Algorithm Analysis and Design

Name: Jiang Bowen **Due Date:** May 30, 2022. Student Number: 202117052 **Assignment:** Number 4

Problem In this assignment, you need to implement Qin's Algorithm (from Lecture 2) using Java or C/C++ (you need to use the package GMP for C/C++).

Solution RSA public key e and private key d is

$$ed \equiv 1 \pmod{\phi(n)}$$

There exists a positive integer k, the equation

$$ed - 1 = k \cdot \phi(n)$$

is satisfied. By dividing both sides by $d\phi(n)$, it turns

$$\frac{e}{\phi(n)} - \frac{k}{d} = \frac{1}{d\phi(n)}$$

If e is too big, approximately equal to n and $n \approx \phi(n)$, $e \approx \phi(n)$. Applying Wiener's results to the RSA Cryptosystem, we are certain to find one of the convergents of the continued fraction expansion of e/n, which is exactly equal to k/d.

Now the first step is that how to modify Qin's algorithm to calculate the convergents of continued fraction. To achieve this, a simple method is to add a new column in state matrix:

$$\begin{bmatrix} x_{11} & x_{12} & x_{13} \\ x_{21} & x_{22} & x_{23} \end{bmatrix}$$

 $\begin{bmatrix} x_{11} & x_{12} & x_{13} \\ x_{21} & x_{22} & x_{23} \end{bmatrix}$ where the initial state is $\begin{bmatrix} 1 & e & 1 \\ 0 & n & n/e \end{bmatrix}$. The entries x_{13} and x_{23} update as the red part of modify_qin_fun() (Algorithm 1, Appendix A). The numerator of continued fractions is x_{21} and the denominator is $x_{13} \mod x_{21}$.

The second step is to find the correct d from the entry x_{21} of Qin's state matrix. We test each denominator k'. If $k'|e \cdot x_{21} - 1$, it means that possible $\phi(n)$ can be calculated. Once n and $\phi(n)$ are known, we can factor n by solving the quadratic equation $x^2 - (n - 1)^2$ $\phi(n) + 1$ x + n which two roots are p and q. The result is possibly more than one, so we need to try each one until pq = n is satisfied. Blue part(Algorithm 1, Appendix A) and function $qet_pq()$ (Algorithm 2, Appendix A) implements this step.

The algorithm requires $O((\log_2 n)^4)$ steps so it is efficient. Finally, the procedure output d, p, q if attack successfully. The results are following and the full code is in Appendix B.

p = 1486713299424019561461244889335894217927531801918421416974272319123048047842301139582064854327076454620356056462712742806057865452106752623867 8306000706225490710033813203994580538080520612702355239318915736363373202 999835780729080942467

 $\begin{array}{l} \mathbf{q} = 160577100375614564018341557641357114288004878328768285449283860863028\\ 8567283929477948775814275770975955370361826425824326324412656409110830668\\ 3888187675774743571970459026047807501138224282313456944860989266461891195\\ 8144028778491530154987321607508754456707395614577095775911210216108867790\\ 257353074868685910047 \end{array}$

 $\begin{array}{lll} \textbf{d} &=& 187456565214517214244104553102686867789547551527992287744974433496534\\ 2842680392447827670891418993436667017805286135357879048826764659115334044\\ 918773405437 \end{array}$

Appendix A

```
Algorithm 1 modify_qin_fun(e, n)
```

```
Input: RSA public key e, n with 1 < e < n, \gcd(e, n) = 1
Output: return Attack Success if find RSA private key d, otherwise return Attack Fail
    \begin{bmatrix} x_{11} & x_{12} & x_{13} \\ x_{21} & x_{22} & x_{23} \end{bmatrix} \leftarrow
                                    1 e
                                    \begin{bmatrix} 0 & n & n/e \end{bmatrix}
   while x_{12} > 1 \text{ do}
         if x_{22} > x_{12} then
               q \leftarrow \lfloor \tfrac{x_{22}-1}{x_{12}} \rfloor
               r \leftarrow x_{22} - qx_{12}
               x_{21} \leftarrow qx_{11} + x_{21}
               x_{13} \leftarrow qx_{23} + x_{13}
               x_{22} \leftarrow r
         end if
         if x_{12} > x_{22} then
               q \leftarrow \lfloor \tfrac{x_{12}-1}{x_{22}} \rfloor
               r \leftarrow x_{12} - qx_{22}
               x_{11} \leftarrow qx_{21} + x_{11}
               x_{23} \leftarrow qx_{13} + x_{23}
               x_{12} \leftarrow r
         end if
         k' \leftarrow x_{13} \mod x_{21}
         d' \leftarrow x_{21}
         if k' is not zero and k'|(d'e-1) then
               \phi'(n) \leftarrow \frac{(d'e-1)}{k'}
               (p,q) \leftarrow get_{-}pq(\phi'(n),n)
               if p is not None and q is not None then
                    d \leftarrow d'
                    return Attack Success
               else
                    d \leftarrow None
               end if
         end if
   end while
```

Algorithm 2 get_pq($\phi'(n), n$)

```
Input: a possible \phi'(n) and RSA public key n

Output: decompose n into p and q. Return (p,q) if success, return (None, None)

b \leftarrow n - \phi'(n) + 1

\Delta \leftarrow b^2 - 4n

if \Delta \geq 0 then

p \leftarrow \frac{-b + \sqrt{\Delta}}{2}

q \leftarrow \frac{-b - \sqrt{\Delta}}{2}

if pq == n then

return (p,q)

end if

end if

return (None, None)
```

Appendix B

```
#include <stdio.h>
#include <stdlib.h>
#include <gmp.h>
#define ATTACK_SUCCESS 0
#define ATTACK_FAIL 1
int get_pq(mpz_t phi, mpz_t n){
   mpz_t b, par, t;
   mpz_t p, q;
   mpz_inits(b, par, t, NULL);
   mpz_inits(p, q, NULL);
    // b = n - phi + 1
   mpz_sub(b, n, phi);
   mpz_add_ui(b, b, 1);
    // par = b * b - 4 * n
   mpz_mul(par, b, b);
   mpz_mul_ui(t, n, 4);
   mpz_sub(par, par, t);
    if(mpz_cmp_ui(par, 0) >= 0){
        mpz_sqrt(par, par);
        // p = (-b + par) / 2
        mpz_sub(t, par, b);
        mpz_fdiv_q_ui(p, t, 2);
```

```
// q = (-b - par) / 2
        mpz_add(t, par, b);
        mpz_fdiv_q_ui(q, t, 2);
        mpz_mul_si(q, q, -1);
        mpz_abs(p, p);
        mpz_abs(q, q);
        mpz_mul(t, p, q);
        if(mpz_cmp(t, n) == 0) {
            gmp_printf("p = %Zd\n", p);
            gmp_printf("q = \frac{Zd}{n}", q);
            return ATTACK_SUCCESS;
        }
    }
    mpz_clears(b, par, t, NULL);
    mpz_clears(p, q, NULL);
    return ATTACK_FAIL;
}
int qin_fun(mpz_t a, mpz_t m){
    mpz_t x11, x12, x21, x22;
    mpz_t x13, x23;
    mpz_inits(x11, x12, x21, x22, x13, x23, NULL);
    mpz_set_str(x11, "1", 10);
    mpz_set(x12, a);
    mpz_set_str(x21, "0", 10);
    mpz_set(x22, m);
    mpz_set_str(x13, "1", 10);
    mpz_fdiv_q(x23, m, a);
    mpz_t t, q, r, k, phi;
    mpz_inits(t, q, r, k, phi, NULL);
    while(mpz_cmp_ui(x12, 1)) {
        if (mpz_cmp(x22, x12)) {
            // q = (x22 - 1) // x12
            // r = x22 - q * x12
            mpz_fdiv_qr(q, r, x22, x12);
            // x21 = q * x11 + x21
            mpz_mul(t, q, x11);
```

```
mpz_add(x21, t, x21);
            // c13 = q * x23 + x13
            mpz_mul(t, q, x23);
            mpz_add(x13, t, x13);
            // x22 = r
            mpz_set(x22, r);
        }
        if (mpz_cmp(x12, x22)) {
            mpz_fdiv_qr(q, r, x12, x22);
            // x21 = q * x21 + x11
            mpz_mul(t, q, x21);
            mpz_add(x11, t, x11);
            // x23 = q * x13 + x23
            mpz_mul(t, q, x13);
            mpz_add(x23, t, x23);
            // x12 = r
            mpz_set(x12, r);
        }
        mpz_mul(t, a, x21);
        mpz_sub_ui(t, t, 1);
        mpz_mod(k, x13, x21);
        if (mpz_cmp_ui(k, 0) == 0) {
            continue;
        }
        mpz_fdiv_qr(phi, t, t, k);
        if (mpz_cmp_ui(t, 0) == 0) {
            int ret = get_pq(phi, m);
            if (!ret) {
                gmp_printf("d = %Zd\n", x21);
                printf("attack success");
                return ATTACK_SUCCESS;
            }
        }
    }
    mpz_clears(x11, x12, x21, x22, x13, x23, NULL);
    mpz_clears(t, q, r, k, phi, NULL);
    return ATTACK_FAIL;
}
int main(int argc, char *argv[]) {
```

```
mpz_t e, n;
int ret = mpz_init_set_str(e,
"15126654297964752013098236737142792343328885821043850717591212219416
121876442886718732758175450185439250097909940465459448006358513001188
315695951467846403260337198869714837909361928327989272855362432071137
603123766827717346921751885958867720015317748270921597161083490383278
302493913553369743780995203398626939715733755155009681289940104874444
139325427474016608879888482717790115539038236875583889703891126293623
009944822130924615253652912319543400303803395910987309654804989665931
463997217772013338972458030372517101804468582533624114654185261016834
363449414265885477703830170220100277420183737256953449921432442769",
10):
if (ret != 0){
    perror("mpz for e error");
    exit(EXIT_FAILURE);
}
ret = mpz_init_set_str(n,
"23873211071137189930614166310267339487454538068422443657969041262460
444429964816431426710183695755476413937115540614428835263204196063562
204136440332086975595534127489616794634958117262920797357760383121216
527578368541040720097880511281265751148514970585852544357025892510419
412343969413669659546355293839984987174937768419702376847760976489365
521587044048554452410130309174187881781817343014915775915442567796307
120118050011988970522886279315724310652960260354317732408237949885912
477148109782713365688772833978483650246761301470938166706596058470226
807335420535020271139433490249384246034110447894017222870344265949",
10);
if (ret != 0){
    perror("mpz for n error");
    exit(EXIT_FAILURE);
}
ret = qin_fun(e, n);
if (ret != 0){
    perror("attack fail");
    exit(EXIT_FAILURE);
}
mpz_clears(e, n, NULL);
return EXIT_SUCCESS;
```

}