Exercise 3.55 [Boyd & Vandenberghe, 2004]

Ans: (a) Differentiate $f(x) = \int_{-\infty}^{x} e^{-h(t)} dt$ and get

$$f'(x) = e^{-h(x)}, f''(x) = -h'(x)e^{-h(x)}$$

Therefore for $x \in \operatorname{dom} h$, $h'(x) \ge 0$,

$$f''(x)f(x) = -h'(x)e^{-h(x)} \int_{-\infty}^{x} e^{-h(t)} dt \le 0 \le (f'(x))^{2}$$

since exponentials and their integrals are always positive.

(b) Apply the inequality $h(t) \ge h(x) + h'(x)(t-x)$ to the integral,

$$\begin{split} \int_{-\infty}^x e^{-h(t)} dx &\leq \int_{-\infty}^x e^{-h(x) - h'(x)(t-x)} dt \\ &= e^{-h(x)} \int_{-\infty}^x e^{-h'(x)(t-x)} dt \\ &= e^{-h(x)} \left(\frac{1}{-h'(x)}\right) \int_{-\infty}^0 e^s ds \text{ (change of variable with } s = -h'(x)(t-x)) \\ &= \frac{e^{-h(x)}}{-h'(x)} \cdot 1 \end{split}$$

Now that h'(x) < 0, we get

$$f''(x)f(x) \le -h'(x)e^{-h(x)}\left(\frac{e^{-h(x)}}{-h'(x)}\right) = e^{-2h(x)} = (f'(x))^2$$

Additional Exercise 3.5 [Boyd & Vandenberghe, 2017]

Ans: First observe the objective can be written as an over all maximization among $i \times j$ quotients,

$$f_0(x) = \frac{\max_{i=1,\dots,m} a_i^T x + b_i}{\min_{j=1,\dots,p} c_j^T x + d_j} = \max_{\substack{i=1,\dots,m\\j=1,\dots,p}} \frac{a_i^T x + b_i}{c_j^T x + d_j}$$

since for fixed i_0 and j_0 ,

$$\frac{\max_{i=1,\cdots,m} a_i^T x + b_i}{\min_{j=1,\cdots,p} c_j^T x + d_j} \ge \frac{a_{i_0}^T x + b_{i_0}}{\min_{j=1,\cdots,p} c_j^T x + d_j} \ge \frac{a_{i_0}^T x + b_{i_0}}{c_{j_0}^T x + d_{j_0}}$$

(the reverse inequality follows from the combination of maximizing index of numerator and minimizing index of denominator is within the range of the $i \times j$ index pairs.) Let

$$y_j = \frac{x}{c_j^T x + d_j}, z_j = \frac{1}{c_j^T x + d_j}$$

then the following 2 problems are equivalent (identical trick as in textbook §4.3.2),

$$\begin{cases} \text{minimize} & \frac{a_i^T x + b_i}{c_j^T x + d_j} \\ \text{subject to} & Fx \prec q \end{cases} \qquad \begin{cases} \text{minimize} & a_i^T y_j + b_i z_j \\ \text{subject to} & Fy_j - g z_j \leq 0 \end{cases}$$

Therefore, combine them all together, we can get the equivalent LP for the original problem,

$$\begin{cases} \text{minimize} & t \\ \text{subject to} & a_i^T y_j - b_i z_j \leq t, \text{ for } i = 1, \dots, m, j = 1, \dots, p \\ & F y_j - g z_j \leq 0, \text{ for } j = 1, \dots, p \\ \text{with variables} & y_j, z_j, t \end{cases}$$

Exercise 4.21 [Boyd & Vandenberghe, 2004]

Ans: Here only does part (b). Consider Lagrange multiplier,

$$c = \lambda 2A(x^* - x_c)$$
$$x^* = \frac{1}{2\lambda}A^{-1}c + x_c$$

(Note that $A \in \mathbf{S}_{++}^n$ must be invertible.) Since x^* must be on the boundary of the ellipsoid $\{x \mid (x-x_c)^T A (x-x_c) \leq 1\}$,

$$1 = (x^* - x_c)^T A (x^* - x_c)$$
$$= \frac{1}{(2\lambda)^2} c^T A^{-1} c$$

Solve for λ ,

$$\lambda = \frac{1}{2} \sqrt{c^T A^{-1} c}$$

and

$$x^* = \frac{1}{\sqrt{c^T A^{-1} c}} A^{-1} c + x_c$$

(Note how the final formula assumes $c \neq 0$.)

Exercise 4.25 [Boyd & Vandenberghe, 2004]

Ans: Without loss of generality, first look into the constraint induced by the first ellipsoid; we are looking for $a \in \mathbb{R}^n, b \in R$ such that

$$a^{T}x + b > 0$$
 for $x = P_{1}u + q_{1}$ with $||u||_{2} \le 1$

or,

$$0 < a^{T}(P_{1}u + q_{1}) + b = u^{T}P_{1}a + q_{1}^{T}a + b$$

Rewrite the inequality we get

$$-b - q_1^T a < u^T P_1 a \le ||P_1 a||_2$$

where the last inequality is from Cauchy's. Combine such inequalities for all ellipsoids, we get the SOCP feasibility problem,

$$\begin{cases} \text{find} & a, b \\ \text{subject to} & -b - q_i^T a \le ||P_i a||_2, \text{ for } i = 1, \dots, K \\ & b + q_i^T a \le ||P_i a||_2, \text{ for } i = K + 1, \dots, L \end{cases}$$

Additional Exercise 7.9 [Boyd & Vandenberghe, 2017]

Ans: (a) It is to show the objective function

$$g(x) = \max_{k=1,\dots,N} ||f_k(x) - y_k||_2$$

is a quasiconvex function. Take $t \in \mathbb{R}$, the t-sublevelset

$$\{x \mid g(x) \leq t\} = \{x \mid ||f_k(x) - y_k||_2 \leq t, k = 1, \dots N\}$$

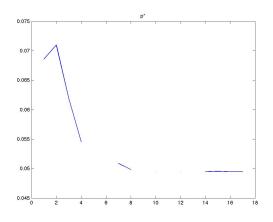
$$= \bigcap_{i=1}^N \{x \mid ||f_k(x) - y_k||_2 \leq t\}$$

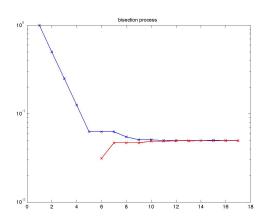
$$= \bigcap_{i=1}^N \left\{x \mid \left\| \frac{1}{c_k^T x + d_k} (A_k x + b_k) - y_k \right\|_2 \leq t\right\}$$

$$= \bigcap_{i=1}^N \left\{x \mid \left\| (A_k x + b_k) - y_k (c_k^T x + d_k) \right\|_2 \leq t (c_k^T x + d_k)\right\}$$

is the intersection of N convex cones and therefore convex. (Note the assumption $c_k^T x + d > 0$ is used in the last equality.)

(b) Here's the graph of p^* and the upper/lower bound during the bisection process. Note the p^* curve is missing some parts because some subproblems were infeasible; there was no solution and CVX outputted NaN.





%% Problem A7.9
P = zeros(3,4,4);
P(:,:,1) = [1 0 0 0; 0 1 0 0; 0 0 1 0];
P(:,:,2) = [1 0 0 0; 0 0 1 0; 0 -1 0 10];
P(:,:,3) = [1 1 1 -10; -1 1 1 0; -1 -1 1 10];
P(:,:,4) = [0 1 1 0; 0 -1 1 0; -1 0 0 10];

```
y = zeros(2,4);
y(:,1) = [.98; .93];
y(:,2) = [1.01; 1.01];
y(:,3) = [.95; 1.05];
y(:,4) = [2.04; 0];
f = Q(x,k) (P(1:2,:,k)*[x;1])/(P(3,:,k)*[x;1]);
g = Q(x) \max([norm(f(x,1)-y(:,1)), norm(f(x,2)-y(:,2)), norm(f(x,3)-y(:,3)), norm(f(x,3)-y(:,3)))
phi_k = Q(t,x,k) \text{ norm}(P(1:2,:,k)*[x;1]-y(:,k)*P(3,:,k)*[x;1]) - t*P(3,:,k)*[x;1];
phi = 0(t,x) \max([phi_k(t,x,1), phi_k(t,x,2), phi_k(t,x,3), phi_k(t,x,4)]);
TOL = 1e-4;
1 = 0; % objective function is always nonnegative
u = 1; % random upper bound
num_iter = 0;
while u-1 > TOL
    num_iter = num_iter + 1;
    t = 0.5 * (u + 1);
    cvx_begin
        variable x(3);
        minimize 1;
        subject to
            phi(t,x) \ll 0;
    cvx_end
    p_star(num_iter) = g(x);
    u_data(num_iter) = u;
    l_data(num_iter) = 1;
    if cvx_optval == 1
        u = t;
    else
        1 = t;
    end
end
h1 = plot(p_star,'LineWidth', 1.1);
title('p*');
saveas(h1,'A79_pstar','jpg');
figure;
semilogy(u_data,'b-x','LineWidth',1.05);
h2 = semilogy(l_data, 'r-x', 'LineWidth', 1.05);
title('bisection process');
```

saveas(h2, 'A79_bisection', 'jpg');

Additional Exercise 14.8 [Boyd & Vandenberghe, 2017]

Ans:

(a) First write all the constraints in discretized form. The glide slope constraint can be viewed as

$$[0,0,1]p_k \ge \alpha \left\| \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} p_k \right\|_2$$
 for $k = 0, 1, \dots, K$

The position and velocity requires more work.

$$\begin{aligned} v_k &= v_{k-1} + \left(\frac{h}{m}\right) f_{k-1} - hge_3 \\ &= v_{k-2} + \left(\frac{h}{m}\right) f_{k-2} - hge_3 + \left(\frac{h}{m}\right) f_{k-1} - hge_3 \\ &= \cdots \\ &= v_1 + \frac{h}{m} \sum_{i=1}^{k-1} f_i - (k-1)hge_3 \\ p_{k+1} &= p_k + \frac{h}{2} (v_k + v_{k+1}) \\ &= p_{k-1} + \frac{h}{2} (v_{k-1} + v_k) + \frac{h}{2} (v_k + v_{k+1}) \\ &= \cdots \\ &= p_1 + \frac{h}{2} v_1 + h(v_2 + \cdots + v_k) + \frac{h}{2} v_{k+1} \\ &= p_1 + \frac{h}{2} v_1 + (k-1)hv_1 + h \sum_{j=2}^{k} \left(\frac{h}{m} \sum_{i=1}^{j-1} f_i - (j-1)hge_3\right) + \frac{h}{2} v_1 + \frac{h}{2} \left(\frac{h}{m} \sum_{i=1}^{k} f_i - khge_3\right) \\ &= p_1 + khv_1 - \frac{k(k+1)}{2} h^2 ge_3 + \frac{h^2}{m} \sum_{i=2}^{k} \sum_{j=1}^{j-1} f_i + \frac{h^2}{2m} \sum_{i=1}^{k} f_i \end{aligned}$$

Therefore the constraints for touch down is

$$0 = v_{K+1}$$

$$= v_1 + \frac{h}{m} \sum_{i=1}^{K} f_i - Khge_3$$

$$= v_0 + \frac{h}{m} f \mathbf{1} - Khge_3$$

$$0 = p_{K+1}$$

$$= p_1 + Khv_1 - \frac{K(K+1)}{2} h^2 ge_3 + \frac{h^2}{m} \sum_{j=2}^{K} \sum_{i=1}^{j-1} f_i + \frac{h^2}{2m} \sum_{i=1}^{K} f_i$$

$$= p_1 + Khv_1 - \frac{K(K+1)}{2} h^2 ge_3 + \frac{h^2}{m} f \begin{bmatrix} K-1 \\ \vdots \\ 2 \\ 1 \\ 0 \end{bmatrix} + \frac{h^2}{2m} f \mathbf{1}$$

$$= p_1 + Khv_1 - \frac{K(K+1)}{2} h^2 ge_3 + \frac{h^2}{m} f \begin{bmatrix} K-0.5 \\ \vdots \\ 2.5 \\ 1.5 \\ 0.5 \end{bmatrix}$$

where $\mathbf{1} = \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix}$ and $f = [f_1, f_2, \dots, f_K] \in \mathbb{R}^{3 \times K}$. Aside of the constraints, the total fuel

use can also be expressed in discretized form since

$$\int_{0}^{T_{t}d} \gamma \|f(t)\|_{2} dt \approx \sum_{i=1}^{K} \gamma \|f_{k}\|_{2} h$$

The fuel descent problem can be now formalized
$$\begin{cases} & \text{minimize} \quad \sum_{i=1}^{K} \gamma h \| f_k \|_2 \\ & \text{subject to} \quad v_1 + \frac{h}{m} \sum_{i=1}^{K} f_i - K h g e_3 = 0 \\ & p_1 + K h v_1 - \frac{K(K+1)}{2} h^2 g e_3 + \frac{h^2}{m} \sum_{j=2}^{K} \sum_{i=1}^{j-1} f_i + \frac{h^2}{2m} \sum_{i=1}^{K} f_i = 0 \\ & [0,0,1] p_k \geq \alpha \left\| \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} p_k \right\|_2 \text{ for } k = 0, 1, \cdots, K \end{cases}$$

$$p_{k+1} = p_1 + k h v_1 - \frac{k(k+1)}{2} h^2 g e_3 + \frac{h^2}{m} \sum_{j=2}^{k} \sum_{i=1}^{j-1} f_i + \frac{h^2}{2m} \sum_{i=1}^{k} f_i \text{ for } k = 0, 1, \cdots, K \end{cases}$$

(b) The problem for solving minimum time descent is

$$\begin{cases} \text{minimize} & K \\ \text{subject to} & v_1 + \frac{h}{m} \sum_{i=1}^{K} f_i - Khge_3 = 0 \\ & p_1 + Khv_1 - \frac{K(K+1)}{2} h^2 ge_3 + \frac{h^2}{m} \sum_{j=2}^{K} \sum_{i=1}^{j-1} f_i + \frac{h^2}{2m} \sum_{i=1}^{K} f_i = 0 \\ & \left[[0,0,1] p_k \ge \alpha \left\| \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} p_k \right\|_2 \text{ for } k = 0, 1, \cdots, K \\ & p_{k+1} = p_1 + khv_1 - \frac{k(k+1)}{2} h^2 ge_3 + \frac{h^2}{m} \sum_{j=2}^{k} \sum_{i=1}^{j-1} f_i + \frac{h^2}{2m} \sum_{i=1}^{k} f_i \text{ for } k = 0, 1, \cdots, K \end{cases}$$

Without any thoughts, this problem can be solved by solving at most \aleph_0 (jk, but also not) feasibility problems with $K = 1, 2, \cdots$.

TOO MUCH WORK

NOT ENOUGH TIME