



# STONE DUALITY FOR COMPLETENESS OF BOOLEAN ALGEBRA

---

*Nov 12th ,2019 Jialu Bao*



# DUALITY FOR COMPLETENESS OF

---

Nov 12th ,2019 Jialu Bao  
Based on works by Simon Docherty

# HIGH-LEVEL

---

# (KRIPKE) FRAME

---

- It's from Kripke semantics, which are first conceived for modal logic and then adapted to other logics.
- In our context, a frame is a triple  $\mathcal{K} = (X, \geq)$  where
  - $X$  is a set of worlds
  - $\geq$  is a preorder (reflexive and transitive) on  $X$
- A frame and a valuation determining whether a world satisfies a formula —> a model

# (KRIPKE) FRAME

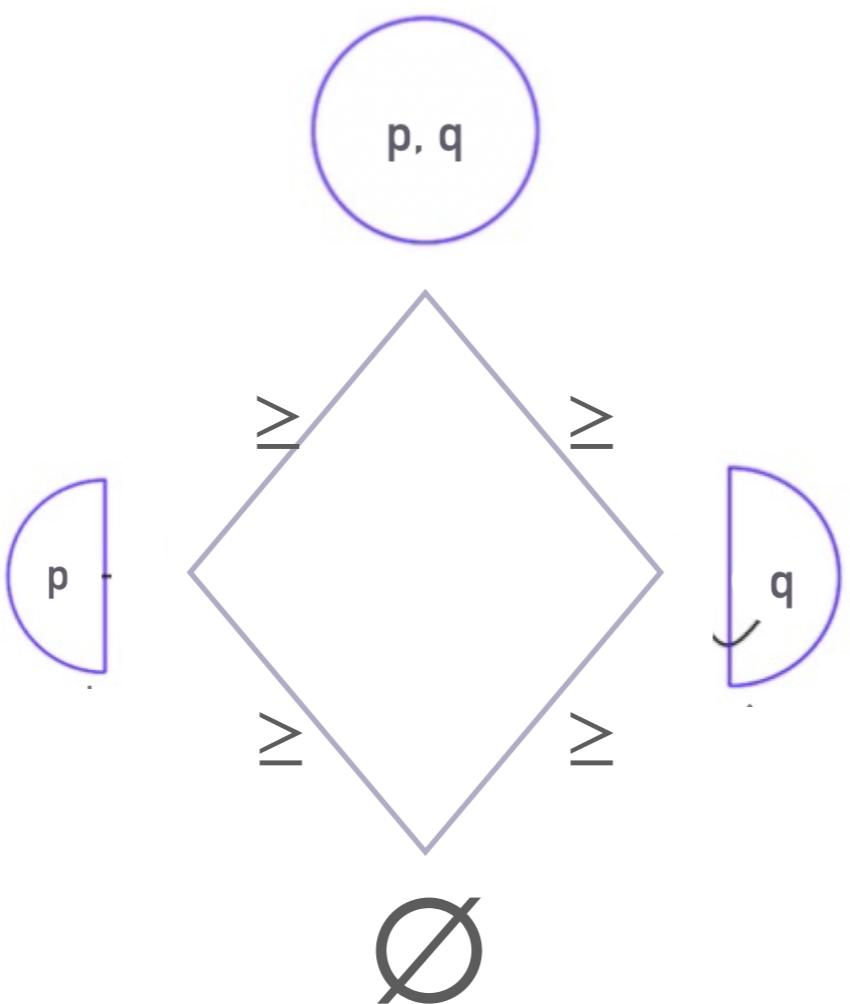
---

- It's from Kripke semantics, which are first conceived for modal logic and then adapted to other logics.
- More generally, a frame is a triple  $\mathcal{K} = (X, \geq, \circ)$  where
  - $X$  is a set of worlds
  - $\geq$  is a preorder (reflexive and transitive) on  $X$
  - For bunch implication  $\circ : X \times X \rightarrow 2^X$  is a binary operation*
  - For modal logic,  $\circ$  is a binary relation  $X \times X$*
- A frame and a valuation determining whether a world satisfies a formula —> a model

# EXAMPLE. (INFORMAL)

---

- We can define semantics in a frame
- eg.  $p \wedge q$



# BOOLEAN ALGEBRA

---

- Algebra: a set with some operators, satisfying some equations
- A Boolean Algebra is an algebra  $\mathbb{A} = (A, \wedge, \vee, \rightarrow, \top, \perp)$  such that  $(A, \wedge, \vee, \top, \perp)$  is a bounded distributive lattice and  $\rightarrow$  is a binary operation satisfying for all  $a, b, c$

$$a \vee (a \rightarrow \perp)$$

$$a \wedge b \leq c \text{ iff } a \leq b \rightarrow c$$

# BOOLEAN ALGEBRA

---

- A Boolean Algebra is an algebra  $\mathbb{A} = (A, \wedge, \vee, \rightarrow, \top, \perp)$  on some partially ordered set  $A$  such that

$\wedge, \vee$  are commutative and associative binary operators

$\wedge$  Absorption :  $a \vee (a \wedge b) = a$

$\vee$  Absorbtion :  $a \wedge (a \vee b) = a$

Distributive :  $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c), a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$

$\top \neq \perp$

$\top$  is the identity of  $\wedge$  :  $\top \wedge a = a$

$\perp$  is the identity of  $\vee$  :  $\perp \vee a = a$

$a \vee (a \rightarrow \perp)$

$a \wedge b \leq c$  iff  $a \leq b \rightarrow c$

Distributive  
Lattice

Bounded  
Distributive  
Lattice

Exclude the middle

# EXAMPLES:

---

## ► Two-Element Boolean Algebra

$\mathbb{A} = (\{0,1\}, \wedge, \vee, \rightarrow, 1, 0)$  where  $\vee, \wedge$  are defined by

$\wedge$	0	1
0	0	0
1	0	1

$\vee$	0	1
0	0	1
1	1	1

## ► Power set $\mathbb{A} = (2^X, \cap, \cup, \rightarrow, X, \emptyset)$

# STONE REPRESENTATION THEOREM

---

- Two-Element Boolean Algebra

$\mathbb{A} = (\{0,1\}, \wedge, \vee, \rightarrow, 1, 0)$  where  $\vee, \wedge$  are defined by

$\wedge$	0	1
0	0	0
1	0	1

$\vee$	0	1
0	0	1
1	1	1

- Power set  $\mathbb{A} = (2^X, \cap, \cup, \rightarrow, X, \emptyset)$
- Every Boolean algebra is isomorphic to a sub algebra of a power set algebra.

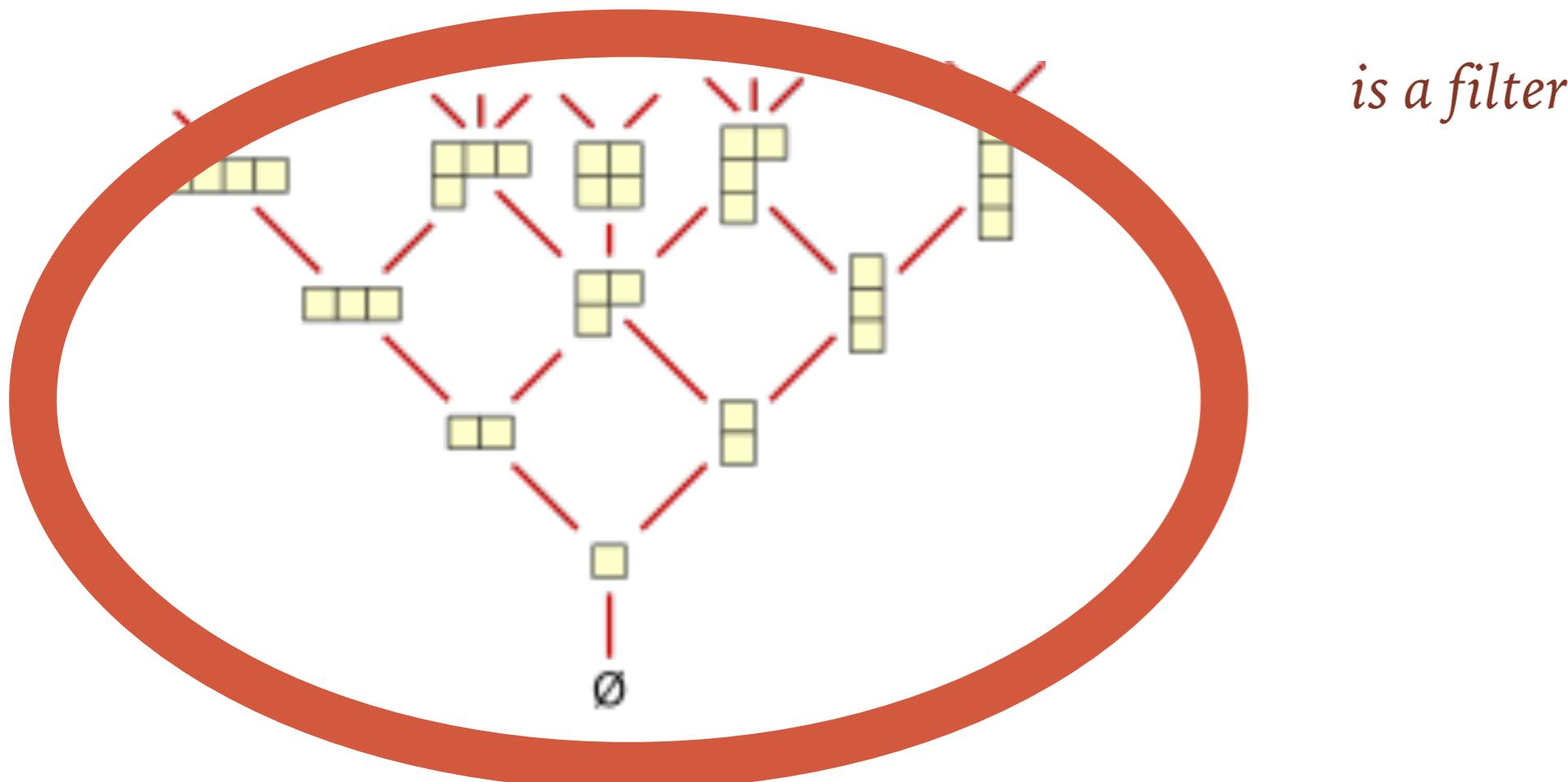
# FILTERS

---

- A filter on a distributive lattice  $A$  is a non-empty subset  $F$  of  $A$  satisfying

*Upwards Closure:  $a \in F$  and  $b \geq a$  implies  $b \in F$*

*Meet Closure:  $a, b \in F$  implies  $a \wedge b \in F$*



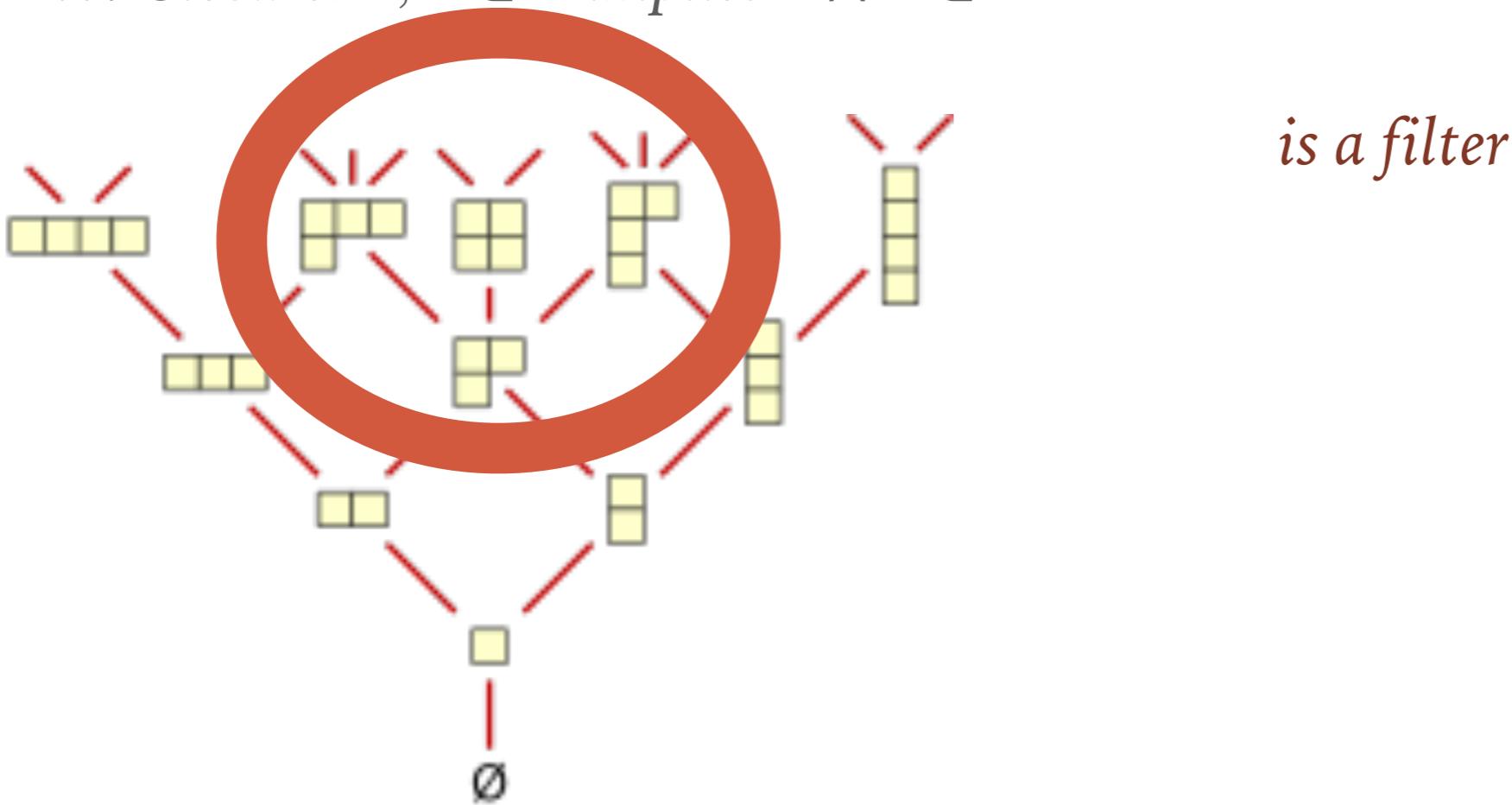
# FILTERS

---

- A filter on a distributive lattice  $A$  is a non-empty subset  $F$  of  $A$  satisfying

*Upwards Closure:  $a \in F$  and  $b \geq a$  implies  $b \in F$*

*Meet Closure:  $a, b \in F$  implies  $a \wedge b \in F$*



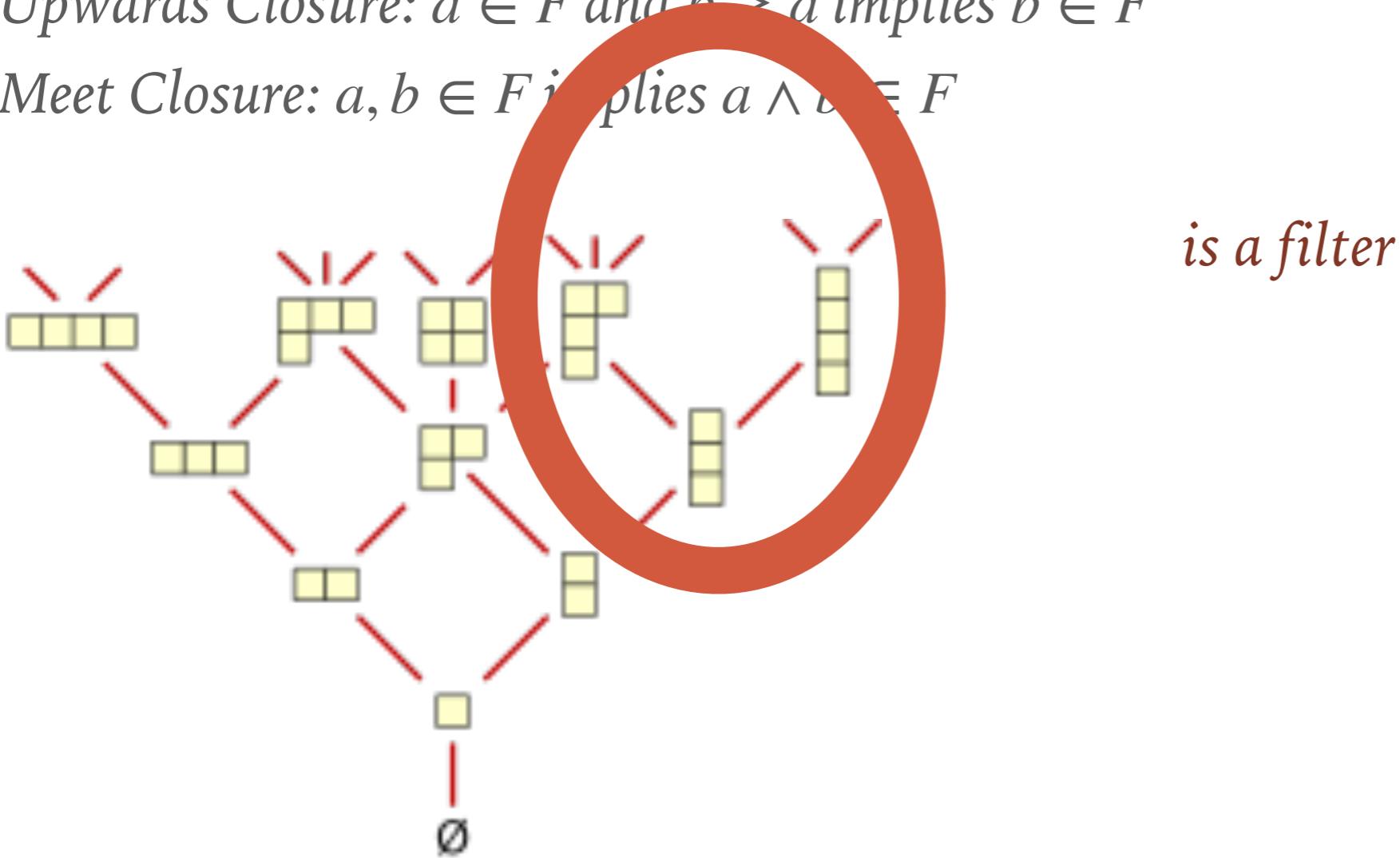
# FILTERS

---

- A filter on a distributive lattice  $A$  is a non-empty subset  $F$  of  $A$  satisfying

*Upwards Closure:  $a \in F$  and  $b > a$  implies  $b \in F$*

*Meet Closure:  $a, b \in F$  implies  $a \wedge b \in F$*



*is a filter*

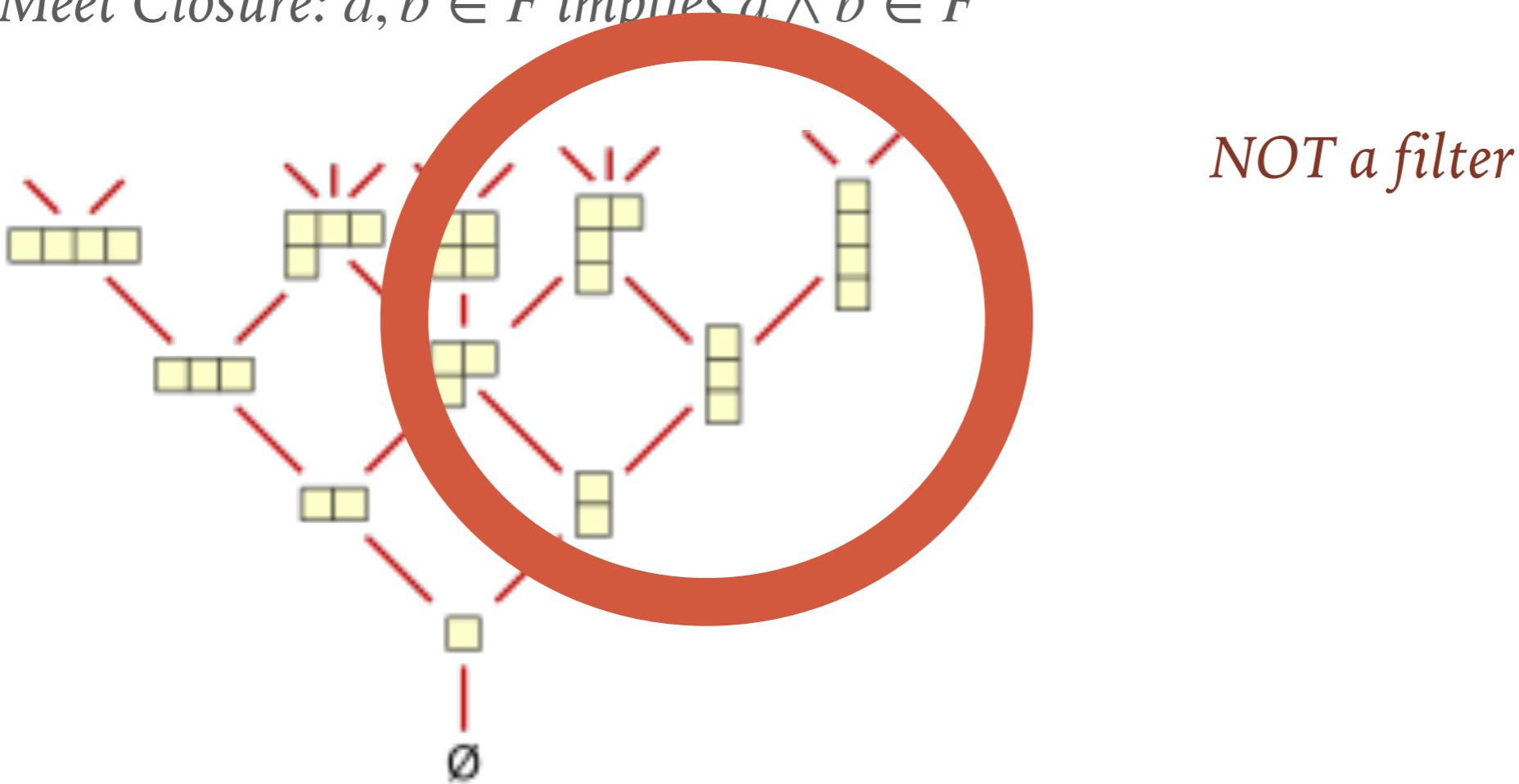
# FILTERS

---

- A filter on a distributive lattice  $A$  is a non-empty subset  $F$  of  $A$  satisfying

*Upwards Closure:  $a \in F$  and  $b \geq a$  implies  $b \in F$*

*Meet Closure:  $a, b \in F$  implies  $a \wedge b \in F$*



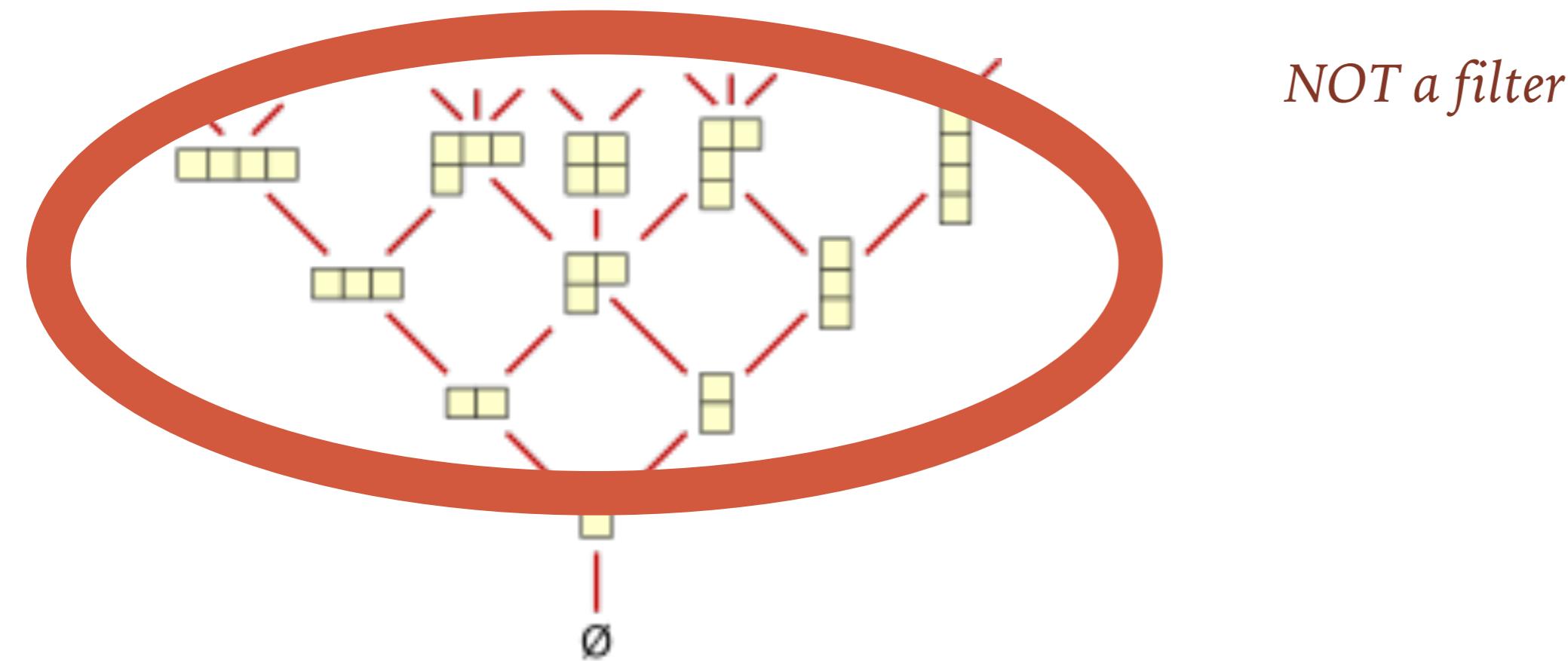
# FILTERS

---

- A filter on a distributive lattice  $A$  is a non-empty subset  $F$  of  $A$  satisfying

*Upwards Closure:  $a \in F$  and  $b \geq a$  implies  $b \in F$*

*Meet Closure:  $a, b \in F$  implies  $a \wedge b \in F$*



# FILTERS

---

- A filter on a distributive lattice  $A$  is a non-empty subset  $F$  of  $A$  satisfying

*Upwards Closure:  $a \in F$  and  $b \geq a$  implies  $b \in F$*

*Meet Closure:  $a, b \in F$  implies  $a \wedge b \in F$*

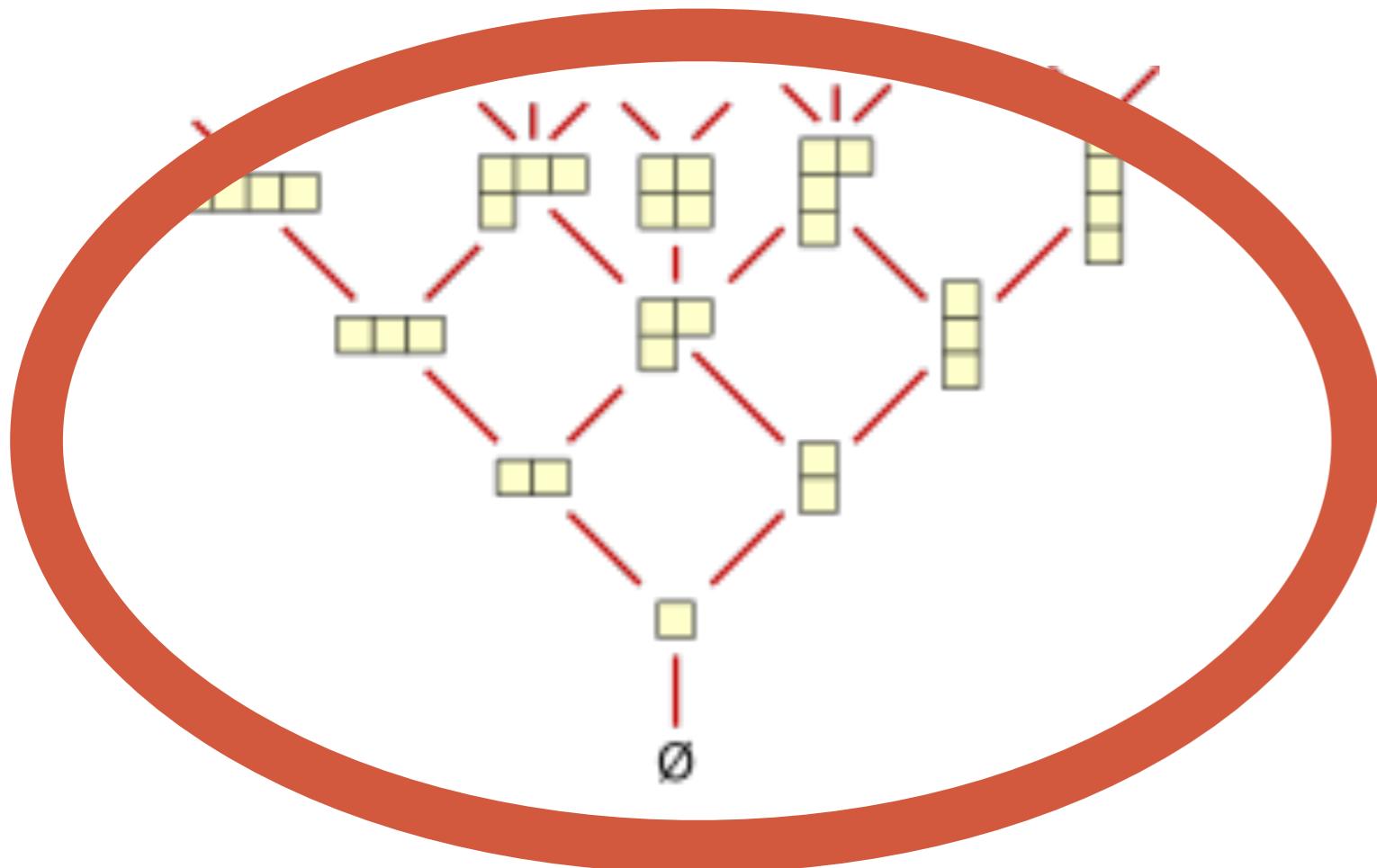
- A filter is proper if it additionally satisfies  $\perp \notin F$
- A proper filter is prime if it further satisfies

*$a \vee b \in F$  implies  $a \in F$  or  $b \in F$*

# FILTERS

---

- A filter is proper if it additionally satisfies  $\perp \notin F$
- A proper filter is prime if it further satisfies  
 $a \vee b \in F$  implies  $a \in F$  or  $b \in F$

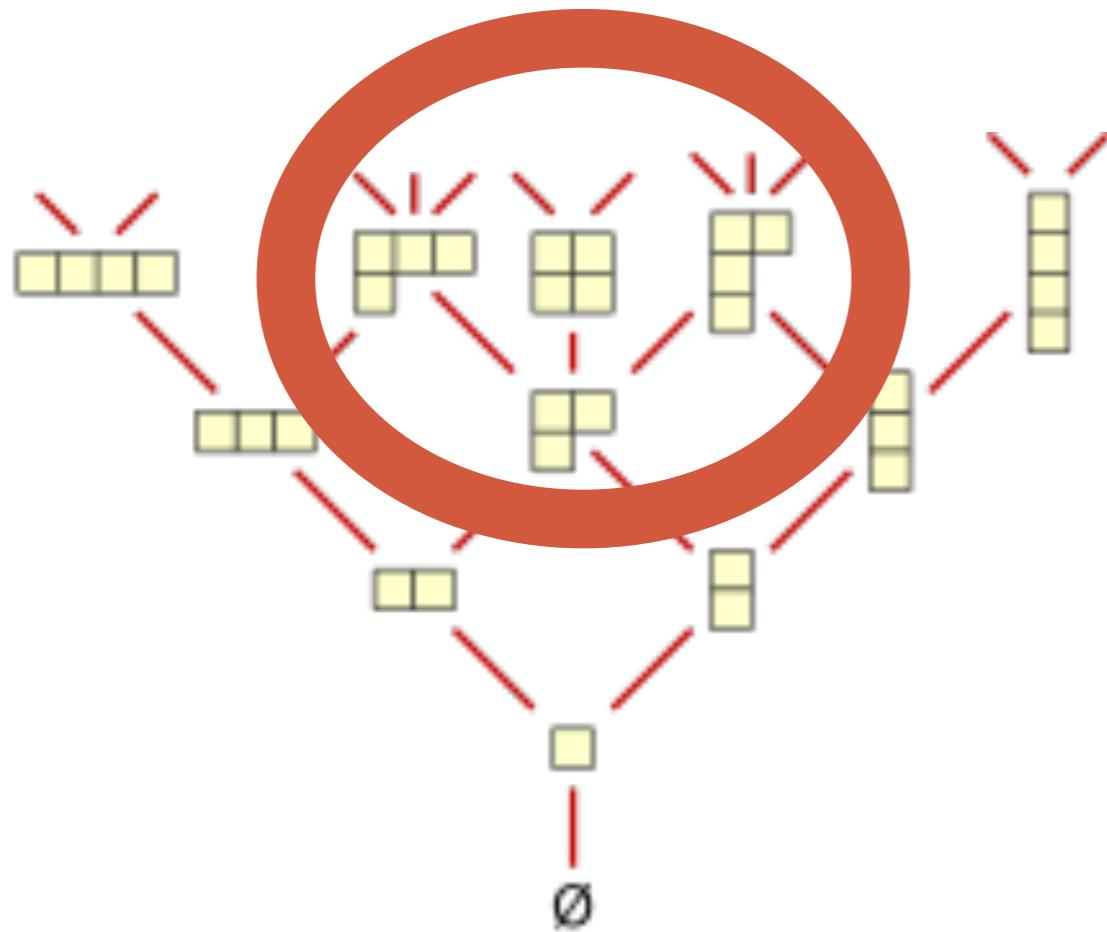


*is a filter but is not proper*

# FILTERS

---

- A filter is proper if it additionally satisfies  $\perp \notin F$
- A proper filter is prime if it further satisfies  
 $a \vee b \in F$  implies  $a \in F$  or  $b \in F$

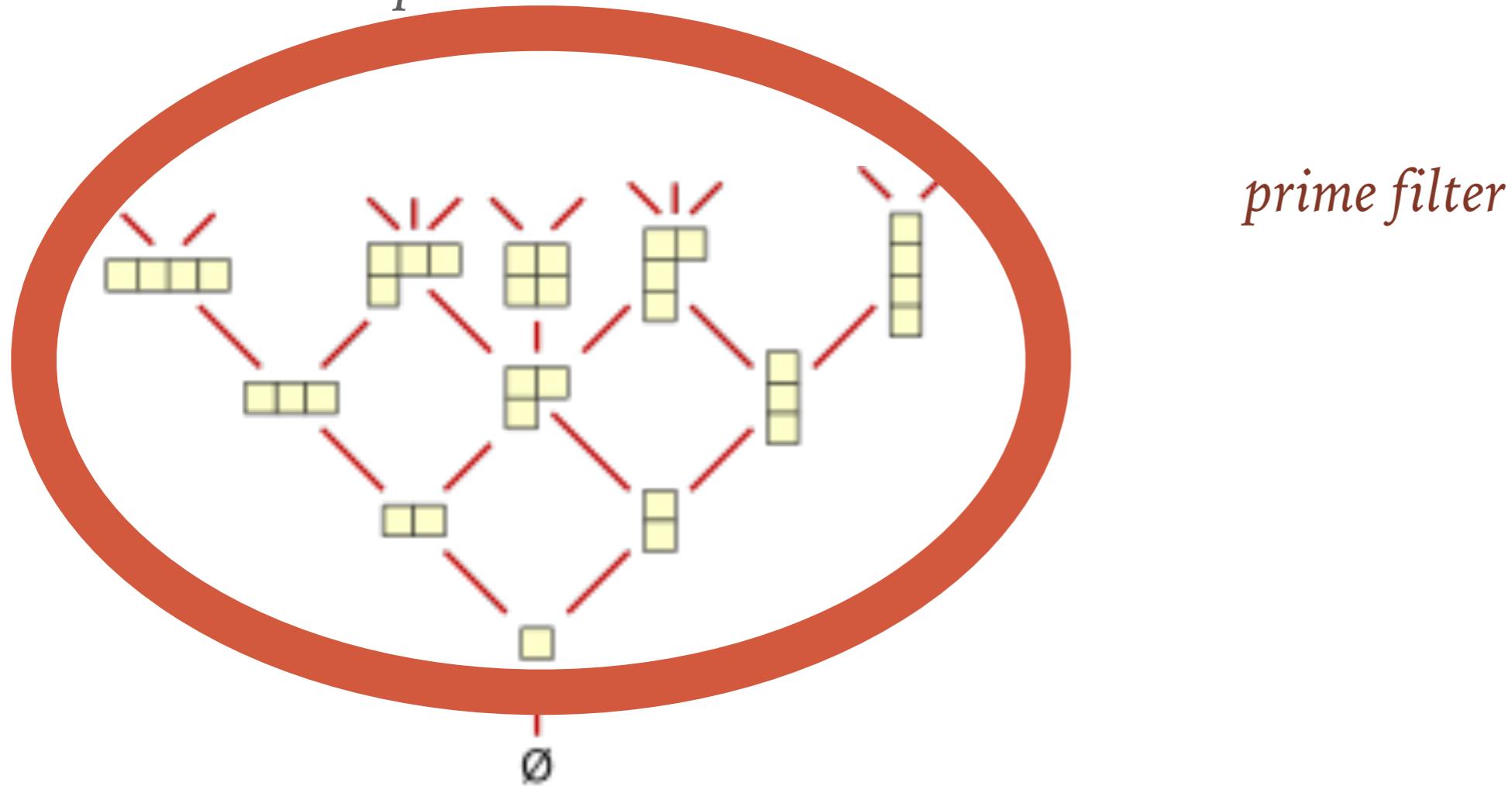


*is proper but is not prime*

# FILTERS

---

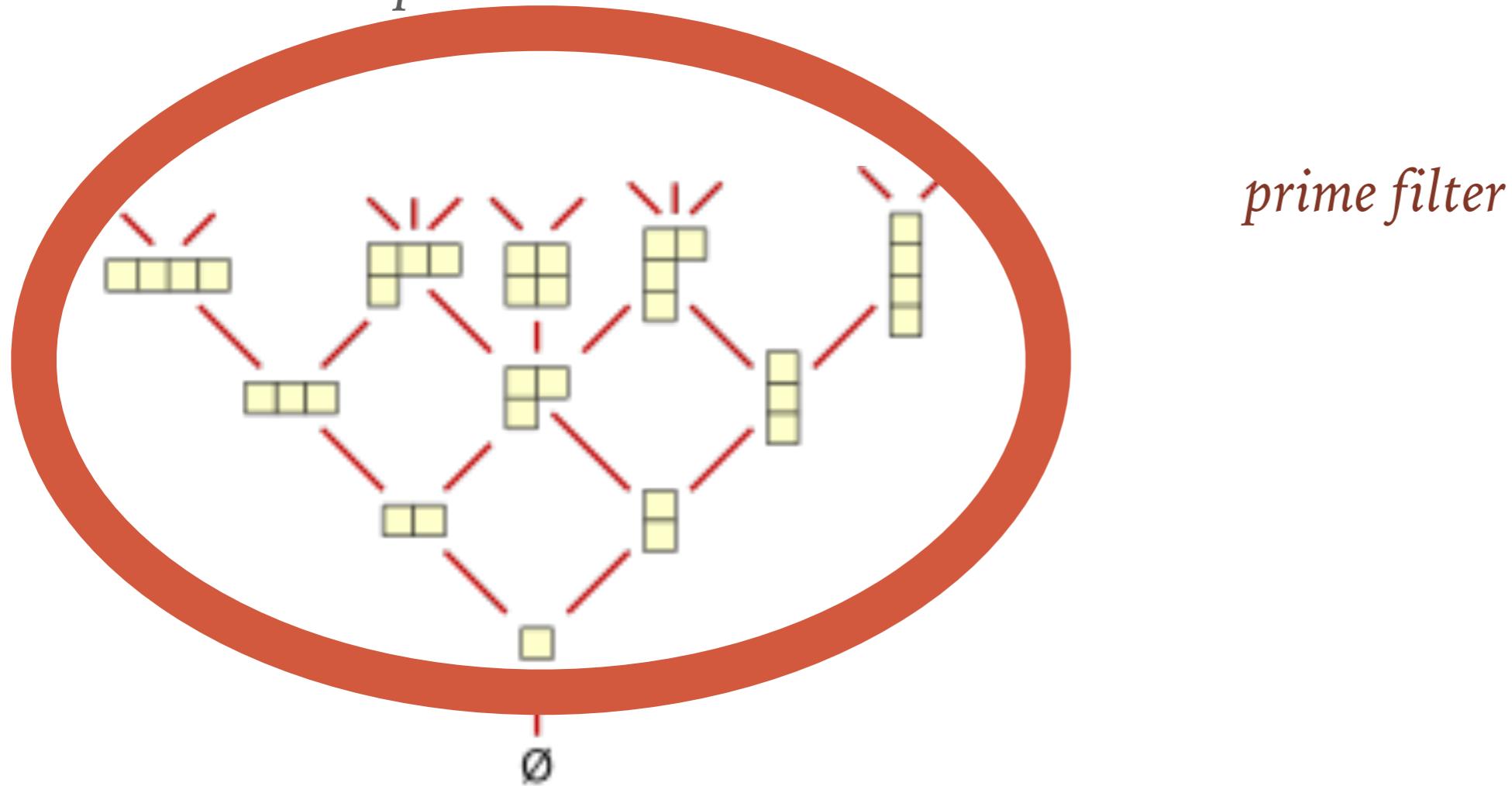
- A filter is proper if it additionally satisfies  $\perp \notin F$
- A proper filter is prime if it further satisfies
$$a \vee b \in F \text{ implies } a \in F \text{ or } b \in F$$



# FILTERS

---

- A filter is proper if it additionally satisfies  $\perp \notin F$
- A proper filter is prime if it further satisfies
$$a \vee b \in F \text{ implies } a \in F \text{ or } b \in F$$



# FILTERS

---

- A filter can be seen as a propositional theory, closed under a simple form of logical consequence.
- If  $\perp$  is in the theory then that theory is ‘improper’ as it is inconsistent and thus contains every proposition.
- Given a boolean algebra  $A$ , we can define a frame,  
 $Pr(A) := (\text{the set of prime filters of } A, \supseteq)$ 
  - we call it prime filter frame

# STONE REPRESENTATION THEOREM

---

- Every Boolean algebra is isomorphic to a sub algebra of a power set algebra.
- Given a frame semantic  $(X, \geq, \circ)$ 
  - define  $Com(\mathcal{K}) = (\mathcal{P}_{\geq}(X), \cap, \cup, \Rightarrow_{\mathcal{K}}, X, \emptyset)$  where
    - $\mathcal{P}_{\geq}(X) = \{A \subseteq X \mid \text{if } a \in A \text{ and } b \geq a \text{ then } b \in A\}$
    - $A \Rightarrow_B \mathcal{K} = \{x \mid x' \geq x \text{ and } x' \in A \text{ implies } x' \in B\}$

# STONE REPRESENTATION THEOREM

---

- Every Boolean algebra is isomorphic to a sub algebra of a power set algebra.
- Given a frame semantic  $(X, \geq, \circ)$ 
  - define  $Com(\mathcal{K}) = (\mathcal{P}_{\geq}(X), \cap, \cup, \Rightarrow_{\mathcal{K}}, X, \emptyset)$  where
    - $\mathcal{P}_{\geq}(X) = \{A \subseteq X \mid \text{if } a \in A \text{ and } b \geq a \text{ then } b \in A\}$
    - $A \Rightarrow_B \mathcal{K} = \{x \mid x' \geq x \text{ and } x' \in A \text{ implies } x' \in B\}$
  - A boolean algebra  $A$  is isomorphic to  $Com(Pr(A))$ 
    - The isomorphism  $h$  can also be written as
$$h(a) = \{F \in Pr(A) \mid a \in F\}$$

# STONE REPRESENTATION THEOREM

---

- Every Boolean algebra is isomorphic to a sub algebra of a power set algebra.
- Given a frame semantic  $(X, \geq, \circ)$ 
  - define  $Com(\mathcal{K}) = (\mathcal{P}_{\geq}(X), \cap, \cup, \Rightarrow_{\mathcal{K}}, X, \emptyset)$  where
    - $\mathcal{P}_{\geq}(X) = \{A \subseteq X \mid \text{if } a \in A \text{ and } b \geq a \text{ then } b \in A\}$
    - $A \Rightarrow_B \mathcal{K} = \{x \mid x' \geq x \text{ and } x' \in A \text{ implies } x' \in B\}$
  - A boolean algebra  $A$  is isomorphic to  $Com(Pr(A))$ 
    - The isomorphism  $h$  can also be written as
$$h(a) = \{F \in Pr(A) \mid a \in F\}$$

# PROOF SKETCH

---

- A boolean algebra  $A$  is isomorphic to  $\text{Com}(\text{Pr}(A))$

- The isomorphism  $h$  can also be written as

$$h(a) = \{F \in \text{Pr}(A) \mid a \in F\}$$

- $h$  is injective

$$h(a \rightarrow b) = h(a) \Rightarrow_{\text{Pr}(A)} h(b)$$

- $\text{Pr}(A)$  and topology on it forms a stone space
- With some topology, it's possible to show that the dual equivalence categories between Boolean algebra and stone space  $=>$  Equivalence of Algebraic and Frame Semantics.

# TOPOLOGICAL SPACE

---

- A topologocial space is a pair  $\mathcal{K} = (X, \mathcal{O})$  where  $X$  is a set and  $\mathcal{O} \subseteq 2^X$ , satisfying
  - $\emptyset, X \in \mathcal{O}$ ;
  - Given any  $O_i \in \mathcal{O}$  indexed by  $I$ ,  $\cup_{i \in I} O_i \in \mathcal{O}$ ;
  - $O_0, \dots, O_n \in \mathcal{O}$  implies  $\cap_{i=0}^n O_i \in \mathcal{O}$

# PROOF SKETCH

---

- A boolean algebra  $A$  is isomorphic to  $\text{Com}(\text{Pr}(A))$

- The isomorphism  $h$  can also be written as

$$h(a) = \{F \in \text{Pr}(A) \mid a \in F\}$$

- $h$  is injective

$$h(a \rightarrow b) = h(a) \Rightarrow_{\text{Pr}(A)} h(b)$$

- $\text{Pr}(A)$  and topology on it forms a stone space

# EQUIVALENCE OF ALGEBRAIC AND FRAME SEMANTICS

---

- Equivalence of Algebraic and Frame Semantics.
- Given algebra A and frame K. Let  $[-]$  be an algebraic interpretation on A

*valuation on  $\mathcal{X}$ . Define  $\mathcal{V}_{[-]} : \text{Prop} \rightarrow \mathcal{P}(\text{Pr}(A))$  by  $\mathcal{V}_{[-]}(p) = \theta_A(\llbracket p \rrbracket)$  and  $\llbracket - \rrbracket_{\mathcal{V}}$  as the algebraic interpretation on  $\text{Com}^{\mathcal{L}}(\mathcal{X})$  generated by  $\mathcal{V}$ . For all  $\mathcal{L}$  formulae  $\varphi$  the following hold.*

1.  $x \models_{\mathcal{V}} \varphi$  iff  $x \in \llbracket \varphi \rrbracket_{\mathcal{V}}$ ;
2.  $\llbracket \varphi \rrbracket \in F$  iff  $F \models_{\mathcal{V}_{[-]}} \varphi$ .

# ALGEBRAIC COMPLETENESS

---

**Theorem 7.2** (Algebraic Completeness (cf. [189])). *For all bunched logics  $\mathcal{L}$  and  $\mathcal{L}$  formulae  $\varphi$  and  $\psi$ , if  $\llbracket \varphi \rrbracket \leq \llbracket \psi \rrbracket$  holds for all algebraic interpretations then  $\varphi \vdash \psi$  is provable in the Hilbert system for  $\mathcal{L}$ .*

*Proof.* We reason contrapositively: suppose  $\varphi \vdash \psi$  is not provable in the Hilbert system for  $\mathcal{L}$ . Consider the Lindenbaum-Tarski algebra over the propositional variables occurring in  $\varphi$  and  $\psi$  with the interpretation given by sending each formula to its equivalence class.  $[\varphi]_{\equiv} \leq [\psi]_{\equiv}$  iff  $[\varphi \rightarrow \psi]_{\equiv} = [\top]_{\equiv}$  iff  $\top \vdash \varphi \rightarrow \psi$  provable iff  $\varphi \vdash \psi$  (using the definition of Heyting implication and the deduction theorem for bunched logics), and so  $[\varphi]_{\equiv} \not\leq [\psi]_{\equiv}$ .  $\square$

# THE TECHNIQUE

---

- Given a set  $L$  of inference rules and semantic frame  $F$ 
  - Find a algebra  $A$  that is algebraically sound and complete of  $L$
  - Define morphisms on  $F$  to mirror the operators in  $A$
  - Show  $\text{Com}(F)$  is a algebra  $A$
  - Show  $\text{Pr}(A)$  is a semantic frame  $F$
  - $\Rightarrow$  Equivalence of Algebraic and Frame Semantics.
  - $A$  is algebraically sound and complete w.r.t  $L$  implies that  $F$  is sound and complete w.r.t.  $L$  too

# GENERALIZATION

---

- Boolean algebra — Stone space

# THANKS