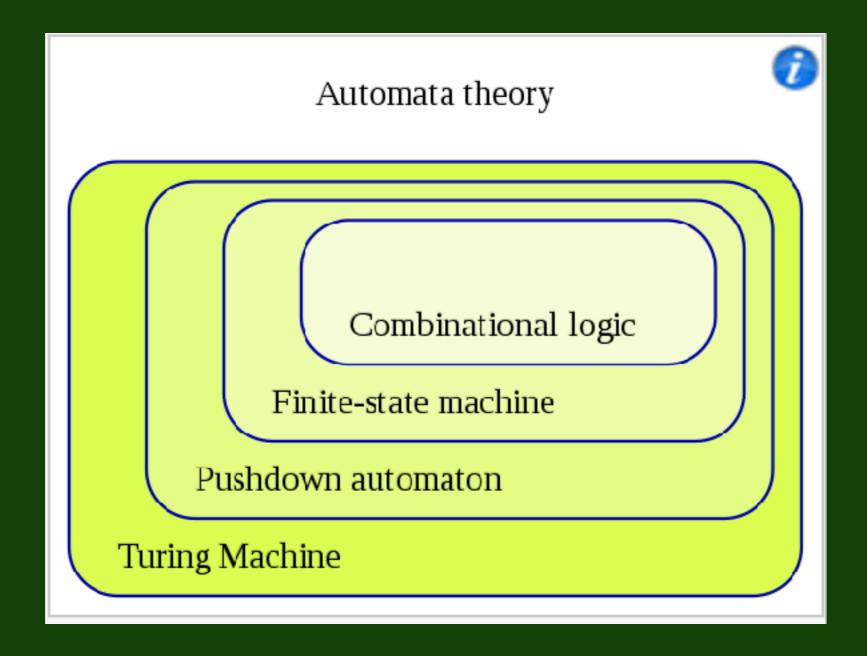
Finite Automata, Regular Expressions

What are automata?

Automata theory is the study of abstract machines and automata, as well as the computational problems that can be solved using them. It is a theory in theoretical computer science and discrete mathematics (a subject of study in both mathematics and computer science). The word automata (the plural of automaton) comes from the Greek word αὐτόματα, which means "self-acting".

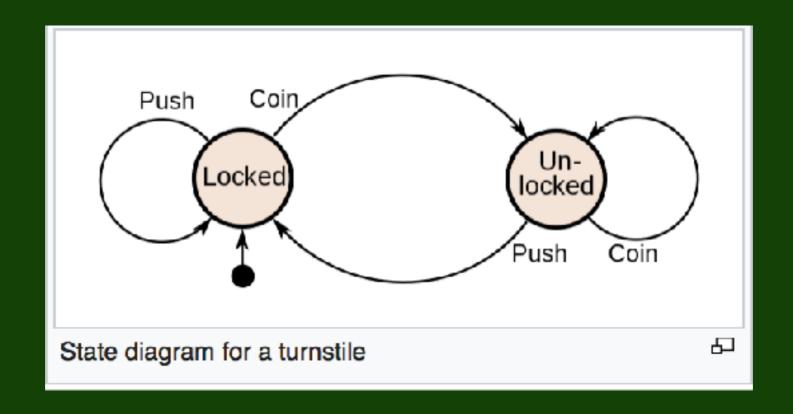
Automata



Automata are further divided into deterministic and non-deterministic

Finite state machine

 A simple two-state FSM: a turnstile (like the gates going on to the "T")—substitute "Charlie Card" for coin:



FSMs, etc.

- Finite state machines have limited "memory"—only the various states. In order to read realistic regular expressions, we need "push-down" automata.
- Push-down automata essentially give us a stack where we can manipulate the stack in addition to the various transitions (but we can't look inside the stack).
- Turing machine gives us essentially unlimited memory to store results/states.

Regular Expressions

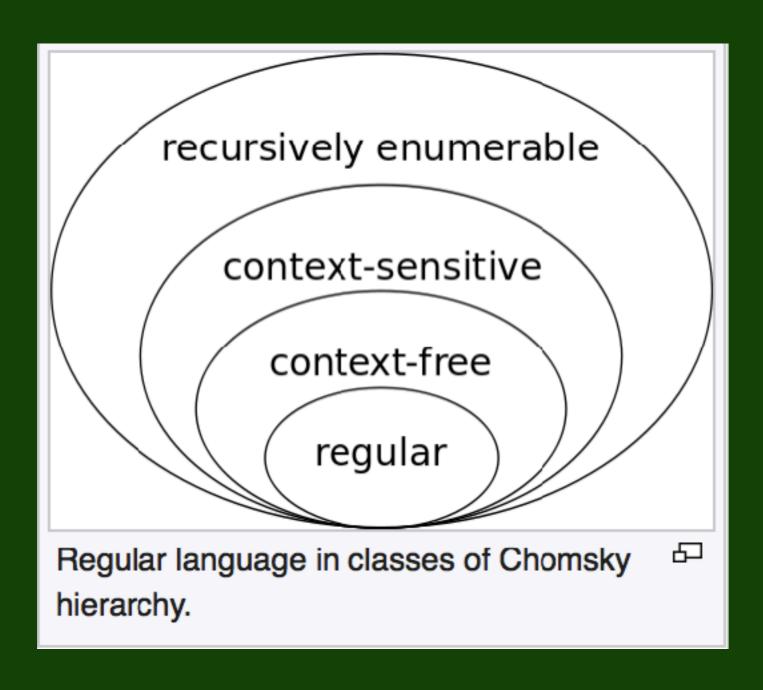
- Regular Expressions derive from the concept of a Regular Language;
- A regular language can be recognized by a finite automaton.
- Kleene's theorem: regular expressions and regular language are equivalent.

Regular Language

Defined by "productions," for example:

```
expr ::= term { "+" term | "-" term }.
term ::= factor { "*" factor | "/" factor}.
factor ::= floatingPointNumber | "(" expr ")".
```

Chomsky Hierarchy



Using regular expressions

- Regular expressions come in many different flavors (unfortunately):
 - Perl (Python?)
 - Java/POSIX (also used by other languages such as Scala)
- They are the basis of parsing (used for DSLs)
- Best resource for testing/understanding regex: https://regex101.com/ (but not Java-specific)
- RegexBuddy (but not free).

Untyped parsing (Scala)

- OK, now we just need to be able to define the grammar that our parser can operate on:
 - Take a look at this set of "productions" in BNF (Backus-Naur form) followed by examples:

```
expr ::= term { "+" term | "-" term }.
term ::= factor { "*" factor | "/" factor}.
factor ::= floatingPointNumber | "(" expr ")".
```

- 1+4.5-3 is an *expr*; 2*3.14/5 is a *term*; 3.1415927 is a *factor*; (7-5) is also a *factor*.
- The **Scala Parser Combinator** library allows us to code this parser with only a few substitutions: "~" replaces " "; "rep(" replaces "{"; ")"

Typed parsing

```
package edu.neu.coe.scala.parse
import scala.util.parsing.combinator._
class Arith extends JavaTokenParsers {
 trait Expression { def eval: Double }
  abstract class Factor extends Expression
  case class Expr(t: Term, ts: List[String~Term]) extends Expression {
    def term(t: String~Term): Double = t match {case "+"~x => x.eval; case "-"~x => -x.eval }
    def eval = ts.foldLeft(t.eval)(_ + term(_))
  case class Term(f: Factor, fs: List[String~Factor]) extends Expression {
    def factor(t: String~Factor): Double = t match {case "*"~x => x.eval; case "/"~x => 1/x.eval }
    def eval = fs.foldLeft(f.eval)(_ * factor(_))
  case class FloatingPoint(x: Any) extends Factor { def eval = x.toString.toDouble  }
  case class Parentheses(e: Expr) extends Factor { def eval = e.eval }
  def expr: Parser[Expr] = term~rep("+"~term | "-"~term) ^^ { case t~r => r match {case x:
List[String~Term] => Expr(t,x)}}
  def term: Parser[Term] = factor\sim rep("*"\sim factor | "/"\sim factor) \land \land \{ case f\sim r \Rightarrow r match \{ case x : r match \} \}
List[String~Factor] => Term(f,x)}
  def factor: Parser[Factor] = (floatingPointNumber | "("~expr~")") ^^ { case "("~e~")" => e match
{case x: Expr => Parentheses(x)}; case s => FloatingPoint(s) }
```