

A Bilevel Edge Computing Architecture and Collaborative Offloading Mechanism for Offshore Buoys Network (Appendix)

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III. PROBLEM SOLVING

B. Joint Iterative Mechanism

2) *Solving for Decision Variables* $\{\mathcal{P}^b\}$: In this subsection, the derivation of the decision variable $\mathcal{P}^b = \{p_{ij}\}_{i \in \mathcal{N}, j \in \mathcal{N}}$ is proposed, when other decision variables are fixed, and the optimization problem P_3 is given as follows.

$$\mathbf{P}_3^{\text{Sub}} : \min_{\mathcal{P}^b} \sum_{i=1}^N \sum_{q=1}^Q E_{iq}^{c1}(p_{i\tilde{\xi}q}) \quad (33a)$$

$$\text{s.t. } C_3, C'_4, C_6 \quad (33b)$$

The objective function in (33a) is can be regarded as the product of two convex functions.

$$E_{iq}^{c1}(p_{i\tilde{\xi}q}) = \left(\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q p_{i\tilde{\xi}q} \right) \cdot \frac{1}{B_0 \log_2 \left(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right)} \quad (34)$$

where $B_0 \left(\log_2 \left(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right) \right)^{-1}$ is a convex function. That's because the outer function of that can be regarded as an inverse

proportional function, and the inner function of that is a monotonically increasing function. According to the composition rules, it is easy to know that such function is a convex function with respect to $\{\mathcal{P}^b\}$, which implies that the objective function in (33a) is can be regarded as the product of two convex functions. Next, the convex upper bound $\tilde{E}_{iq}^{c1}(p_{i\tilde{\xi}q}; \tilde{p}_{i\tilde{\xi}q})$ for non-convex function $E_{iq}^{c1}(p_{i\tilde{\xi}q})$ is derived in (32) based on SCA, where $H_q(p_{i\tilde{\xi}q}) = \frac{1}{2} (\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q p_{i\tilde{\xi}q} + (B_0 \log_2(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2}))^{-1})^2$ holds.

Then, C'_4 constraint, as mentioned in (18), is also a convex function. Similarly, C_3 and C_7 constraints can also be combined as follows.

$$C'_3, C''_6 : \sum_{n=1}^N p_{in} \leq c_6, \quad \forall i \in \mathcal{N} = \mathcal{N}^M \cup \mathcal{N}^Q \quad (35)$$

where constant c_6 is calculated as follows.

$$c_6 = \begin{cases} P_0, & \text{for } C_3 \text{ constraint} \\ P_0 - \sum_{\omega=1}^W g_{ms}^\omega, & \text{for } C_7 \text{ constraint} \end{cases} \quad (36)$$

$$\begin{aligned} E_{iq}^{c1}(p_{i\tilde{\xi}q}) &= \frac{\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q \cdot p_{i\tilde{\xi}q}}{B_0 \log_2 \left(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right)} = \left(\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q \cdot p_{i\tilde{\xi}q} \right) \cdot \left(\frac{1}{B_0 \log_2 \left(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right)} \right) \\ &= \frac{1}{2} \left(\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q p_{i\tilde{\xi}q} + \frac{1}{B_0 \log_2 \left(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right)} \right)^2 - \frac{1}{2} \left(\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q p_{i\tilde{\xi}q} \right)^2 - \frac{1}{2} \left[B_0 \log_2 \left(1 + \frac{p_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right) \right]^{-2} \\ &\leq H_q(p_{i\tilde{\xi}q}) - \frac{1}{2} \left(\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q \tilde{p}_{i\tilde{\xi}q} \right)^2 - \frac{1}{2} \left[B_0 \log_2 \left(1 + \frac{\tilde{p}_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right) \right]^{-2} \\ &\quad - \ln 2 \cdot B \cdot \frac{\tilde{p}_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2 + \tilde{p}_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2} - \frac{(\mathbb{I}_{\{i \in \Gamma_{q\bar{m}}\}} I_q)^2 \cdot \tilde{p}_{i\tilde{\xi}q} \cdot (p_{i\tilde{\xi}q} - \tilde{p}_{i\tilde{\xi}q})}{\left(B \log_2 \left(1 + \frac{\tilde{p}_{i\tilde{\xi}q} |h_{i\tilde{\xi}q}^{\text{IoT}}|^2}{\sigma_{\tilde{\xi}q}^2} \right) \right)^3} \cdot (p_{i\tilde{\xi}q} - \tilde{p}_{i\tilde{\xi}q}) \\ &\triangleq \tilde{E}_{iq}^{c1}(p_{i\tilde{\xi}q}; \tilde{p}_{i\tilde{\xi}q}) \end{aligned} \quad (32)$$

Next, the Lagrangian function of P_3^{Sub} is expressed as follows.

$$\begin{aligned}
& L_3(\{\mathcal{P}^b\}, \{\mu_i\}_{i \in \mathcal{N}}, \{\delta_q\}_{q \in \mathcal{N}^Q}) \\
&= \sum_{i=1}^N \sum_{q=1}^Q \tilde{E}_{iq}^{c1}(p_{i\tilde{\xi}q}; \tilde{p}_{i\tilde{\xi}q}) + \sum_{i=1}^N \mu_i \left(\sum_{n=1}^N p_{in} - c_6 \right) \\
&+ \sum_{q=1}^Q \left[\delta_q^1 \cdot (-R_{q \rightarrow \xi_1} + R_{\min}) + \sum_{k \in \{1, \dots, K_{q\tilde{m}}-1\}} \delta_q^k \cdot (-R_{\xi_k \rightarrow \xi_{k+1}} + R_{\min}) \right. \\
&\left. + \delta_q^{K_{q\tilde{m}}} \cdot (-R_{\xi_{K_{q\tilde{m}}} \rightarrow \tilde{m}} + R_{\min}) \right] \\
&\triangleq \sum_{m=1}^M L_3^m(\{p_{m\tilde{\xi}m}\}_{m \in \mathcal{N}^M}) + \sum_{q=1}^Q L_3^q(\{p_{q\tilde{\xi}q}\}_{q \in \mathcal{N}^Q}) + c_7 \quad (37)
\end{aligned}$$

where $L_3^m(\{p_{m\tilde{\xi}m}\}_{m \in \mathcal{N}^M})$ is defined as the function of decision variables $\{p_{m\tilde{\xi}m}\}_{m \in \mathcal{N}^M}$ for gateway buoys \mathcal{N}^M , $L_3^q(\{p_{q\tilde{\xi}q}\}_{q \in \mathcal{N}^Q})$ is defined as the function of decision variables $\{p_{q\tilde{\xi}q}\}_{q \in \mathcal{N}^Q}$ for detection buoys \mathcal{N}^Q , and c_7 is the constant that is independent of \mathcal{P}^b . We use the symbol $\tilde{\xi}^q$ to refer to the next hop of buoy q , and the symbol $\tilde{\xi}^m$ to refer to the next hop of buoy m .

Then, take the partial derivative of the Lagrangian function $L_3^m(\{p_{m\tilde{\xi}m}\}_{m \in \mathcal{N}^M})$ with respect to $p_{m\tilde{\xi}m}$ for any $m \in \mathcal{N}^M$, take the partial derivative of the Lagrangian function $L_3^q(\{p_{q\tilde{\xi}q}\}_{q \in \mathcal{N}^Q})$ with respect to $p_{q\tilde{\xi}q}$ for any $q \in \mathcal{N}^Q$. We represent all power variables in terms of $\{p_{ij}\}_{i \in \mathcal{N}, j \in \mathcal{N}}$, and a unified representation for updates is formulated as follows:

$$p_{ij}^{(\kappa+1)} = \begin{cases} p_{m\tilde{\xi}m}^{(\kappa)} - \zeta_{m\tilde{\xi}m}^{(\kappa)} \cdot \frac{\partial L_3^m(\{p_{m\tilde{\xi}m}\}_{m \in \mathcal{N}^M})}{\partial p_{m\tilde{\xi}m}^{(\kappa)}}, & \text{if } i \in \mathcal{N}^M \wedge i \in \tilde{\Gamma} \\ p_{q\tilde{\xi}q}^{(\kappa)} - \zeta_{q\tilde{\xi}q}^{(\kappa)} \cdot \frac{\partial L_3^q(\{p_{q\tilde{\xi}q}\}_{q \in \mathcal{N}^Q})}{\partial p_{q\tilde{\xi}q}^{(\kappa)}}, & \text{if } i \in \mathcal{N}^Q \wedge i \in \tilde{\Gamma} \\ 0, & \text{if } i \notin \tilde{\Gamma} \end{cases} \quad (38)$$

where the calculate of variable $\{p_{ij}\}_{i \in \mathcal{N}, j \in \mathcal{N}}$ is divided into three cases. If the node i serves as a gateway buoy and has an associated forwarding task, (i.e., $i \in \mathcal{N}^M \wedge i \in \tilde{\Gamma}$) holds, then we can use $L_3^m(\{p_{m\tilde{\xi}m}\}_{m \in \mathcal{N}^M})$ to update $p_{m\tilde{\xi}m}$. In addition, if the node i serves as a detection buoy and has an associated forwarding task, (i.e., $i \in \mathcal{N}^Q \wedge i \in \tilde{\Gamma}$) holds, then we can use $L_3^q(\{p_{q\tilde{\xi}q}\}_{q \in \mathcal{N}^Q})$ to update $p_{q\tilde{\xi}q}$. Otherwise, $p_{ij} = 0$ when the node i does not have an associated forwarding task.

Let κ represent the index of dual ascent algorithm for outer loop, and let $\tilde{\kappa}$ represent the index of SCA algorithm for inner loop. Then, the update of $\tilde{p}_{ij}^{(\tilde{\kappa}+1)}$ for $\tilde{\kappa}+1$ -th iteration is given as follows.

$$\tilde{p}_{ij}^{(\kappa, \tilde{\kappa}+1)} = \arg \min_{\mathcal{P}^b} \sum_{i=1}^N \sum_{q=1}^Q \tilde{E}_{iq}^{c1}(p_{i\tilde{\xi}q}; \tilde{p}_{i\tilde{\xi}q}) \quad (39a)$$

$$s.t. \begin{cases} C'_3, C''_6 : \sum_{n=1}^N p_{in} \leq c_6, \quad \forall i \in \mathcal{N} \\ C'_4 \text{ (refer to (16))} \end{cases} \quad (39b)$$

Algorithm 3 Solving for Computing Resources $\{\mathcal{P}^b\}$ for Gateway Guoys

Input: Optimization Problem P_3^{Sub} .

Output: $\{\mathcal{P}^b\} = \{p_{ij}\}_{i \in \mathcal{N}, j \in \mathcal{N}}$

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1: Initialize  $\{\{p_{ij}^{(0)}\}_{j \in \mathcal{N}}, \mu_i^{(0)}, \gamma_i^{(0)}\}_{i \in \mathcal{N}}, \{\delta_q^{k(0)}\}_{q \in \mathcal{N}^Q, \xi_k \in \Gamma_{q\tilde{m}}}$ .
2: while (1) do
3:   // Part(a): Dual-Ascent-based Variable Update
4:   for  $m = \{1, 2, \dots, M\}$  do
5:     Calculate the Lagrangian function based on (37).
6:     Calculate the partial derivatives and update decision variable  $p_{ij}^{(\kappa+1)}$  for  $\forall i, j$ , based on (38).
7:     Update Lagrange multipliers  $\mu_i^{(\kappa+1)}, \gamma_i^{(\kappa+1)}, \delta_q^{k(\kappa+1)}$  for  $\forall i, q$  based on (40).
8:   end for
9:   if  $\max_i \{\mu_i^{(\kappa+1)} - \mu_i^{(\kappa)}\} < \varepsilon_\mu \wedge \max_{q,k} \{\delta_q^{k(\kappa+1)} - \delta_q^{k(\kappa)}\} < \varepsilon_\delta$  then
10:    BreakWhile
11:   end if
12:   Update  $\kappa = \kappa + 1$ .
13:   // Part(b): SCA-based Power Iteration
14:   Initialize  $\{\tilde{p}_{ij}^{(\kappa, 0)}\}_{i \in \mathcal{N}, j \in \mathcal{N}}$  variables.
15:   while (1) do
16:     Update  $\tilde{p}_{ij}^{(\kappa, \tilde{\kappa}+1)}$  by solving the convex optimization problem (39) with substitute variable  $\tilde{h}_q(p_{ij}^{(\kappa)}; \tilde{p}_{ij}^{(\kappa, \tilde{\kappa})})$  mentioned in (32).
17:     if  $\{\tilde{p}_{i\tilde{\xi}q}^{(\kappa, \tilde{\kappa}+1)}\}_{i \in \mathcal{N}, j \in \mathcal{N}}$  converge then
18:       Break While
19:     end if
20:     Update  $\tilde{\kappa} = \tilde{\kappa} + 1$ .
21:   end while
22: end while
23: return  $\{\mathcal{P}^b\}$ 

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where the above formula can be used to update the auxiliary variables in SCA.

Then, the Lagrange multipliers are calculated as follows.

$$\mu_i^{(\kappa+1)} = \mu_i^{(\kappa)} + \zeta_{m\mu}^{(\kappa)} \cdot \left(\sum_{n=1}^N p_{in}^{(\kappa)} - c_6 \right) \quad (40)$$

$$\delta_q^{k(\kappa+1)} = \delta_q^{k(\kappa)} + \zeta_{qk\delta}^{(\kappa)} \cdot (-R_{\xi_k \rightarrow \xi_{k+1}} + R_{\min})$$

where $\zeta_{m\mu}^{(\kappa)}, \zeta_{qk\delta}^{(\kappa)}$ are the step sizes for Lagrange multipliers. Finally, the solution algorithm for decision variables $\{\mathcal{P}^b\}$ is given in Alg. 3.

3) *Solving for Decision Variables $\{\mathcal{C}\}$* : In this subsection, the derivation of the decision variable \mathcal{C} is proposed, when other decision variables are fixed, and the optimization problem P_4^{Sub} is given as follows.

$$P_4^{\text{Sub}} : \min_{\mathcal{C}} \sum_{m=1}^M E_m^s(\mathcal{C}) \quad (41a)$$

$$s.t. \quad C_2 \quad (41b)$$

where the objective function is a quadratic functions with respect to $\mathcal{C} = \{c_{qm}\}_{q \in \mathcal{N}^Q, m \in \mathcal{N}^M}$, which is also a convex function.

Algorithm 4 Solving for Computing Resources $\{\mathcal{C}\}$ for Gateway Guoys

Input: Optimization Problem P_4^{Sub} .

Output: $\{\mathcal{C}\} = \{c_{qm}\}_{q \in \mathcal{N}^Q, m \in \mathcal{N}^M}$

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1: Initialize  $\{\{c_{qm}^{(0)}\}_{q \in \mathcal{N}^Q, m \in \mathcal{N}^M}, \lambda_m^{(0)}, \theta_m^{(0)}\}_{m \in \mathcal{N}^M}$ 
2: while (1) do
3:   for  $m = \{1, 2, \dots, M\}$  do
4:     Calculate the Lagrangian function based on (42).
5:     Calculate the partial derivatives based on (43).
6:     Update decision variable  $c_{qm}^{(\kappa+1)}$  for  $\forall q, m$ , based on (44).
7:     Update Lagrange multipliers  $\lambda_m^{(\kappa+1)}$  and  $\theta_m^{(\kappa+1)}$  for  $\forall m$  based on (45).
8:   end for
9:   if  $\max_m \{\lambda_m^{(\kappa+1)} - \lambda_m^{(\kappa)}\} < \varepsilon_\lambda \wedge \max_m \{\theta_m^{(\kappa+1)} - \theta_m^{(\kappa)}\} < \varepsilon_\theta$  then
10:    BreakWhile
11:   end if
12: end while
13: return  $\{\mathcal{C}\}$ 

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Then, C_2 constraint is linear functions with respect to \mathcal{C} , which satisfies the definition of convex function. Thus, the Lagrangian function of P_4^{Sub} is expressed as follows.

$$\begin{aligned}
& L_4(\{\mathcal{C}\}, \{\lambda_m\}_{m \in \mathcal{N}^M}) \\
&= \sum_{m=1}^M E_m^s(\mathcal{C}) + \sum_{m=1}^M \lambda_m \left(\sum_{q=1}^Q c_{qm} - C_0 \right) \\
&\triangleq \sum_{m=1}^M L_4^m(\{c_{qm}\}_{m \in \mathcal{N}^M}) + c_9
\end{aligned} \quad (42)$$

where $L_4^m(\{c_{qm}\}_{m \in \mathcal{N}^M})$ is defined as the function of decision variables $\{\mathcal{C}\}$, and c_9 is the constant that is independent of $\{\mathcal{C}\}$.

Then, take the partial derivative of the Lagrangian function $L_4^m(\{c_{qm}\}_{m \in \mathcal{N}^M})$ with respect to c_{qm} for any $m \in \mathcal{N}^M$ as

$$\begin{aligned}
& \frac{\partial L_4^m(\{c_{qm}\}_{m \in \mathcal{N}^M})}{\partial c_{qm}} \\
&= \left(2\theta_m \epsilon_m x_{qm} \phi_q(1 - y_{qm}) \right) \cdot c_{qm} + \lambda_m
\end{aligned} \quad (43)$$

let $\frac{\partial L_4^m(\{c_{qm}\}_{m \in \mathcal{N}^M})}{\partial c_{qm}} = 0$, we have

$$c_{qm}^{(\kappa+1)} = \max \left\{ \frac{\lambda_m}{2\theta_m \epsilon_m x_{qm} \phi_q(y_{qm} - 1)}, 0 \right\} \quad (44)$$

Then, the Lagrange multipliers are calculated as follows.

$$\lambda_m^{(\kappa+1)} = \lambda_m^{(\kappa)} + \zeta_{m\lambda}^{(\kappa)} \cdot \left(\sum_{q=1}^Q c_{qm} - C_0 \right) \quad (45)$$

Finally, the solution algorithm for decision variables $\{\mathcal{C}\}$ is given in Alg. 4.

4) *Solving for Decision Variables $\{\mathcal{Y}\}$* : In this subsection, the solution of the decision variable \mathcal{Y} is proposed, when other

Algorithm 5 Joint Iteration Algorithm (Top Algorithm)

Input: Optimization Problems $P_2^{\text{Sub}}, P_3^{\text{Sub}}$ and P_4^{Sub} .

Output: $\{\varrho, \mathcal{G}, \mathcal{P}^b, \mathcal{C}, \mathcal{Y}\}$.

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1: Initialize  $\varrho^{(0)}, \mathcal{G}^{(0)}, \mathcal{P}^{b(0)}, \mathcal{C}^{(0)}, \mathcal{Y}^{(0)}, \tau = 0$ .
2: while (1) do
3:   Fix the decision variables  $\{\mathcal{P}^{b(\tau)}, \mathcal{C}^{(\tau)}, \mathcal{Y}^{(\tau)}\}$ , and solve  $\{\varrho^{(\tau+1)}, \mathcal{G}^{(\tau+1)}\}$  by Alg. 2.
4:   Fix the decision variables  $\{\varrho^{(\tau+1)}, \mathcal{G}^{(\tau+1)}, \mathcal{C}^{(\tau)}, \mathcal{Y}^{(\tau)}\}$ , and solve  $\{\mathcal{P}^{b(\tau+1)}\}$  by Alg. 3.
5:   Fix the decision variables  $\{\varrho^{(\tau+1)}, \mathcal{G}^{(\tau+1)}, \mathcal{P}^{b(\tau+1)}, \mathcal{Y}^{(\tau)}\}$ , and solve  $\{\mathcal{C}^{(\tau+1)}\}$  by Alg. 4.
6:   Fix the decision variables  $\{\varrho^{(\tau+1)}, \mathcal{G}^{(\tau+1)}, \mathcal{P}^{b(\tau+1)}, \mathcal{C}^{(\tau+1)}\}$ , and solve  $\{\mathcal{Y}^{(\tau+1)}\}$  by barrier function interior point method.
7:   if  $\{\mathcal{P}^{b(\tau+1)}, \mathcal{C}^{(\tau+1)}, \mathcal{Y}^{(\tau+1)}, \varrho^{(\tau+1)}, \mathcal{G}^{(\tau+1)}\}$  converge then
8:     BreakWhile
9:   end if
10:  Update  $\tau = \tau + 1$ .
11: end while
12: return  $\{\varrho, \mathcal{G}\}$ 

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decision variables are fixed, and the optimization problem P_5^{Sub} is given as follows.

$$\begin{aligned}
P_5^{\text{Sub}} : \min_{\mathcal{Y}} & \sum_{m=1}^M E_m^s + \sum_{m=1}^M \sum_{\omega=1}^W E_{m\omega}^{c2} \\
s.t. & C_5
\end{aligned} \quad (46a)$$

where both the objective function and the constraints are linear function with respect to $\mathcal{Y} = \{y_{qm}\}_{q \in \mathcal{N}^Q, m \in \mathcal{N}^M}$. Thus, we can solve the optimization problem P_5^{Sub} using a similar method mentioned above (or using the barrier function interior point method, CVX, and so on).

5) *Outer Iteration Algorithm and Complexity Analysis*: It is worth noting that Alg. 2, Alg. 3, and Alg. 4 serve as inner-layer iterative algorithms, each solving one variable while keeping the others fixed. Based on Fig. 3, a joint iterative algorithm based on BCD method for the outer layer is proposed. This algorithm iteratively polls and resolves the different variables in sequence, as detailed in Alg. 5.

The computational complexity is analyzed in this paper. In particular, the BCD method, as mentioned in Alg. 5, has a computational complexity of $\mathcal{O}(\tau \cdot C(M, Q, W))$, where τ represents the iteration count for BCD, and $C(M, Q, W)$ denotes the maximum complexity for solving each optimization problem (Alg. 2, Alg. 3, and Alg. 4) when other decision variables are fixed, which implies that $C(M, Q, W) = \max\{C_2(M, W), C_3(M, Q), C_4(M, Q), LP\}$.

Next, we analyze each individual sub-optimization problem. It is worth mentioning that each sub-optimization problem mentioned in this article can be decomposed in the form of $L^m(\cdot)$ or $L^q(\cdot)$, as mentioned in Fig. 7, which implies that distributed parallel computing is feasible, and the algorithm complexity can be effectively reduced. Specifically, for the sub-optimization problem P_2 , the computational complexity $C_2(M, W) = k'_1 \cdot W$, where M is the number of nodes

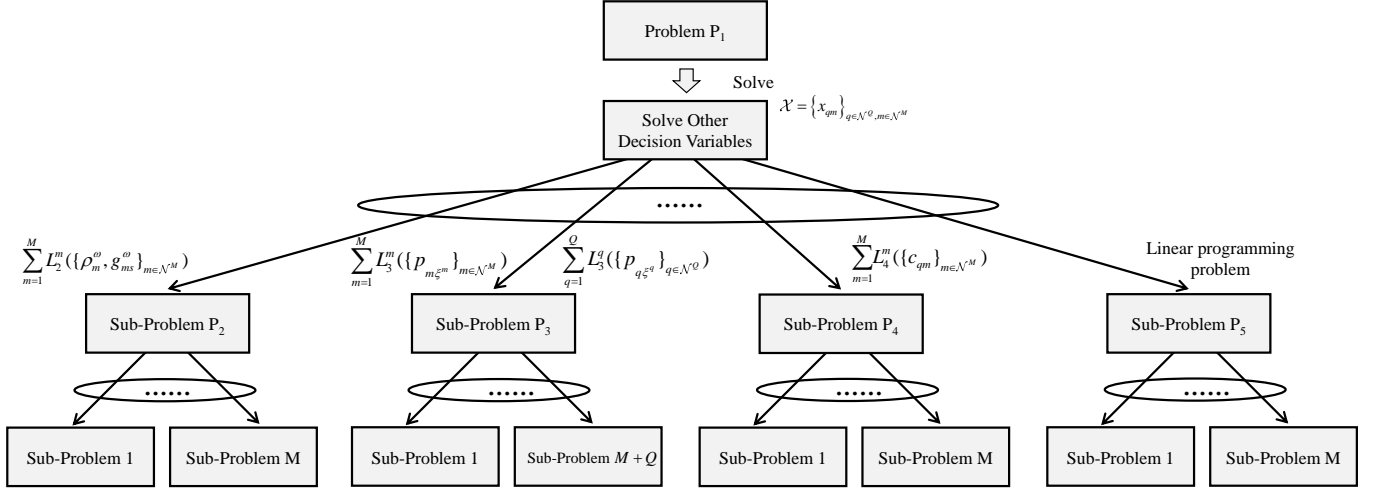


Fig. 7: Detailed decomposition architecture for the optimization problem P_1 .

for parallel computing, and k'_1 is the number of iteration index. For the sub-optimization problem P_3 , the computational complexity $C_3(M, W) = \max\{C_3(M), C_3(Q)\} = \max\{k'_2 \cdot M, k'_3 \cdot Q\}$. In addition, $C_4(M, Q) = k'_4 \cdot Q$, and

LP is the computational complexity of linear programming problem. Thus, based on the decomposition architecture for the optimization problem P_1 , the computational complexity can be effectively reduced