Programming with Message Passing PART II: MPI

HPC Fall 2012

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Overview

- MPI overview
- MPI process creation
- MPI point-to-point communication
- MPI collective communications
- MPI groups and communicators
- MPI virtual topologies
- MPE and Jumpshot
- Further reading



MPI

- Message Passing Interface (MPI) is an API and protocol standard with portable implementations
 - □ MPICH, LAM-MPI, OpenMPI, ...
- Hardware platforms:
 - □ Distributed Memory
 - ☐ Shared Memory
 - Particularly SMP / NUMA architectures
 - □ Hybrid
 - SMP clusters, workstation clusters, and heterogeneous networks
- Parallelism is explicit
 - Programmer is responsible for correctly identifying parallelism and implementing parallel algorithms using MPI constructs
- SPMD model with static process creation
 - □ MPI-2 allows dynamic process creation: MPI_Comm_spawn()



MPI Process Creation

- Static process creation
 - Start a N processes running the same program prog1
 mpirun prog1 -np N
 but this does not specify where the prog1 copies run
 - □ Use batch processing tools, such as SGE, to run MPI programs on a cluster
- MPI-2 supports dynamic process creation:
 MPI Comm spawn()

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MPI SPMD Computational Model

```
main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);
    doWork();

    MPI_Finalize();
}
```

All processes execute this work

All processes belong to the

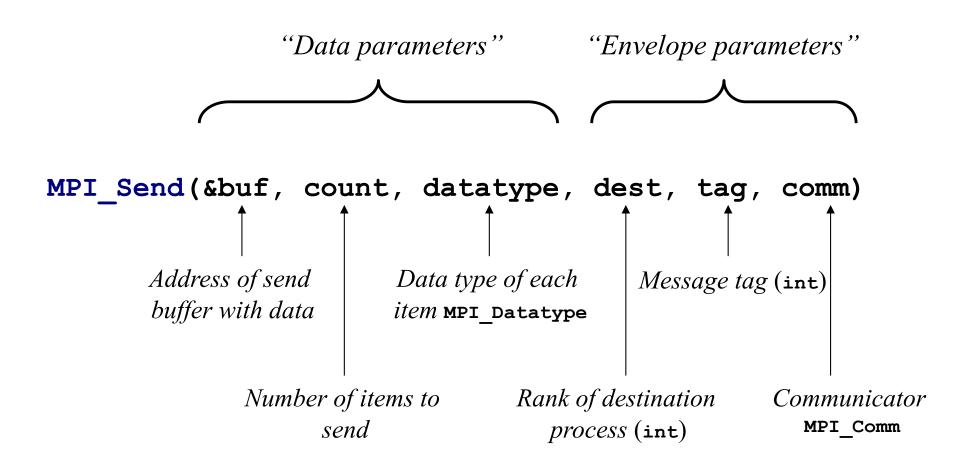
MPI_COMM_WORLD communicator,
and each process has a rank from 0
to P-1 in the communicator

```
doWork()
{
  int myrank;
  MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
  if (myrank == 0)
    printf("I'm the master process");
}
```

Use rank numbers to differentiate work

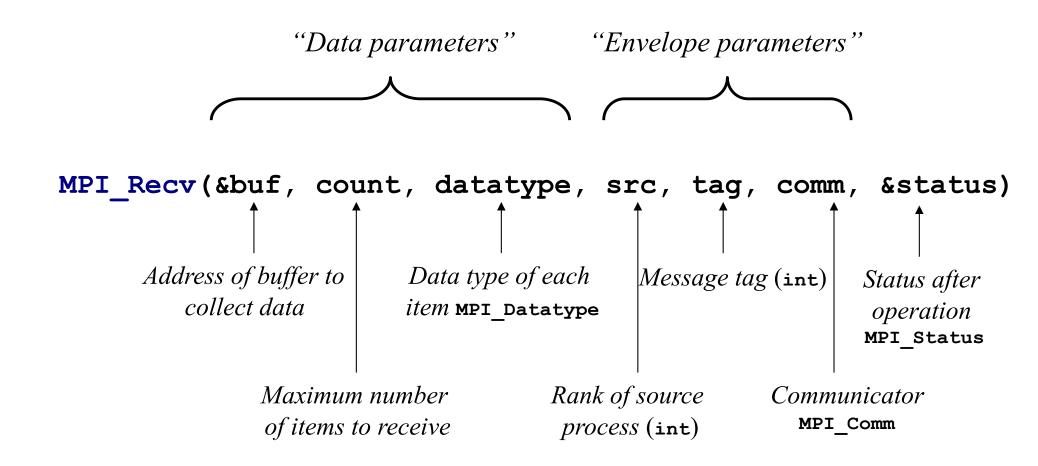


MPI Point-to-point Send Format





MPI Point-to-point Recv Format





Example: Send-Recv Between two Processes

```
main(int argc, char *argv[])
  int myrank;
  int value = 123;
  MPI Status status;
 MPI Init(&argc, &argv);
  MPI Comm Rank(MPI COMM WORLD, &myrank);
  if (myrank == 0)
    MPI Send(&value, 1, MPI INT, 1, MPI ANY TAG, MPI COMM WORLD);
  else if (myrank == 1)
    MPI Recv(&value, 1, MPI INT, 0, MPI ANY TAG, MPI COMM WORLD, &status);
  MPI Finalize();
```



MPI Point-to-point Communication Modes

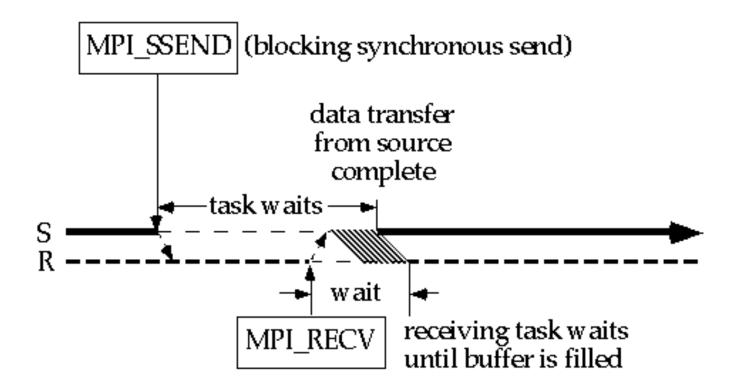
Communication mode	Blocking routines	Nonblocking routines
Synchronous	MPI_Ssend	MPI_Issend
Ready	MPI_Rsend	MPI_Irsend
Buffered	MPI_Bsend	MPI_Ibsend
Standard	MPI_Send	MPI_Isend
	MPI_Recv	MPI_Irecv
	MPI_Sendrecv	
	MPI_Sendrecv_replace	

Roughly speaking: MPI_Xsend = MPI_Ixsend + MPI_Wait

MPI_Recv = MPI_Irecv + MPI_Wait



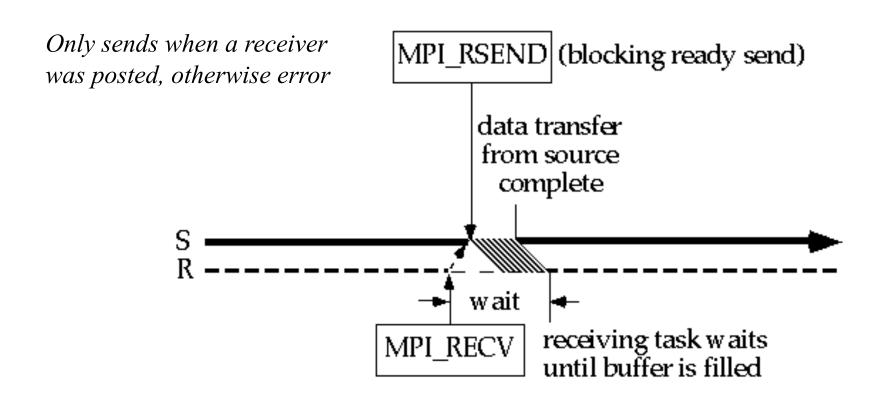
MPI Synchronous Send



MPI_Ssend(void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)



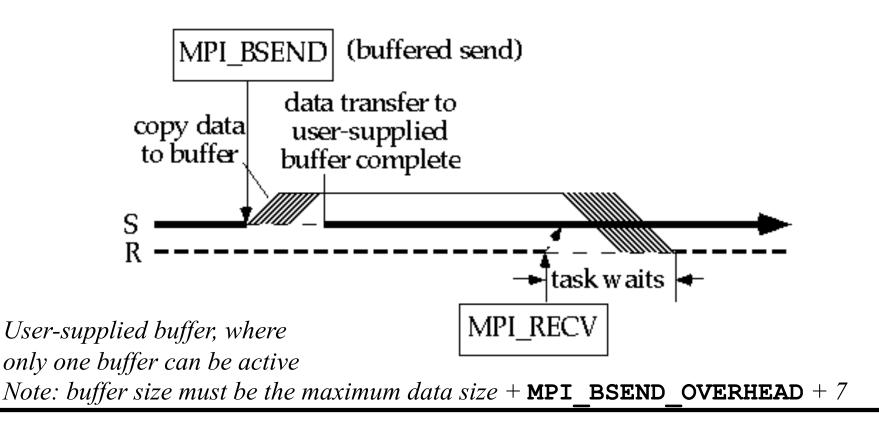
MPI Blocking Ready Send



MPI_Rsend(void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)



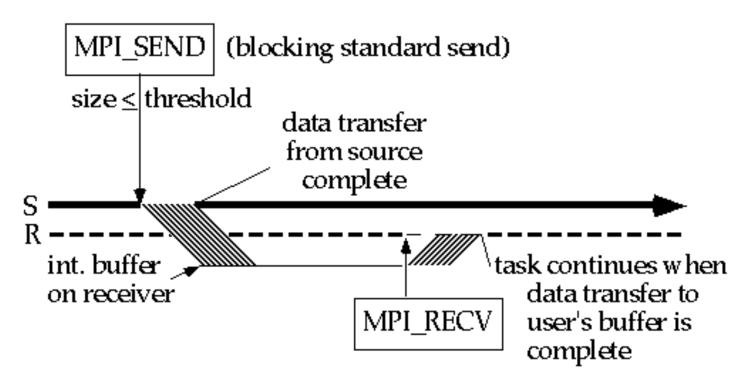
MPI Buffered Send



MPI_Buffer_attach(void *buf, int size)
MPI_Bsend(void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)



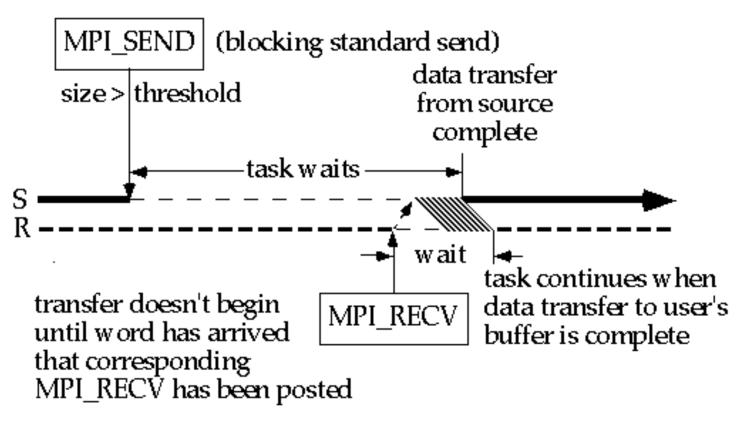
MPI Blocking Standard Send Small Message Size



The threshold value (also called the "eager limit") differs between systems



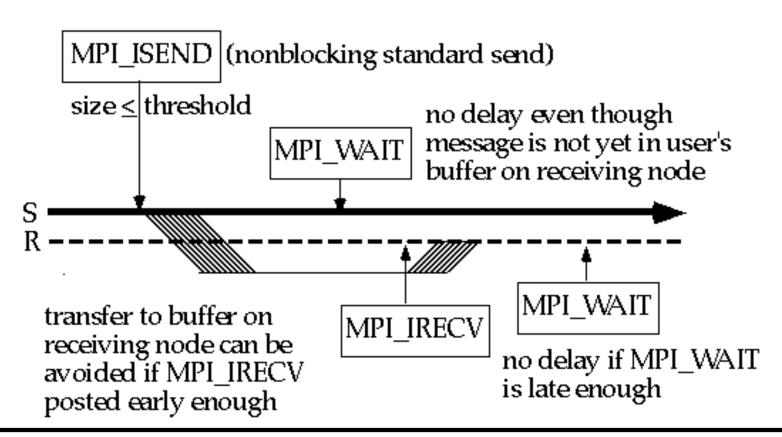
MPI Blocking Standard Send Large Message Size



The threshold value (also called the "eager limit") differs between systems



MPI Nonblocking Standard Send Small Message Size



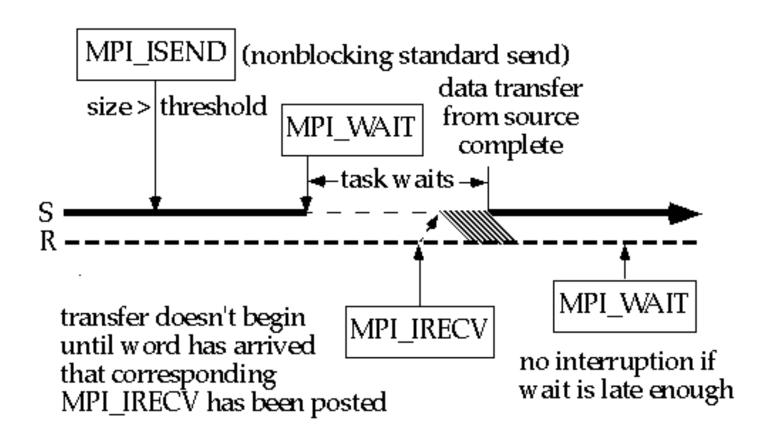
MPI_Isend(void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm, MPI_Request *request)

MPI_Irecv(void* buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Request *request)

MPI_Wait(MPI_Request *request, MPI_Status *status)



MPI Nonblocking Standard Send Large Message Size





MPI Communication Modes

Communication mode	Advantages	Disadvantages
Synchronous	Safest, and therefore most portable SEND/RECV order not critical Amount of buffer space irrelevant	Can incur substantial synchronization overhead
Ready	Lowest total overhead SEND/RECV handshake not required	RECV must precede SEND
Buffered	Decouples SEND from RECV No sync overhead on SEND Order of SEND/RECV irrelevant Programmer can control size of buffer space	Additional system overhead incurred by copy to buffer
Standard	Good for many cases	May not be suitable for your program



Example 1: Nonblocking Send and Blocking Recv

```
main(int argc, char *argv[])
  int myrank;
  int value = 123;
 MPI Status status;
 MPI Request req;
 MPI Init(&argc, &argv);
 MPI Comm Rank (MPI COMM WORLD, &myrank);
  if (myrank == 0)
  { MPI Isend(&value, 1, MPI INT, 1, MPI ANY TAG, MPI COMM WORLD, req);
    doWork(); /* do not modify value */
    MPI Wait(&req, &status);
  else if (myrank == 1)
    MPI Recv(&value, 1, MPI INT, 0, MPI ANY TAG, MPI COMM WORLD, &status);
  MPI Finalize();
```



Example 2: Nonblocking Send/Recv

```
main(int argc, char *argv[])
  int myrank;
  int value = 123;
 MPI Status status;
 MPI Request req;
  MPI Init(&argc, &argv);
  MPI Comm Rank (MPI COMM WORLD, &myrank);
  if (myrank == 0)
  { MPI Isend(&value, 1, MPI INT, 1, MPI ANY TAG, MPI COMM WORLD, &req);
    doWork(); /* do not modify value */
    MPI Wait(&req, &status);
  else if (myrank == 1)
  { MPI Irecv(&value, 1, MPI INT, 0, MPI ANY TAG, MPI COMM WORLD, &req);
    doWork(); /* do not read or modify value */
   MPI Wait(&req, &status);
    /* using value is OK now */
 MPI Finalize();
```



MPI Wait



MPI Test



MPI Combined Sendrecy



Example 3: Combined Sendrecv

```
main(int argc, char *argv[])
  int myrank, hisrank;
  float value;
 MPI Status status;
 MPI Init(&argc, &argv);
  MPI Comm Rank(MPI COMM WORLD, &myrank);
  hisrank = 1-myrank;
  if (myrank == 0)
   value = 3.14;
  else
    value = 1.41;
 MPI Sendrecv replace (&value, 1, MPI FLOAT, hisrank, MPI ANY TAG,
                       hisrank, MPI ANY TAG, MPI COMM WORLD, &status);
  printf("Process %d got value %g\n", myrank, value);
  MPI Finalize();
```



MPI Basic Datatypes

MPI CHAR MPI SHORT MPI INT MPI LONG MPI UNSIGNED CHAR MPI UNSIGNED SHORT MPI UNSIGNED MPI UNSIGNED LONG MPI FLOAT MPI DOUBLE MPI LONG DOUBLE MPI BYTE MPI PACKED

signed char
signed short int
signed int
signed long int
unsigned char
unsigned short int
unsigned int
unsigned long int
float
double
long double



MPI Collective Communications

- Coordinated communication within a group of processes identified by an MPI communicator
- Collective communication routines block until locally complete
- Amount of data sent must exactly match amount of data specified by receiver
- No message tags are needed



MPI Collective Communications Performance Considerations

- Communications are hidden from user
 - Communication patterns depend on implementation and platform on which MPI runs
 - ☐ In some cases, the root process originates or receives all data
 - □ Performance depends on implementation of MPI
- Communications may, or may not, be synchronized (implementation dependent)
 - Not always the best choice to use collective communication
 - There may be forced synchronization, which can be avoided with nonblocking point-to-point communications



MPI Collective Communications Classes

Three classes:

- 1. Synchronization
 - Barrier synchronization
- 2. Data movement
 - Broadcast
 - Scatter
 - Gather
 - □ All-to-all
- Global computation
 - Reduce
 - □ Scan



MPI Barrier

MPI_Barrier(MPI_Comm comm)

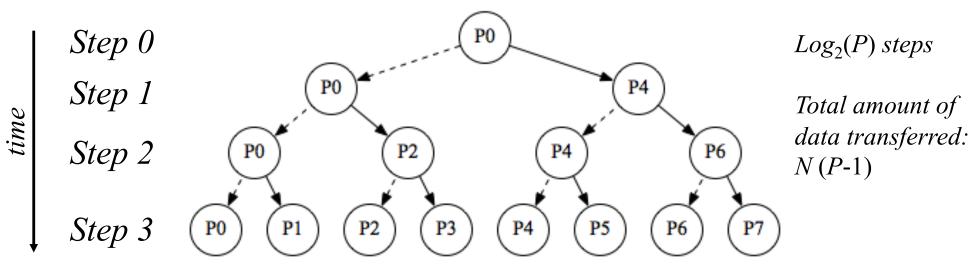
A node invoking the barrier routine will be blocked until all the nodes within the group (communicator) have invoked it



MPI Broadcast

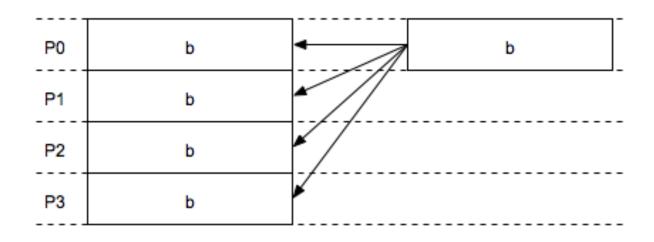
Simple broadcast implementation: root sends data to all processes, which is efficient on a bus-based parallel machine

More efficient on a network: broadcast as a tree operation





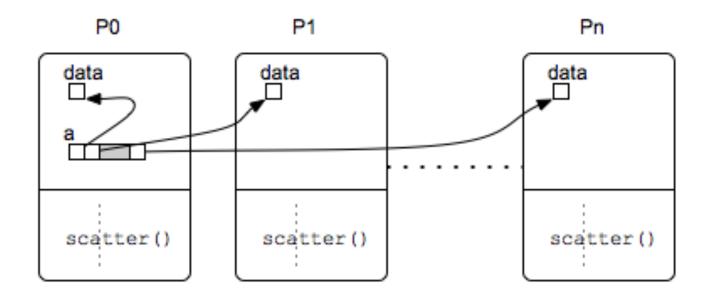
MPI Broadcast Example



```
float A[N][N], Ap[N/P][N], b[N], c[N], cp[N/P];
...
root = 0;
MPI_Broadcast(b, N, MPI_Float, root, MPI_COMM_WORLD);
...
```

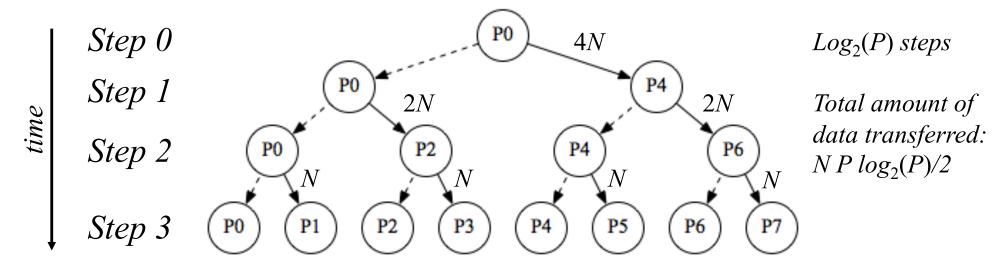


MPI Scatter



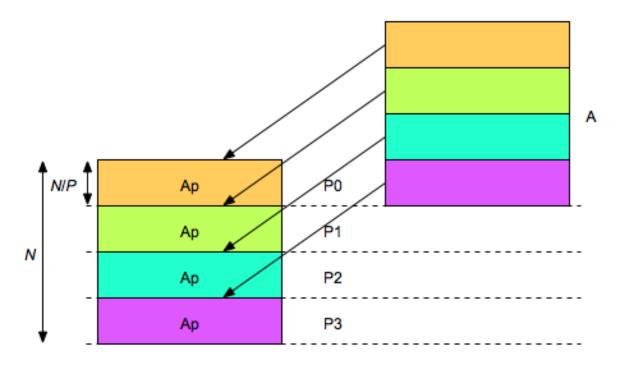


MPI Scatter



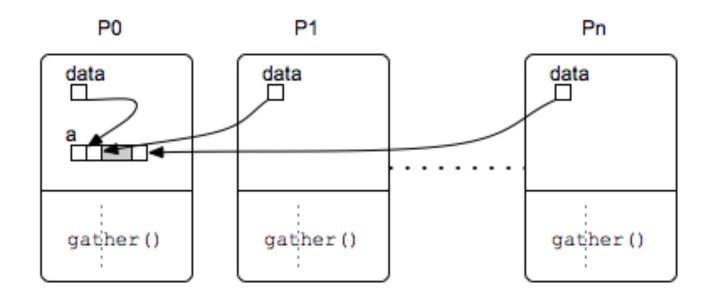


MPI Scatter Example



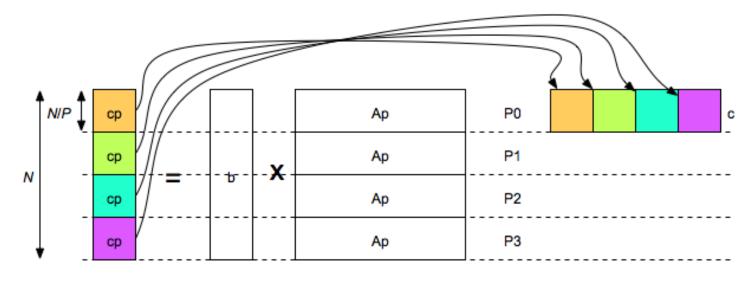


MPI Gather



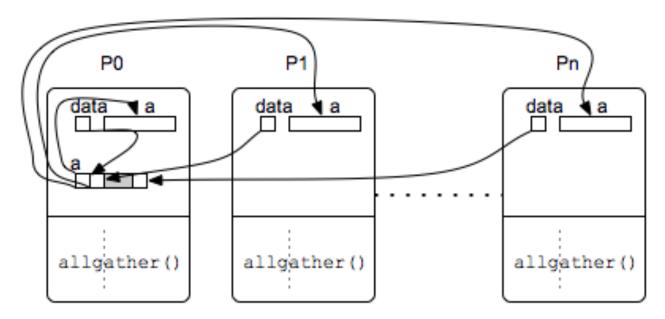


MPI Gather Example



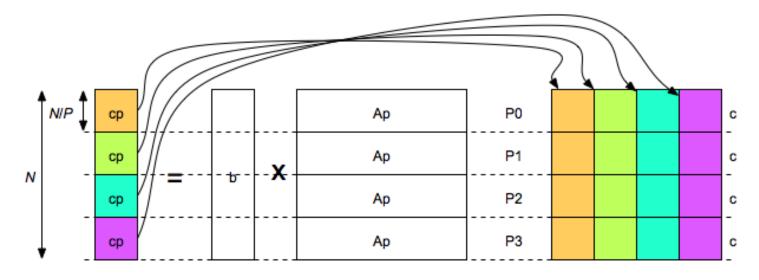


MPI AllGather





MPI AllGather Example



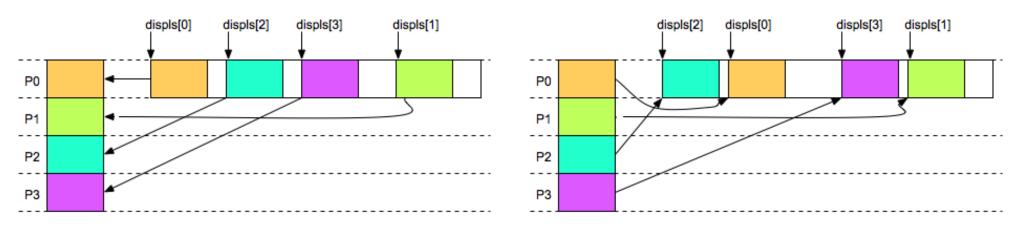
```
float A[N][N], Ap[N/P][N], b[N], c[N], cp[N/P];
...
for (i = 1; i < N/P; i++)
{
    cp[i] = 0;
    for (k = 0; k < N; k++)
        cp[i] = cp[i] + Ap[i][k] * b[k];
}
MPI_AllGather(cp, N/P, MPI_Float, c, N/P, MPI_Float, MPI_COMM_WORLD);</pre>
```



MPI Scattery and (All)Gathery

```
MPI Gatherv(void *sbuf,
MPI Scatterv(void *sbuf,
                                              int scount,
             int *scounts,
             int *displs,
                                              MPI Datatype stype,
             MPI Datatype stype,
                                              void *rbuf,
             void *rbuf,
                                              int *rcounts,
                                              int *displs,
             int rcount,
                                              MPI Datatype rtype,
             MPI Datatype rtype,
             int root,
                                              int root,
                                              MPI Comm comm)
             MPI Comm comm)
```

Adds displacements to scatter/gather operations and counts may vary per process



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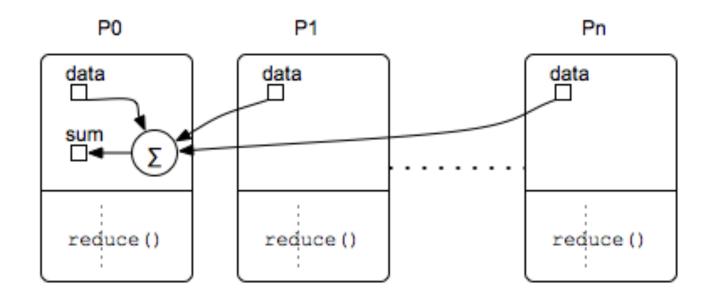
MPI All to All

Global transpose: the jth block from processor i is received by processor j and stored in ith block

P0	a0	b0	c0	d0		a0	a1	a2	a3
P1	a1	b1	c1	d1		b0	b1	b2	b3
P2	a2	b2	c2	d2		c0	c1	c2	сЗ
Р3	a3	b3	сЗ	d3	γ	d0	d1	d2	d3



MPI Reduce



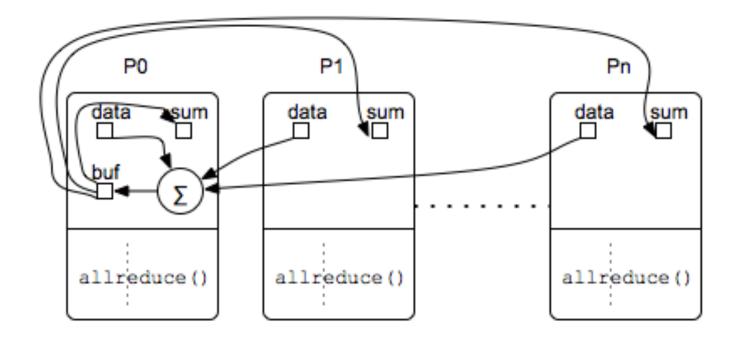


MPI Reduce Example

Р0	a0	b0	c0	d0	a0+a1+a2+a3	b0+b1+b2+b3	c0+c1+c2+c3	d0+d1+d2+d3
P1	a1	b1	c1	d1				
P2	a2	b2	c2	d2				
P3	а3	b3	с3	d3				



MPI AllReduce





MPI AllReduce Example

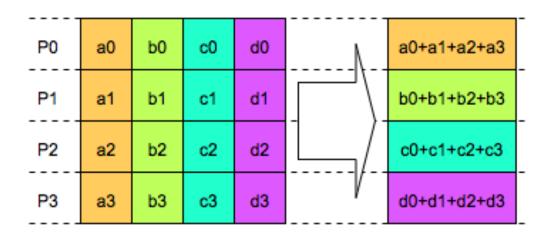
P0	a0	b0	c0	d0		a0+a1+a2+a3	b0+b1+b2+b3	c0+c1+c2+c3	d0+d1+d2+d3
P1	a1	b1	c1	d1		a0+a1+a2+a3	b0+b1+b2+b3	c0+c1+c2+c3	d0+d1+d2+d3
P2	a2	b2	c2	d2		a0+a1+a2+a3	b0+b1+b2+b3	c0+c1+c2+c3	d0+d1+d2+d3
P3	a3	b3	сЗ	d3	<i> </i>	a0+a1+a2+a3	b0+b1+b2+b3	c0+c1+c2+c3	d0+d1+d2+d3



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MPI Reduce_scatter

Same as Reduce followed by Scatter



Note: rcounts = number of elements received, which is >1 when N>P

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MPI Scan

Р0	a0	b0	c0	d0		a0	ь0	c0	d0
P1	a1	b1	c1	d1		a0+a1	b0+b1	c0+c1	d0+d1
P2	a2	b2	c2	d2		a0+a1+a2	b0+b1+b2	c0+c1+c2	d0+d1+d2
Р3	a3	b3	сЗ	d3	J	a0+a1+a2+a3	b0+b1+b2+b3	c0+c1+c2+c3	d0+d1+d2+d3



MPI Reduce and Scan Reduction Operators

MPI_OP	Operation	С	Fortran
MPI_MAX	maximum	integer, float	integer, real, complex
MPI_MIN	minimum	integer, float	integer, real, complex
MPI_SUM	sum	integer, float	integer, real, complex
MPI_PROD	product	integer, float	integer, real, complex
MPI_LAND	logical and	integer	logical
MPI_BAND	bit-wise and	integer, MPI_BYTE	integer, MPI_BYTE
MPI_LOR	logical or	integer	logical
MPI_BOR	bit-wise or	integer, MPI_BYTE	integer, MPI_BYTE
MPI_LXOR	logical xor	integer	logical
MPI_BXOR	bit-wise xor	integer, MPI_BYTE	integer, MPI_BYTE
MPI_MAXLOC	max val and loc	float, double, long double	real, complex, double precision
MPI_MINLOC	min val and loc	float, double, long double	real, complex, double precision

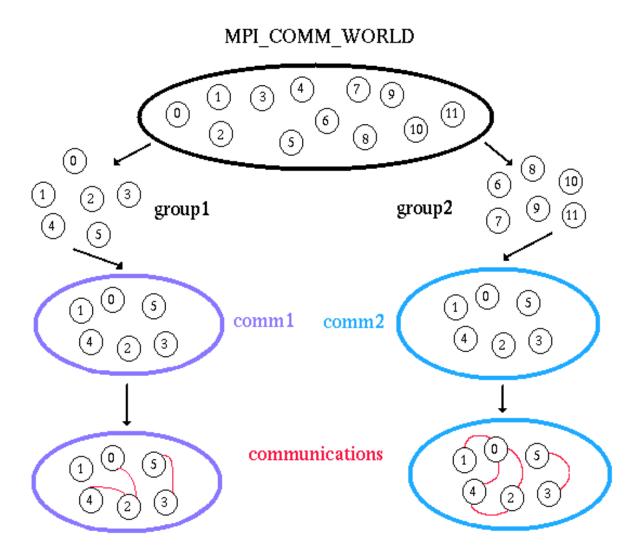


MPI Groups and Communicators

- A group is an ordered set of processes
 - □ Each process in a group is associated with a unique integer rank between 0 and *P*-1, with *P* the number of processes in the group
- A communicator encompasses a group of processes that may communicate with each other
 - □ Communicators can be created for specific groups
 - □ Processes may be in more than one group/communicator
- Groups/communicators are dynamic and can be setup and removed at any time
- From the programmer's perspective, a group and a communicator are the same



MPI Groups and Communicators



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MPI Group Operations

```
MPI Comm group
returns the group associated with a communicator
MPI Group union
creates a group by combining two groups
MPI Group intersection
creates a group from the intersection of two groups
MPI Group difference
creates a group from the difference between two groups
MPI Group incl
creates a group from listed members of an existing group
MPI Group excl
creates a group excluding listed members of an existing group
MPI Group range incl
creates a group according to first rank, stride, last rank
MPI Group range excl
creates a group by deleting according to first rank, stride, last rank
MPI Group free
marks a group for deallocation
```



MPI Communicator Operations

```
MPI Comm size
returns number of processes in communicator's group
MPI Comm rank
returns rank of calling process in communicator's group
MPI Comm_compare
compares two communicators
MPI Comm dup
duplicates a communicator
MPI Comm create
creates a new communicator for a group
MPI Comm split
splits a communicator into multiple, non-overlapping communicators
MPI Comm free
marks a communicator for deallocation
```



MPI Virtual Topologies

- A virtual topology can be associated with a communicator
- Two types of topologies supported by MPI
 - ☐ Cartesian (grid)
 - ☐ Graph
- Increases efficiency of communications
 - □ Some MPI implementations may use the physical characteristics of a given parallel machine
 - ☐ Introduces locality: low communication overhead with nodes that are "near" (few message hops), while distant nodes impose communication penalties (many hops)



MPE and Jumpshot

```
MPE Log get state eventIDs(int *startID, int *finalID)
Get a pair of event numbers to describe a state (see next)
MPE Describe state(int startID, int finalID, const char *name
                       const char *color)
Describe the state with a name and color (e.g. "red", "blue", "green")
MPE Log get solo eventID(int *eventID)
Get an event number to describe an event (see next)
MPE Describe event(int eventID, const char *name, const char *color)
Describe the event with a name and color (e.g. "red", "blue", "green")
MPE Log event(int event, int data, const char *bytebuf)
Record event in the log, data is unused and bytebuf carries informational data are should be NULL.
Use two calls, one with the startID and the other with the finalID to log the time of a state.
```



MPE and Jumpshot

```
int event1, state1s, state1f;
int myrank;
MPI Comm rank(MPI COMM WORLD, &myrank);
MPE Log get solo eventIDs(&event1);
MPE Log get state eventIDs(&state1s, &state1f);
if (myrank == 0)
{ MPE Describe event(event1, "Start", "green");
  MPE Describe state(state1s, state1f, "Computing", "red");
MPE Log event(event1, 0, NULL);
MPI Bcast(...);
MPE Log event(state1s, 0, NULL);
doComputation();
MPE Log event(state1f, 0, NULL);
MPI Barrier();
```



MPI+MPE Compilation, Linking, and Run

- Compilation with MPE requires mpe and 1mpe libs:
 - □ mpicc -o myprog myprog.c -lmpe -llmpe
 - □ Note: without -llmpe you must init and finalize the logs
- Run directly...
 - □ mpirun -np 4 ./myprog
- ... or in batch mode (e.g. with SGE)
 - □ qsub runmyprog.sh
- This generates the log file (typically in home dir)
 - □ myprog.clog2
- Display log file with
 - □ jumpshot

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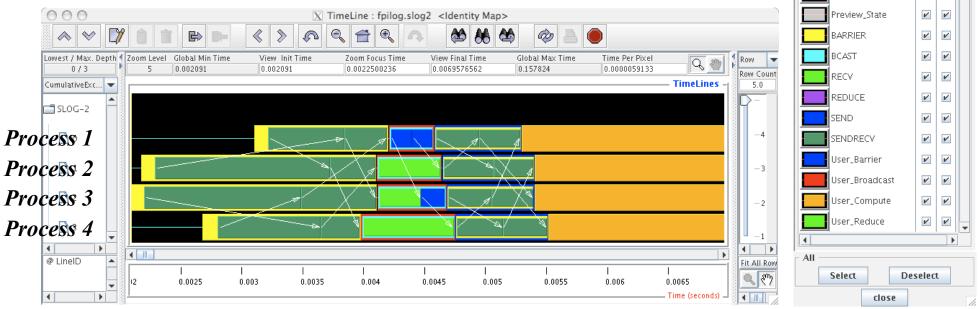
```
#!/bin/bash
# All environment variables active within the qsub
# utility to be exported to the context of the job:
#$ -V
# use openmpi with 4 nodes
#$ -pe openmpi_4 4
# use current directory
#$ -cwd
mpirun -np $NSLOTS ./myprog
```



Jumpshot

- Jumpshot generates a space-time diagram from a log file
 - □ Supports clog2 (through conversion) and newer slog2 formats
 - □ Shows user-defined states (start-end) and single point events
 - www.cs.fsu.edu/~engelen/courses/HPC/cpilog.c
 - www.cs.fsu.edu/~engelen/courses/HPC/fpilog.f

Space-time diagram



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O O X Legend : fpilog.slog2

review_Arrow

VV

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Further Reading

- [PP2] pages 52-61
- Optional:

Lawrence Livermore National Laboratories MPI Tutorial http://www.llnl.gov/computing/tutorials/mpi/