

# Algorithm Design and Analysis - Assignment 4

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## Problem1

**Analysis:** Assume  $x_{ij} = 1$  means j-th job is assigned to i-th worker. Vice versa.

then we have:

$$\begin{aligned} \min \max_i \sum_{j=1}^n C_{ij}x_{ij} \\ s. t. \sum_{i=1}^m x_{ij} = 1 \text{ for all } j \\ x_{ij} = 0/1 \text{ for all } i \text{ and } j \end{aligned}$$

## Problem2

**Analysis:** Assume  $x_i = 1$  means i-th food is bought. Vice versa. At the same time, we say the nutrition is satisfied when it is greater than 0.

then we have:

$$\begin{aligned} \min \sum_{i=1}^m C_i x_i \\ s. t. \sum_{i=1}^m x_i W_{ij} > 0 \text{ for all } j \\ x_i = 0/1 \text{ for all } i \end{aligned}$$

## Problem3

**Analysis:** Assume  $x_{ij}$  is the amount of A of i-th place that is delivered to j-th market. Assume that the balance is achieved if and only if all A delivered is below the sum-output and all A accepted is below the sum-demand.

then we have:

$$\begin{aligned} \min \sum_{i=1}^n \sum_{j=1}^m C_{ij}x_{ij} \\ s. t. \sum_{j=1}^m x_{ij} \leq a_i \text{ for all } i \end{aligned}$$

$$\sum_{i=1}^n x_{ij} \leq b_i \text{ for all } j$$

$$x_{ij} \geq 0 \text{ for all } i \text{ and } j$$

## Problem4

**Analysis:** Assume  $x_{ij} = 1$  means  $i$ -th stuff is packed into  $j$ -th box. Vice versa.

then we have:

$$\min \sum_{j=1}^n \sum_{i=1}^m x_{ij}$$

$$s. t. \sum_{j=1}^n x_{ij} \leq 1 \text{ for all } i$$

$$\sum_{i=1}^n x_{ij} C_i \leq S_j \text{ for all } j$$

$$x_{ij} = 0/1 \text{ for all } i \text{ and } j$$

## Problem6

**Analysis:** Assume  $x_i = 1$  means  $i$ -th project is fund. Vice versa.

then we have:

$$\max \sum_{i=1}^n (c_i - b_i) x_i$$

$$s. t. \sum_{i=1}^n x_i b_i \leq B$$

$$x_i = 0/1 \text{ for all } i$$

## Problem 7

**Analysis:** Assume we build  $x$  A-dorms and  $y$  B-dorms. Assume that profit is sum of  $c_i$ .

then we have:

$$\max x c_i + y c_j$$

$$s. t. x \leq n_i$$

$$y \leq n_j$$

$$x s_i + y s_j \leq s$$

$$x, y \geq 0$$

## Problem 8

**Analysis:** Assume  $x_i = 1$  means place  $A_i$  is chosen. Vice versa.

then we have:

$$\max \sum_{i=1}^7 (c_i - b_i)x_i$$

$$s. t. \sum_{i=1}^7 x_i b_i \leq B$$

$$x_1 + x_2 + x_3 \leq 2$$

$$x_4 + x_5 \geq 1$$

$$x_6 + x_7 \leq 1$$

$$x_i = 0/1 \text{ for all } i$$

## Problem 9

**Analysis:** Assume  $x_i = 1$  means i-th item is packed. Vice versa. Assume there are n items in total.

then we have:

$$\max \sum_{i=1}^n c_i x_i$$

$$s. t. \sum_{i=1}^n x_i a_i \leq b$$

$$x_i = 0/1 \text{ for all } i$$

## Problem 10

**Analysis:** Assume  $x_{ij} = 1$  means j-th job is assigned to i-th worker. Vice versa. Assume that num of jobs is less than num of workers, and there are m workers and n jobs. So  $m \leq n$ .

then we have:

$$\max \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij}$$

$$s. t. \sum_{i=1}^m x_{ij} = 1 \text{ for all } j$$

$$\sum_{j=1}^n x_{ij} = 1 \text{ for all } i$$

$$x_{ij} = 0/1 \text{ for all } i \text{ and } j$$