#### CS 491 CAP Intro to Competitive Algorithmic Programming

# Lecture 4 Ad Hoc and Simulation

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### **Tryouts**

Our first tryouts will be on Saturday, Sept 23



#### What are Ad Hoc and Simulation?

- Simulation: Do exactly what the problem statement tells you.
  - E.g. simulate a board game, such as UNO, Blockus, and King of Tokyo, after the statement tells you the rules and strategies (may be simplified).
- Ad Hoc (a.k.a. Bruteforce): The algorithm is very straightforward.
  - E.g. just enumerate all possible solutions and figure out the answer.
- Backtracking is also very useful in bruteforce search problem.
  - E.g. Exact Cover Problems, such as 8-Queens and Sudoku.



### Josephus problem: Statement

- ♦ *N* people standing in a circle, numbered from 1 to *N*
- $\diamond$  The next person of *i* is i % N + 1
  - E.g. the next person of 1 is 2
  - E.g. the next person of *N* is 1
- ♦ The counting out begins at *s*-th person
- ♦ In each step, K 1 people are skipped and the next person is executed.
  - If we start with 1, the first one to be executed is *K*.
- $\diamond$  Given *N*, *K*, and *s*, who is the last person?



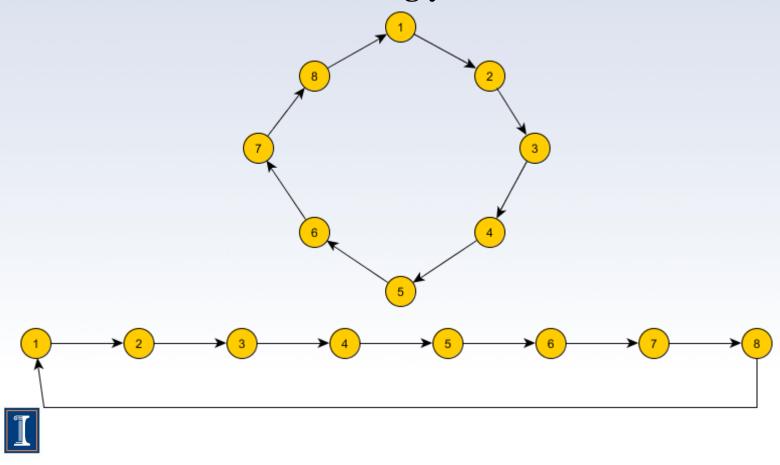
### Josephus problem: Example

- ♦ Example, N = 7, K = 2, s = 1
- ♦ 0<sup>th</sup> round: **1** 2 3 4 5 6 7
- ♦ 1<sup>st</sup> round: 1 **3** 4 5 6 7
- ♦ 2<sup>nd</sup> round: 1 3 **5** 6 7
- ♦ 3<sup>rd</sup> round: 1 3 5 **7**
- ♦ 4<sup>th</sup> round: **3** 5 7
- ♦ 5<sup>th</sup> round: 3 **7**
- $\diamond$  6<sup>th</sup> round: 7



### Josephus problem: Solution

♦ Data Structure: Circular Singly Linked List



#### Josephus problem: Solution

- ♦ Data Structure: Circular Singly Linked List
- ♦ Two for-loops to simulate the execution process

```
current_person = S-th person

for (int iteration = 1; iteration < N; ++
iteration) {
   for (int i = 1; i < K; ++ i) {
      current_person = move to next person
   }
   kill current_person
}</pre>
```



### Josephus problem: Solution

- ♦ Data Structure: Circular Singly Linked List
- ♦ Two for-loops to simulate the execution process
- $\diamond$  Time Complexity: O(NK)



#### Josephus problem: Practice

- ♦ POJ 1012
- ♦ <a href="http://poj.org/problem?id=1012">http://poj.org/problem?id=1012</a>
- ♦ Tips: There might be some duplicated test cases in a single run. You can store the answers in your program and print them when you see the same test cases again.



### Josephus problem: Future

 $\diamond$  You will learn more efficient solutions in Dynamic Programming, which are O(N) and O(KlogN).



#### Social Constraints: Statement

- $\diamond$  There are  $n \leq 8$  movie goers
- ♦ They will sit in the front row with *n* consecutive open seats
- ♦ There are  $m \le 20$  seating constraints among them, i.e. a and b must be at most (or at least) c seats apart
- How many possible seating arrangements are there?



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### Social Constraints: Example

- ♦ 3 people in total.
- ♦ 1 constraint: 1 and 2 must sit adjacently.
- All possible assignments:
  - **1** 2 3
  - **2** 1 3
  - **3** 1 2
  - **321**



#### **Social Constraints: Solution**

♦ Try all possible seats assignment

```
vector<int> perm;
for (int i = 0; i < n; ++ i) {
    perm.push_back(i);
}
do {
    // check the conditions
} while (next_permutation(perm.begin(), perm.end());</pre>
```



#### **Social Constraints: Solution**

- ♦ Try all possible seats assignment
- $\diamond$  Time Complexity: O(n! m)
- $\diamond$  Worse case:  $8! \times 20 = 806,400$



#### Tips

- ♦ Usually, within 1 second running time, you can apply 10<sup>7</sup> multiplication operations. Sometimes, on some efficient machine, 10<sup>8</sup> is also fine.
- \$\langle\$ /, sqrt, cos, sin, atan2, and so on. These operations are a little slower.
- ⋄ |, &, ^, ~ are much faster.



#### **Subset Enumeration**

- ♦ How many subsets of {1..n} are "good"?
- "good" is very efficient and easy to judge after you have the subset.
- $\diamond$   $n \leq 20$



## Backtracking

- ♦ Save the states before recursion
- ♦ Restore the states after recursion



#### **Subset Enumeration**

♦ Backtracking:

```
vector<int> subset;
int n;
void search(int i)
{
    if (i == n) {
        // do something for subset ...
        return;
    }
    // choose i
    subset.push_back(i);
    dfs(i + 1);
    subset.pop_back(i);

    // not choose i
    dfs(i + 1);
}
```



#### **Subset Enumeration**

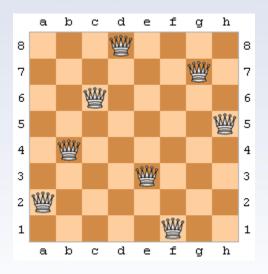
♦ Let's use bitmask instead of backtracking.

```
for (int mask = 0; mask < 1 << n; ++ mask) {
   vector<int> subset;
   for (int i = 0; i < n; ++ i) {
      if (mask >> i & 1) {
        subset.push_back(i);
      }
   }
   // do something for subset ...
}
```



### 8 Queens (UVA 750)

♦ Print ALL possible solution for 8 Queens





#### 8 Queens

- ♦ *i*-th step
  - For all 1 <= x <= 8
    - Try to put a new queen on (i, x)
    - Update the "attackness" of each grid
    - Recursion to (i+1)-th step
    - Remove (i, x)
    - · Re-update the "attackness" of each grid



### Sudoku (POJ 3074)

- ♦ Given a partially filled 9 x 9 Sudoku, find a possible solution.
- ♦ Every row/column/square contains 1-9 exactly once.

	2	7	3	8			1	
	1				б	7	3	5
					6		2	9
3		5	6	9	2		8	
					. 4			
	6		1	7	4	5		3
6	4				5			· ·
9	5	1	8				7	
	8			6	5	3	4	



#### Sudoku

- $\diamond$  (i,j)-th step
  - For all 1 <= *d* <= 9
    - Try to put d on (i, j)
    - Update the "conflict table" of each grid
    - Recursion to (i', j')-th step
    - Reset (i, j) to unknown
    - · Re-update the "conflict table" of each grid



#### DuLL: ICPC Mid-Central 2009

http://www.icpc midcentral.us/archives/2009/mcpc2009/dull/dull.html



#### DuLL: ICPC Mid-Central 2009

- ♦ Discrete Event Simulation
- Use frequency dictionary to keep track of how many programs rely on a DLL
- Iterate through *event list* (each program entry and exit) instead of over each time step
- At each event:
  - Update the frequency dictionary
  - Check if DLL memory usage is greater than previously seen



#### Summary

- In the following situations, it will be likely a simulation or ad hoc problem:
  - The statement tells what you need to do
    - E.g. tells you the detailed rules of a poker game
  - Ask for almost all (intermediate) results/solutions
    - E.g. print all primes between 1 to n
  - The size of data is quite small
    - Permutation: O(n!) algorithm for  $n \sim 10$
    - Subsets:  $O(2^n)$  algorithm for  $n \sim 20$
  - Backtracking when they ask for only ONE possible solutions
    - Exponential Time Complexity!



## Questions?

