

A Distributed OpenCL Framework using Redundant Computation and Data Replication

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Outline

- OpenCL programming model
- Previous approaches for clusters
- Overview of SnuCL-D
- Correctness problems
- Optimization techniques
- Limitation
- Conclusion





Introduction

- Heterogeneous systems
 - Different types of processors
 - E.g., CPUs+GPUs

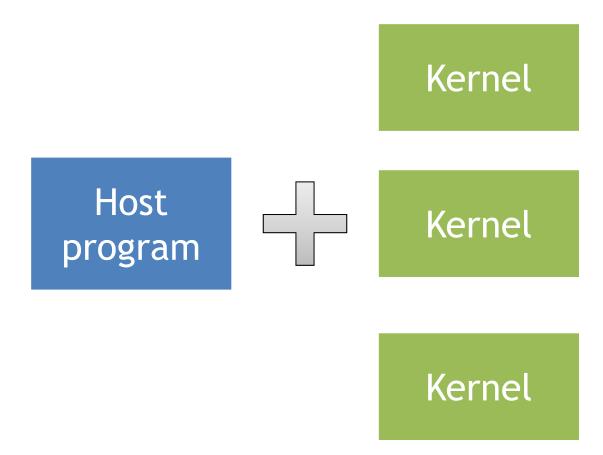
- Major programming models
 - CUDA and OpenCL







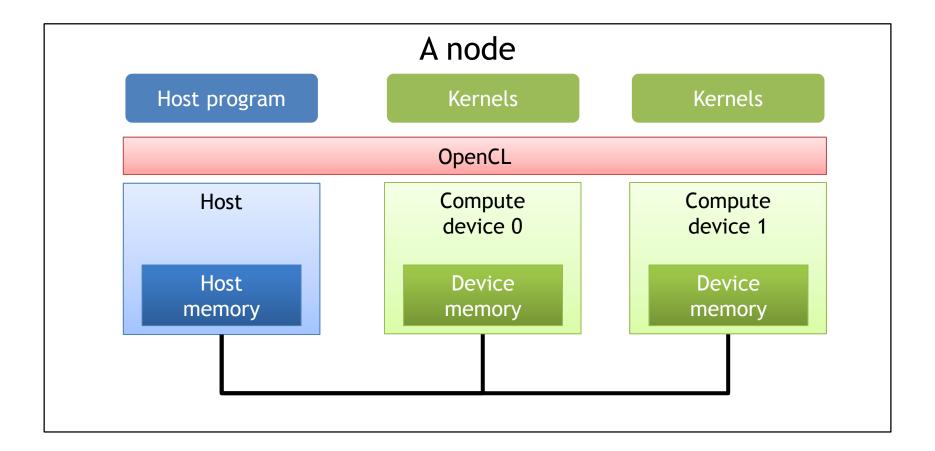
OpenCL Program





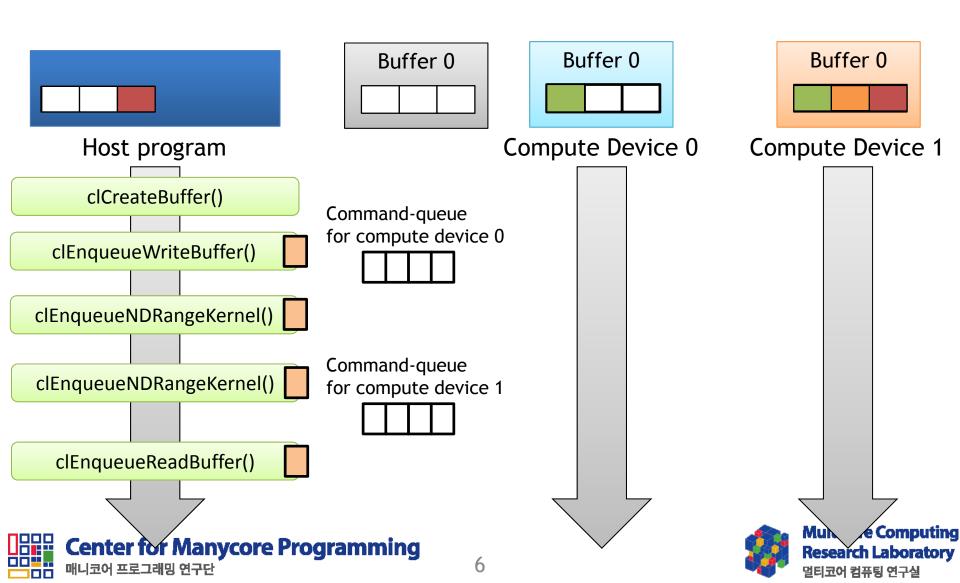


OpenCL Platform Model





OpenCL Programming Model





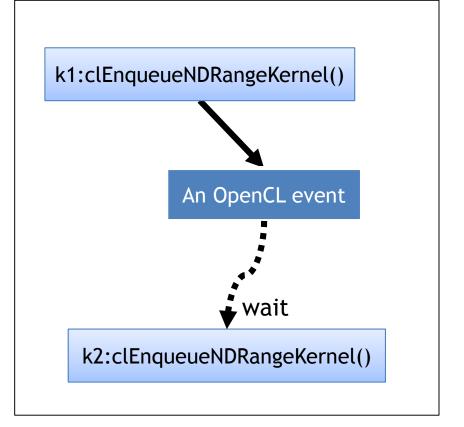
Order of Commands

Unordered

k1:clEnqueueNDRangeKernel()

k2:clEnqueueNDRangeKernel(

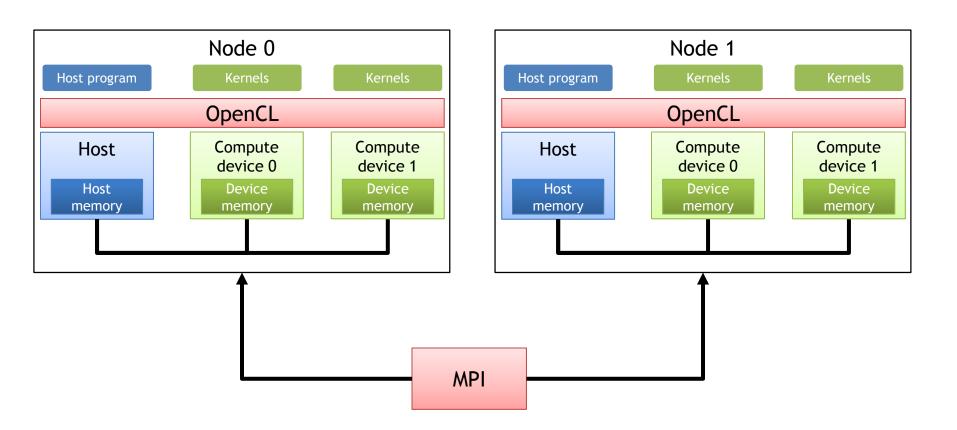
Ordered







OpenCL for Heterogeneous Clusters



MPI+OpenCL: cumbersome and error-prone







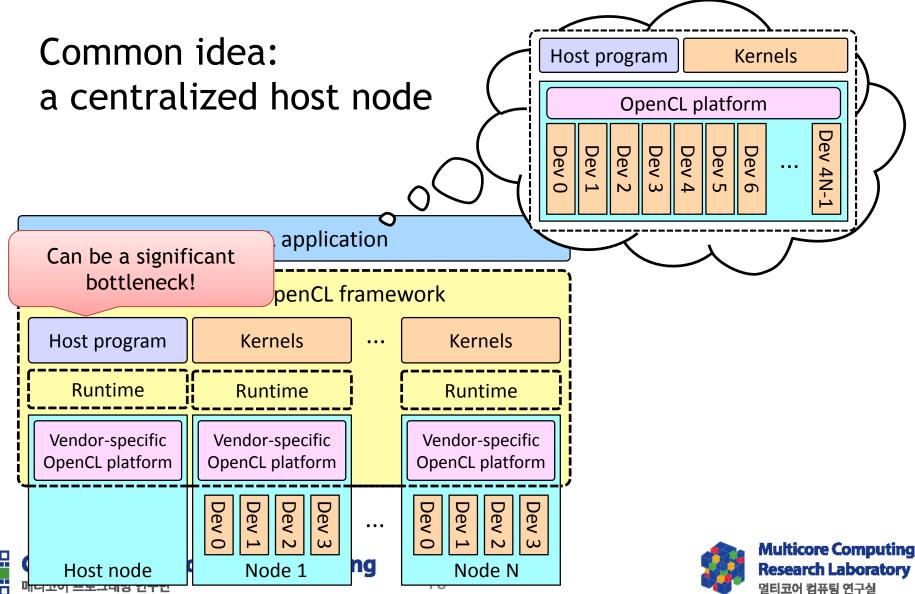
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Previous Approaches for Clusters





Centralized Approaches

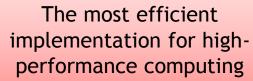
- clOpenCL
- Hybrid OpenCL
- SocketCL
- dOpenCL
- CLara
- SnuCL

- DistributedCL
- CLuMPI
- rCUDA
- DS-CUDA

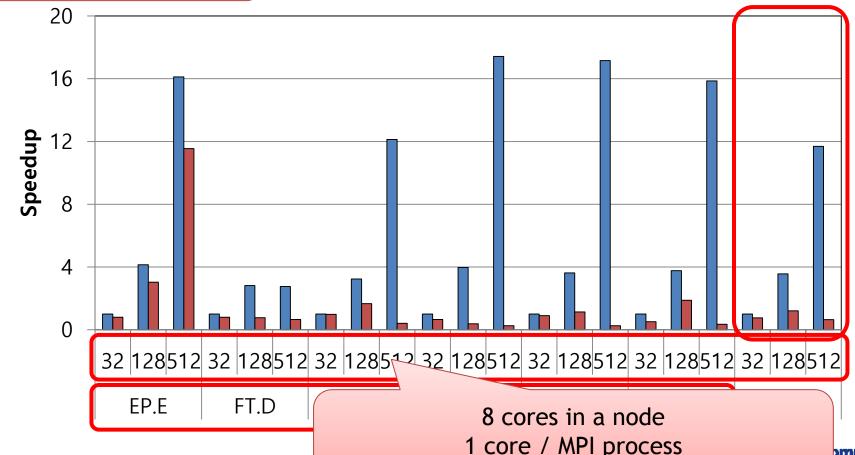
Evaluates it on a large number of nodes (256 nodes)



MPI-Fortran vs. SnuCL









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1 core / MPI process A set of 4 cores / OpenCL compute device

omputing boratory 연구실



Goals

- Develop a scalable OpenCL framework for clusters
 - Comparable to MPI-Fortran
 - Achieve ease of programming with high performance
- Key idea
 - Eliminate the centralized host node
 - Redundant host computation
 - Data replication
- SnuCL-D
 - The successor of SnuCL
 - A distributed OpenCL framework







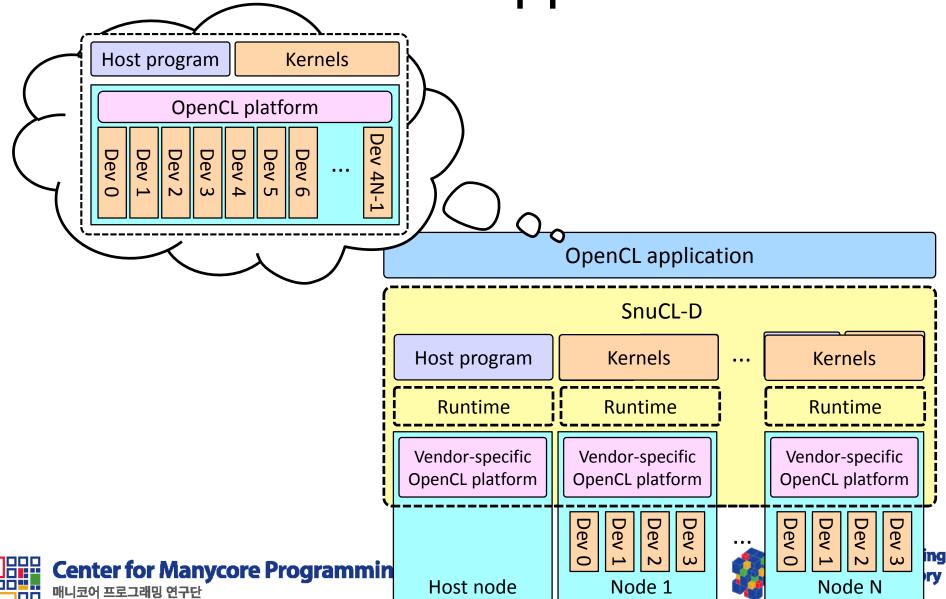
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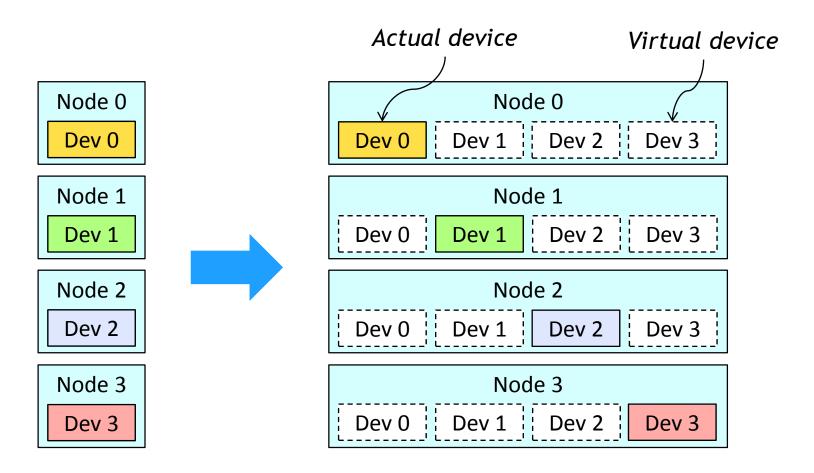


SnuCL-D's Approach





Remote Device Virtualization

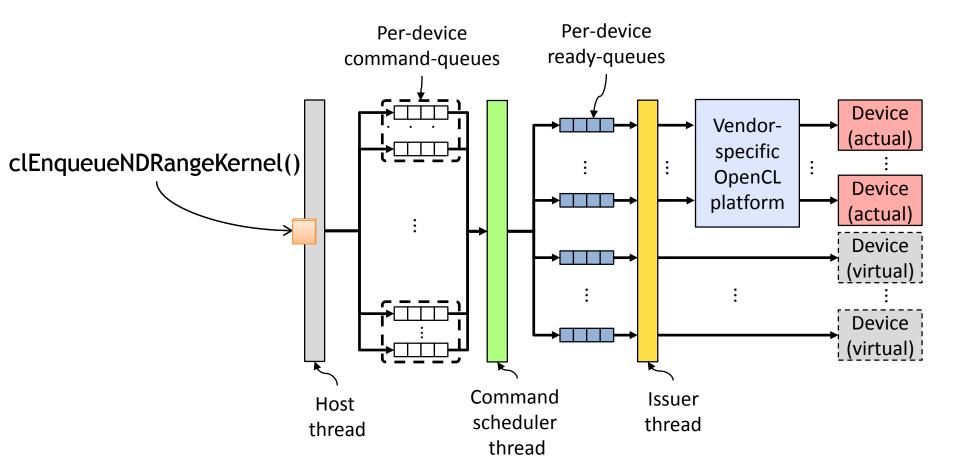








SnuCL-D Runtime









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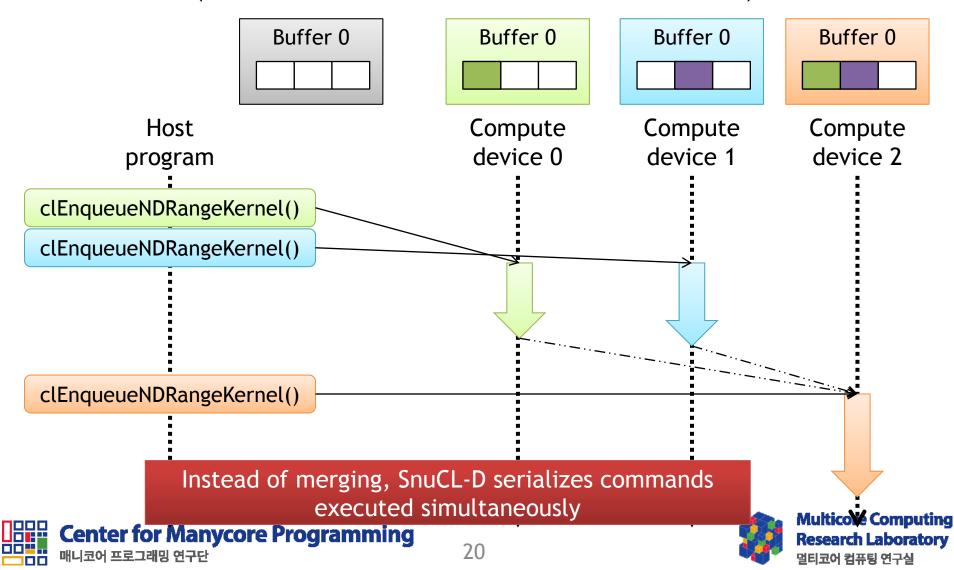
Correctness Problems

Consistency problem

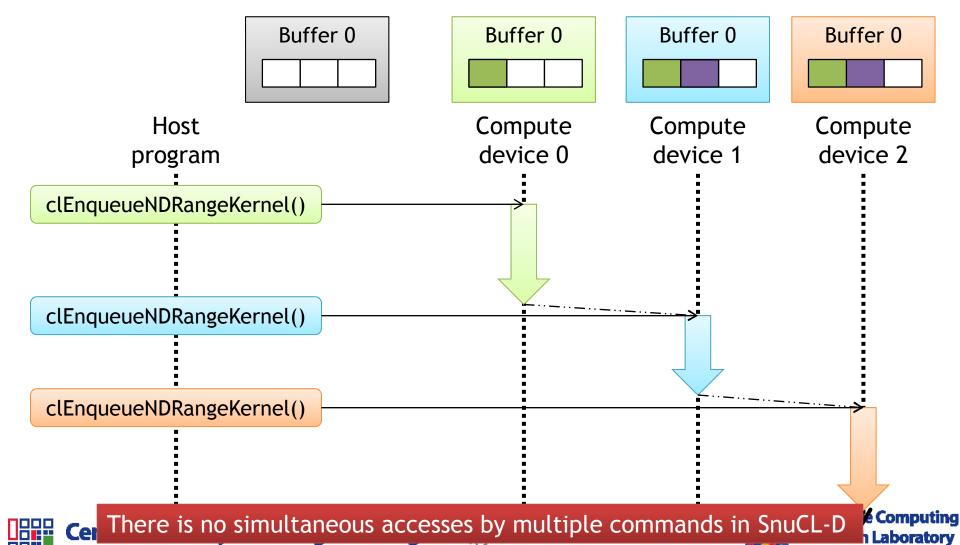
Non-determinacy problem



Consistency Problem (Simultaneous Accesses)



Consistency Problem (Simultaneous Accesses)

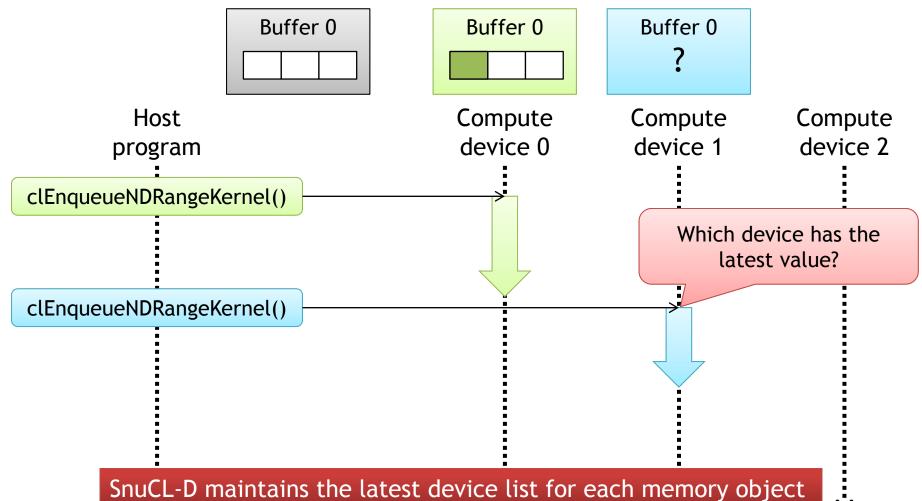


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Consistency Problem (Sequential Accesses)





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22

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Consistency Problem (Sequential Accesses)

Latest device list for buffer 0 Buffer 0 Buffer 0 Buffer 0 Dev 0 Host Compute Compute Compute device 0 device 2 device 1 program clEnqueueNDRangeKernel() Check the latest device list! clEnqueueNDRangeKernel() Multicove Computing **Center for Manycore Programming**



Correctness Problems

- Consistency problem
 - Solved by serialization and latest device lists

Non-determinacy problem



Non-deterministic Command Scheduling (Single-threaded Host Programs)

Problem

- Redundant command scheduling on every node
- If there is no enforced order between commands,
 - Command execution order can be different across nodes
 - May cause a deadlock, data inconsistency, etc.

Solution

- The enqueueing order is fixed
 - The order of clEnqueue...() calls
- SnuCL-D enforces the enqueueing order



Non-deterministic Command Scheduling (Multi-threaded Host Programs)

- Problem
 - Even the enqueueing order is not guaranteed
- Solution
 - Can be solved by deterministic multithreading
 - E. D. Berger [OOPSLA 09], C. Bienia [PACT 08], T. Liu [SOSP 11], M. Olszewski [ASPLOS 09]
 - Using this, we can make the enqueueing order deterministic
- Single-threaded host programs are more common
 - SnuCL-D assumes single-threaded host programs





Non-deterministic Result of a Function Call

Problem

- The result of a function call can be different across nodes
- E.g., file I/O and srand()

Solution

- Use global synchronization between nodes
- After designating a root node, the root node performs the call in the host program
- The others receive the result from the root



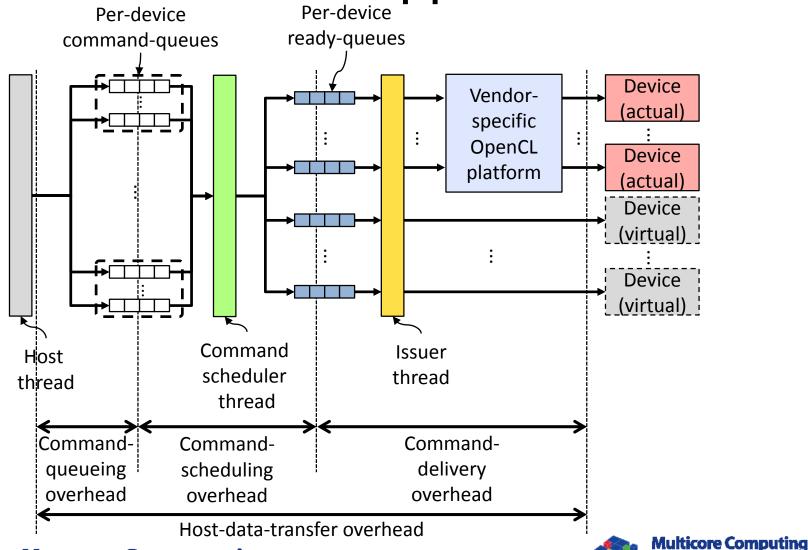


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Runtime Overheads **SIGN** of Centralized Approaches





29

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Decentralization Technique

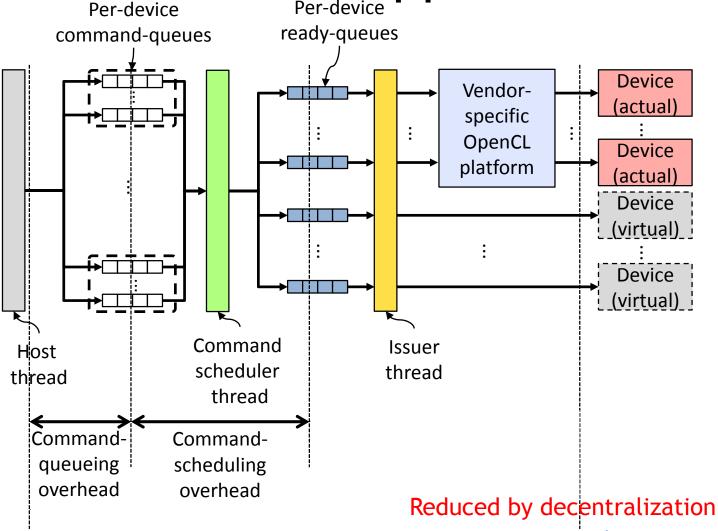
- To reduce
 - Command-delivery overhead
 - Host-data-transfer overhead
- The host program is executed on every node
 - Redundant computation
 - Data replication
- Remote device virtualization
 - Deliver commands to only actual devices



Runtime Overheads Signature



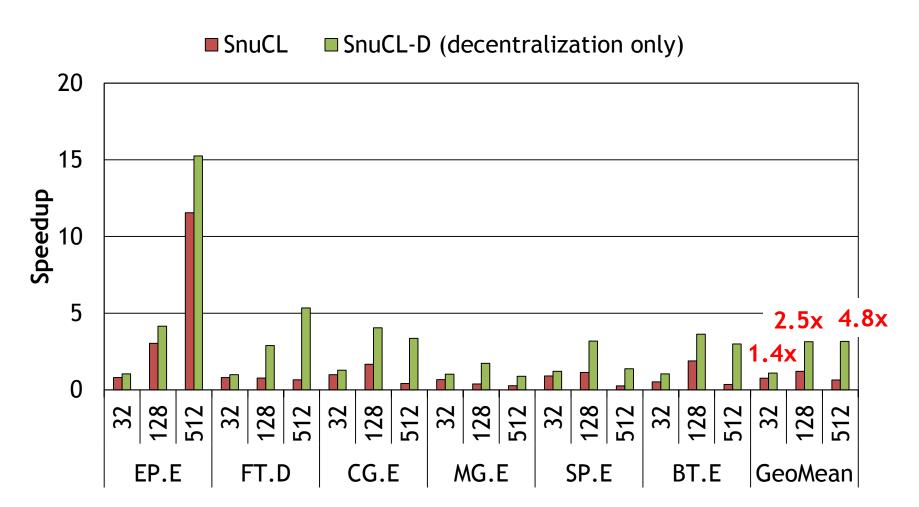
of Centralized Approaches
Per-device







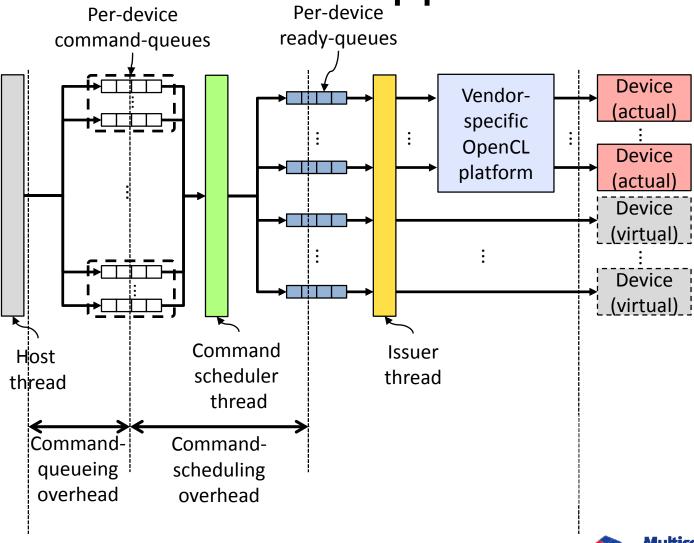
Performance







Runtime Overheads FILL of Centralized Approaches









New API Function

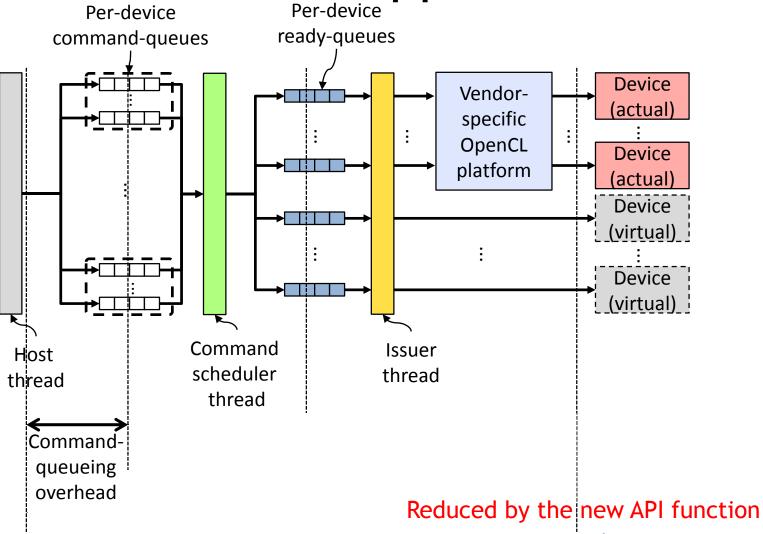
void clAttachBufferToDevice(cl_mem m, cl_device_id d);

- In OpenCL,
 - Memory object m is not bound to any devices
 - Scheduling overhead
 - Need to maintain latest device lists (consistency management)
- If this function is called,
 - SnuCL-D assumes
 - d always has the latest copy of m
 - Scheduling overhead reduced
 - No need to maintain latest device lists





Runtime Overheads **SIGN** of Centralized Approaches



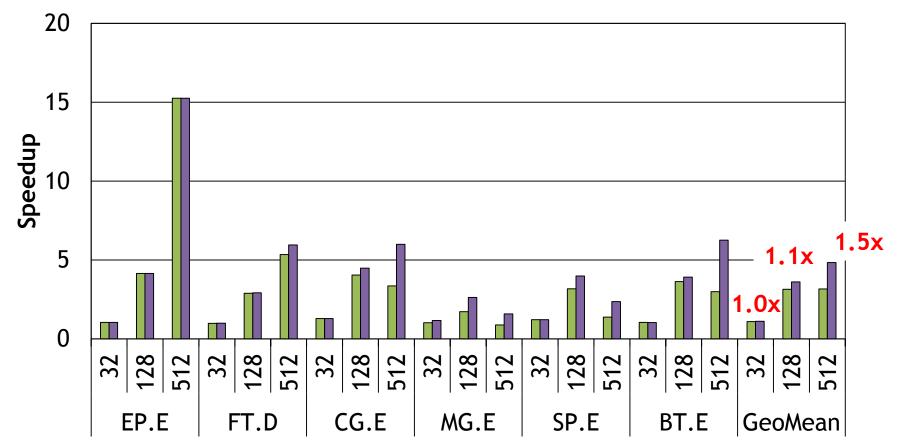






Performance

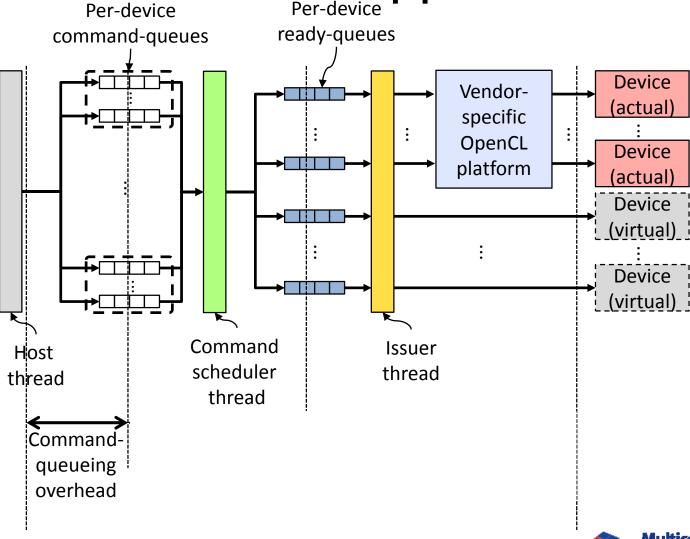
- SnuCL-D (decentralization only)
- SnuCL-D (decentralization + clAttachBufferToDevice)







Runtime Overheads FILE of Centralized Approaches







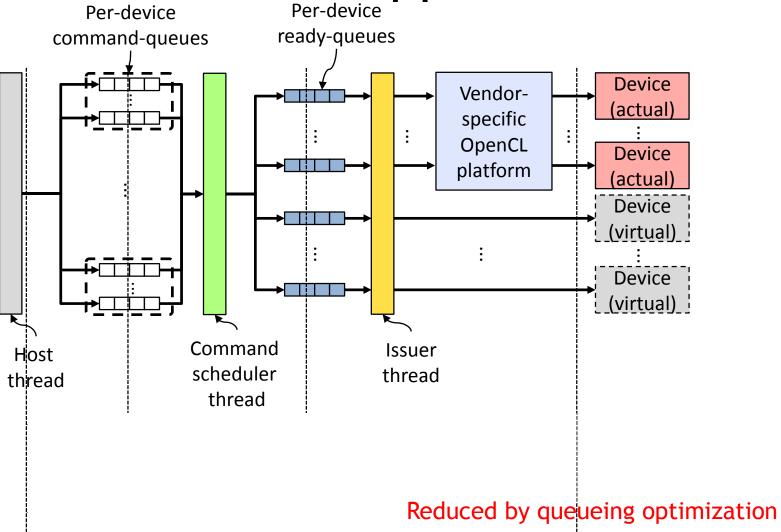


Queueing Optimization

- Two conditions under which commands for a virtual device do not need to be enqueued
 - No events in the event wait list
 - Each memory object is attached
- If the two conditions are met by a command
 - Discarding it does not affect correctness
- The commands can be safely discarded when enqueued



Runtime Overheads FILL of Centralized Approaches



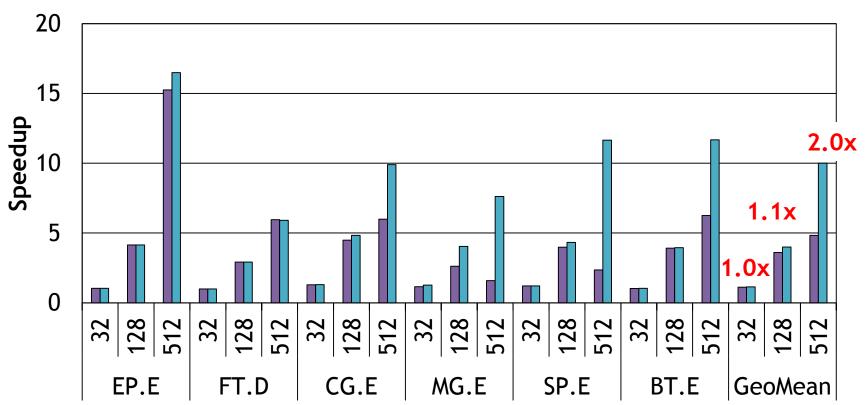






Performance

- SnuCL-D (decentralization + clAttachBufferToDevice)
- SnuCL-D (decentralization + clAttachBufferToDevice + queueing optimization)



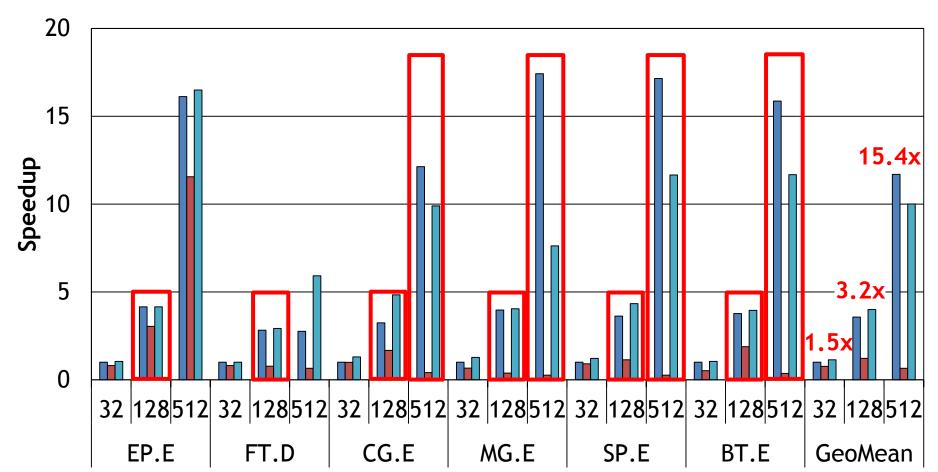






Performance

■ MPI-Fortran ■ SnuCL ■ SnuCL-D









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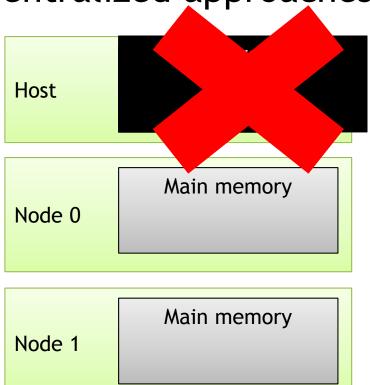
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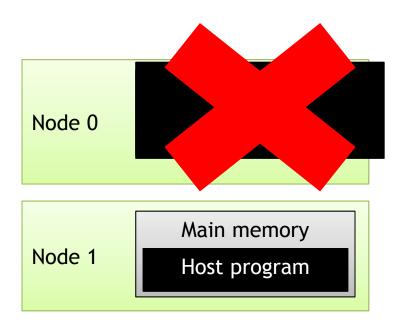


Limitation: Memory Footprint

Centralized approaches



SnuCL-D







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Conclusion

- SnuCL-D
 - A scalable and distributed OpenCL framework for clusters
 - OpenCL programs for multiple devices
 - Efficiently executed on a large-scale cluster
- Correctness Problems
 - Consistency problem and non-determinacy problem
- Three optimization techniques
 - Decentralization, new API function, and queueing optimization
- Available at http://snucl.snu.ac.kr
 - July 11, 2016







Thank you



