
Differentially Private Bayesian Inference

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Abstract

1. Setting up

The Bayesian inference process is denoted as $\text{BI}(x, \text{prior})$ taking an observed data set $x \in \mathcal{X}^n$ and a prior distribution as input, outputting a posterior distribution *posterior*. For conciseness, when prior is given, we use $\text{BI}(x)$.

For now, we already have a prior distribution *prior*, an observed data set x .

1.1. Exponential Mechanism with Global Sensitivity

1.1.1. MECHANISM SET UP

In exponential mechanism, candidate set R can be obtained by enumerating $y \in \mathcal{X}^n$, i.e.

$$R = \{\text{BI}(y) \mid y \in \mathcal{X}^n\}.$$

Hellinger distance H is used here to score these candidates. The utility function:

$$u(x, r) = -H(\text{BI}(x), r); r \in R. \quad (1)$$

Exponential mechanism with global sensitivity selects and outputs a candidate $r \in R$ with probability proportional to $\exp(\frac{\epsilon u(x, r)}{2\Delta_g u})$:

$$P[r] = \frac{\exp(\frac{\epsilon u(x, r)}{2\Delta_g u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_g u})},$$

where global sensitivity is calculated by:

$$\Delta_g u = H(\text{BI}(x'), r) - H(\text{BI}(y'), r) \max_{\{x', y' \leq 1; x', y' \in \mathcal{X}^n\}} \max_{\{r \in R\}}.$$

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1.1.2. SECURITY ANALYSIS

It can be proved that exponential mechanism with global sensitivity is ϵ -differentially private. We denote the BI with privacy mechanism as PrivInfer . For adjacent data set $\|x, y\|_1 = 1$:

$$\begin{aligned} & \frac{P[\text{PrivInfer}(x, u, R) = r]}{P[\text{PrivInfer}(y, u, R) = r]} \\ &= \frac{\exp(\frac{\epsilon u(x, r)}{2\Delta_g u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_g u})} \\ &= \frac{\exp(\frac{\epsilon u(y, r)}{2\Delta_g u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(y, r')}{2\Delta_g u})} \\ &= \left(\frac{\exp(\frac{\epsilon u(x, r)}{2\Delta_g u})}{\exp(\frac{\epsilon u(y, r)}{2\Delta_g u})} \right) \cdot \left(\frac{\sum_{r' \in R} \exp(\frac{\epsilon u(y, r')}{2\Delta_g u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_g u})} \right) \\ &= \exp\left(\frac{\epsilon(u(x, r) - u(y, r))}{2\Delta_g u}\right) \\ &\quad \cdot \left(\frac{\sum_{r' \in R} \exp(\frac{\epsilon u(y, r')}{2\Delta_g u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_g u})} \right) \\ &\leq \exp\left(\frac{\epsilon}{2}\right) \cdot \exp\left(\frac{\epsilon}{2}\right) \cdot \left(\frac{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_g u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_g u})} \right) \\ &= \exp(\epsilon). \end{aligned}$$

Then, $\frac{P[\text{PrivInfer}(x, u, R) = r]}{P[\text{PrivInfer}(y, u, R) = r]} \geq \exp(-\epsilon)$ can be obtained by symmetry.

1.2. Exponential Mechanism with Local Sensitivity

1.2.1. MECHANISM SET UP

Exponential mechanism with local sensitivity share the same candidate set and utility function as it with global sensitivity. This outputs a candidate $r \in R$ with probability proportional to $\exp(\frac{\epsilon u(x, r)}{2\Delta_l u})$:

$$P[r] = \frac{\exp(\frac{\epsilon u(x, r)}{2\Delta_l u})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_l u})},$$

where local sensitivity is calculated by:

$$\Delta_l u(x) = H(\text{BI}(x), r) - H(\text{BI}(y'), r) |$$

$$\max_{\{x, y' | \leq 1; y' \in \mathcal{X}^n\}} \max_{\{r \in R\}}.$$

1.2.2. SECURITY ANALYSIS

We will then prove that exponential mechanism with local sensitivity is non-differentially private.

$$\frac{P[\text{PrivInfer}(x, u, R) = r]}{P[\text{PrivInfer}(y, u, R) = r]}$$

$$= \exp\left(\frac{\epsilon u(x, r)}{2\Delta_l u(x)} - \frac{\epsilon u(y, r)}{2\Delta_l u(y)}\right) \cdot \left(\frac{\sum_{r' \in R} \exp(\frac{\epsilon u(y, r')}{2\Delta_l u(y)})}{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r')}{2\Delta_l u(x)})}\right)$$

$$= \frac{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r)}{2\Delta_l u(x)} + \frac{\epsilon u(y, r')}{2\Delta_l u(y)})}{\sum_{r' \in R} \exp(\frac{\epsilon u(y, r)}{2\Delta_l u(y)} + \frac{\epsilon u(x, r')}{2\Delta_l u(x)})}.$$

Without loss of generality, we consider the case that $\Delta_l u(y) < \Delta_l u(x)$, $r = \arg(\max_{r' \in R} \{u(x, r')\}) = \arg(\min_{r' \in R} \{u(y, r')\})$ and $\Delta_l u(y) = u(x, r) - u(y, r)$. We have:

$$\frac{\sum_{r' \in R} \exp(\frac{\epsilon u(x, r)}{2\Delta_l u(x)} + \frac{\epsilon u(y, r')}{2\Delta_l u(y)})}{\sum_{r' \in R} \exp(\frac{\epsilon u(y, r)}{2\Delta_l u(y)} + \frac{\epsilon u(x, r')}{2\Delta_l u(x)})}$$

$$> \frac{\sum_{r' \in R} \exp(\frac{\epsilon(u(x, r) + u(y, r'))}{2\Delta_l u(x)})}{\sum_{r' \in R} \exp(\frac{\epsilon(u(y, r) + u(x, r'))}{2\Delta_l u(y)})}$$

$$> \frac{|R| \exp(\frac{\epsilon(u(x, r) + u(y, r))}{2\Delta_l u(x)})}{|R| \exp(\frac{\epsilon(u(y, r) + u(x, r))}{2\Delta_l u(y)})}$$

$$= \exp(\frac{\epsilon}{2}(\frac{u(x, r) + u(y, r)}{\Delta_l u(x)} - \frac{u(x, r) + u(y, r)}{\Delta_l u(y)})).$$

From Eq. 1, $\{u(x, r') \leq 0 | r' \in R\}$ and $\{u(y, r') \leq 0 | r' \in R\}$, we can infer that $r = \arg(\max_{r' \in R} \{u(x, r')\}) = \text{BI}(x)$ and $u(x, r) = 0$. From $\Delta_l u(y) = u(x, r) - u(y, r)$, we can also infer that $\Delta_l u(y) = -u(y, r)$. Then, the following relationship between $u(x, r)$, $u(y, r)$, $\Delta_l u(x)$ and $\Delta_l u(y)$:

$$-\Delta_l u(x) < \Delta_l u(y)$$

$$\Delta_l u(x) - \Delta_l u(y) < 2\Delta_l u(x)$$

$$-\Delta_l u(y)(\Delta_l u(y) - \Delta_l u(x)) < 2\Delta_l u(x)\Delta_l u(y)$$

$$u(y, r)(\Delta_l u(y) - \Delta_l u(x)) < 2\Delta_l u(x)\Delta_l u(y)$$

$$\frac{u(x, r) + u(y, r)}{\Delta_l u(x)} - \frac{u(x, r) + u(y, r)}{\Delta_l u(y)} > 2.$$

holds.

Then we can have:

$$\exp(\frac{\epsilon}{2}(\frac{u(x, r) + u(y, r)}{\Delta_l u(y)} - \frac{u(x, r) + u(y, r)}{\Delta_l u(x)}))$$

$$> \exp(\frac{\epsilon}{2} * 2)$$

$$= \exp(\epsilon),$$

i.e.

$$\frac{P[\text{PrivInfer}(x, u, R) = r]}{P[\text{PrivInfer}(y, u, R) = r]} > \exp(\epsilon).$$

Since there are cases where exponential mechanism with local sensitivity's privacy loss is greater than e^ϵ . we can say it is non-differentially private.

1.3. Exponential Mechanism of Varying Sensitivity

1.3.1. MECHANISM SETTING UP

1.3.2. SECURITY ANALYSIS

1.4. Exponential Mechanism of Smooth Sensitivity

1.4.1. MECHANISM SETTING UP

1.4.2. SECURITY ANALYSIS

2. Privacy Fix

2.1. Propositions

Assume we have a prior distribution $\text{beta}(1, 1)$, an observed data set $x \in \{0, 1\}^n$, $n > 0$. We use the $x + 1$ and $x - 1$ to denote:

if $\text{BayesInfer}(x) = \text{beta}(a_1 + 1, b_1 + 1)$

then $\text{BayesInfer}(x + 1) = \text{beta}((a_1 + 1) + 1, (b_1 - 1) + 1)$

$\text{BayesInfer}(x - 1) = \text{beta}((a_1 - 1) + 1, (b_1 + 1) + 1)$,

x_0 to denote:

if n is even

then $\text{BI}(x_0) = \text{beta}(\frac{n}{2} + 1, \frac{n}{2} + 1)$

else $\text{BI}(x_0) = \{\text{beta}(\frac{n+1}{2} + 1, \frac{n-1}{2} + 1), \text{beta}(\frac{n-1}{2} + 1, \frac{n+1}{2} + 1)\}$

$\text{beta}(\alpha, \beta)$ is the beta function with two arguments α and β .

Then, we have the following three statements, and proofs of the statements.

I $H(\text{BI}(x), \text{BI}(x + 1)) < H(\text{BI}(x + 1), \text{BI}(x + 2)) \forall x \geq x_0$;

or $H(\text{BI}(x), \text{BI}(x+1)) > H(\text{BI}(x+1), \text{BI}(x+2)) \forall x \leq x_0$.

II $\Delta_l u(x) = H(\text{BI}(x), \text{BI}(x+1)), \forall x \geq x_0$;

$\Delta_l u(x) = H(\text{BI}(x), \text{BI}(x-1)), \forall x \leq x_0$.

III $\forall x \neq x_0 : \Delta_l u(x) > \Delta_l u(x_0)$.

2.2. proof

2.2.1. STATEMENT I

We use the MI (Mathematical Induction) method to prove the first statement.

Proof. Since the Hellinger distance is symmetric, if we prove the $H(\text{BI}(x), \text{BI}(x+1)) < H(\text{BI}(x+1), \text{BI}(x+2)) \forall x \geq x_0$, the other part when $\forall x \leq x_0$ also holds.

1. if $x = x_0$, $H(\text{BI}(x_0), \text{BI}(x_0+1)) < H(\text{BI}(x_0+1), \text{BI}(x_0+2))$ holds:

$$\begin{aligned} & \sqrt{1 - \frac{\text{beta}(\frac{\frac{n}{2}+1+m+\frac{n}{2}+1+m+1}{2}, \frac{\frac{n}{2}+1-m+\frac{n}{2}+1-m-1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+1+m, \frac{n}{2}+1-m)\text{beta}(\frac{n}{2}+2+m, \frac{n}{2}-m)}}} \\ & < \sqrt{1 - \frac{\text{beta}(\frac{\frac{n}{2}+1+m+1+\frac{n}{2}+1+m+2}{2}, \frac{\frac{n}{2}+1-m-1+\frac{n}{2}+1-m-2}{2})}{\sqrt{\text{beta}(\frac{n}{2}+2+m, \frac{n}{2}-m)\text{beta}(\frac{n}{2}+3+m, \frac{n}{2}-m-1)}}} \\ & \frac{\text{beta}(\frac{n+2m+3}{2}, \frac{n-2m+1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+1+m, \frac{n}{2}+1-m)\text{beta}(\frac{n}{2}+2+m, \frac{n}{2}-m)}} \\ & > \frac{\text{beta}(\frac{n+2m+5}{2}, \frac{n-2m-1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+2+m, \frac{n}{2}-m)\text{beta}(\frac{n}{2}+3+m, \frac{n}{2}-m-1)}} \end{aligned}$$

Now, we need to proof $H(\text{beta}(\frac{n}{2}+1+m+1, \frac{n}{2}+1-m-1), \text{beta}(\frac{n}{2}+1+m+2, \frac{n}{2}+1-m-2)) < H(\text{beta}(\frac{n}{2}+1+m+3, \frac{n}{2}+1-m-3))$ by using what we know.

From $x = x_0 + m$ and property of $\text{beta}(\alpha, \beta)$ function, we know:

$$\begin{aligned} & H(\text{beta}(\frac{n}{2}+1, \frac{n}{2}+1), \text{beta}(\frac{n}{2}+1+1, \frac{n}{2}+1-1)) < H(\text{beta}(\frac{n}{2}+1+1, \frac{n}{2}+1-1), \text{beta}(\frac{n}{2}+1+2, \frac{n}{2}+1-2)) \\ & \sqrt{1 - \frac{\text{beta}(\frac{\frac{n}{2}+1+\frac{n}{2}+1+1}{2}, \frac{\frac{n}{2}+1+\frac{n}{2}+1-1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+1, \frac{n}{2}+1)\text{beta}(\frac{n}{2}+1+1, \frac{n}{2}+1-1)}}} < \sqrt{1 - \frac{\text{beta}(\frac{\frac{n}{2}+1+1+\frac{n}{2}+1+2}{2}, \frac{\frac{n}{2}+1-1+\frac{n}{2}+1-2}{2})}{\sqrt{\text{beta}(\frac{n}{2}+1+1, \frac{n}{2}+1-1)\text{beta}(\frac{n}{2}+1+2, \frac{n}{2}+1-2)}}} \\ & \sqrt{1 - \frac{\text{beta}(\frac{n+3}{2}, \frac{n+1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+1, \frac{n}{2}+1)\text{beta}(\frac{n}{2}+2, \frac{n}{2})}}} < \sqrt{1 - \frac{\text{beta}(\frac{n+5}{2}, \frac{n-1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+2, \frac{n}{2})\text{beta}(\frac{n}{2}+3, \frac{n}{2}-1)}}} \\ & \frac{\text{beta}(\frac{n+3}{2}, \frac{n+1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+1, \frac{n}{2}+1)\text{beta}(\frac{n}{2}+2, \frac{n}{2})}} > \frac{\text{beta}(\frac{n+5}{2}, \frac{n-1}{2})}{\sqrt{\text{beta}(\frac{n}{2}+2, \frac{n}{2})\text{beta}(\frac{n}{2}+3, \frac{n}{2}-1)}} \\ & \frac{\text{beta}(\frac{n+3}{2}, \frac{n-1}{2})^{\frac{\frac{n-1}{2}+\frac{n+3}{2}}{\frac{n}{2}-1+\frac{n}{2}+1}}}{\sqrt{\text{beta}(\frac{n}{2}+1, \frac{n}{2}-1)^{\frac{\frac{n}{2}-1}{\frac{n}{2}-1+\frac{n}{2}+1}}\text{beta}(\frac{n}{2}+2, \frac{n}{2}+1)^{\frac{\frac{n}{2}}{\frac{n}{2}-1+\frac{n}{2}+1}}}} > \frac{\text{beta}(\frac{n+3}{2}, \frac{n-1}{2})^{\frac{\frac{n+3}{2}}{\frac{n+3}{2}+\frac{n-1}{2}}}}{\sqrt{\text{beta}(\frac{n}{2}+1, \frac{n}{2}-1)^{\frac{\frac{n}{2}+1}{\frac{n}{2}+1+\frac{n}{2}-1}}\text{beta}(\frac{n}{2}+2, \frac{n}{2}-1)^{\frac{\frac{n}{2}+2}{\frac{n}{2}+1+\frac{n}{2}-1}}}} \\ & \frac{\frac{n-1}{2}}{\sqrt{(\frac{n}{2}-1)(\frac{n}{2})}} > \frac{\frac{n+3}{2}}{\sqrt{(\frac{n}{2}+1)(\frac{n}{2}+2)}} \\ & (n-1)^2(n+2)(n+4) > (n+3)^2n(n-2) \\ & n > -1. \end{aligned}$$

Since $n > 0$, it always holds.

2. if $x = x_0 + m$ holds, then also $x = x_0 + m + 1$ holds:

i.e $H(\text{beta}(\frac{n}{2}+1+m, \frac{n}{2}+1-m), \text{beta}(\frac{n}{2}+1+m+1, \frac{n}{2}+1-m-1)) < H(\text{beta}(\frac{n}{2}+1+m+1, \frac{n}{2}+1-m-1), \text{beta}(\frac{n}{2}+1+m+2, \frac{n}{2}+1-m-2))$ is what we know:

$$\begin{aligned} & \frac{\text{beta}(\frac{n+2m+5}{2}, \frac{n-2m-1}{2})^{\frac{n-2m-1}{n+2m+3}}}{\sqrt{\text{beta}(\frac{n}{2}+2+m, \frac{n}{2}-m)\text{beta}(\frac{n}{2}+3+m, \frac{n}{2}-m-1)^{\frac{n-2m}{n+2m+2}}}} \\ & > \frac{\text{beta}(\frac{n+2m+7}{2}, \frac{n-2m-3}{2})^{\frac{n-2m-3}{n+2m+5}}}{\sqrt{\text{beta}(\frac{n}{2}+2+m, \frac{n}{2}-m)\text{beta}(\frac{n}{2}+3+m, \frac{n}{2}-m-1)^{\frac{n-2m-2}{n+2m+6}}}} \end{aligned}$$

2.2.3. STATEMENT III

Proof. From Statement I and Statement II, we can conclude that:

$$\begin{aligned} \text{when } x > x_0 \\ & H(\text{BI}(x), \text{BI}(x+1)) \\ & > H(\text{BI}(x_0), \text{BI}(x_0+1)); \\ & \text{i.e. } \Delta_l u(x) > \Delta_l u(x_0) \\ \text{when } x < x_0 \\ & H(\text{BI}(x), \text{BI}(x-1)) \\ & > H(\text{BI}(x_0), \text{BI}(x_0-1)); \\ & \text{i.e. } \Delta_l u(x) > \Delta_l u(x_0). \end{aligned}$$

□

i.e. $x = x_0 + m + 1$ also holds when $x = x_0 + m$ is valid. □

2.2.2. STATEMENT II

Proof.

$$\begin{aligned} \because \Delta_l u(x) &= |H(\text{BI}(x), r) - H(\text{BI}(y'), r)|, \\ & \max_{\{|x, y'| \leq 1; y' \in \mathcal{X}^n\}} \max_{\{r \in R\}}; \\ \because H(\text{BI}(x), r) - H(\text{BI}(y'), r) \\ & \leq H(\text{BI}(x), \text{BI}(y')); \\ \therefore \Delta_l u(x) &= H(\text{BI}(x), \text{BI}(y')), \\ & \max_{\{|x, y'| \leq 1; y' \in \mathcal{X}^n\}}; \\ \therefore \Delta_l u(x) &= \max\{H(\text{BI}(x), \text{BI}(x+1)), \\ & H(\text{BI}(x), \text{BI}(x-1))\}; \end{aligned}$$

According to Statement I :

$$\begin{aligned} \text{if } x > x_0 \\ \text{then } & H(\text{BI}(x), \text{BI}(x-1)) \\ & < H(\text{BI}(x), \text{BI}(x+1)); \\ \text{then } & \Delta_l u(x) = H(\text{BI}(x), \text{BI}(x+1)); \\ \text{if } x < x_0 \\ \text{then } & H(\text{BI}(x), \text{BI}(x-1)) \\ & > H(\text{BI}(x), \text{BI}(x+1)); \\ \text{then } & \Delta_l u(x) = H(\text{BI}(x), \text{BI}(x-1)); \\ \text{else } & \Delta_l u(x_0) = H(\text{BI}(x_0), \text{BI}(x_0-1)) \\ & = H(\text{BI}(x_0), \text{BI}(x_0+1)). \end{aligned}$$

From above, we can conclude the Statement II. □

3. Experimental Evaluations

We got some results from these mechanisms.

4. Dilation Property of Laplace Noise

Proof. We take 1-dimensional Laplace distribution, $h(z) = \frac{1}{2}e^{-|z|}$. The dilation property is:

$$Pr[z \in S] \leq e^{\frac{\epsilon}{2}} Pr[z \in e^\lambda S] + \frac{\delta}{2}$$

In this case, we have $\alpha = \frac{\epsilon}{2}$, $\beta = \frac{\epsilon}{2\rho_{\delta/3}(|z|)}$ or $\frac{\epsilon}{2ln(2/\delta)}$. We have some prior knowledge that $|\lambda| \leq \beta$.

 • case 1: $\lambda > 0$

$$\begin{aligned} \because h(e^\lambda z) &= \frac{1}{2}e^{-|e^\lambda z|} < \frac{1}{2}e^{-|z|} = h(z) \\ \therefore \frac{Pr[z \in e^\lambda S]}{Pr[z \in S]} &= \frac{\int_{e^\lambda S} \frac{1}{2}e^{-|z|} dz}{\int_S \frac{1}{2}e^{-|z|} dz} = \frac{\int_S \frac{1}{2}e^{-|e^\lambda z|} e^\lambda dz}{\int_S \frac{1}{2}e^{-|z|} dz} \\ &= \frac{e^{-|e^\lambda z|} e^\lambda}{e^{-|z|}} = \frac{e^\lambda h(e^\lambda z)}{h(z)} \leq e^\lambda \\ \therefore \ln\left(\frac{e^\lambda h(e^\lambda z)}{h(z)}\right) &\leq \lambda \\ \because \lambda \leq \beta &= \frac{\epsilon}{2ln(3/\delta)}, \delta < 1 \\ \therefore \lambda &\leq \frac{\epsilon}{2} \\ \therefore \frac{Pr[z \in e^\lambda S]}{Pr[z \in S]} &\leq \frac{\epsilon}{2} \end{aligned}$$

- **case 2:** $\lambda < 0$

$$\because \frac{h(e^\lambda z)}{h(z)} = \exp(|z|(1 - e^\lambda)) \leq |\lambda|$$

$$\therefore \ln\left(\frac{e^\lambda h(e^\lambda z)}{h(z)}\right) \leq |z||\lambda|$$

We consider an event $G = \{z \mid |z| \leq \log(\frac{1}{\delta})\}$. Under this event, we have:

$$|z||\lambda| \leq \log\left(\frac{1}{\delta}\right)|\lambda| \leq \log\left(\frac{1}{\delta}\right) \frac{\epsilon}{2\log(\frac{3}{\delta})} \leq \frac{\epsilon}{2}.$$

Also, from the Laplace distribution property we have:

$$Pr[G] \geq 1 - \delta;$$

$$Pr[\overline{G}] \leq \delta.$$

We start from the thing we want to proof:

$$Pr[z \in S] \leq e^{\frac{\epsilon}{2}} Pr[z \in e^\lambda S] + \frac{\delta}{2};$$

$$\begin{aligned} Pr_{z \sim Lap(1)}[z \in S] &\leq Pr[z \in S \cap G] + Pr[z \in \overline{G}] \\ &\leq Pr[z \in S \cap G] + \delta \\ &\leq e^{|z||\lambda|} Pr_{z \sim Lap(e^\lambda)}[z \in S \cap G] + \delta \\ &\leq e^{|z||\lambda|} Pr_{z \sim Lap(e^\lambda)}[z \in S] + \delta \\ &\leq e^{\frac{\epsilon}{2}} Pr_{z \sim Lap(e^\lambda)}[z \in S] + \delta \\ &= e^{\frac{\epsilon}{2}} Pr_{z \sim Lap(1)}[z \in e^\lambda S] + \delta \end{aligned}$$

□