

Computational Genomics

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גנומיקה חישובית

פרופ' רון שמיר ופרופ' רודד שרן

ביה"ס למדעי המחשב, אוניברסיטת תל אביב

Lecture 4:

Multiple alignment

13,15/11/12

Multiple Sequence Alignment

Reference:

Gusfield, Algorithms on Strings, Trees & Sequences

Some slides from:

- Jones, Pevzner, USC *Intro to Bioinformatics Algorithms*
<http://www.bioalgorithms.info/>
- S. Batzoglou, Stanford http://ai.stanford.edu/~serafim/CS262_2006/
- Geiger, Wexler, Technion <http://www.cs.technion.ac.il/~cs236522/>
- Ruzzo, Tompa U. Washington CSE 590bi
- Poch, Strasbourg www.inra.fr/internet/Projets/agroBI/PHYLO/Poch.ppt
- A. Drummond, Auckland, NZ



Multiple Alignment vs. Pairwise Alignment

- Up until now we have only tried to align two sequences.
- What about more than two? And what for?
- A faint similarity between two sequences becomes significant if present in many
- Multiple alignments can reveal subtle similarities that pairwise alignments do not reveal



Multiple Alignment vs. Pairwise Alignment

- "Pairwise alignment whispers ...
multiple alignment shouts out loud"

Hubbard, Lesk, Tramontano, Nature Structural
Biology 1996.



Multiple Alignment Definition

Input: Sequences S_1, S_2, \dots, S_k over the same alphabet

Output: Gapped sequences S'_1, S'_2, \dots, S'_k of equal length

1. $|S'_1| = |S'_2| = \dots = |S'_k|$
2. Removal of spaces from S'_i gives S_i for all i

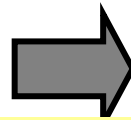


Example

$S_1 = \text{AGGTC}$

$S_2 = \text{GTTCG}$

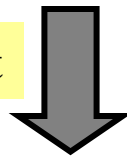
$S_3 = \text{TGAAC}$



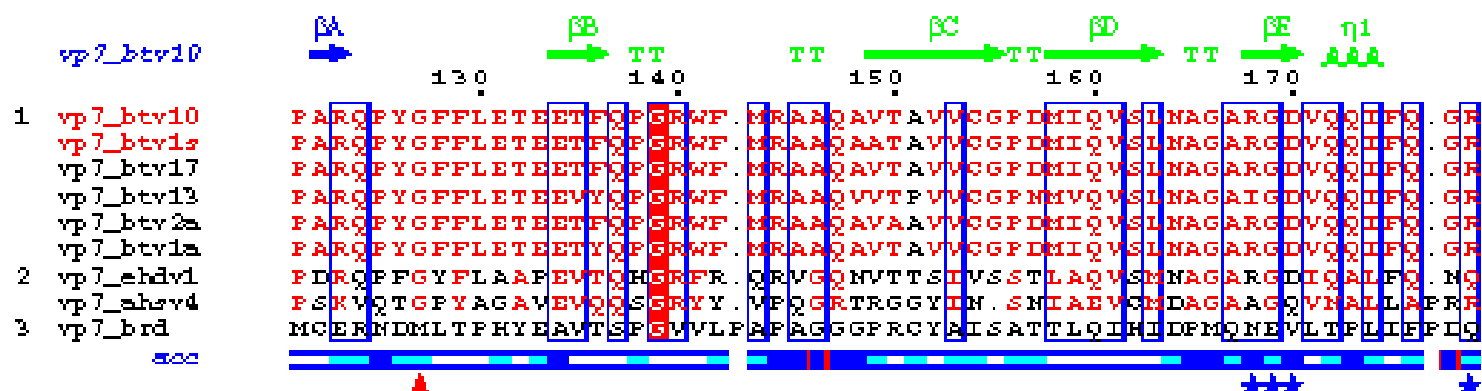
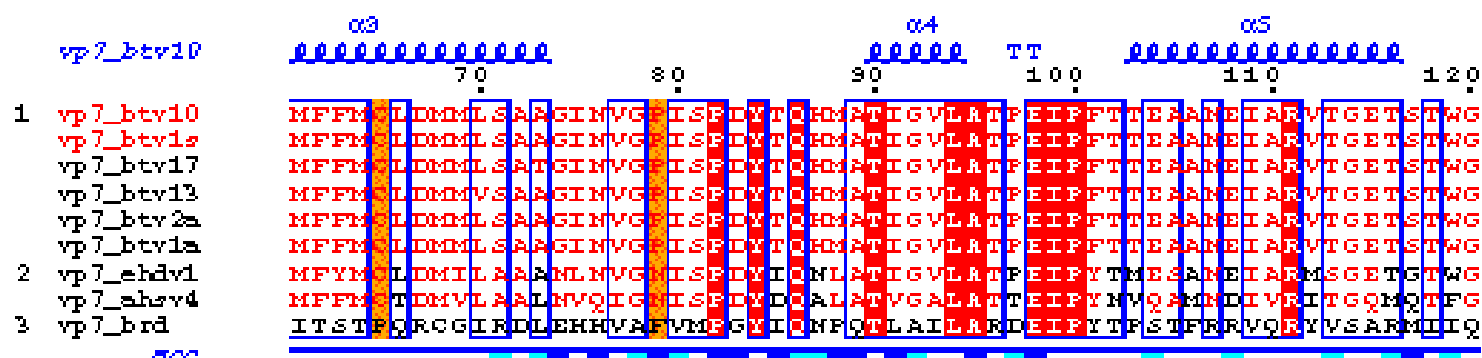
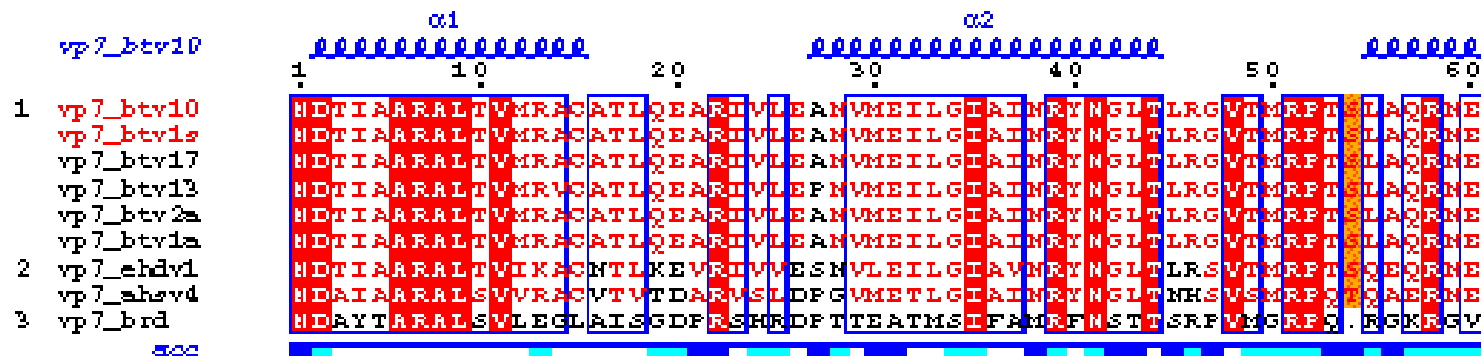
Possible alignment

A	G	G	T	-	C	-
-	G	-	T	T	C	G
T	G	-	A	A	C	-

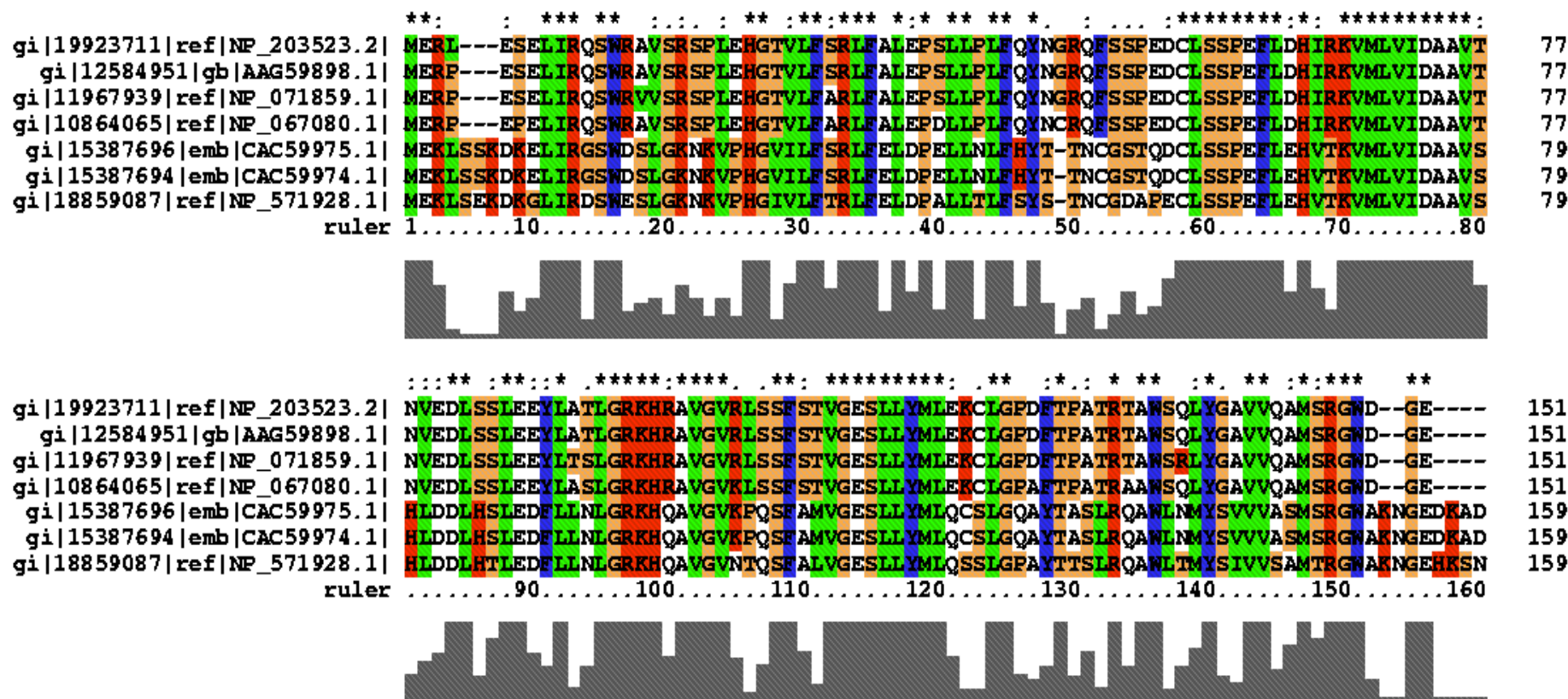
Possible alignment



A	G	G	T	-	C	-
G	T	T	-	-	C	G
-	T	G	A	A	A	C



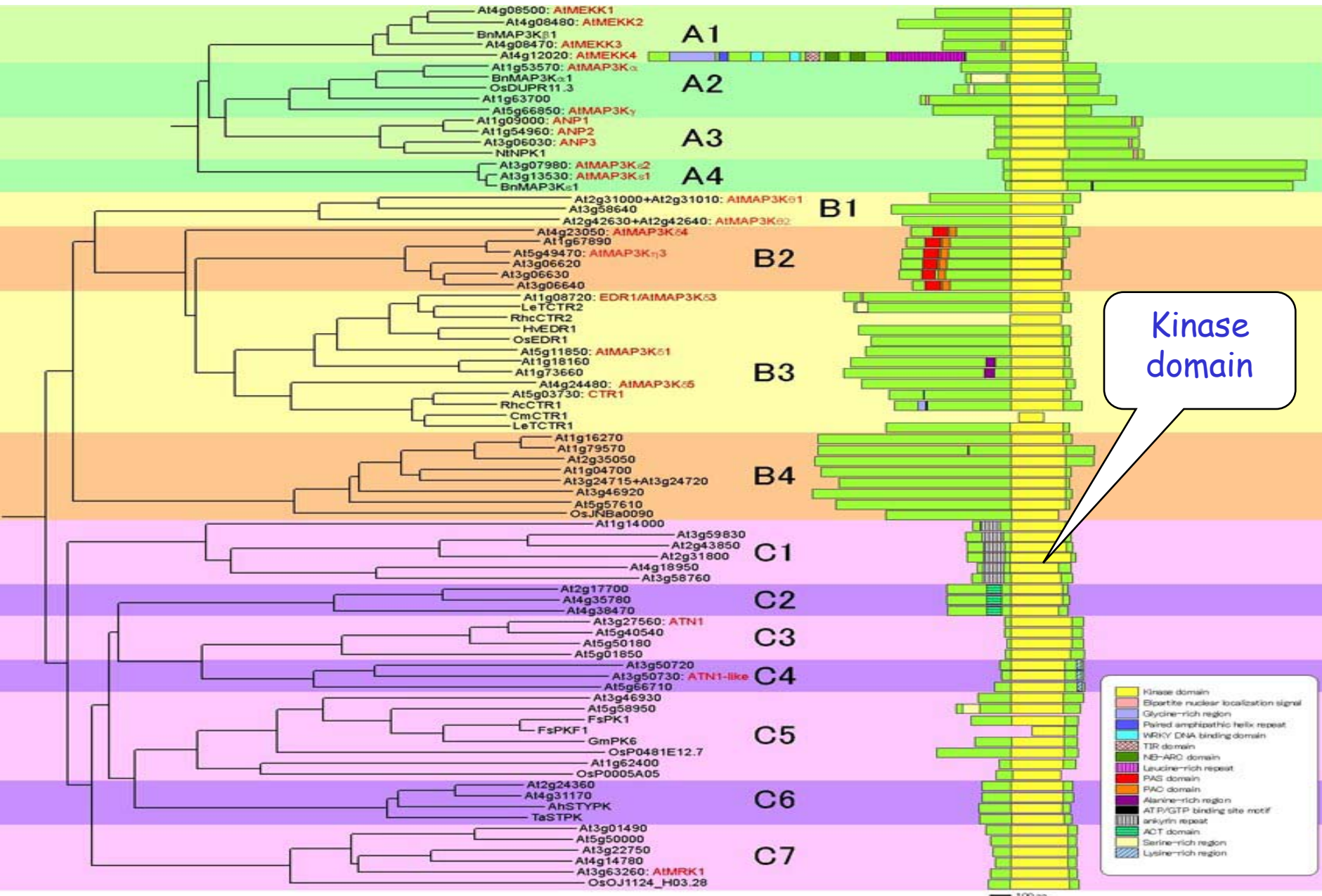
Example



Multiple sequence alignment of 7 neuroglobins using clustalx



Protein Phylogenies - Example



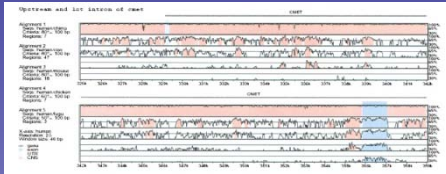
Motivation again

- Common structure, function or origin may be only weakly reflected in sequence - multiple comparisons may highlight weak signals
- Major uses:
 - Identify and represent protein families
 - Identify and represent conserved sequence or structure elements (e.g. domains)
 - Deduce evolutionary history

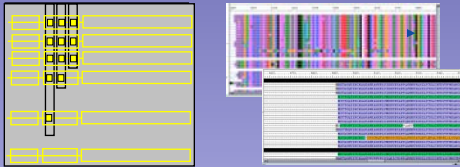


MSA : central role in biology

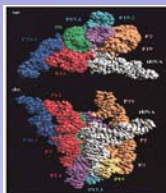
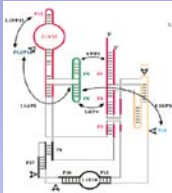
Comparative genomics



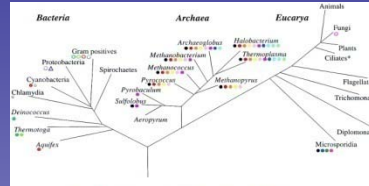
Gene identification, validation



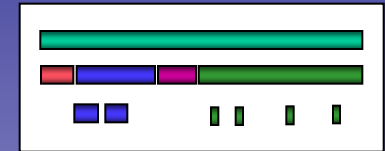
RNA sequence, structure, function



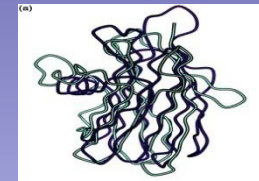
Phylogenetic studies



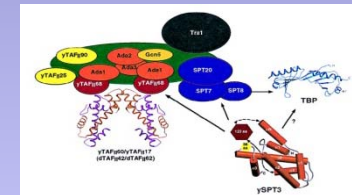
Hierarchical function annotation: homologs, domains, motifs



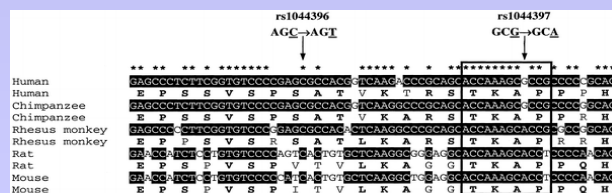
Structure comparison, modelling



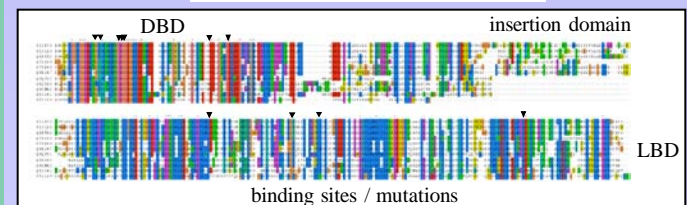
Interaction networks



Human genetics, SNPs



Therapeutics, drug design



Scoring alignments

- Given input seqs. S_1, S_2, \dots, S_k find a multiple alignment of optimal score
- Scores preview:
 - Sum of pairs
 - Consensus
 - Tree



Sum of Pairs score

Def: Induced pairwise alignment

A pairwise alignment induced by the multiple alignment

Example:

x: AC-GCGG-C
y: AC-GC-GAG
z: GCCGC-GAG

Induces:

x: ACGCGG-C; **x:** AC-GCGG-C; **y:** AC-GCGAG
y: ACGC-GAC; **z:** GCCGC-GAG; **z:** GCCGCGAG

$$S(M) = \sum_{k < l} \sigma(S'_k, S'_l)$$



SOP Score Example

Consider the following alignment:

A	C	-	C	D	B	-
-	C	-	A	D	B	D
A	-	B	C	D	A	D

Scoring scheme: match - 0
 mismatch/indel - -1

SP score: -3 -5 -4 = -12



Alignments = Paths

- Align 3 sequences: *ATGC*, *AATC*, *ATGC*

	A	--	T	G	C
--	---	----	---	---	---

	A	A	T	--	C
--	---	---	---	----	---

	--	A	T	G	C
--	----	---	---	---	---



Alignment Paths

0	1	1	2	3	4
	A	--	T	G	C

x coordinate

	A	A	T	--	C
--	---	---	---	----	---

	--	A	T	G	C
--	----	---	---	---	---



Alignment Paths

- Align 3 sequences: *ATGC*, *AATC*, *ATGC*

0	1	1	2	3	4
	A	--	T	G	C

x coordinate

0	1	2	3	3	4
	A	A	T	--	C

y coordinate

	--	A	T	G	C
--	----	---	---	---	---



Alignment Paths

0	1	1	2	3	4
	A	--	T	G	C
0	1	2	3	3	4
	A	A	T	--	C
0	0	1	2	3	4
	--	A	T	G	C

x coordinate

y coordinate

z coordinate

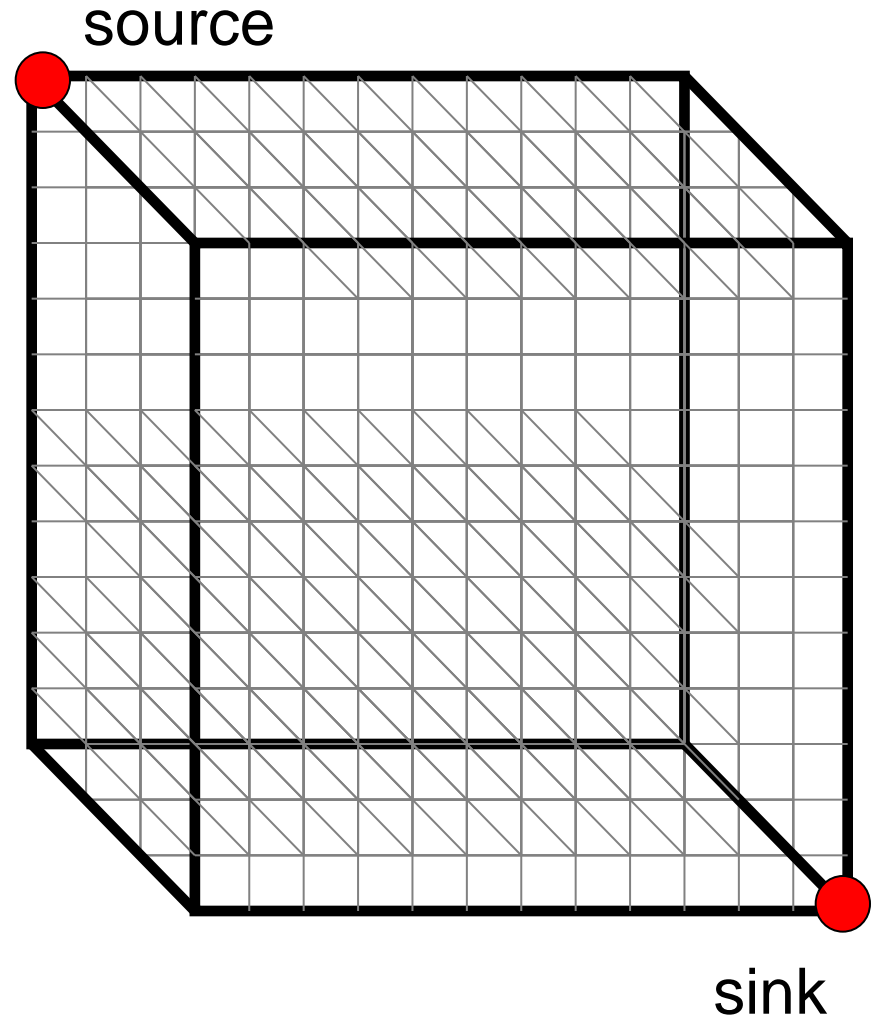
- Resulting path in (x,y,z) space:

$(0,0,0) \rightarrow (1,1,0) \rightarrow (1,2,1) \rightarrow (2,3,2) \rightarrow (3,3,3) \rightarrow (4,4,4)$

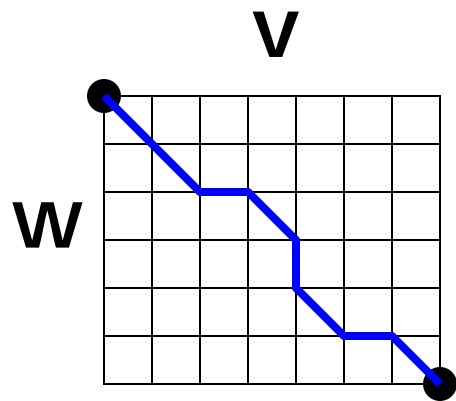


Aligning Three Sequences

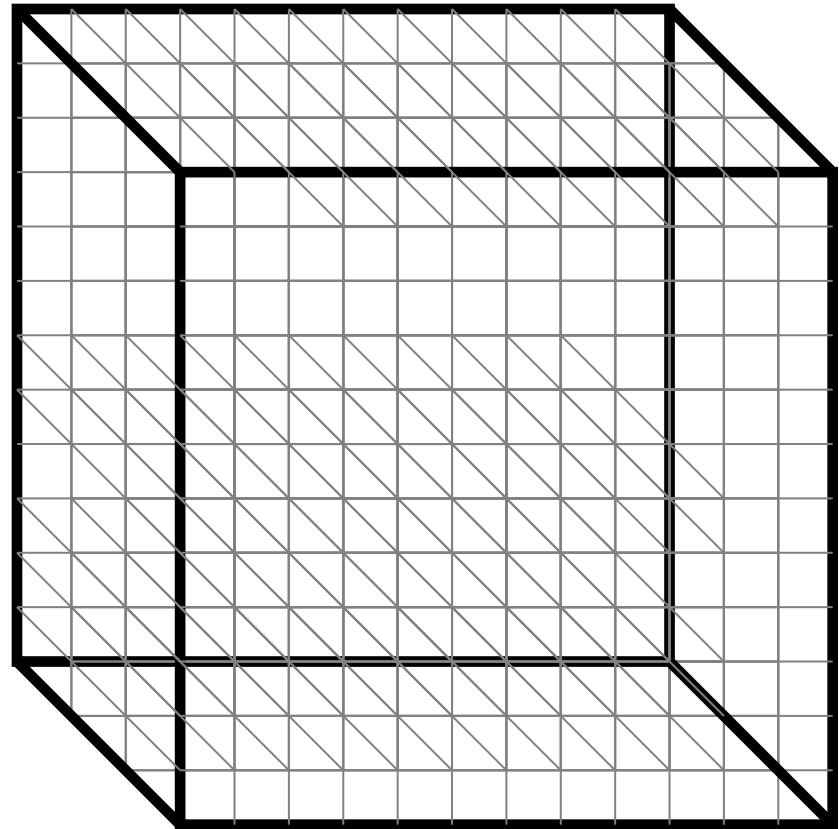
- Same strategy as aligning two sequences
- Use a 3-D "Manhattan Cube", with each axis representing a sequence to align
- For global alignments, go from source to sink



2-D vs 3-D Alignment Grid



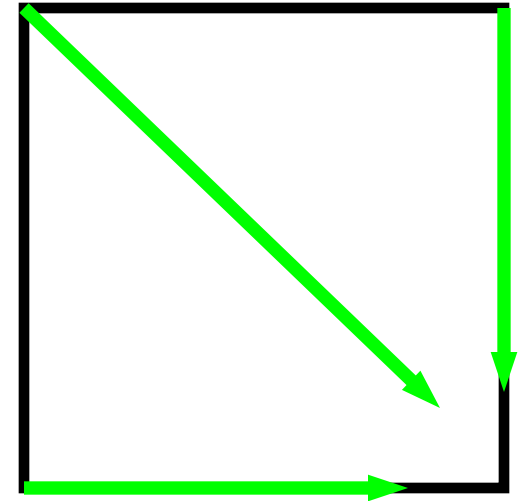
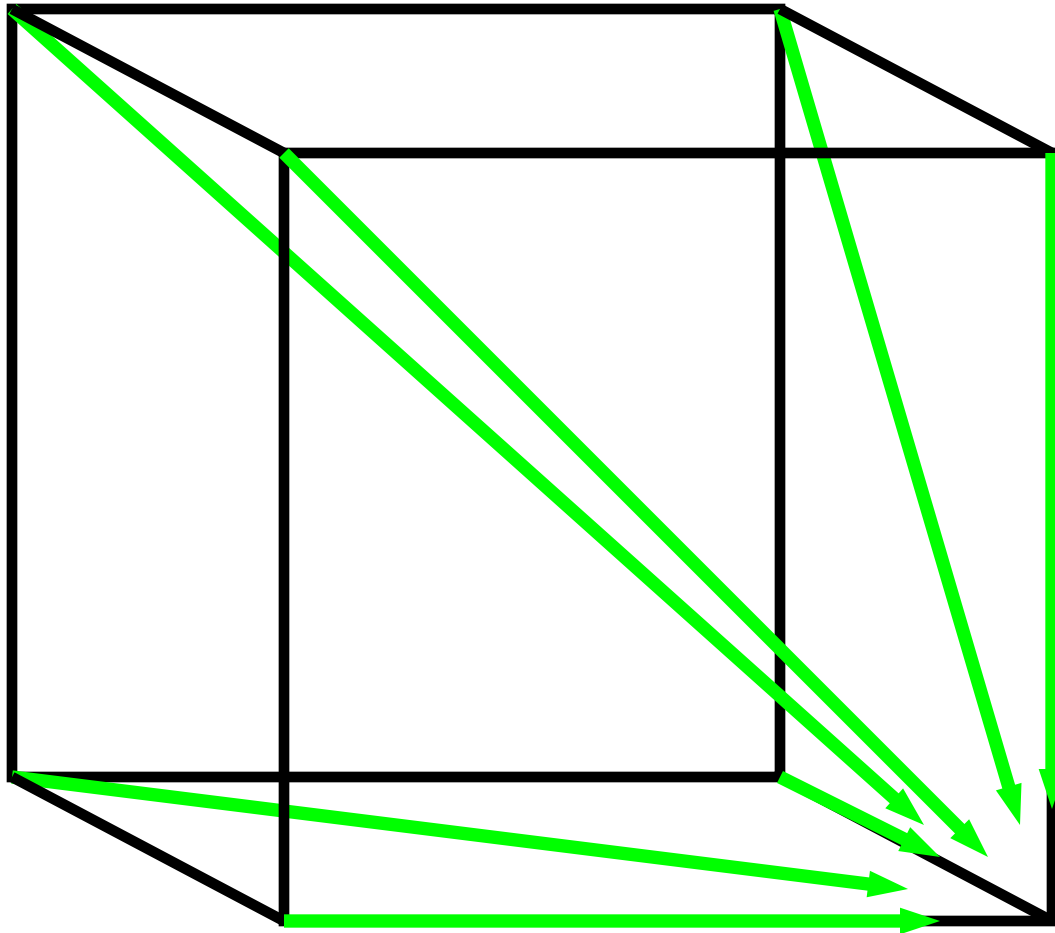
2-D edit graph



3-D edit graph



2-D cell versus 2-D Alignment Cell

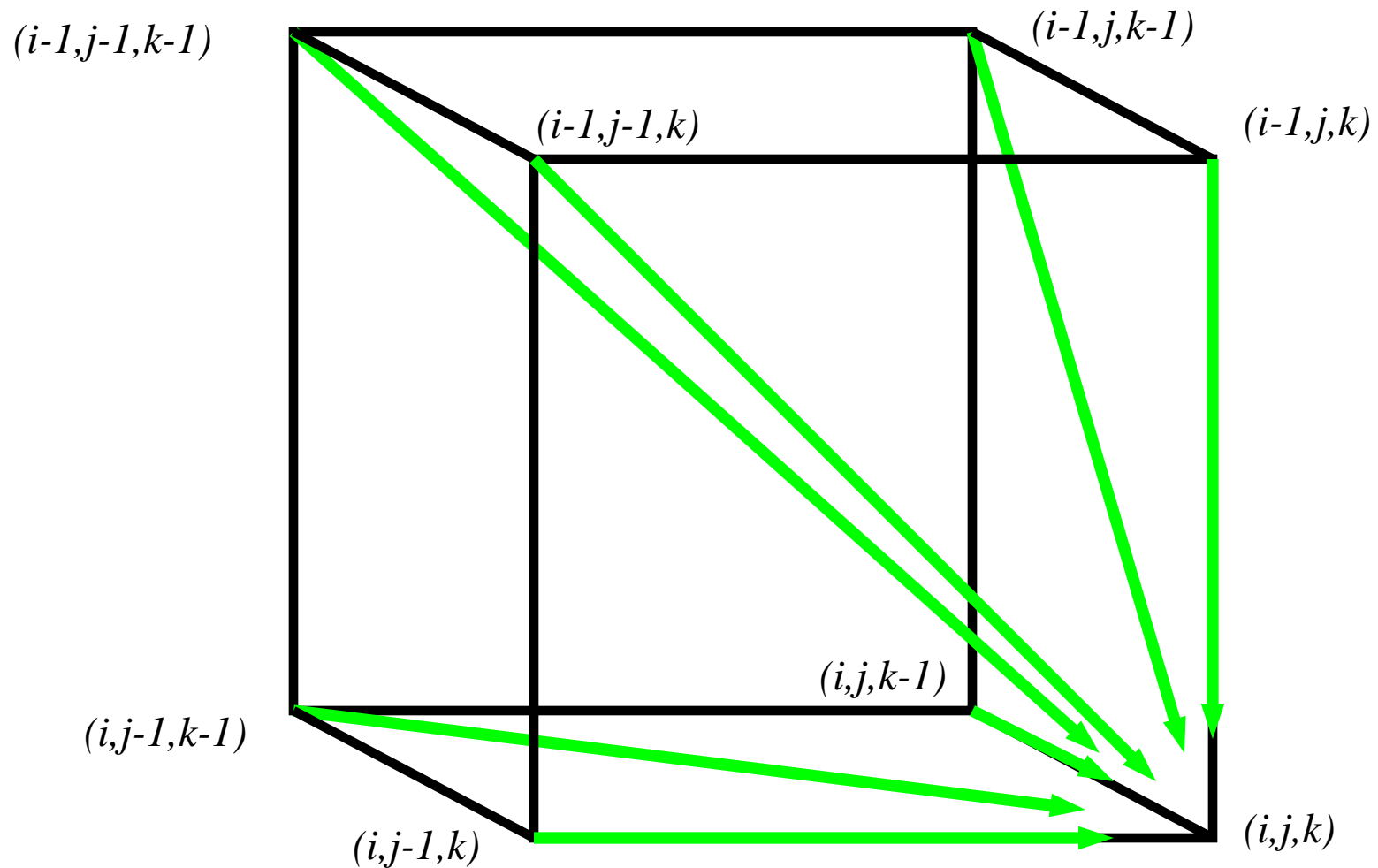


In **2-D**, 3 edges
in each unit
square

In **3-D**, 7 edges
in each unit cube



Architecture of 3-D Alignment Cell



Multiple Alignment: Dynamic Programming

- $s_{i,j,k} = \max \left\{ \begin{array}{ll} s_{i-1,j-1,k-1} + \delta(v_i, w_j, u_k) & \text{cube diagonal: no indels} \\ s_{i-1,j-1,k} + \delta(v_i, w_j, _) & \\ s_{i-1,j,k-1} + \delta(v_i, _, u_k) & \\ s_{i,j-1,k-1} + \delta(_, w_j, u_k) & \text{face diagonal: one indel} \\ s_{i-1,j,k} + \delta(v_i, _, _) & \\ s_{i,j-1,k} + \delta(_, w_j, _) & \\ s_{i,j,k-1} + \delta(_, _, u_k) & \text{edge diagonal: two indels} \end{array} \right\}$
- $\delta(x, y, z)$ is an entry in the 3-D scoring matrix



Running Time

- For 3 sequences of length n , the run time is $O(n^3)$
- For k sequences, build a k -dimensional cube, with run time $O(2^k n^k)$ [another 2^k factor for affine gaps]
- Impractical for most realistic cases
- NP-hard (Elias'03 for general matrices)



Minimum cost - SOP

We use min **cost** instead of max **score**

→ Find alignment of minimal cost

Assumption (for the approx. algs to follow):

the cost function **δ** is a **distance function**

- $\delta(x,x) = 0$ (also for gaps)
- $\delta(x,y) = \delta(y,x) \geq 0$
- $\delta(x,y) + \delta(y,z) \geq \delta(x,z)$ (triangle inequality)



Forward Dynamic Programming

- An alternative approach to DP for pairwise (and multiple) alignment:
- $D(v)$ - opt value of path $\text{source} \rightarrow v$
- $p(w)$ - best-yet solution of path $\text{source} \rightarrow w$
- When $D(v)$ is computed, send its value forward on the arcs from $v \rightarrow w$:
$$p(w) = \min\{p(w), D(v) + \text{cost}(v, w)\}$$
- Once $p(w)$ has been updated by all incoming edges - that value is optimal; set as $D(w)$



Forward Dynamic Programming (2)

- Maintain a queue of nodes whose D is not set yet
- For the node w at the head of the queue: Set $D(w) \leftarrow p(w)$ and remove
- \forall out-neighbor x of w - update p ; if x is not in the queue - add it at the end
 - Breaking ties lexicographically
 - Only x -s with some forward transmission are added to the queue
- Same complexity as the regular (backwards) DP



Faster DP Algorithm for MultiAlign

Carillo-Lipman 88

- Use forward DP.
- We'll demonstrate on three sequences
- $f_{12}(i,j)$ = opt pairwise alignment score of suffixes $S_1(i+1,..n_1)$, $S_2(j+1,..n_2)$, etc.
- Key idea: if \exists a known soln of cost z , if $D(i,j,k) + f_{12}(i,j) + f_{13}(i,k) + f_{23}(j,k) > z$
→ Do not send $D(i,j,k)$ forward
- Guarantees opt soln - no improved time bound, but often saves a lot in practice.



An approximation algorithm

Gusfield 93

- Compute a *center* string, minimizing the sum of pairwise distances to the other strings
- Use it as a “pivot” for the alignment

$D(S,T)$ - cost of minimum global alignment between S and T



The Center Star algorithm

Input: Γ - set of k strings S_1, \dots, S_k .

1. Find the string $S^* \in \Gamma$ (**center**) that minimizes $\sum_{S \in \Gamma \setminus \{S^*\}} D(S^*, S)$
2. Denote $S_1 = S^*$ and the rest of the strings as S_2, \dots, S_k
3. Iteratively add S_2, \dots, S_k to the alignment as follows:
 - a. Suppose S_1, \dots, S_{i-1} are already aligned as S'_1, \dots, S'_{i-1}
 - b. Optimally align S_i to S'_1 to produce S'_i and S''_1 aligned
 - c. Adjust S'_2, \dots, S'_{i-1} by adding spaces where spaces were added to S''_1
 - d. Replace S'_1 by S''_1



Inheriting gaps

x: AGAC
y: ATGA
z: ATGGA

← center

1: y: ATGA-
x: A-GAC

2: y: ATG-A
z: ATGGA

y: ATG-A-
x: A-G-AC
z: ATGGA-



Running time

- Choosing S_1 - execute DP for all sequence-pairs - $O(k^2 n^2)$
- Adding S_i to the alignment - execute DP for S_i, S'_1 - $O(i \cdot n^2)$.
(In the i^{th} stage the length of S'_1 can be up-to $i \cdot n$)

$$\sum_{i=1}^{k-1} O(i \cdot n^2) = O(k^2 n^2)$$

↑
total complexity



Approximation ratio

- M^* - An optimal alignment
- M - The alignment produced by this algorithm
- $d(i,j)$ - The distance M induces on the pair S_i, S_j
-

$$v(M) = \sum_{i=1}^k \sum_{\substack{j=1 \\ j \neq i}}^k d(i, j) = 2 \sum_{i < j} d(i, j)$$

- recall $D(S, T)$ - min cost of alignment between S and T

For all i : $d(1, i) = D(S_1, S_i)$

(we perform optimal alignment between S'_1 and S_i and $\delta(-, -) = 0$)



Approximation ratio (2)

Triangle
inequality +
symmetry

$$v(M) = \sum_{i=1}^k \sum_{\substack{j=1 \\ j \neq i}}^k d(i, j) \leq \sum_{i=1}^k \sum_{\substack{j=1 \\ j \neq i}}^k (d(1, i) + d(1, j))$$

$$= 2(k-1) \sum_{l=2}^k d(1, l) = 2(k-1) \sum_{l=2}^k D(S_1, S_l)$$

$$v(M^*) = \sum_{i=1}^k \sum_{\substack{j=1 \\ j \neq i}}^k d^*(i, j) \geq \sum_{i=1}^k \sum_{\substack{j=1 \\ j \neq i}}^k D(S_i, S_j) \geq$$

$$\geq \sum_{i=1}^k \sum_{j=2}^k D(S_1, S_j) = k \sum_{j=2}^k D(S_1, S_j)$$

$$\frac{v(M)}{v(M^*)} \leq \frac{2(k-1)}{k} \leq 2$$

Definition of S_1 :

$$\forall i: \sum_{j=2}^k D(S_1, S_j) \leq \sum_{\substack{j=1 \\ j \neq i}}^k D(S_i, S_j)$$

Theorem (Gusfield 93)



- We have proved:
- The center star algorithm is a polynomial algorithm that guarantees a solution at most twice the optimum.
- "a 2-approximation"
- "an approximation ratio of 2"



Steiner/consensus string

- Input: Γ - set of k strings S_1, \dots, S_k .
- $D(X, Y)$ - score of aligning X, Y .
- S - arbitrary sequence (unrelated to Γ)
- The **consensus error** of S relative to Γ :
$$E(S) = \sum_{i \leq k} D(S, S_i)$$
- S^* is an optimal **Steiner string** for Γ if it minimizes $E(S)$
- Different objective function - linear no of terms
- No direct relation to multialign! (for now)

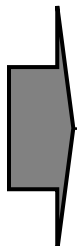


Optimal Steiner String: Approximation

Thm: Assume D satisfies triangle ineq. Then
 $\exists S \in \Gamma$ that gives a 2-approximation.

Pf: Pick $S \in \Gamma$

$$\begin{aligned} E(S) &= \sum_{S \neq S_i} D(S, S_i) \leq \sum_{S \neq S_i} (D(S, S^*) + D(S^*, S_i)) \\ &= (k-2)D(S, S^*) + D(S, S^*) + \sum_{S_i \neq S} D(S^*, S_i) \\ &= (k-2)D(S, S^*) + E(S^*) \end{aligned}$$

Pick $S \in \Gamma$ closest to S^* (not constructively)  $\frac{E(S)}{E(S^*)} \leq \frac{(k-2)}{k} + 1 < 2$

$$E(S^*) = \sum_{S_i \in \Gamma} D(S^*, S_i) \geq k \cdot D(S, S^*)$$



Consensus string from MA

- The **consensus string** of a MA is obtained by taking the most frequent character in each position

- S^* : AC-GC-GAG
- x: AC-GCGG-C
- y: AC-GC-GAG
- z: GCCGA-GAG
- u: AC-T-GGCA
- v: -CAGT-GAG
- w: AC-GC-GAG

Alignment error: $S(M) = \sum_k \sigma(S'_k, S^*)$

The **opt consensus MA** of the set of input sequences Γ : MA with least alignment error



Consensus multiple alignment

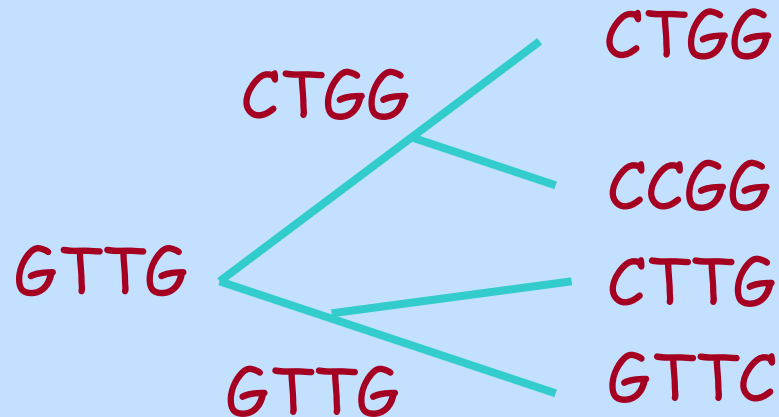
Thm: opt soln of consensus MA =
Steiner string (up to spaces)

.



Tree Multiple Alignment

- Input: Tree T , a string for each leaf
- **Phylogenetic (also Tree) alignment** for T :
Assignment of a string to each internal node



- Score - (weighted) sum of scores along edges
- Goal: find tree alignment of optimal score
- Consensus = tree alignment where T is a star



Tree MA - complexity

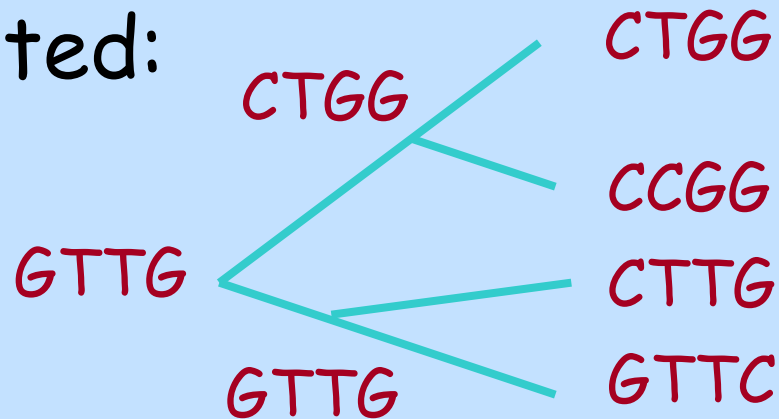
- NP-hard
- Poly time approximations:
 - 2-approximation
 - Better approximation with more time (PTAS)



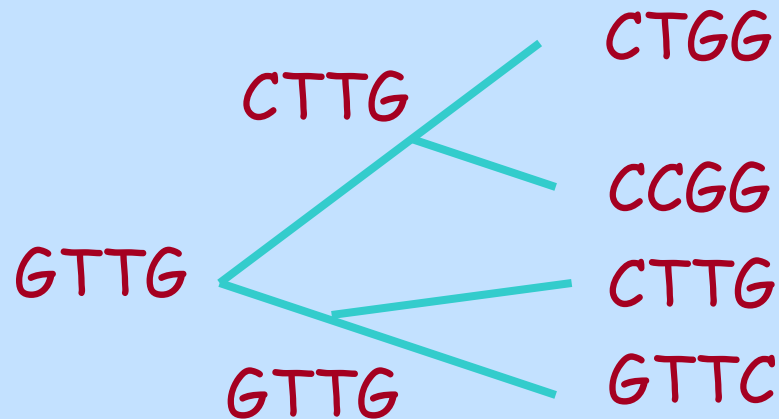
Lifted alignment

- The seq. label at every internal node is lifted from one of its children

- Lifted:



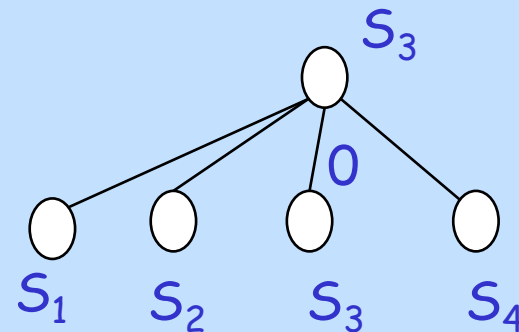
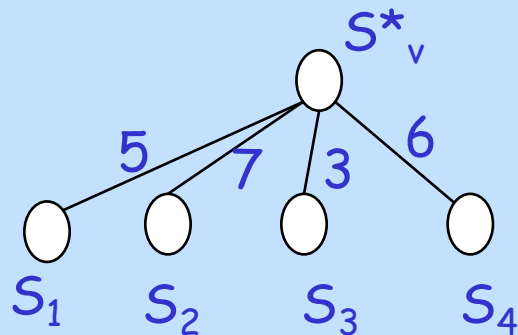
- Not lifted



A 2-approximation to Tree MSA

[Jiang, Wang, Lawler 1996]

- Assumes triangle inequality
- (non constructively) Transform an optimal tree T^* to a lifted alignment T^L :
 - At each internal node v , assign seq. of a child that is closest to the optimal label of v



- Claim: T^L has at most twice the distance of T^*



Pf sketch: $\text{cost}(T^L) \leq 2 \text{cost}(T^*)$

- In T^L , take $e=(v,w)$, $v=\text{Pa}(w)$ with labels S_j for v , S_i for w , $S_i \neq S_j$
 - $D(S_j, S_i) \leq D(S_j, S_v^*) + D(S_v^*, S_i) \leq 2D(S_i, S_v^*)$
(why?)
 - Path P_e from leaf labeled S_i up to v has cost:
 - $D(S_j, S_i)$ in T^L
 - At least $D(S_v^*, S_i)$ in T^*
- Paths $\{P_e\}$ are edge disjoint and cover all nonzero edges in T^L



Dynamic Programming alg for optimal lifted alignment

- $d(v, S)$ - distance of the best lifted alignment of T_v s.t. string S is assigned to node v

$$d(v, S) = \sum_w \min_T [D(S, T) + d(w, T)]$$

here w - child of v , T - string at a leaf of T_w

- Complexity: k = no of leaves, tot length N
 - Compute all pairwise leaf distances in $O(N^2)$
 - Computation per internal node: $O(k^2)$
 - $\rightarrow O(N^2 + k^3)$ (can do $O(N^2 + k^2)$)



Wrapping up lifted alignment

- \exists a lifted alignment that is $\leq 2 \text{ OPT}$
- We can find the min cost lifted alignment in poly time
- That alignment is also $\leq 2 \text{ OPT}$
- \Rightarrow Thm: lifted alignment alg gives a poly-time 2-approximation to Tree Alignment



Profile Representation of MA

-	A	G	G	C	T	A	T	C	A	C	C	T	G
T	A	G	-	C	T	A	C	C	A	-	-	-	G
C	A	G	-	C	T	A	C	C	A	-	-	-	G
C	A	G	-	C	T	A	T	C	A	C	-	G	G
C	A	G	-	C	T	A	T	C	G	C	-	G	G

A		1				1			.8				
C	.6			1			.4	1		.6	.2		
G			1	.2					.2			.4	1
T	.2				1		.6					.2	
-	.2		.8							.4	.8	.4	

- Alternatively, use log odds:
- $p_i(a)$ = fraction of a's in col i
- $p(a)$ = fraction of a's overall
- $\log p_i(a)/p(a)$



Aligning a sequence to a profile

- Key in pairwise alignment is scoring two positions x, y : $\sigma(x, y)$
- For a letter x and a column C in a profile, $\sigma(x, C)$ = value of x in col. C
- Invent a score for $\sigma(x, -)$
- Run the DP alg for pairwise alignment



Aligning alignments

- Given two alignments, how can we align them?
- Hint: use DP on the corresponding profiles.

x	GGGCACTGCAT	
y	GGTTACGTC--	
z	GGGAACTGCAG	

Alignment 1

w	GGACGTACC--	
v	GGACCT-----	

Alignment 2

x	GGGCACTGCAT
y	GGTTACGTC--
z	GGGAACTGCAG
w	GGACGTACC--
v	GGACCT-----



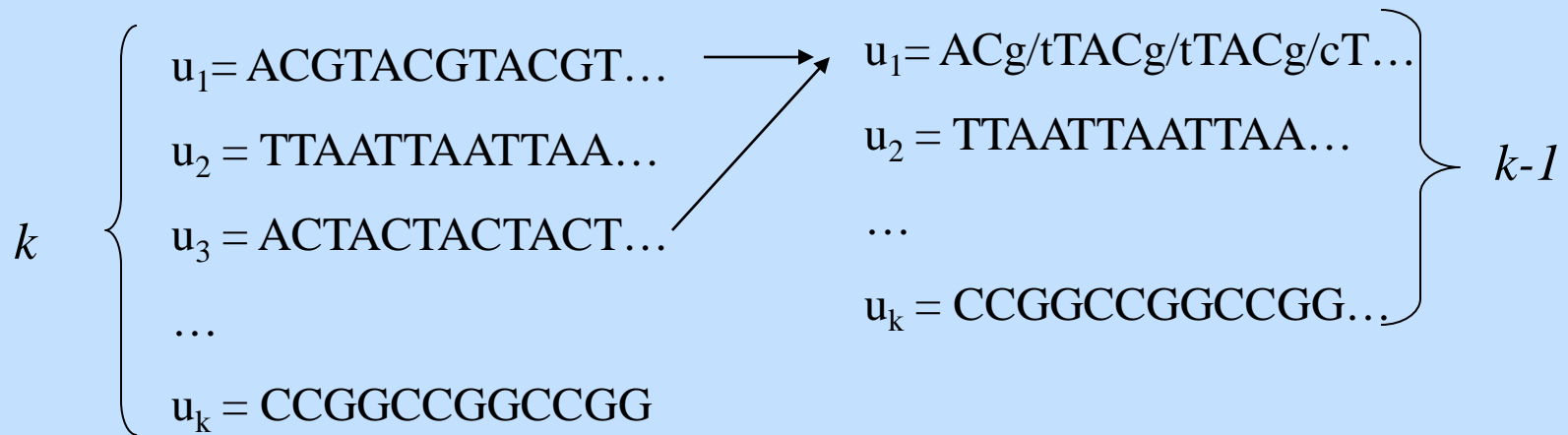
Profile-profile scoring

- Fix a position in the alignment
 - p_i - prob (i in 1st profile); q_i - in 2nd profile
- Expected score: $\sum_{ij} p_i q_j s_{ij}$
- Euclidean distance
- Pearson correlation
- KL-divergence (relative entropy)
- ...



Multiple Alignment: Greedy Heuristic

- Choose most similar pair of sequences and combine into a profile, thereby reducing alignment of k sequences to an alignment of $k-1$ sequences/profiles. **Repeat**

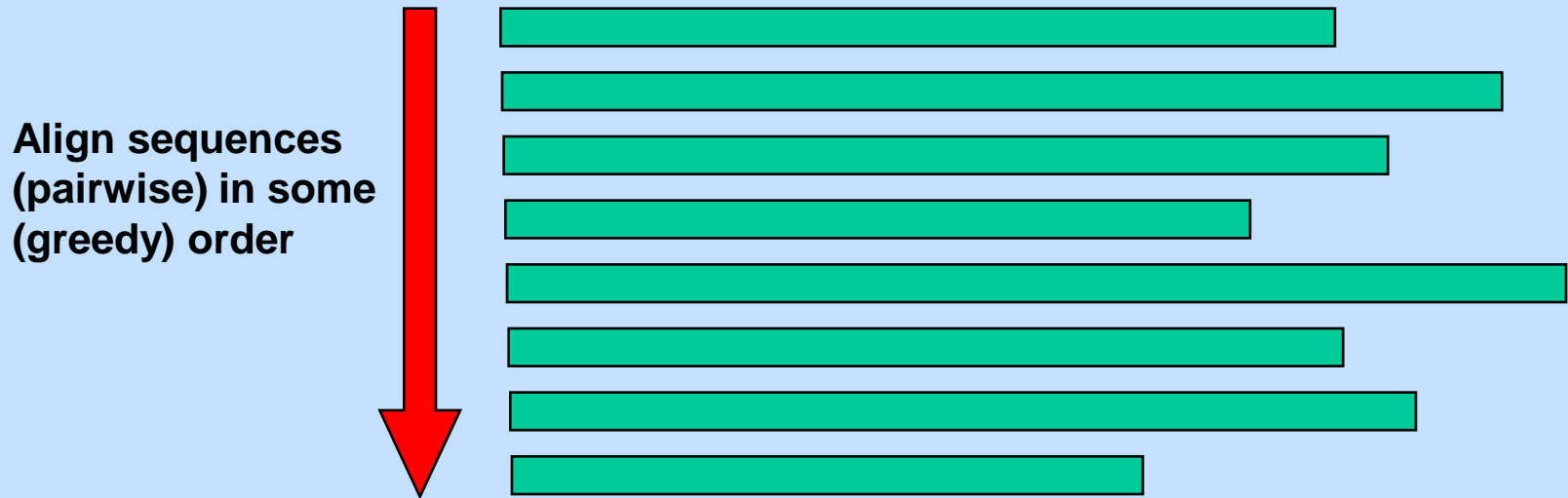


Progressive Alignment

- *A variation of greedy algorithm with a somewhat more intelligent strategy for choosing the order of alignments.*



Progressive alignment

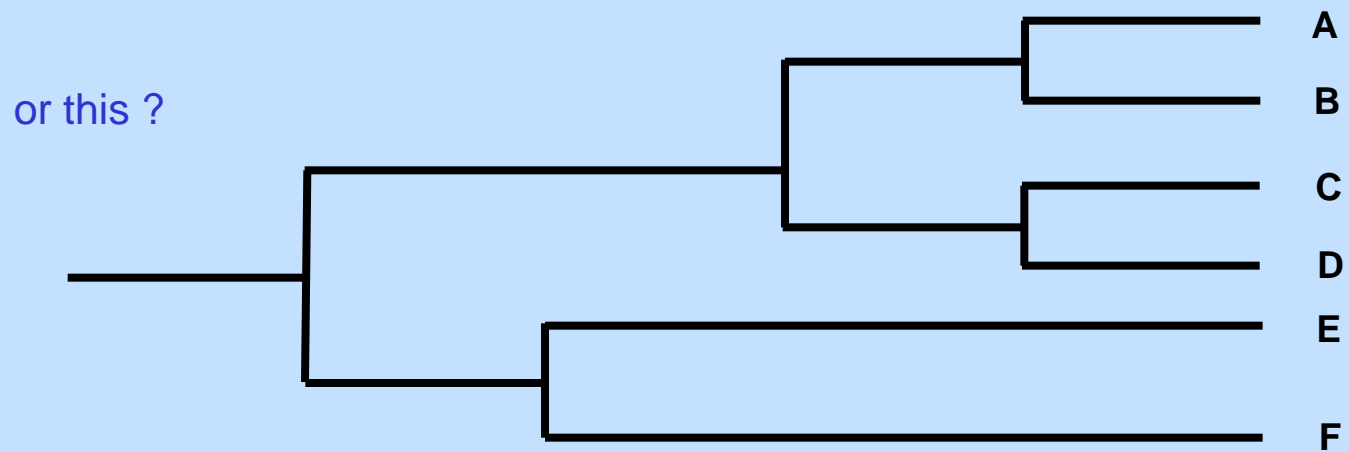
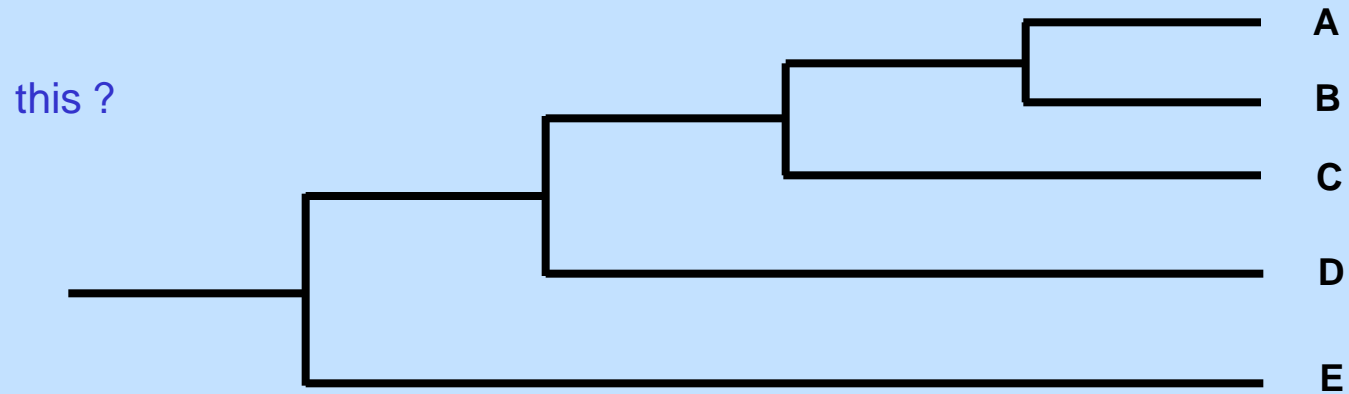


Decisions

- (1) Order of alignments
- (2) Alignment of sequence to group (only), or allow group to group
- (3) Method of alignment, and scoring function



Guide tree



Feng & Doolittle (1987)

Overview

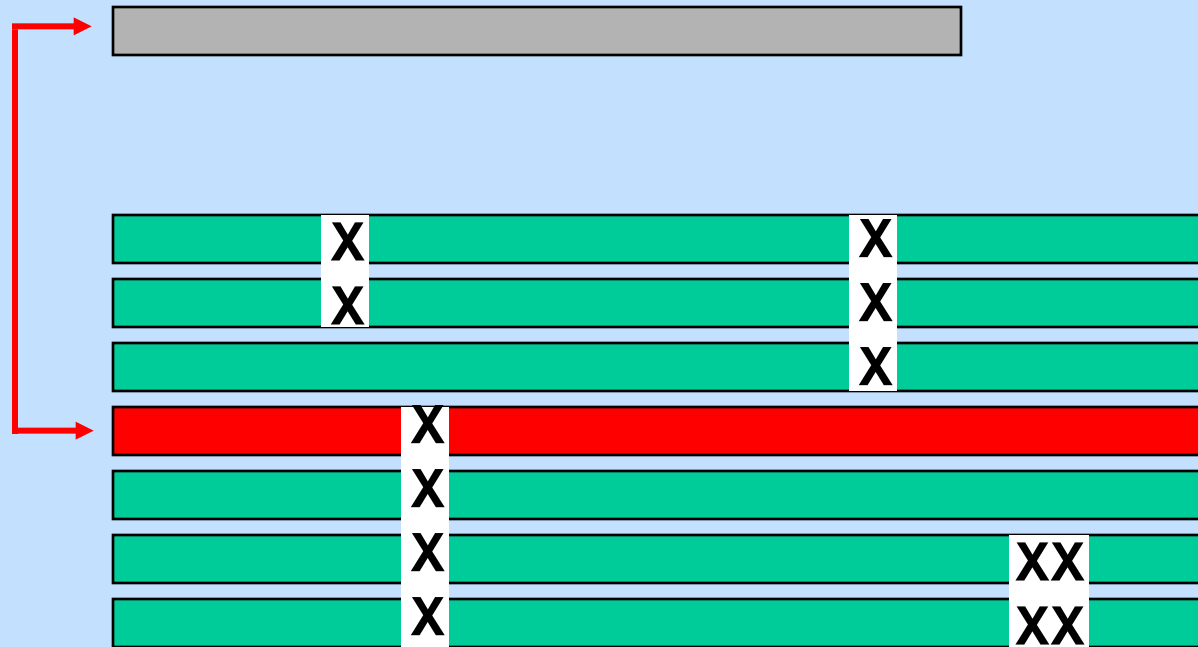
- (1) Calculate distances between all pairs of sequences by pairwise alignment; convert raw alignment scores to approximate **pairwise "distances"**.
- (2) Construct **guide tree** from the distance matrix by using a tree building (clustering) alg.
- (3) Starting from first node added to the tree, align the child nodes (which may be two sequences, a sequence and an alignment, or two alignments). Repeat for all other nodes in the order that they were added to tree, until all sequences have been aligned.



Feng & Doolittle (1987)

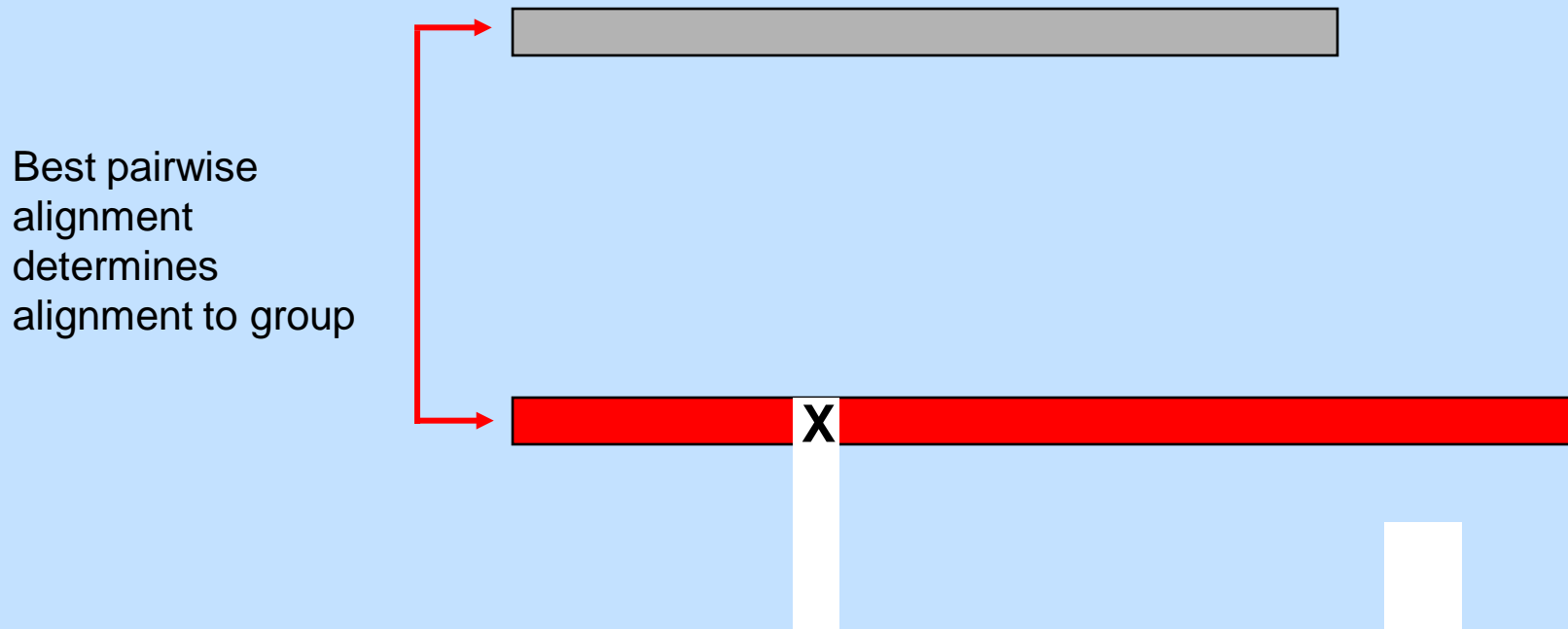
sequence-to-group

Best pairwise
alignment
determines
alignment to group



Feng & Doolittle (1987)

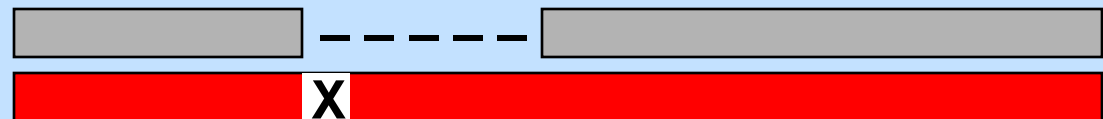
sequence-to-group



Feng & Doolittle (1987)

sequence-to-group

Best pairwise
alignment
determines
alignment to group



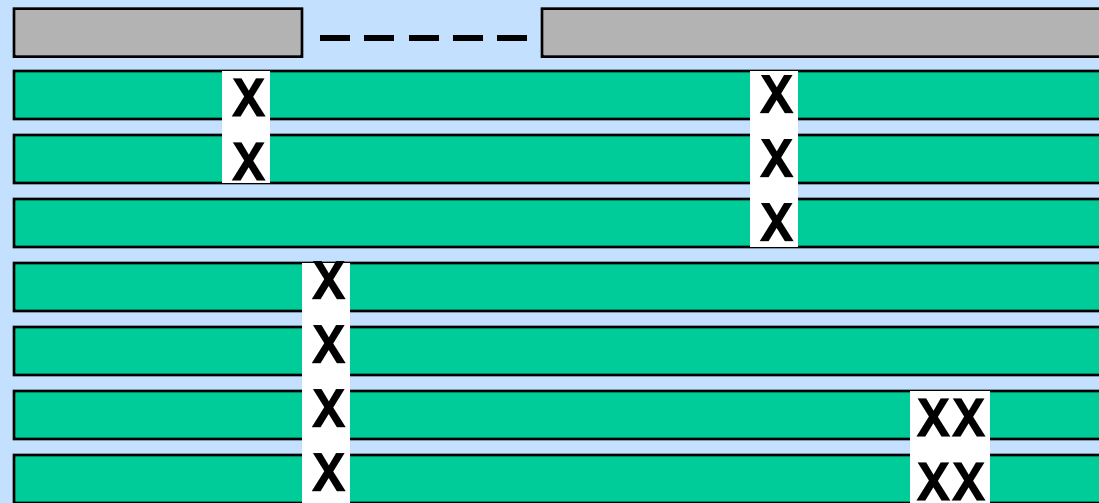
A space in the grey seq in
this column is encouraged
because it has no cost



Feng & Doolittle (1987)

sequence-to-group

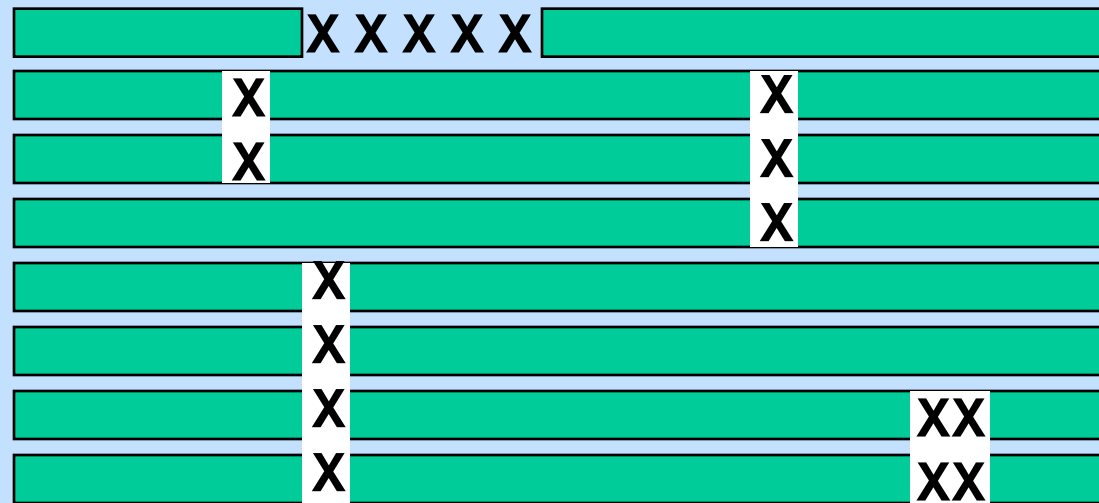
Best pairwise
alignment
determines
alignment to group



Feng & Doolittle (1987)

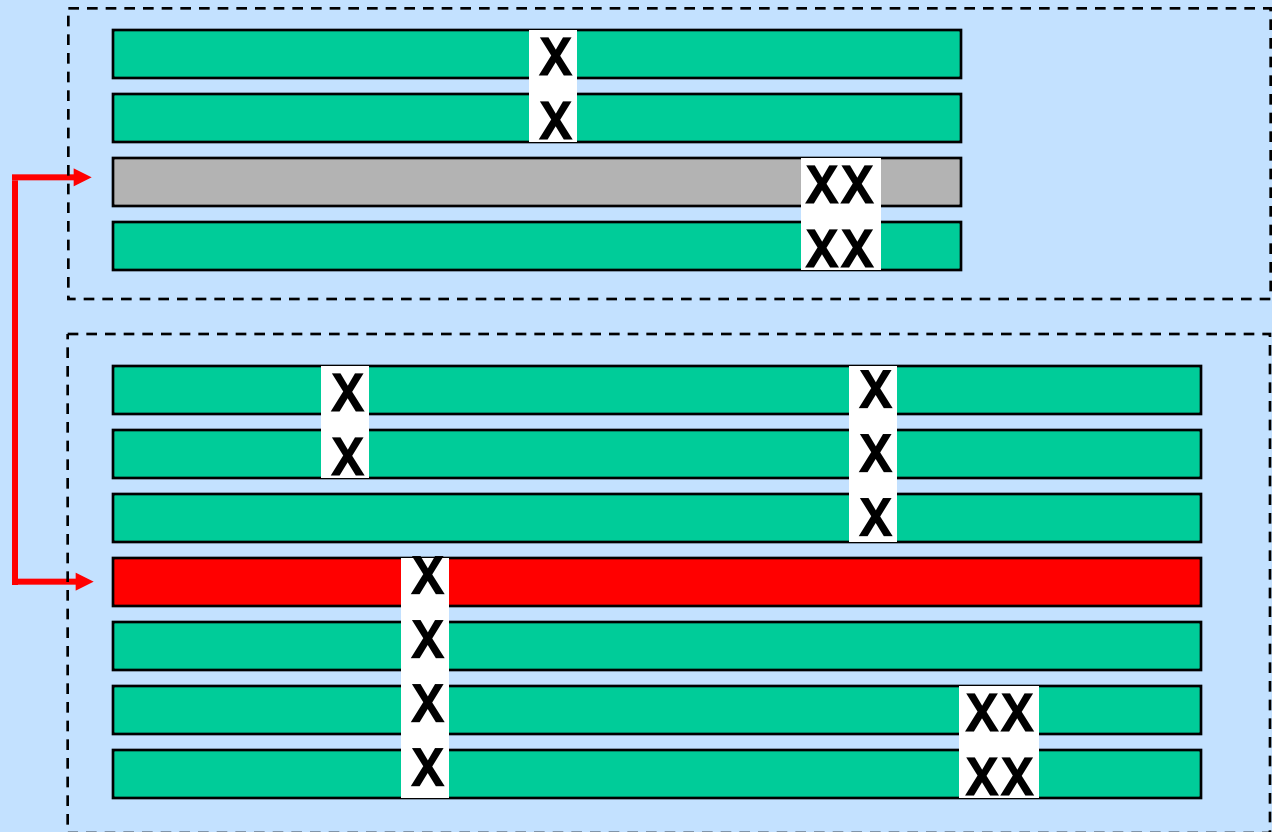
sequence-to-group

Best pairwise alignment determines alignment to group



Feng & Doolittle (1987)

group-to-group



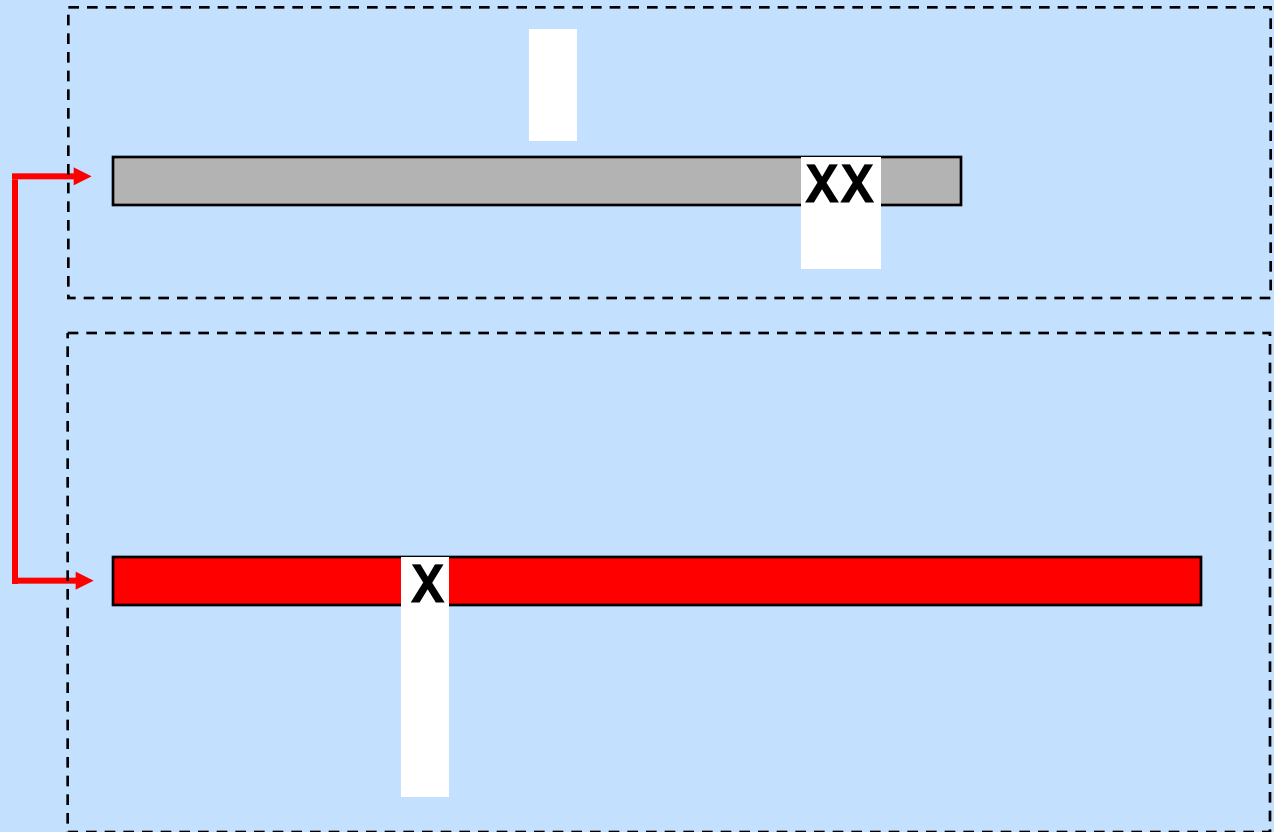
Best pairwise
alignment
determines
alignment of
groups



Feng & Doolittle (1987)

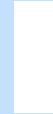
group-to-group

Best pairwise
alignment
determines
alignment of
groups

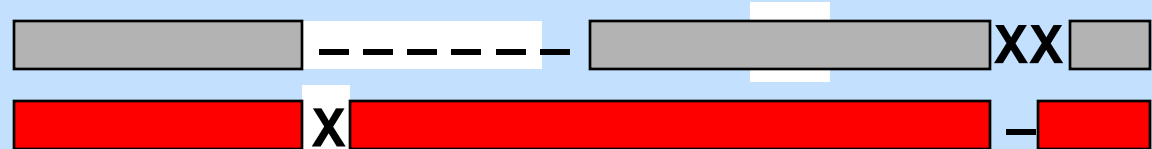


Feng & Doolittle (1987)

group-to-group

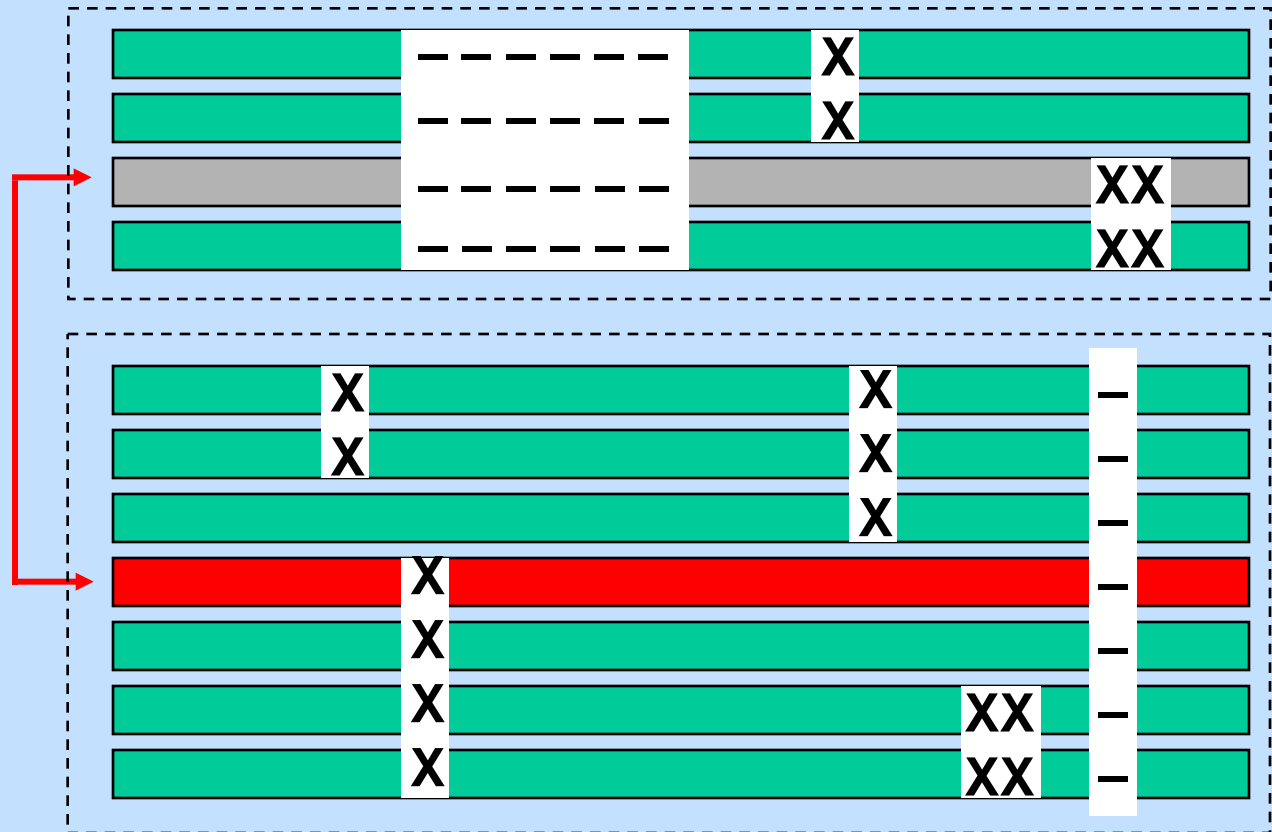


Best pairwise
alignment
determines
alignment of
groups



Feng & Doolittle (1987)

group-to-group



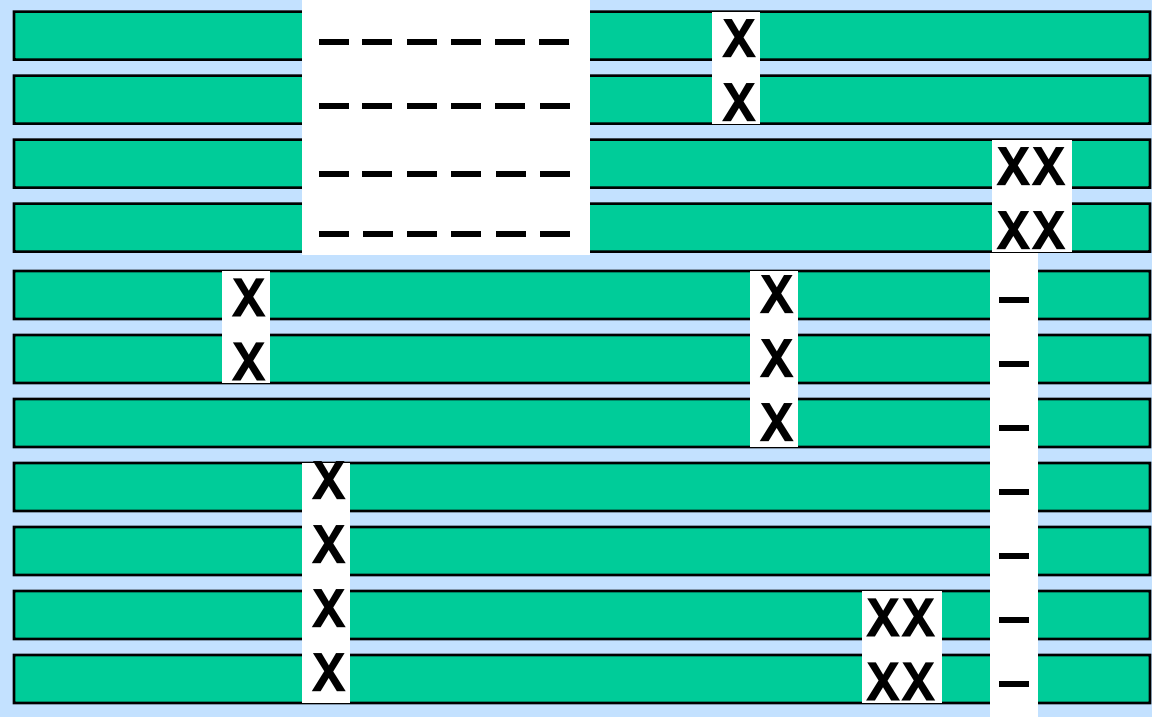
Best pairwise
alignment
determines
alignment of
groups



Feng & Doolittle (1987)

group-to-group

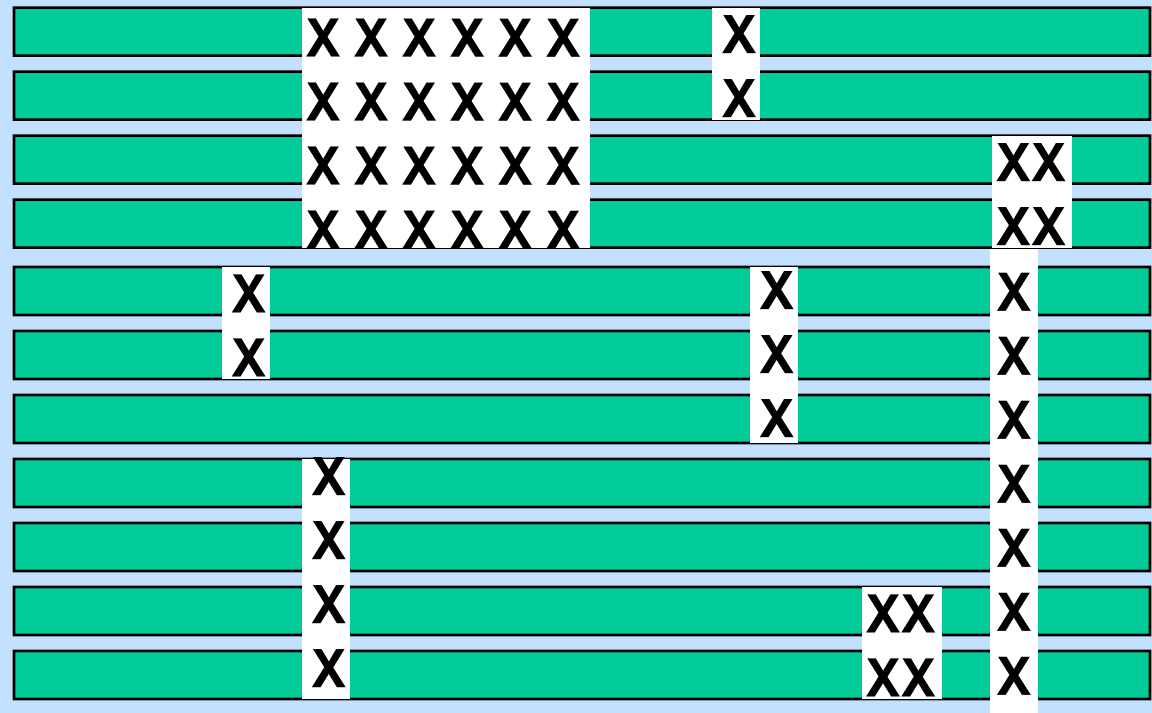
Best pairwise
alignment
determines
alignment of
groups



Feng & Doolittle (1987)

group-to-group

Best pairwise
alignment
determines
alignment of
groups



Feng & Doolittle (1987)

"Once a gap, always a gap".

Encourages gaps to occur in same columns in subsequent alignments.

This early version of progressive alignment is based on pairwise alignments. Later methods use profile-profile alignments.



ClustalW

Thompson, Higgins, Gibson 94

- Popular multiple alignment tool today
- Three-step process
 - 1.) Construct pairwise alignments
 - 2.) Build guide tree
 - 3.) Progressive alignment guided by the tree



Step 1: Pairwise Alignment

- Aligns each pair of sequences, giving a similarity matrix
- Similarity = exact matches / sequence length (percent identity)

	v_1	v_2	v_3	v_4
v_1	—			
v_2	.17	—		
v_3	.87	.28	—	
v_4	.59	.33	.62	—

(.17 means 17 % identical)



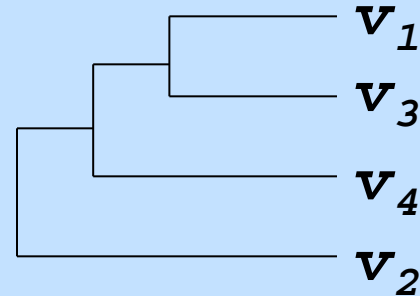
Step 2: Guide Tree

- Use the similarity method to create a guide tree by applying some clustering method*
- Guide tree roughly reflects evolutionary relations
 - *ClustalW uses the neighbor-joining method (to be described later in the course)



Step 2: Guide Tree (cont'd)

	v_1	v_2	v_3	v_4
v_1	—			
v_2	.17	—		
v_3	.87	.28	—	
v_4	.59	.33	.62	—



Calculate:

$$\begin{aligned}
 & \mathbf{v}_{1,3} &= \text{alignment } (v_1, v_3) \\
 & \mathbf{v}_{1,3,4} &= \text{alignment } ((v_{1,3}), v_4) \\
 & \mathbf{v}_{1,2,3,4} &= \text{alignment } ((\mathbf{v}_{1,3,4}), v_2)
 \end{aligned}$$



Step 3: Progressive Alignment

- Start by aligning the two most similar sequences
- Using the guide tree, add in the most similar pair (seq-seq, seq-prof or prof-prof)
- Insert gaps as necessary
- Many ad-hoc rules: weighting, different matrices, special gap scores....

FOS_RAT	PEEMSVTS-LDLTGGLPEATTPESEEAFTLPLLNDPEPK-PSLEPVKNISNMELKAEPFD
FOS_MOUSE	PEEMSVAS-LDLTGGLPEASTPESEEAFTLPLLNDPEPK-PSLEPVKSISNVELKAEPFD
FOS_CHICK	SEELAAATALDLG-----APSPAAAEFAFALPLMTEAPPAVPPKEPSG--SGLELKAEPFD
FOSB_MOUSE	PGPGPLAEVRDLPG-----STSAKEDGFGWLLPPPPPPP-----LPPFQ
FOSB_HUMAN	PGPGPLAEVRDLPG-----SAPAKEDGFSWLLPPPPPPP-----LPPFQ
	. . : ** . :.. *:.* * . *



Dots and stars show how well-conserved a column is.



Multiple Alignment: History

1975 Sankoff

*Formulated multiple alignment problem
and gave dynamic programming solution*

1988 Carrillo-Lipman

Branch and Bound approach for MSA

1990 Feng-Doolittle

Progressive alignment

1994 Thompson-Higgins-Gibson-ClustalW ← *> 40K citations!*

Most popular multiple alignment program

1998 Morgenstern et al. - DIALIGN

Segment-based multiple alignment

2000 Notredame-Higgins-Heringa-T-coffee

Using the library of pairwise alignments

2004 MUSCLE

2005 ProbCons

2011 Clustal Omega

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