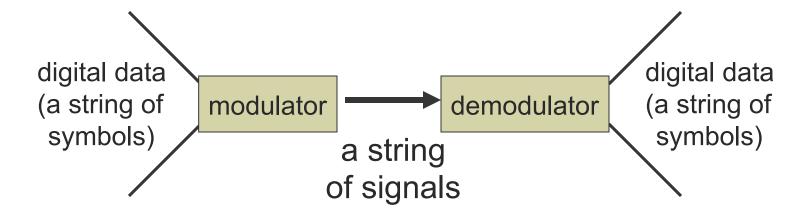
Direct Link Networks - Framing

Reading: Peterson and Davie, Chapter 2

Framing



- Encoding translates symbols to signals
- Framing demarcates units of transfer
 - Separates continuous stream of bits into frames
 - Marks start and end of each frame



Benefits of framing

- Synchronization recovery
 - Breaks up continuous streams of unframed bytes
 - Recall RS-232 start and stop bits
- Link multiplexing
 - Multiple hosts on shared medium
 - Simplifies multiplexing of logical channels
- Efficient error detection
 - Per-frame error checking and recovery



Framing

- Demarcates units of transfer
- Goal
 - Enable nodes to exchange blocks of data
- Challenge
 - How can we determine exactly what set of bits constitute a frame?
 - How do we determine the beginning and end of a frame?



Framing

Approaches

- Sentinel: delimiter at end of frame (like C strings)
- Length-based: length field in header (like Pascal strings)
- Clock-based: periodic, time-based

Characteristics

- Bit- or byte-oriented
- Fixed or variable length
- Data-dependent or data-independent length



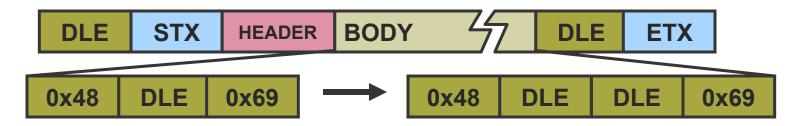
Sentinel-Based Framing

- End of Frame
 - Marked with a special byte or bit pattern
 - Frame length is data-dependent
 - Challenge
 - Frame marker may exist in data
 - Requires stuffing
- Examples
 - BISYNC, HDLC, PPP, IEEE 802.4 (token bus)



ARPANET IMP-IMP

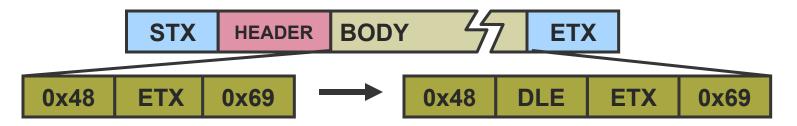
- Interface Message processors (IMPs)
 - Packet switching nodes in the original ARPANET
 - Byte oriented, Variable length, Data dependent
 - Frame marker bytes
 - STX/ETX start of text/end of text
 - DLE data link escape
 - Byte Stuffing
 - DLE byte in data sent as two DLE bytes back-to-back





BISYNC

- Blnary SYNchronous Communication
 - Developed by IBM in late 1960's
 - Byte oriented, Variable length, Data dependent
 - Frame marker bytes:
 - STX/ETX start of text/end of text
 - DLE data link escape
 - Byte Stuffing
 - ETX/DLE bytes in data prefixed with DLE's





Byte Stuffing: BISYNC

0000 0011 1110 0111 1111 1110 0001 0000 0001 1111

ETX/DLE bytes in data prefixed with DLE's

DLE = 16; STX = 2; ETX = 3

Ans: 0000 0010 0001 0000 0000 0011 1110 0111 1111 1110 0001 0000 0001 0000 0001 1111 0000 0011



Byte Stuffing: Efficiency

0000 0011 1110 0111 1111 1110 0001 0000 0001 1111

Frame: 0000 0010 0001 0000 0000 0011

1110 0111 1111 1110 **0001 0000**

0001 0000 0001 1111 0000 0011

Efficiency:

- 72 bits were sent for 40 bits of data
- Efficiency is 40/72 = 55.6%



High-Level Data Link Control Protocol (HDLC)

- Bit oriented, Variable length, Datadependent
- Frame Marker
 - O1111110
- Bit Stuffing
 - Insert 0 after pattern 011111 in data
 - Example
 - 01111110 end of frame
 - 01111111 error! lose one or two frames



Bit Stuffing: HDLC

0000 0011 1110 0111 1111 1110 0001 0000 0001 1111

Insert 0 after pattern 011111 in data

Ans: 0111 1110 0000 0011 11100 0111 11011 1110 0001 0000 0001 11110 0111 1110



Bit Stuffing: Efficiency

0000 0011 1110 0111 1111 1110 0001 0000 0001 1111

- Frame: 0111 1110 0000 0011 11100 0111 11011 1110 0001 0000 0001 11110
 0111 1110
- Efficiency
 - 59 bits were sent for 40 bits of data
 - Efficiency = 67.8%



IEEE 802.4 (token bus)

- Alternative to Ethernet (802.3) with fairer arbitration
- End of frame marked by encoding violation,
 - i.e., physical signal not used by valid data symbol
 - Recall Manchester encoding
 - low-high means "0"
 - high-low means "1"
 - low-low and high-high are invalid
- IEEE 802.4
 - byte-oriented, variable-length, data-independent
- Another example
 - Fiber Distributed Data Interface (FDDI) uses 4B/5B
- Technique also applicable to bit-oriented framing



Length-Based Framing

- End of frame
 - Calculated from length sent at start of frame
 - Challenge
 - Corrupt length markers
- Examples
 - DECNET's DDCMP
 - Byte-oriented, variable-length
 - RS-232 framing
 - Bit-oriented, implicit fixed-length



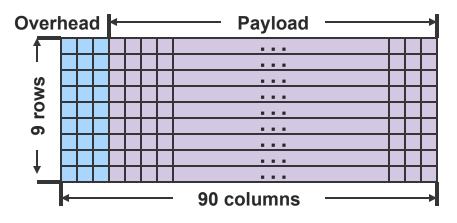


Clock-Based Framing

- Continuous stream of fixed-length frames
 - Clocks must remain synchronized
- STS-1 frames 125μs long
 - No bit or byte stuffing
- Example
 - Synchronous Optical Network (SONET)
- Problems
 - Frame synchronization
 - Clock synchronization



- Frames (all STS formats) are 125 µsec long
 - Ex: STS-1 51.84 Mbps = 90 bytes
- Frame Synchronization
 - 2-byte synchronization pattern at start of each frame





SONET: Challenges

- How to recover frame synchronization
 - Synchronization pattern unlikely to occur in data
 - Wait until pattern appears in same place repeatedly
- How to maintain clock synchronization
 - NRZ encoding
 - Data scrambled (XOR' d) with127-bit pattern
 - Creates transitions
 - Also reduces chance of finding false sync. pattern

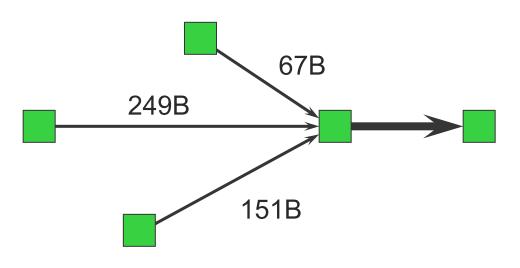


- A single SONET frame may contain multiple smaller SONET frames
- Bytes from multiple SONET frames are interleaved to ensure pacing



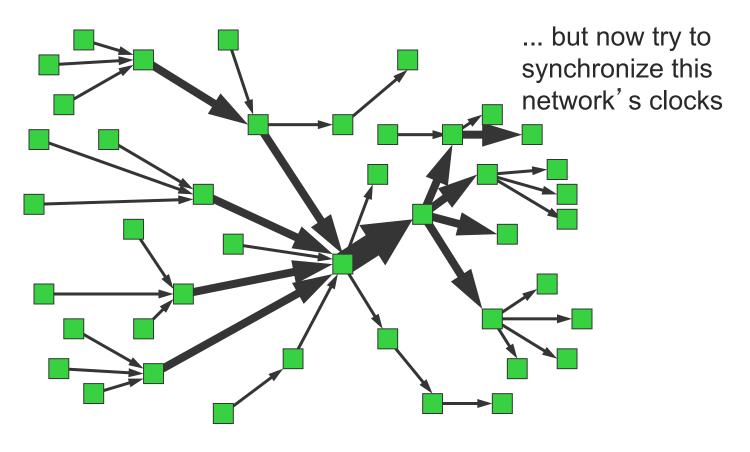


- STS-1 merged bytewise round-robin into STS-3
- Unmerged (single-source) format called STS-3c
- Problem: simultaneous synchronization of many distributed clocks

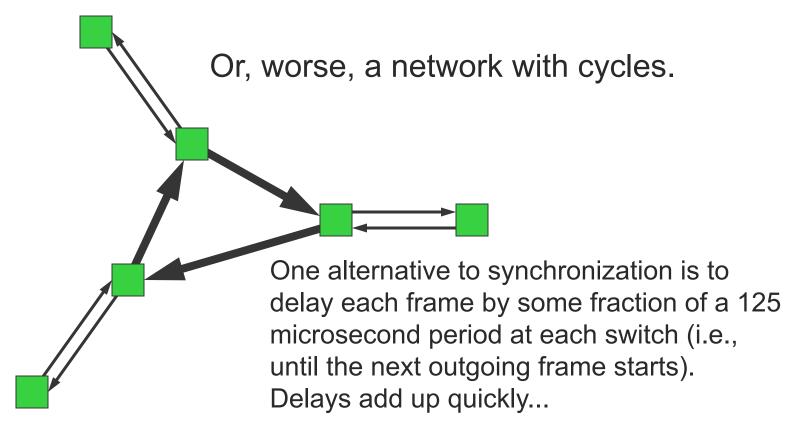


not too difficult to synchronize clocks such that first byte of all incoming flows arrives just before sending first 3 bytes of outgoing flow



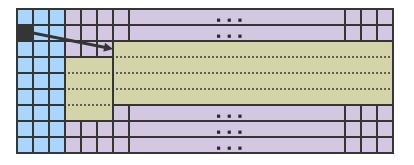








- Problem
 - Clock synchronization across multiple machines
- Solution
 - Allow payload to float across frame boundaries
 - Part of overhead specifies first byte of payload





Framing Summary

- Technique
 - Demarcate units of transfer
- Benefits
 - Synchronization recovery
 - Link multiplexing
 - Efficient error detection

- Approaches
 - Sentinel
 - Length-based
 - Clock based
- Characteristics
 - Bit- or byte-oriented
 - Fixed or variable length
 - Data-dependent or data-independent length

