Verifying Snapping Mechanism

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In order to verify the differential privacy proeprty of snapping mechanism, we follow the logic rules designed from [1].

Some new rules are added into this logic in Figure 1 following with correctness proof. Then we formalized the snapping mechanism and verified its differential privacy property under these logic rules.

1 Logic Rules

Definition 1 (Laplce mechanism [2])

Let $\epsilon > 0$. The Laplace mechanism $\mathcal{L}_{\epsilon} : \mathbb{R} \to \mathsf{Distr}(\mathbb{R})$ is defined by $\mathcal{L}(t) = t + v$, where $v \in \mathbb{R}$ is drawn from the Laplace distribution $\mathsf{Laplce}(\frac{1}{\epsilon})$.

Definition 2

Let $\epsilon \leq 0$. The ϵ -DP divergence $\Delta_{\epsilon}(\mu_1, \mu_2)$ between two sub-distributions $\mu_1 \in \mathsf{Distr}(U)$, $\mu_2 \in \mathsf{Distr}(U)$ is defined as:

$$\sup_{E \in U} \left(\Pr_{x \leftarrow \mu_1} [x \in E] - \exp(\epsilon) \Pr_{x \leftarrow \mu_2} [x \in E] \right)$$

Definition 3 $((\epsilon, \delta)$ - **lifting [1]**)

Two sub-distributions $\mu_1 \in \mathsf{Distr}(U_1)$, $\mu_2 \in \mathsf{Distr}(U_2)$ are related by the (ϵ, δ) - dilation lifting of $\Psi \subseteq U_1 \times U_2$, written $\mu_1 \Psi^{\#(\epsilon, \delta)} \mu_2$, if there exist two witness sub-distributions $\mu_L \in \mathsf{Distr}(U_1 \times U_2)$ and $\mu_R \in \mathsf{Distr}(U_1, U_2)$ s.t.:

- 1. $\pi_1(\mu_L) = \mu_1$ and $\pi_2(\mu_R) = \mu_2$;
- 2. $supp(\mu_L) \subseteq \Psi$ and $supp(\mu_R) \subseteq \Psi$; and
- 3. $\Delta_{\epsilon}(\mu_L, \mu_R) \leq \delta$.

The logic rules we are using in our work is presented in Figure 1. The correctness of rules is proved in Theorem 1 and Theorem 2

Theorem 1

Let $\mu_1 \in \mathsf{Distr}(\mathbb{R})$, $\mu_2 \in \mathsf{Distr}(\mathbb{R})$ are defined:

$$\mu_1(x) = \text{unif}(x)$$

$$\mu_2(y) = \mathsf{unif}(y)$$

$$\frac{}{\vdash u_{1} \stackrel{\$}{\leftarrow} \mu \sim_{\epsilon,0} u_{2} \stackrel{\$}{\leftarrow} \mu : \top \Rightarrow e^{-\epsilon} u_{2} \leq u_{1} \leq e^{\epsilon} u_{2}} \xrightarrow{\text{AXUNIF}} \\
\frac{}{\vdash y_{1} \stackrel{\$}{\leftarrow} \mathcal{L}_{\epsilon}(e_{1}) \sim_{k' \cdot \epsilon, 0} y_{2} \stackrel{\$}{\leftarrow} \mathcal{L}_{\epsilon}(e_{2}) : |k + e_{1} - e_{2}| \leq k' \Rightarrow y_{1} + k = y_{2}} \xrightarrow{\text{LAPGEN}} \\
\frac{}{\vdash y_{1} \stackrel{\$}{\leftarrow} \mathcal{L}_{\epsilon}(e_{1}) \sim_{0,0} y_{2} \stackrel{\$}{\leftarrow} \mathcal{L}_{\epsilon}(e_{2}) : \top \Rightarrow y_{1} - y_{2} = e_{1} - e_{2}} \xrightarrow{\text{LAPNULL}} \\
\frac{}{\vdash y_{1} \stackrel{\$}{\leftarrow} \mu \sim_{0,0} y_{2} \stackrel{\$}{\leftarrow} \mu : \top \Rightarrow y_{1} = y_{2}} \xrightarrow{\text{AXNULL}} \\
\frac{}{\vdash y_{1} \stackrel{\$}{\leftarrow} \mu \sim_{0,0} y_{2} \stackrel{\$}{\leftarrow} \mu : \top \Rightarrow y_{1} = y_{2}} \xrightarrow{\text{COMP}} \\
\frac{}{\vdash p_{1}; c_{1} \sim_{k + k', 0} p_{2}; c_{2} : \Phi_{1} \Rightarrow \Phi_{2}} \xrightarrow{\text{COMP}}$$

Figure 1: Logic Rules from [1]

where unif is uniform distribution over [0,1) whoes pdf. is defined as:

$$\mathsf{pdf}_{\mathsf{unif}}(x) = \begin{cases} 1 & x \in [0,1) \\ 0 & o.w. \end{cases}.$$

Then, $\mu_1 \Psi^{\#(\epsilon,0)} \mu_2$, where

$$\Psi = \{(x, y) \in \mathbb{R} \times \mathbb{R} | x \cdot e^{-\epsilon} \le y \le x \cdot e^{\epsilon} \}$$

Theorem 2

For any distributions $\mu_1 \in \mathsf{Distr}(\mathbb{R}), \, \mu_2 \in \mathsf{Distr}(\mathbb{R}), \, \mu_1 \Psi^{\#(0,0)} \mu_2$, where

$$\Psi = \{(x, y) \in \mathbb{R} \times \mathbb{R} | x = y\}$$

Proof of Theorem 1. Existing $\mu_L, \mu_R \in \mathsf{Distr}(\mathbb{R} \times \mathbb{R})$:

$$\mu_L(x,y) = \begin{cases} \operatorname{unif}(x) & x \cdot e^{-\epsilon} = y \wedge x \in [0,1) \\ 0 & o.w. \end{cases} \\ \mu_R(x,y) = \begin{cases} \operatorname{unif}(y) & x \cdot e^{-\epsilon} = y \wedge y \in [0,1) \\ 0 & o.w. \end{cases}.$$

Their pdf. are defined:

$$\mathsf{pdf}_{\mu_L}(x,y) = \begin{cases} \mathsf{pdf}_{\mathsf{unif}}(x) & x \cdot e^{-\epsilon} = y \land x \in [0,1) \\ 0 & o.w. \end{cases}$$

$$\mathsf{pdf}_{\mu_R}(x,y) = \begin{cases} \mathsf{pdf}_{\mathsf{unif}}(y) & x \cdot e^{-\epsilon} = y \land y \in [0,1) \\ 0 & o.w. \end{cases}.$$

- $supp(\mu_L) \in \Psi \land supp(\mu_R) \in \Psi$
 - $\sup(\mu_L) \subseteq \Psi$ By definition of the pdf of μ_L , we have: $\Pr_{(x,y) \overset{\$}{\leftarrow} \mu_L} [(x,y) \notin \Psi] = 0.$ Then we can derive $\sup(\mu_L) \in \Psi$

- supp(μ_R) ⊆ Ψ

By definition of the pdf of μ_R , we have: $\Pr_{(x,y) \stackrel{\$}{\leftarrow} \mu_R} [(x,y) \notin \Psi] = 0.$

Then we can derive $supp(\mu_L) \in \Psi$

•
$$\pi_1(\mu_L) = \mu_1 \wedge \pi_2(\mu_R) = \mu_2$$

$$- \pi_1(\mu_L) = \mu_1$$

By definition of the π_1 and pdf of μ_L , we have $\forall x \in \mathbb{R}$:

$$\mathsf{pdf}_{\pi_1(\mu_L)}(x) = \begin{cases} \int_{\mathcal{Y}} \mathsf{pdf}_{\mathsf{unif}}(x) & (x,y) \in \Psi \land x \in [0,1) \\ 0 & o.w. \end{cases} = \begin{cases} \mathsf{pdf}_{\mathsf{unif}}(x) & x \in [0,1) \\ 0 & o.w. \end{cases} = \mathsf{pdf}_{\mu_1}(x)$$

 $- \pi_1(\mu_R) = \mu_2$

Equivalent to showpdf_{$\pi_2(\mu_R)$} = pdf_{μ_2}.

By definition of the π_2 and pdf of μ_R , we have $\forall y \in \mathbb{R}$:

$$\mathsf{pdf}_{\pi_2(\mu_R)}(y) = \begin{cases} \int_{\mathcal{X}} \mathsf{pdf}_{\mathsf{unif}}(y) & (x,y) \in \Psi \land y \in [0,1) \\ 0 & o.w. \end{cases} = \begin{cases} \mathsf{pdf}_{\mathsf{unif}}(y) & y \in [0,1) \\ 0 & o.w. \end{cases} = \mathsf{pdf}_{\mu_2}(y)$$

• $\Delta_{\epsilon}(\mu_L, \mu_R) \leq 0$

By definition of ϵ -DP divergence, we have:

$$\Delta_{\epsilon}(\mu_{L}, \mu_{R}) = \sup_{S} \left(\Pr_{(x,y) \stackrel{\$}{\leftarrow} \mu_{L}} [(x,y) \in S] - e^{\epsilon} \Pr_{(x,y) \stackrel{\$}{\leftarrow} \mu_{R}} [(x,y) \in S] \right)$$
$$= \sup_{S} \left(\int_{(x,y) \in S} \mathsf{pdf}_{\mu_{L}}(x,y) - e^{\epsilon} \int_{(x,y) \in S} \mathsf{pdf}_{\mu_{R}}(x,y) \right)$$

case $S \subseteq \{(x, y) | x \in [0, 1) \land x \cdot e^{-\epsilon} = y\}$:

$$\begin{array}{ll} \Delta_{\epsilon}(\mu_L,\mu_R) &= \int_{(x,y)\in S} \mathsf{pdf}_{\mathsf{unif}}(x) - e^{\epsilon} \int_{(x,y)\in S} \mathsf{pdf}_{\mathsf{unif}}(y) \\ &= \int_{(x,y)\in S} \mathsf{pdf}_{\mathsf{unif}}(x) - e^{\epsilon} \int_{(x,y)\in S} \mathsf{pdf}_{\mathsf{unif}}(x*e^{-\epsilon}) \\ &= \int_{(x,y)\in S} \mathsf{pdf}_{\mathsf{unif}}(x) - e^{\epsilon} * e^{-\epsilon} \int_{(x,y)\in S} \mathsf{pdf}_{\mathsf{unif}}(x) \\ &= 0 \end{array}$$

case $S \subseteq \{(x, y) | x \in [1, e^{\epsilon}) \land x \cdot e^{-\epsilon} = y\}$:

$$\begin{array}{ll} \Delta_{\epsilon}(\mu_L, \mu_R) &= 0 - e^{\epsilon} \int_{(x,y) \in S} \mathsf{pdf}_{\mathsf{unif}}(y) \\ &< 0 \end{array}$$

case o.w.

$$\Delta_{\epsilon}(\mu_L, \mu_R) = 0 - 0 = 0$$

$$\overline{u_{1} \stackrel{\$}{\rightleftharpoons} \mu \sim_{\epsilon,0} u_{2} \stackrel{\$}{\rightleftharpoons} \mu : \top \Rightarrow e^{-\epsilon} u_{2} \leq u_{1} \leq e^{\epsilon} u_{2}} \xrightarrow{\text{AXUNIF}}$$

$$\overline{u_{1} \stackrel{$\rightleftharpoons}{\rightleftharpoons} \mu \sim_{\epsilon,0} u_{2} \stackrel{\$}{\rightleftharpoons} \mu : \top \Rightarrow e^{-\epsilon} u_{2} \leq u_{1} \leq e^{\epsilon} u_{2} \Rightarrow y_{2} - 1 \leq y_{1} \leq 1 + y_{2}} \xrightarrow{\text{AXNULL}}$$

$$\overline{u_{1} \stackrel{$\rightleftharpoons}{\rightleftharpoons} \{-1,1\} \sim_{0,0} u_{2} \stackrel{\$}{\rightleftharpoons} \{-1,1\} : \top \Rightarrow u_{1} = u_{2} \text{AXNULL}} \xrightarrow{\text{AXNULL}}$$

$$\overline{u_{1} = u_{1} \vee u_{1} \vee u_{2} = u_{2} \vee u_{2} : u_{1} = u_{2} \wedge u_{2} \wedge u_{2} = u_{2} \vee u_{2} = u_{2} \vee u_{2} \wedge u_{2} \wedge u_{2} \wedge u_{2} = u_{2} \vee u_{2} \wedge u_{2} \wedge$$

Figure 2: Coupling Derivation of two Snap mechanisms: $Snap(a_1)$, $Snap(a_2)$

2 Formalization

Definition 4 (Snap(a): $A \rightarrow Distr(B)$)

The ideal Snapping mechanism Snap(a) is defined as:

$$u \overset{\$}{\leftarrow} \mu; y = \frac{\ln(u)}{\epsilon}; s \overset{\$}{\leftarrow} \{-1, 1\}; z = s * y; x = f(a); w = x + z; w' = \lfloor w \rfloor_{\Lambda}; r = \mathsf{clamp}_{B}(w')$$

where f is the query function over input $a \in A$, ϵ is the privacy budget, B is the clampping bound and Λ is the rounding argument satisfying $\lambda = 2^k$ where 2^k is the smallest power of 2 greater or equal to the $\frac{1}{\epsilon}$.

Theorem 3

The Snap mechanism is differentially praivate by following derivation in Figure 2.

Definition 5 (ϵ - dilation)

Let $\epsilon \ge 0$. The ϵ -dilation $D_{\epsilon}(\mu_1, \mu_2)$ between two sub-distributions $\mu_1 \in \mathsf{Distr}(U)$, $\mu_2 \in \mathsf{Distr}(U)$ is defined as:

$$\sup_{E \in U} \left(\Pr_{x \leftarrow \mu_1} [x \in E] - \exp(\epsilon) \Pr_{x \leftarrow \mu_2} [x \in \exp(-\epsilon) \cdot E] \right)$$

Proposition 4 $((\epsilon, \delta)$ -differential privacy)

For every pair of sub-distributions $\mu_1 \in \mathsf{Distr}(U), \, \mu_2 \in \mathsf{Distr}(U), \, \mathrm{s.t.}$

$$D_{\epsilon}(\mu_1, \mu_2) \leq \delta$$
,

The snapping mechanism $\mathsf{Snap}(\mu, a) : \mathsf{Distr}(U) \to A \to \mathsf{Distr}(B)$ is (ϵ, δ) - differentially private w.r.t. an adjacency relation Φ for every two adjacent inputs a, a' and μ_1, μ_2

Proof. Followed directly by unfolding the Snap mechanism.

$$\begin{array}{lll} \Pr_{x \leftarrow \mathsf{Snap}(\mu_1,a)}[x=e] & = & \Pr_{u \leftarrow \mu_1}[\lfloor f(a) + \frac{S \cdot \log(u)}{\varepsilon} \rfloor_{\Lambda} = e] \\ & = & \Pr_{u \leftarrow \mu_1}[u \in [\frac{\exp((e - \frac{\lambda}{2} - f(a))\varepsilon)}{S}, \frac{\exp((e + \frac{\lambda}{2} - f(a))\varepsilon)}{S})] \\ & \leq & \exp(\varepsilon) \Pr_{u \leftarrow \mu_2}[u \in \exp(-\varepsilon)[\frac{\exp((e - \frac{\lambda}{2} - f(a))\varepsilon)}{S}, \frac{\exp((e + \frac{\lambda}{2} - f(a))\varepsilon)}{S})] \\ & = & \exp(\varepsilon) \Pr_{u \leftarrow \mu_2}[\lfloor f(a') + \frac{S \cdot \log(u)}{\varepsilon} \rfloor_{\Lambda} = e] \\ & = & \exp(\varepsilon) \Pr_{x \leftarrow \mathsf{Snap}(\mu_2,a')}[x=e] \end{array}$$

References

- [1] Gilles Barthe, Marco Gaboardi, Benjamin Grégoire, Justin Hsu, and Pierre-Yves Strub. Proving differential privacy via probabilistic couplings. In *LICS 2016*.
- [2] Cynthia Dwork, Frank McSherry, Kobbi Nissim, and Adam Smith. Calibrating Noise to Sensitivity in Private Data Analysis. In *TCC*, 2016.