

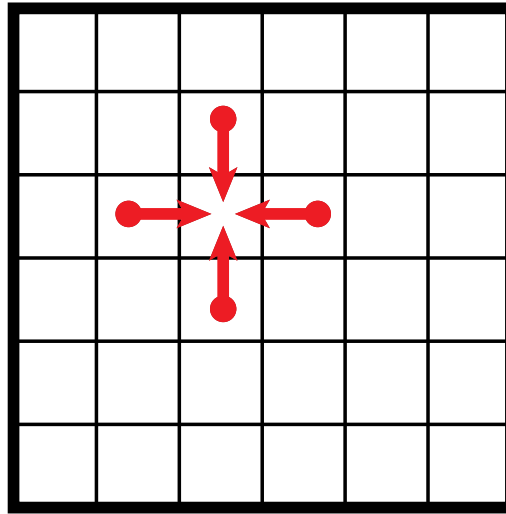
Verified Lifting of Stencil Computations

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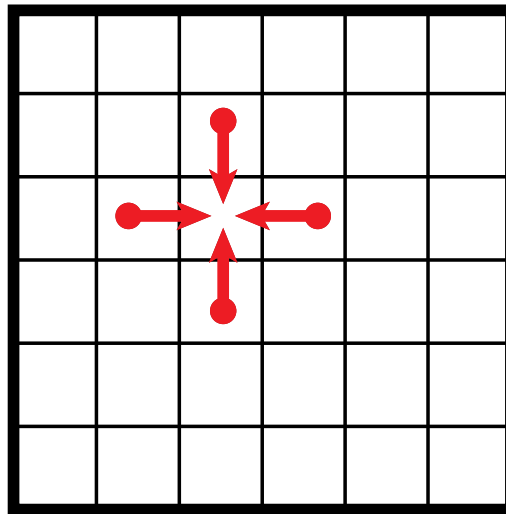
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Stencil computations update each point in a multidimensional regular array with a function of a subset of its neighbors.



$$out(i, j) = in(i, j) + 0.25 \times (in(i + 1, j) + in(i - 1, j) + in(i, j + 1) + in(i, j - 1))$$



```

do ib=istart,iend,2*ibsize
do i=ib, min(ib+ibsize,iend)
do j=jstart,jend,2
    out(i,j) = in(i,j) + (1.0/4.0)*(in(i+1,j)+in(i-1,j)+in(i,j+1)+in(i,j-1))
    out(i,(j+1)) = in(i,(j+1)) + (1.0/4.0)*(in(i+1,(j+1))+ &
        in(i-1,(j+1))+in(i,(j+1)+1)+in(i,(j+1)-1))
enddo
do i=ib+ibsize, min(ib+2*ibsize,iend)
do j=jstart,jend,2
    out(i,j) = in(i,j) + (1.0/4.0)*(in(i+1,j)+in(i-1,j)+in(i,j+1)+in(i,j-1))
    out(i,(j+1)) = in(i,(j+1)) + (1.0/4.0)*(in(i+1,(j+1))+ &
        in(i-1,(j+1))+in(i,(j+1)+1)+in(i,(j+1)-1))
enddo
enddo
enddo

```

Verified Lifting

Goal: Given a snippet of stencil code, *lift* it into a high-level summary that fully describes the algorithm, without optimizations.

Overview

- Introduction
- Verifying A Summary
- Synthesizing Loop Invariants & Postconditions
- Evaluation
- Summary

Verifying a summary

```
procedure sten(imin,imax,jmin,jmax,a,b)
  real (kind=8), dimension(imin:imax,jmin:jmax) :: a
  real (kind=8), dimension(imin:imax,jmin:jmax) :: b
  do j=jmin,jmax
    t = b(imin, j)
    do i=imin+1,imax
      q = b(i,j)
      a(i,j) = q + t
      t = q
    enddo
  enddo
end procedure
```



$$\forall i, j \in [imin + 1, imax] \times [jmin, jmax]$$
$$a(i, j) = b(i - 1, j) + b(i, j)$$

Verifying a summary

```
do j=jmin,jmax  
  t = b(imin, j)  
  do i=imin+1,imax  
    q = b(i,j)  
    a(i,j) = q + t  
    t = q  
  enddo  
enddo
```

Loop invariant I_j

$$\forall i, j' \in [imin + 1, imax] \times [jmin, j) \\ a(i, j') = b(i - 1, j') + b(i, j')$$

Loop invariant I_i

$$\forall i, j' \in [imin + 1, imax] \times [jmin, j) \\ a(i, j') = b(i - 1, j') + b(i, j') \\ \forall i', j' \in [imin + 1, i) \times [j, j] \\ a(i', j') = b(i' - 1, j') + b(i', j')$$



$$\forall i, j \in [imin + 1, imax] \times [jmin, jmax] \\ a(i, j) = b(i - 1, j) + b(i, j)$$

Verifying a summary

```
do j=jmin,jmax
  t = b(imin, j)
  do i=imin+1,imax
    q = b(i,j)
    a(i,j) = q + t
    t = q
  enddo
enddo
```


$$\forall i, j \in [imin + 1, imax] \times [jmin, jmax]$$
$$a(i, j) = b(i - 1, j) + b(i, j)$$

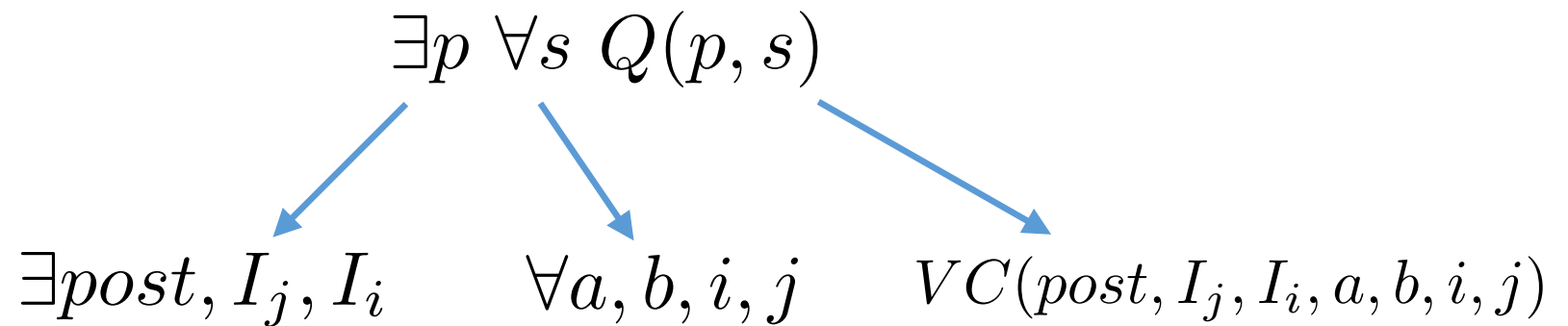
Need loop invariants such that:

- $\forall a, b, j, i . I_j(a, b, jmin)$
- $\forall a, b, j, i . I_j(a, b, j) \wedge (j > jmax) \rightarrow post(a, b)$
- ...
- $\forall a, b, j, i . I_i(a, b, j, i) \wedge (i \leq imax) \rightarrow$
 $I_i(a[(i, j) := b(i - 1, j) + b(i, j)], b, j, i + 1)$

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Program Synthesis



Template Generation

```
do j=jmin,jmax  
  t = b(imin, j)  
  do i=imin+1,imax  
    q = b(i,j)  
    a(i,j) = q + t  
    t = q  
  enddo  
enddo
```

$jmin, jmax, imin, imax = rand()$

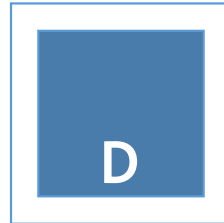
$b = \{b_0, b_1, b_2, \dots\}$

$a = \{\dots, b_9 + b_{10}, b_{10} + b_{11}, \dots\}$

$a(i, j) = b(??) + b(??)$

Template Generation

$$\forall \vec{x} \in D. \text{out}[\vec{x}] = \text{expr}(\vec{x})$$



$$\exists post, I_j, I_i \quad \forall a, b, i, j \quad VC(post, I_j, I_i, a, b, i, j)$$



$$I_j(a, b, j) \wedge (j > jmax) \rightarrow post(a, b)$$

$$\neg I_j(a, b, j) \vee \neg(j > jmax) \vee post(a, b)$$



$$\forall i, j' \in [imin + 1, imax] \times [jmin, j]$$

$$a(i, j') = b(i - 1, j') + b(i, j')$$



EAE

Skolemization

$$\exists x \forall y \exists z$$



$$\exists x \forall y f(y)$$

```
if (j<jmin)
    something
else if (j<jmax)
    something_else
else if...
    another_thing
```


Partial Skolemization

$$\exists x \forall y \exists z$$



$$\exists x \forall y f_p(y)$$

$$\exists x \forall y Q(x, y, p_1) \vee Q(x, y, p_2)$$

Partial Skolemization

```
boolean Inv_j(ref double[sz] a, ref double[sz] b, int jmin, int jmax,
int imin, int imax) {

    int[2] _js = {gen(jmin, jmax, imin, imax), gen(jmin, jmax, imin, imax)};

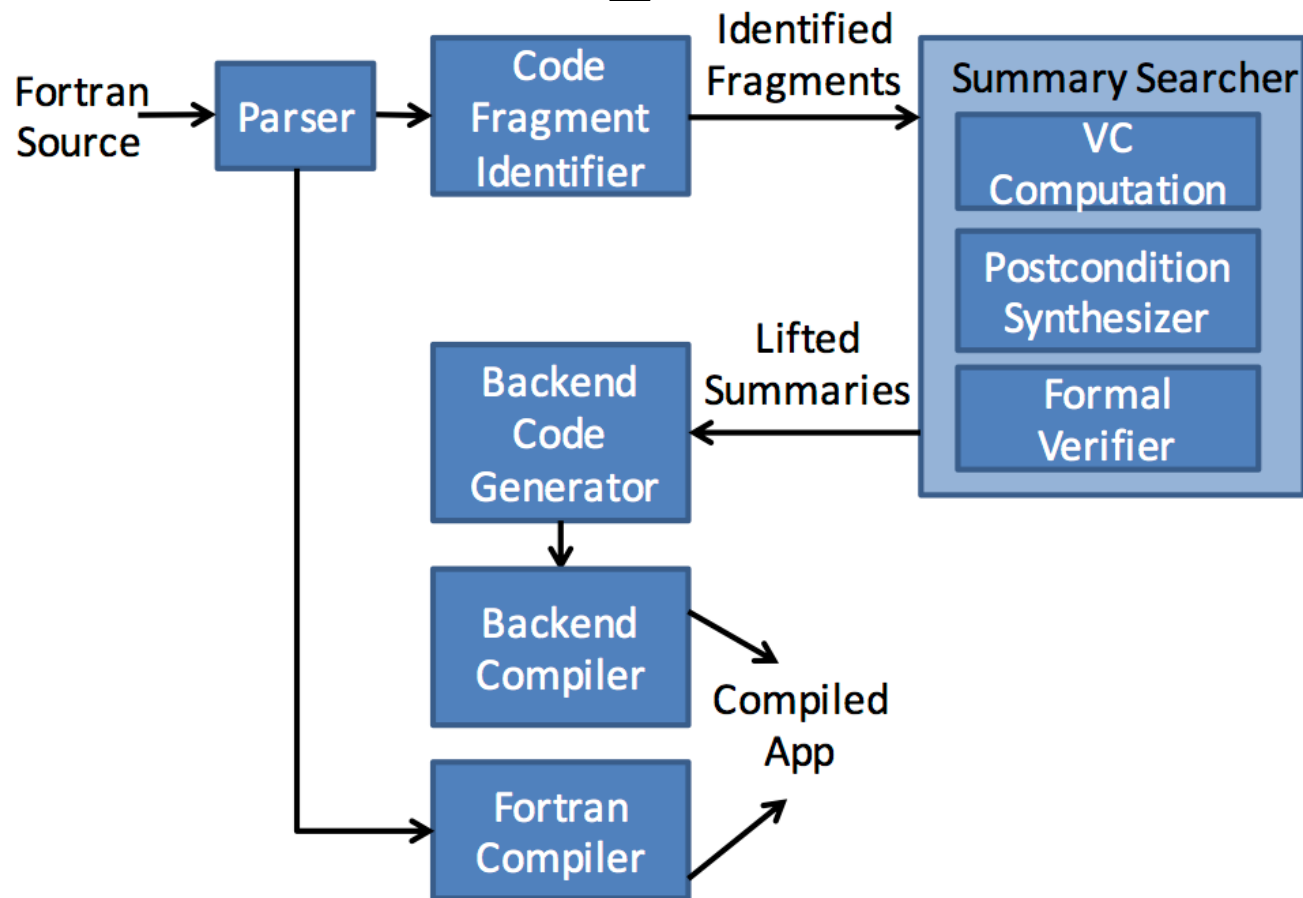
    int[2] _is = {gen(jmin, jmax, imin, imax), gen(jmin, jmax, imin, imax)};

    for (int ji=0; ji<2; ji++) {
        for (int ii=0; ii<2; ii++) {
            i = _is[ii];
            j = _js[ji];
            if a(i,j) != b(gen_pt(i,j)) + b(gen_pt(i,j)) return false;
        }
    }
    return true;
}
```

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STNG: Verified Lifting for Stencils



Limitations

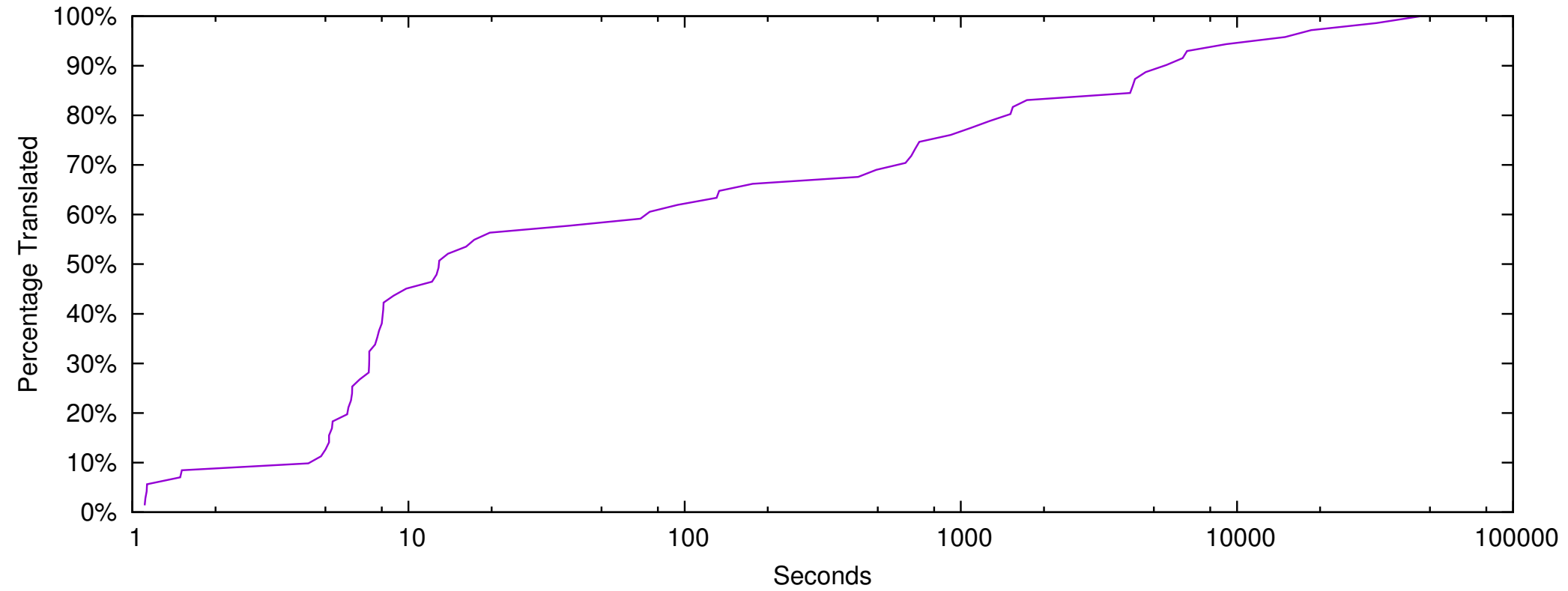
- Concentrate on a specific subset of all stencil computations
 - No conditionals or boundary conditions
 - Distinct input and output arrays
 - Incrementing loops only
- Even with careful generation of Sketches, time to solution can be very large

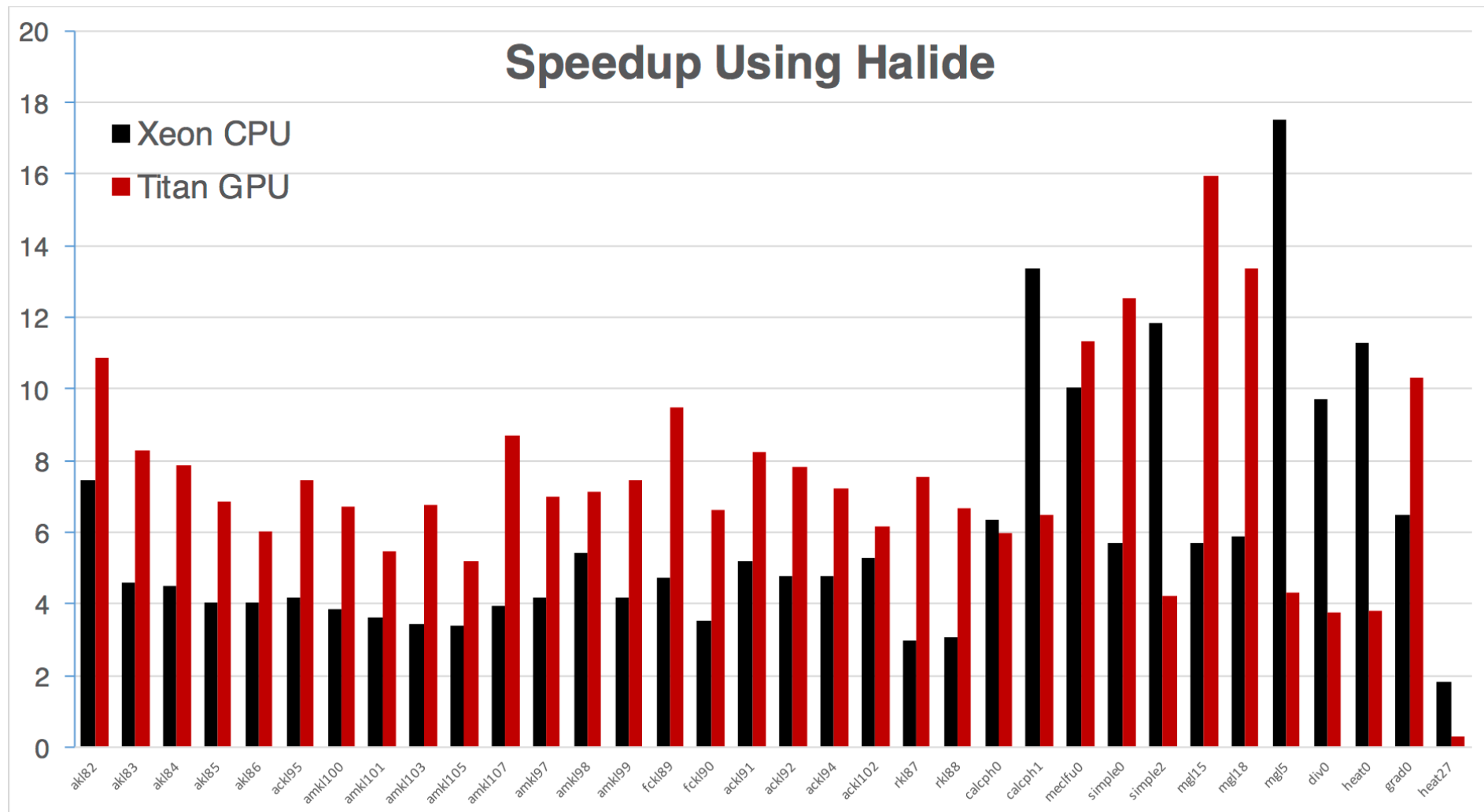
Benchmark Codes

- 4 scientific applications & 2 sets of benchmark codes
- All benchmarks are in (dialects of) Fortran
- Dimensionality from 1D to 6D

	Candidates	Translated	Failed
StencilMark	4	3	1
NAS MG	9	3	5
Cloverleaf	45	40	4
Terra	1	1	0
NFFS-FVM	29	25	1
Challenge	5	5	0
Total	93	77	11

Cumulative Distribution of Times to Synthesize Postcondition





Related Work

- Automatic invariant generation
- Compiler techniques (e.g. polyhedral model)
- QBS: Verified lifting of ORM code

Summary

- STNG performs verified lifting on stencil computations by synthesizing postconditions and loop invariants
- Much of the effort is in generating Sketches that limit the search space of possible postconditions/invariants
- Performance results show that the transformed code can leverage high performance DSLs