# CS 491 CAP Advanced Search & Simulation

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- ♦ Bitmasks for Pruning
- ♦ Search Order
- ♦ Bidirectional Search
- ♦ A\* Search



- **† Bitmasks for Pruning**
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### N Queens

 $\diamond$ Find the total number of layouts of n queens on an  $n \times n$  chessboard, such that any two queens will not attach each other.

#### ♦Bruteforce:

- Every time, put a queen at (x, y)
- Mark (\*, y) and  $(x + \Delta, y + \Delta)$  and  $(x + \Delta, y \Delta)$  as attacked
  - •This step is O(n)



Recursion

#### Bitmask

- $\diamond$  Query the *d*-th digit of *x* 
  - $(x \gg d) \& 1$
- $\diamond$ Set the *d*-th digit of *x* as 1
  - $-x \mid (1 \ll d)$
- $\diamond$ Set the *d*-th digit of *x* as 0
  - $x & (\sim (1 \ll d))$
- $\diamond$ Only keep the last non-zero digit of x
  - -x & (-x)



### N Queens

- ♦ Mark columns, diagonals only
- $\diamond$ It becomes O(1)
- $\diamond$  Find the possible position in *x*-th row
- ♦⇔ Find the non-zero bits



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#### Sudoku

 $\diamond$ 9×9 and 16×16

- ♦No same number in a row
- ♦No same number in a col
- ♦No same number in a sub-square

♦Find a solution

1		3			5		9
		2	1	9	4		
			7	4	9	2	
3			5	2			6
	6				2 2	5	
7		2 - 5	8	3	- 5	8 - 3	4
			4	1	2 00	60 X	
		9	2	5	8		
8		4			1		7



#### Sudoku - Bruteforce

- ♦Find an empty cell
- ♦Enumerate a possible number to fill
- ♦ Mark its row, column, and subsquare
  - Bitmasks could be used again!

1		3			5		9
		2	1	9	4		
			7	4			
3			5	2			6
	6					5	
7		2 0	8	3			4
			4	1		8	
		9	2	5	8		
8		4			1		7



#### Sudoku - Heuristic

- ♦Find an empty cell
- ♦With the **smallest** number of possible numbers



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#### Bidirectional search

- ♦ Find shortest path given an initial node and a target node.
- ♦ Two simultaneous searches
  - Start  $\rightarrow ... \rightarrow$  an overlap!  $\leftarrow ... \leftarrow$  Target
- ♦ Faster!
  - Suppose both searches expand a <u>tree</u> with <u>branching factor</u> b
  - the distance from start to goal is *d*
  - Each of the two searches:  $O(b^{d/2})$
  - A single search:  $O(b^d)$



#### Bidirectional search

♦Requirement: the reversed move is easy to obtain



```
1 2 3
x 4 6
7 5 8
```

- ♦ How many different layouts?
- ♦Bidirectional BFS is much faster than BFS



#### k-sum

- $\diamond$  Given an array a[1..n] and a target sum s
- $\diamond$  Is it possible to find k numbers such that their sum is exactly s?

$$\diamond O(n^{\left\lceil \frac{k}{2} \right\rceil})$$
 is desired



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#### A\* search

- $\diamond g(x)$  is the current cost from start to x
- $\diamond$  Design a heuristic function h(x), which estimates the cost of the cheapest path from x to the goal
  - problem-specific
  - admissible, meaning that it never overestimates the actual cost to get to the nearest goal node

$$\diamond f(x) = g(x) + h(x)$$

 $\diamond$  Use f(x) as the priority



#### A\* search

- $\diamond$  Maintain a priority queue Q
- $\diamond$  Pick x with the highest priority f(x) from Q
- $\diamond$  If x is the goal, f(x) is the answer
- $\diamond$ Expand  $x \rightarrow y_1, y_2, ..., y_k$
- ♦Push all *unseen y*'s into *Q*



### Dijkstra Revisit

- ♦Dijkstra's algorithm can be viewed as a special case of A\*
  - Pick the unseen closet node
  - g(x) is the current shortest distance
  - $\bullet h(x) = 0$



```
♦ Goal:
♦ 1 2 3 4
♦ 5 6 7 8
♦ 9 10 11 12
♦ 13 14 15 X
```

of the tiles with which it shares an edge.





- ♦ How many different layouts?
- \$15!=1,307,674,368,000
- ♦Toooo large for BFS



♦ Any heuristic?

♦Sum of Manhattan Distances to their destinations



### Recommended Readings

- ♦ <u>USACO 1.4.1 Search Techniques</u>
- ♦ <u>Bidirectional Search</u>
- ♦A\* Search
- **♦**Exact Cover
- **♦** Dancing Links



## Q&A

