|x_0 - \bar{x}_k\|_2^2 = \sum_{i=1}^k \left(1 - \frac{i}{k+1}\right)^2 \leq \frac{1}{3} (k-1)$$

(see Nesterov's book for gory details)

Proof of Complexity Lower Bound. Consider fixed  $L > 0, k > 0$ .

Then consider  $f(x) = f_{2k+1}(x)$ ,  $d = 2k+1$ .

$$f(x_k) = f_{2k+1}(x_k) = f_k(x_k) \geq f_k^*$$

$$\begin{aligned} \Rightarrow \frac{f(x_k) - f(x^*)}{\|x_0 - x^*\|_2^2} &\geq \frac{f_k^* - f_{2k+1}^*}{\|x_0 - \bar{x}_{2k+1}\|_2^2} \\ &\geq \frac{\frac{L}{8} \left(1 + \frac{1}{k+1} + 1 - \frac{1}{2k+2}\right)}{\frac{1}{3}(2k+1)} \\ &\geq \frac{3L}{32(k+1)^2} \quad \checkmark \end{aligned}$$

The second part of our theorem holds as

$$\begin{aligned} \|x_k - x^*\|_2^2 &\underset{\substack{\parallel \\ \bar{x}_{2k+1}}}{\geq} \sum_{i=k+1}^{2k+1} (\bar{x}_{2k+1}^{(i)})^2 = \sum_{i=k+1}^{2k+1} \left(1 - \frac{i}{2k+2}\right)^2 \\ &\quad \text{(if } \frac{3}{2}k \text{ is an integer)} \geq \sum_{i=k+1}^{\frac{3}{2}k+1} \left(1 - \frac{i}{2k+2}\right)^2 \\ &\geq \left(\frac{1}{2}k\right) \left(1 - \frac{\frac{3}{2}k+1}{2k+2}\right)^2 \\ &\geq \left(\frac{1}{2}k\right) \left(\frac{1}{4}\right)^2 \\ &\geq \frac{1}{32} k \\ &\quad \left(\text{since } \|x_0 - x_{2k+1}\|_2^2 \geq \frac{1}{3}(2k+1)\right) \geq \frac{3}{64} \|x_0 - x^*\|^2. \quad \square \end{aligned}$$



## Recap of First-Order Smooth Optimization Results

For any function  $f$  with  $L$ -Lipschitz gradient, we found the following guarantees:

Method	Generic Rate	Speed-up from $\geq 0$ .
Gradient Descent (for nonconvex $f$ )	$\frac{1}{T} \sum_{k=0}^{T-1} \ \nabla f(x_k)\ _2^2 \leq O\left(\frac{1}{T}\right)$	Linear Rate Under Strict Positive Definite $\nabla^2 f$
Gradient Descent (for convex $f$ )	$f(x_T) - f(x^*) \leq O\left(\frac{1}{T}\right)$	Linear Rate Under Strong Convexity of $f$
Accelerated Gradient (for convex $f$ )	$f(x_T) - f(x^*) \leq O\left(\frac{1}{T^2}\right)$	Faster Linear Rate Under Restart with Strong Convexity.

↑  
This rate is optimal.  
(No 1<sup>st</sup> order method can  
improve this worstcase)

↑  
Proven in HW2, Q3.

## Structured Constrained/Nonsmooth Optimization

1. Define Proximal Operator
2. Example Orthogonal Projection on Constraints
3. Projected / Proximal Gradient Methods
4. Acceleration
5. More proximal methods (Alternating projections, ADMM).

# 1. Proximal Operator

We want a local improvement without on the existence of a gradient.

$$\text{prox}_f(\bar{x}) := \operatorname{argmin}_{x \in \mathbb{R}^d} \left\{ f(x) + \frac{1}{2} \|x - \bar{x}\|_2^2 \right\}$$

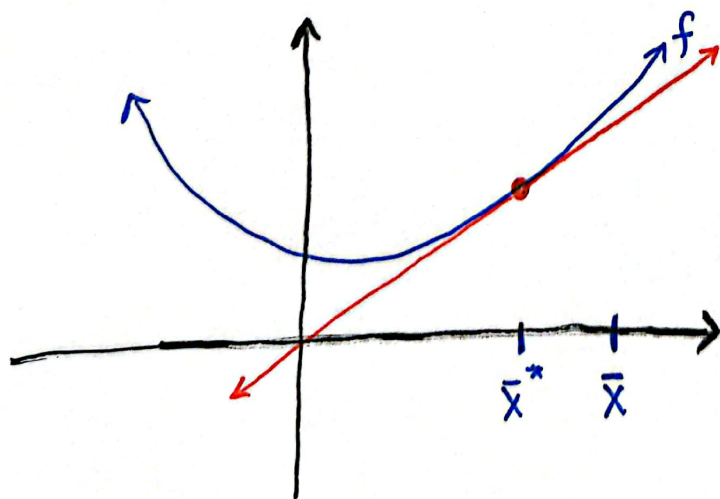
To allow a stepsize, we can rescale by  $\alpha > 0$

$$\begin{aligned} \text{prox}_{\alpha \cdot f}(\bar{x}) &:= \operatorname{argmin} \left\{ \alpha f(x) + \frac{1}{2} \|x - \bar{x}\|_2^2 \right\} \\ &= \operatorname{argmin} \left\{ \underbrace{f(x) + \frac{1}{2\alpha} \|x - \bar{x}\|_2^2}_{\text{looks like HW2, Q2}} \right\} \end{aligned}$$

looks like HW2, Q2.

In particular,  $\bar{x}^* = \text{prox}_{\alpha f}(\bar{x})$ , then  $g = \frac{1}{\alpha}(\bar{x} - \bar{x}^*)$  gives a linearization

$$f(x) \geq f(\bar{x}^*) + g^T(x - \bar{x}^*).$$



"g is a 'subgradient' of f at  $\bar{x}^*$ "

$$\partial f(x) = \left\{ g \in \mathbb{R}^d \mid f(y) \geq f(x) + g^T(y - x) \right. \\ \left. \forall y \in \mathbb{R}^d \right\}$$

"Intro to Convexity" covers these in more depth.