## CS 710: Complexity Theory Lecture 5: Random Walk Instructor: Jin-Yi Cai Scribe: Jialu Bao

## 1 Random Walks

A random walk in graph G = (V, E) starts with a vertex  $s \in V$  at time step 0. At each time step i, the random walk chooses a neighbor v of the current vertex u with some probability  $P_{u,v}$  and jumps to v at the next time step. For simplicity, we work with undirected simple graphs—simple meaning there is no multi-edge and no loop—and assume that the walk picks neighbors with equal probability.

Formally, let  $y_i$  be the vertex that the random walk visits at time step i.

$$P_{u,v} := \Pr (y_{i+1} = v \mid y_i = u)$$

$$= \begin{cases} \frac{1}{deg(u)} & \text{if } (u,v) \in E \\ 0 & \text{Otherwise} \end{cases}$$

## 2 One dimensional Random Walk

Consider the following one dimensional random walk on a chain of n vertices starting from node 1 and will getting absorbed at node n—the random walk stops immediately after it hits node n.

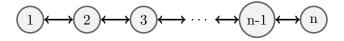


Figure 1: One-dimensional Walk

**Definition 1.** Define **hitting time** of a random walk as the expected number of steps it takes from the starting state to the absorbing state.

**Definition 2.** Define  $h_{i,j} := \mathbb{E}[number\ of\ steps\ the\ random\ walk\ takes\ to\ go\ from\ node\ i\ to\ node\ j$  in Figure 1

Then the hitting time of random walks in Figure 1 equals  $h_{1,n}$ .

Note that although the number of steps between any two states must be integral, the expected value of that number does not have to be integral.

**Lemma 1.** For the random walk in Figure 1, for any i, j, k such that  $1 \le j \le i \le k \le n$ ,

$$h_{i,k} = h_{i,i} + h_{i,k}$$

*Proof.* By the topology of Figure 1, every path going from node j to node k goes through node i. As random walks are memoryless, the number of steps in every path going from node j to node k

is the sum of the number of steps it takes from node j to i and the number of steps it takes from node i to k. Therefore,

$$h_{j,k} = E[\# \text{ of steps to go from } j \text{ to } k]$$
  
=  $E[\# \text{ of steps to go from } j \text{ to } i + \# \text{ of steps to go from } i \text{ to } k]$   
=  $E[\# \text{ of steps to go from } j \text{ to } i] + E[\# \text{ of steps to go from } i \text{ to } k]$   
=  $h_{j,i} + h_{i,k}$ 

The second to the last step follows from the linearity of expectation, and the last step follows from the definition of  $h_{i,i}$  and  $h_{i,k}$ .

Corollary 1. The hitting time,

$$h_{1,n} = h_{1,2} + h_{2,n}$$

$$= h_{1,2} + h_{2,3} + h_{3,n}$$

$$= \dots$$

$$= \sum_{i=1}^{i=n-1} h_{i,i+1}$$

**Corollary 2.** For the random walk in Figure 1, for any i such that  $2 \le i \le n-1$ ,

$$h_{i-1,i+1} = h_{i-1,i} + h_{i,i+1}$$

**Lemma 2.** For  $2 \le i \le n-1$ ,  $h_{i,i+1} = 1 + h_{i-1,i+1}$ , and  $h_{12} = 1$ .

*Proof.* Since we can only go node 2 from node 1, there is only one path from node 1 to 2 and that path has 1 step, so  $h_{1,2} = 1$ .

From node i such that  $2 \le i \le n-1$ , the random walk can either goes to i-1 or i+1 for the next step, each with probability  $\frac{1}{2}$ . If it goes directly to i+1, then it reaches i+1 in 1 step. If it goes to i-1, then in expectation it takes another  $h_{i-1,i+1}$  steps after that steps to reach node i+1. Thus,

# of steps to go from i to  $i+1=\frac{1}{2}\cdot 1+\frac{1}{2}\cdot (1+\# \text{ of steps to go from } i-1 \text{ to } i+1)$ 

Thus, by substituting it into the definition of  $h_{i,j}$ , we have

$$\begin{split} h_{i,i+1} &= \mathrm{E}[ \text{\# of steps to go from } i \text{ to } i+1] \\ &= \mathrm{E}[\frac{1}{2} \cdot 1 + \frac{1}{2} \cdot (1 + \text{\# of steps to go from } i-1 \text{ to } i+1)] \\ &= \mathrm{E}[1 + \frac{1}{2} \text{\# of steps to go from } i-1 \text{ to } i+1)] \\ &= 1 + \frac{1}{2} \mathrm{E}[ \text{\# of steps to go from } i-1 \text{ to } i+1)] \\ &= 1 + \frac{1}{2} h_{i-1,i+1} \end{split}$$

**Theorem 1.** The hitting time,  $h_{1,n}$ , is  $(n-1)^2$ .

*Proof.* Substitute  $h_{i-1,i+1} = h_{i-1,i} + h_{i,i+1}$  from lemma 1 into  $h_{i,i+1} = 1 + \frac{1}{2}h_{i-1,i+1}$ , then we have,

$$h_{i,i+1} = 1 + \frac{1}{2}(h_{i-1,i} + h_{i,i+1}) \tag{1}$$

$$\implies h_{i,i+1} = 2 + h_{i-1,i} \tag{2}$$

Since  $h_{1,2} = 1$ , the recurrence solves to  $h_{i,i+1} = 2i - 1$ . Substitute it into the formula given by Corollary 1, the hitting time equals

$$h_{1,n} = \sum_{i=1}^{n-1} h_{i,i+1} = (n-1)^2$$

## 3 Expected number of times a walk stops visits a node

Now we consider a different problem of random walks, given a undirected graph G = (V, E) and perform the random walk as described in Section 1, then what is the expected number of times a walk visits a node i? In the following, we use  $t_i$  to denote that number and will try to calculate it.

Because node n is the absorbing node and no other node is absorbing state, the walk will eventually visits n and never steps out after that. Thus, without any computations, we can conclude that  $t_n=1$ . Then we consider node n-1, every time the walk reaches node n-1, with probability  $\frac{1}{2}$ , it goes to node n, gets absorbed and never comes back; with probability  $\frac{1}{2}$ , it goes to node n-1, then either directly or after an indefinite number of steps, it comes back and stopst node n-1 again, and repeat the process. Thus, the walk visits a node for only once with probability  $\frac{1}{2}$ , visits a node twice with probability  $\frac{1}{4}$ , visits a node three times with probability  $\frac{1}{8}$ ... That is

$$t_{n-1} = \sum_{k \ge 1} k \cdot \Pr(\text{the walk visits node } n - 1 \text{ for } k \text{ times})$$

$$= \frac{1}{2} \sum_{k \ge 1} k \cdot (\frac{1}{2})^{k-1} = \frac{1}{2} \left[ \sum_{k \ge 1} (x^k) \right]'_{x = \frac{1}{2}}$$

$$= \frac{1}{2} \left[ \frac{x}{1-x} \right]'_{x = \frac{1}{2}} = \frac{1}{2} \left[ \frac{1}{(1-x)^2} \right]_{x = \frac{1}{2}}$$

$$= \frac{1}{2} \left[ \frac{1}{(1-\frac{1}{2})^2} \right] = 2$$

Alternatively, because the random walk is memoryless, every time it visits node n-1, the expected number of upcoming visits to node n-1 again stays the same. Thus, if the walk visits node n-1 once and then goes to node n-2 at the next steps, the expected number of times of additional visits to node n-1 is still  $t_{n-1}$ . Therefore, we have the recurrence  $t_{n-1} = \frac{1}{2} \times 1 + \frac{1}{2} \times (1 + t_{n-1})$ , which also gives  $t_{n-1} = 2$ .

How about  $t_i$  for the rest of nodes?

**Theorem 2.** For  $3 \le i \le n-2$ ,  $t_i = \frac{1}{2}(t_{i-1} + t_{i+1})$ .

*Proof.* We will first present the intuition and then a more rigorous proof.

Because our graph is just a single chain of nodes, for any node i such that  $3 \le i \le n-2$ , the walk must passes either node i-1 or node i+1 en route to other node. Meanwhile, every time the walk visits node i-1, it has  $\frac{1}{2}$  chance of going to node i the next step; and similarly, every time the walk visits node i+1, it has  $\frac{1}{2}$  chance of visiting node i the next step. Thus, the expected number of times that the walk visits node i equals  $\frac{1}{2}(t_{i-1}+t_{i+1})$ .

To prove it formally, we define random variable W to be the number of steps a random walk takes from node 1 to node n, i.e., the length of time that a random walk is alive. We also define 0-1 random variables  $x_{j,i}$  for  $i \geq 1, j \geq 0$ ,

$$x_{j,i} = \begin{cases} 1 & \text{if } W \ge j \text{ and } y_j = i \\ 0 & \text{Otherwise} \end{cases}$$

 $W \geq j$  indicates that a random walk is alive at time j, which guarantees that  $y_j$  is defined.  $y_j$  is the node the random walk visited at the  $j^{th}$  step. So by definition,  $x_{j,i} = 1$  if and only if the walk has not yet ended at the  $j^{th}$  step and the walk is visiting node i at  $j^{th}$  step. Thus,

$$t_i = E[\text{the number of visits to node } i] = E[\sum_{j>0} x_{j,i}]$$

Then, by linearity of expectation and the defintion of expected value,

$$t_i = \sum_{j \ge 0} E[x_{j,i}]$$
$$= \sum_{j \ge 0} \Pr[W \ge j \text{ and } y_j = i]$$

Since the walk starts at node 1 and  $i \ge 3$ , there is no chance to visit node i in the  $0^{th}$  and  $1^{st}$  step, so

$$\sum_{j \ge 0} \Pr[W \ge j \text{ and } y_j = i] = \sum_{j \ge 1} \Pr[W \ge j \text{ and } y_j = i] = \sum_{j \ge 2} \Pr[W \ge j \text{ and } y_j = i]$$
 (3)

Thus,

$$\begin{split} t_i &= \sum_{j \geq 0} \Pr[W \geq j \text{ and } y_j = i] = \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_j = i] \\ &= \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i - 1 \text{ and } y_j = i] + \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i + 1 \text{ and } y_j = i] \\ &= \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i - 1] \cdot \Pr[y_j = i \mid y_{j-1} = i - 1] + \\ &\qquad \qquad \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i + 1] \cdot \Pr[y_j \mid y_{j+1}] \\ &= \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i - 1] \cdot \frac{1}{2} + \sum_{j \geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i + 1] \cdot \frac{1}{2} \end{split}$$

Since  $i \le n-2$ , the random walk had to continue at least one more steps if it visits node i+1 or node i-1 at  $(j-1)^{th}$  step, so

$$\Pr[W \ge j - 1 \text{ and } y_{j-1} = i+1] = \Pr[W \ge j \text{ and } y_{j-1} = i+1]$$
(4)

$$\Pr[W \ge j - 1 \text{ and } y_{j-1} = i - 1] = \Pr[W \ge j \text{ and } y_{j-1} = i - 1]$$
 (5)

Thus,

$$t_{i} = \sum_{j\geq 1} \Pr[W \geq j \text{ and } y_{j-1} = i-1] \cdot \frac{1}{2} + \sum_{j\geq i} \Pr[W \geq j+1 \text{ and } y_{j+1} = i+1] \cdot \frac{1}{2}$$

$$= \sum_{j\geq 1} \Pr[W \geq j-1 \text{ and } y_{j-1} = i-1] \cdot \frac{1}{2} + \sum_{j\geq 1} \Pr[W \geq j-1 \text{ and } y_{j-1} = i+1] \cdot \frac{1}{2} \text{ (By 4,5)}$$

$$= \sum_{j-1\geq 1} \Pr[W \geq j-1 \text{ and } y_{j-1} = i-1] \cdot \frac{1}{2} + \sum_{j-1\geq 1} \Pr[W \geq j-1 \text{ and } y_{j-1} = i+1] \cdot \frac{1}{2} \text{ (By 3)}$$

$$= t_{i-1} \cdot \frac{1}{2} + t_{i+1} \cdot \frac{1}{2}$$

Similarly,  $t_1 = \frac{1}{2}t_2 + 1$  because every walk pays one visit to node 1 at the beginning, and after that, it has to go through node 2 once every time it gets to node 1. For node 2,  $t_2 = \frac{1}{2}t_3 + t_1$  because every visit to node 2 must come from either node 1 or node 3 at the last step, and from node 1 it always goes to node 2 while from node 3 it goes to node 2 with probability  $\frac{1}{2}$ . Substituting  $t_1 = \frac{1}{2}t_2 + 1$  into  $t_2 = \frac{1}{2}t_3 + t_1$ , we get  $t_2 = t_3 + 2$ .

For  $3 \le i \le n-2$ ,  $t_i = \frac{1}{2}[t_{i-1} + t_{i+1}]$  implies that  $t_{i+1} - t_i = t_i - t_{i-1}$ . Thus,  $t_2 - t_3 = t_3 - t_4 = \dots = t_{n-2} - t_{n-1}$ . Since  $t_2 = t_3 + 2$ ,

$$t_2 - t_3 = t_3 - t_4 = \dots = t_{n-2} - t_{n-1} = 2$$

Since  $t_{n-1}=2$ , we can solve that  $t_i=2(n-i)$  for  $2 \le i \le n-1$ . In particular,  $t_2=2(n-2)$ , implying that  $t_1=\frac{1}{2}t_2+1=n-1$ .

We can do a sanity check of our calculated result, the sum of  $t_i$  for all i should be exactly one more than the hitting time.

$$\sum_{i=1}^{n} t_i = t_1 + \sum_{i=2}^{n-1} t_i + t_{n-1} + t_n$$

$$= (n-1) + \left(\sum_{i=2}^{n-2} 2(n-i)\right) + 2 + 1$$

$$= (n-1) + (n^2 - 3n) + 2 + 1$$

$$= n^2 - 2n + 2$$

$$= (n-1)^2 + 1$$

Thus, the sum of our calculated  $t_i$  is exactly the hiting time plus one.