Due Wednesday, January 18th, 4:00 pm in 2131 Kemper

For inductive proofs, you must state when you use your inductive hypothesis in your proof.

- 1. (2 points) Let F_i be the Fibonacci numbers, where $F_0 = 1$, $F_1 = 1$, $F_2 = 2$, $F_3 = 3$, $F_4 = 5$ Prove $\sum_{i=1}^{N-2} F_i = F_N 2$, where N > 2 using induction. (Lipschutz, HW 3)
- 2. (2 points) Use induction to prove that for all natural numbers x > 1 and n, $x^n 1$ is divisible by x 1. (Heileman, p.415)
- 3. (2 points) Prove by induction that $1^2 + 2^2 + 3^2 + ... + n^2 = \frac{n(n+1)(2n+1)}{6}$. (Lipshutz)
- 4. (2 points) Prove by induction that the sum of the first n odd positive integers is n^2 , i.e., $1+3+5+...+(2n-1)=n^2$.
- 5. (2 points) Prove that $\sum_{k=0}^{n} k2^k = (n-1)2^{n+1} + 2$. (Heileman, p. 415)
- 6. (2 points) Assuming a and b are arbitrary constants, and 0 < a < 1 < b, order the following functions by growth rate: $lg\ n,\ log(log\ n),\ n\ log n,\ n^b,\ n^a,\ \frac{1}{n'},\frac{1}{\log n},\ n^n,\ b^n,\ 1,\ n^{\log n},b^{b^n},\ \frac{1}{b^n}$. (Heileman p. 27)
- 7. (2 points) The number of operations executed by algorithms *A* and *B* is $8n\log n$ and $2n^2$, respectively. Determine an n_0 such that A < B for $n > n_0$. (Goodrich, p.185)
- 8. (4 points) Prove that $x^2 + 3x 10$ is $\Theta(N^2)$. This actually involves two proofs. (Lipshutz HW 3)
- 9. (4 points, 2 points each) Assuming that $f_1(n)$ is $O(g_1(n))$ and $f_2(n)$ is $O(g_2(n))$: (Drozdek, p.71)
 - a. Find a counterexample to refute that $f_1(n) f_2(n)$ is $O(g_1(n) g_2(n))$ by supplying f_1 , f_2 , g_1 , and g_2 .
 - b. Find a counterexample to refute that $f_1(n)/f_2(n)$ is $O(g_1(n)/g_2(n))$ by supplying f_1 , f_2 , g_1 , and g_2 .
- 10. (2 points) Show that log(n + 1) = O(log n). (Heileman p.28)
- 11. (12 points, 2 points each) Find the computational complexity for the following code fragments: (a Nyhoff p. 364, b-e Drozdek pp. 72-73, f Weiss, p. 72)

```
a. for(int x = 1, count = 0, i = 0; i < n; i++)
{
    for(int j = 0; j <= x; j++)
        count++;
    x = 2;
}
b. for(int count = 0, i = 0; i < n; i++)
    for(int j = 0; j < n; j++)
        count++;

c. for(int count = 0, i = 0; i < n; i++)
    for(int j = 0; j < i; j++)
        count++;

d. for(int count = 0, i = 1; i < n; i *= 2)
    for(int j = 0; j < n; j++)
        count++;

e. for(int count = 0, i = 1; i < n; i *= 2)</pre>
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for (int j = 0; j < i; j++)

count++;

```
f. for(int count = 0, i = 0; i < n * n; i++)
    if( i % n == 0)
    for(int j = 0; j < i; j++)
        count++;</pre>
```

- 12. (4 points, 2 points each) Let p(x) be a polynomial of degree n, that is, $p(x) = \sum_{i=0}^{n} a_i x^i$. (Goodrich, p. 190)
 - a. Describe a simple $O(n^2)$ time method for computing p(x).
 - b. Now consider a rewriting of p(x) as $p(x) = a_0 + x(a_1 + x(a_2 + x(a_3 + ... + x(a_{n-1} + xa_n) ...)))$, which is known as Horner's method. Using the big-Oh notation, characterize the number of arithmetic operations this method executes.
- 13. (6 points, 2 points each) Evaluate the following sums: (Weiss, p. 47)
 - a. $\sum_{i=0}^{\infty} \frac{1}{4^i}$
 - b. $\sum_{i=0}^{\infty} \frac{i}{4^i}$
 - c. $\sum_{i=0}^{\infty} \frac{i^2}{4^i}$

Sources of questions:

Thomas H. Cormen, Charles E. Leiserson, and Ronald L. Rivest, *Introduction to Algorithms*, New York, New York, McGraw-Hill, 1990.

Adam Drozdek, *Data Structures and Algorithms in C++*, *Second Edition*, Pacific Grove, CA, Brooks/Cole, 2001.

Michael T. Goodrich, Roberto Tamassia, and David Mount, *Data Structures & Algorithms*, *Second Edition*, Hoboken, NJ, John Wiley & Sons, 2011.

Gregory L. Heileman, *Data Structures*, *Algorithms*, and *Object Oriented Programming*, New York, NY, McGraw-Hill, 1996.

Seymour Lipschutz, Schaum's Outline of Theory and Problems of Discrete Mathematics, 3rd ed., NewYork, NY, McGraw-Hill, 2007.

Larry R. Nyhoff, C++: An Introduction to Data Structures, Upper Saddle River, NJ, 1999.

Mark Weiss, *Data Structures and Algorithm Analysis in C++*, *Fourth Edition*, New York, NY, Pearson Education, 2014.