

NORMAL FORMS

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Normalizing Data

- Why not just put it all in one table?

REDUNDANCY!
INCONSISTENCY RISK!

ONE_BIG_TABLE							
Id	SName	GradYear	DName	Title	Prof	YearOffered	Grade
1	Joe	2021	Chemistry	Calculus	Dan	2017	A
2	Alice	2017	Chemistry	Calculus	Dan	2017	A-
3	Bob	2020	Chemistry	Calculus	Dan	2017	A
4	Dan	2018	Chemistry	Calculus	Dan	2017	B+
5	Dean	2023	Math	Thermodynamics	Ed	2017	A-
6	Chuck	2017	Math	Thermodynamics	Ed	2017	B+
7	Ed	2020	Math	Thermodynamics	Ed	2017	A
8	Frank	2018	Math	Thermodynamics	Ed	2017	A-

Schema Tables

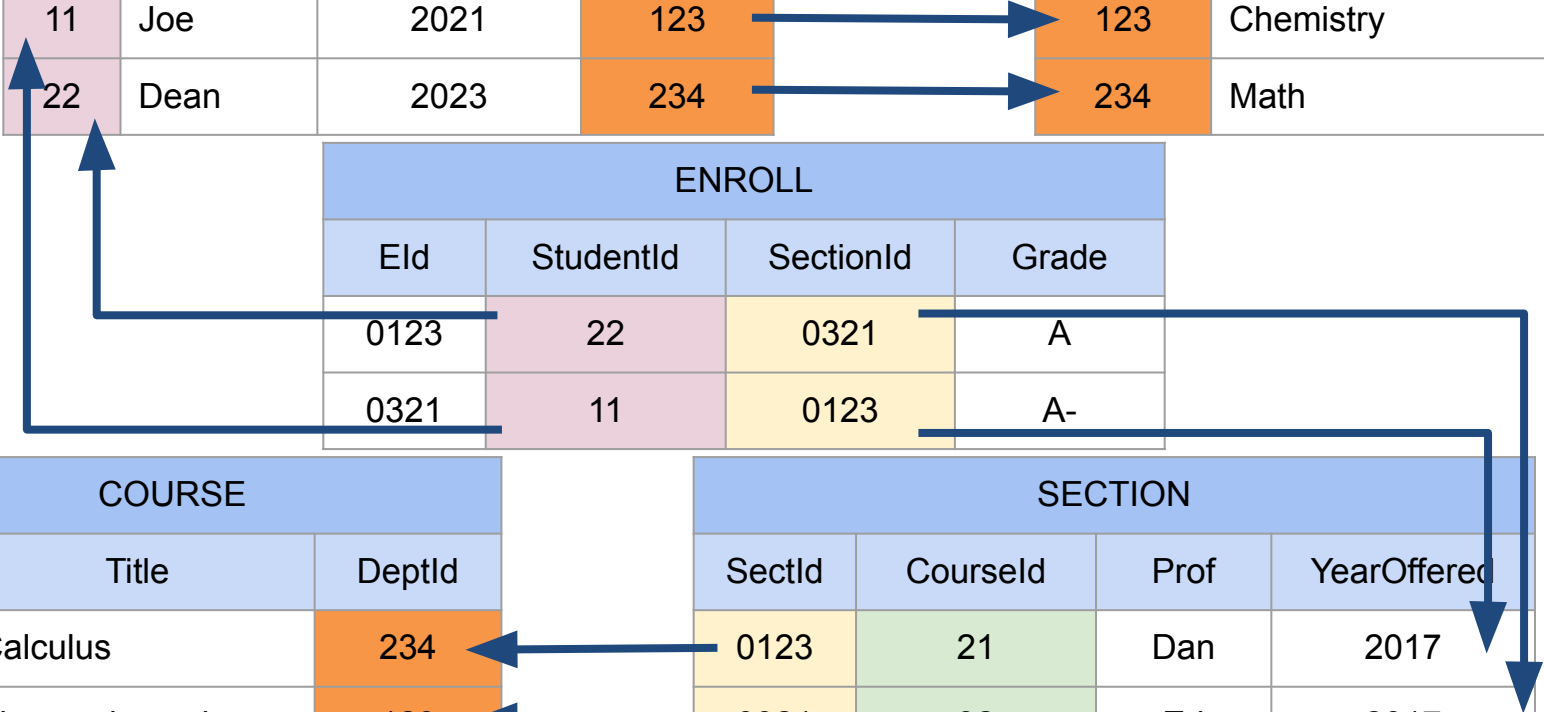
STUDENT			
SId	SName	GradYear	MajorId
11	Joe	2021	123
22	Dean	2023	234

DEPARTMENT	
DId	DName
123	Chemistry
234	Math

ENROLL			
EId	StudentId	SectionId	Grade
0123	22	0321	A
0321	11	0123	A-

COURSE		
CId	Title	DeptId
21	Calculus	234
32	Thermodynamics	123

SECTION			
SectId	CourseId	Prof	YearOffered
0123	21	Dan	2017
0321	32	Ed	2017



Normalization

- Organize fields and tables *minimizing redundancy & dependency*
- Usually *split* large tables *into several* smaller *tables*
- And *define relations* between them

Normalization

- Objective is *decouple* data so *updates* (insertions, deletions, & updates) are to *individual tables*
- Changes propagate through relations

First Normal Form (1NF)

- A table is in first normal form if it contains **no repeating fields / columns**
- Repeating fields
 - **waste disk space**
 - **less flexible**
 - **more difficult to search**
- **1NF – data must depend on the key**

First Normal Form (1NF)

- Consider the following *Customer* table

<u>CID</u>	FirstName	Surname	Telephone
123	Alice	Wonderland	555-123-2345
456	Billy	Squire	555-234-3455
789	Daniel	Craig	555-345-5433

First Normal Form Example (1/3)

- Now consider adding additional phones
- We can put them all in a single field

Telephoness

<u>CID</u>	FirstName	Surname	Telephones
123	Alice	Wonderland	555-123-2345
456	Billy	Squire	555-123-3455, 555-234-3455
789	Daniel	Craig	555-345-5433

First Normal Form Example (1/3)

- Or separate columns: Telephone1, 2.
But what about 3, 4, ...?

<u>CID</u>	FirstName	Surname	Telephone1	Telephone2
123	Alice	Wonderland	555-123-2345	
456	Billy	Squire	555-123-3455	555-234-3455
789	Daniel	Craig	555-345-5433	

First Normal Form Example (2/3)

- Or we can have separate records

<u>CID</u>	FirstName	Surname	Telephone
123	Alice	Wonderland	555-123-2345
456	Billy	Squire	555-123-3455
457	Billy	Squire	555-234-3455
789	Daniel	Craig	555-345-5433

- But this introduces *redundancy* → inconsistency

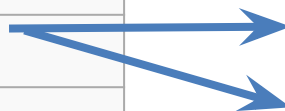
First Normal Form Example (3/3)

- Remove redundant row

<u>CID</u>	FirstName	Surname	Telephone
123	Robert	Ingram	555-861-2025
456	Jane	Wright	555-403-1659
457	Jane	Wright	555-776-4100
789	Maria	Fernandez	555-808-9633

- Revised *Customer* table and new *Telephones* table

<u>CID</u>	FirstName	Surname
123	Robert	Ingram
456	Jane	Wright
789	Maria	Fernandez



<u>CID</u>	<u>Telephone</u>
123	555-861-2025
456	555-403-1659
456	555-776-4100
789	555-808-9633

2NF – The Whole Key

- Second normal form refactors tables that have columns that do not **depend on whole** primary key
- Consider the following table

<u>CID</u>	<u>Semester</u>	#Seats	CourseName
CS1500	2009-1	100	Programming
CS1500	2009-2	150	Programming
CS5200	2009-1	200	Databases
CS5200	2010-1	150	Databases
CS5610	2009-2	120	Web Design

- **CourseName** column does not **depend on whole** PK, only on one of the keys: CID

Achieving 2NF

- We achieve 2NF using the same technique as 1NF

<u>CID</u>	<u>Semester</u>	#Seats
CS1500	2009-1	100
CS1500	2009-2	100
CS5200	2009-1	200
CS5200	2010-1	150
CS5610	2009-2	120

- New Course table referencing table above

CID	CourseName
CS1500	Programming
CS5200	Databases
CS5610	Web Design

Another 2NF Example (1/2)

- Consider the following Skills table

<u>Employee</u>	<u>Skill</u>	Current Location
Alice	QA	Cambridge
Alice	CSS	Cambridge
Alice	HTML	Cambridge
Bob	jQuery	Newton
Charly	Java	Newton
Charly	C#	Newton
Dan	jQuery	Newton

- $\{\underline{\text{Employee}}, \underline{\text{Skill}}\}$ is PK, but *Location* depends only on Employee, not Skill → redundancy

Another 2NF Example (2/2)

- Split into several tables

<u>Employee</u>	<u>Skill</u>
Alice	QA
Alice	CSS
Alice	HTML
Bob	jQuery
Charly	Java
Charly	C#
Dan	jQuery

<u>Employee</u>	<u>Current Location</u>
Alice	Cambridge
Bob	Newton
Charly	Newton
Dan	Newton

But 2NF May Not be Enough

- This 2NF table could have problems

<u>Course</u>	<u>Year</u>	Prof	Prof Date of Birth
Web Dev	2012	Alice Wonder	26 November 1996
<u>Intro to DB</u>	2012	Bob Marley	6 February 1945
C++	2011	Charly Garcia	23 October 1951
Web Dev	2013	Bob Marley	6 February 1945
<u>Intro to DB</u>	2013	Alice Wonder	26 November 1996

But 2NF May Not be Enough

- This 2NF table could have problems

<u>Course</u>	<u>Year</u>	Prof	Prof Date of Birth
Web Dev	2012	Alice Wonder	26 November 1865
<u>Intro to DB</u>	2012	Bob Marley	6 February 1945
C++	2011	Charly Garcia	23 October 1951
Web Dev	2013	Bob Marley	6 February 1945
<u>Intro to DB</u>	2013	Alice Wonder	26 November 1865

- **{Course, Year}** is PK and other fields depend on them → 2NF

But 2NF May Not be Enough

- This 2NF table could have problems

<u>Course</u>	<u>Year</u>	Prof	Prof Date of Birth
Web Dev	2012	Alice Wonder	26 November 1865
<u>Intro to DB</u>	2012	Bob Marley	6 February 1945
C++	2011	Charly Garcia	23 October 1951
Web Dev	2013	Bob Marley	6 February 1945
<u>Intro to DB</u>	2013	Alice Wonder	26 November 1865

- **{Course, Year}** is PK and other fields depend on them → 2NF
- But, we can mistakenly change Prof's birth date

But 2NF May Not be Enough

- This 2NF table could have problems

<u>Course</u>	<u>Year</u>	Prof	Prof Date of Birth
Web Dev	2012	Alice Wonder	26 November 1865
<u>Intro to DB</u>	2012	Bob Marley	6 February 1945
C++	2011	Charly Garcia	23 October 1951
Web Dev	2013	Bob Marley	6 February 1945
<u>Intro to DB</u>	2013	Alice Wonder	26 November 1865

- **{Course, Year}** is PK and other fields depend on them → 2NF
- But, we can mistakenly change Prof's birth date
- Fix with 3NF

3NF – Nothing But the Key

- Every non-key attribute must provide a fact about ***the key, the whole key, and nothing but the key***

<u>Course</u>	<u>Year</u>	Prof	Prof Date of Birth
Web Dev	2012	Alice Wonder	26 November 1865
<u>Intro to DB</u>	2012	Bob Marley	6 February 1945
C++	2011	Charly Garcia	23 October 1951
Web Dev	2013	Bob Marley	6 February 1945
<u>Intro to DB</u>	2013	Alice Wonder	26 November 1865

- To fix, again, split tables into separate related tables

Achieving 3NF

- Removing birth date from table:

<u>Course</u>	<u>Year</u>	Prof
Web Dev	2012	Alice Wonder
<u>Intro to DB</u>	2012	Bob Marley
C++	2011	Charly Garcia
Web Dev	2013	Bob Marley
<u>Intro to DB</u>	2013	Alice Wonder

- And moving it to a dedicated table

Prof	Prof Date of Birth
Alice Wonder	26 November 1865
Bob Marley	6 February 1945
Charly Garcia	23 October 1951