操作系统

Operating Systems

L23 段页结合的实际内存管理

Segmentation & Paging

lizhijun_os@hit.edu.cn 综合楼411室

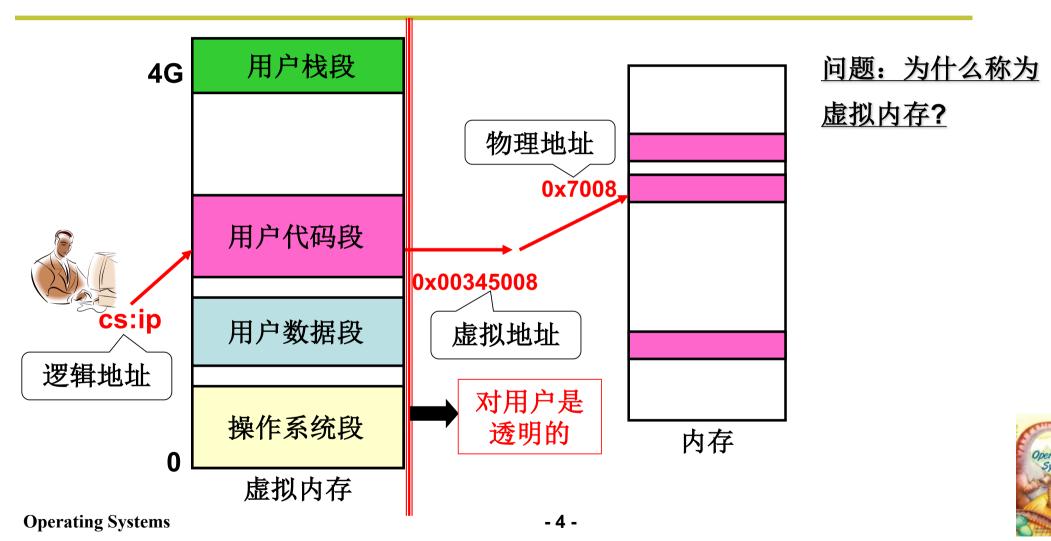
授课教师: 李治军

段、页结合:程序员希望用段,物理内存希望用页,所以....

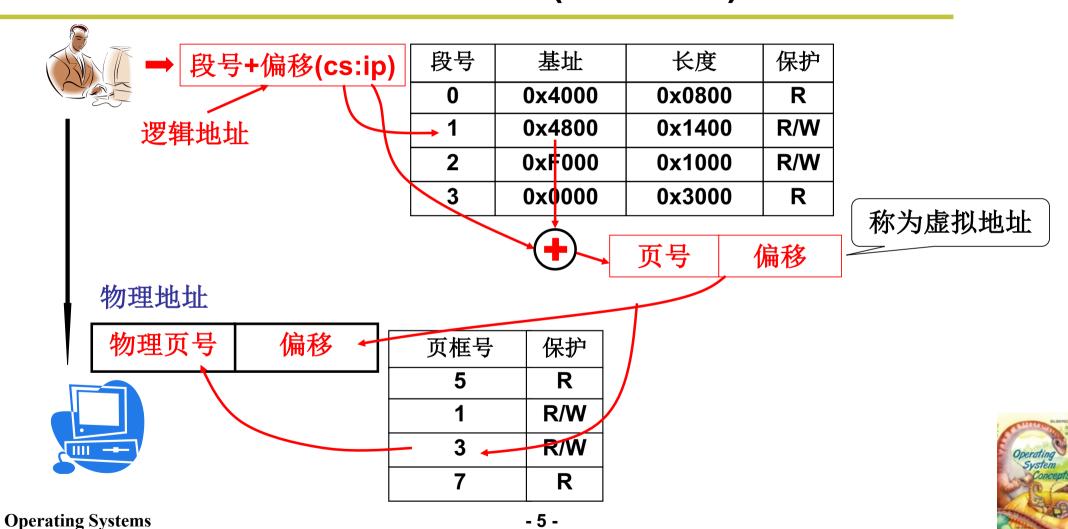




段、页同时存在:段面向用户/页面向硬件



段、页同时存在是的重定位(地址翻译)

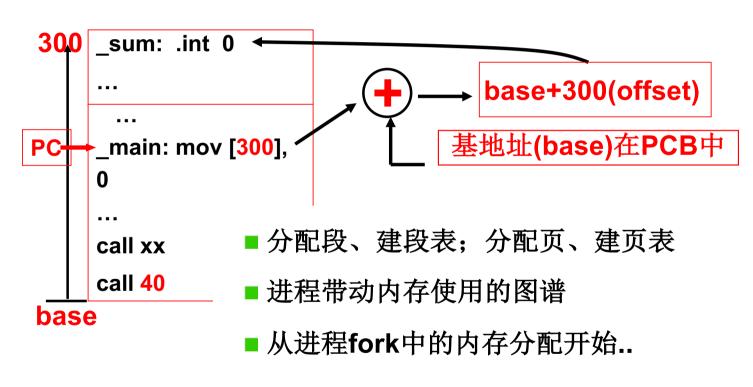


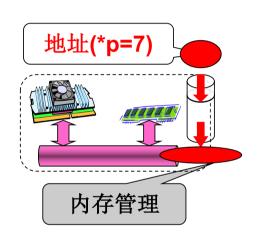
一个实际的段、页式内存管理



这个故事从哪里开始?

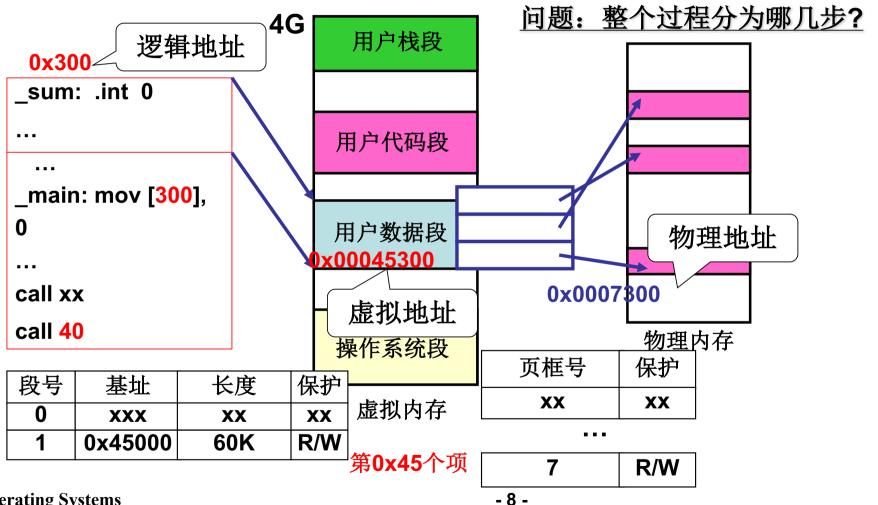
■ 内存管理核心就是内存分配,所以从程序放入内存、使用内存开始...







段、页式内存下程序如何载入内存?





Operating Systems

故事从fork()开始 分配虚存、建段表

■ fork()→sys_fork→copy_process的路都已经走过了

```
在linux/kernel/fork.c中

int copy_process(int nr, long ebp,...)

{ ...

copy_mem(nr, p); ...

的确是进程带动内存!
```

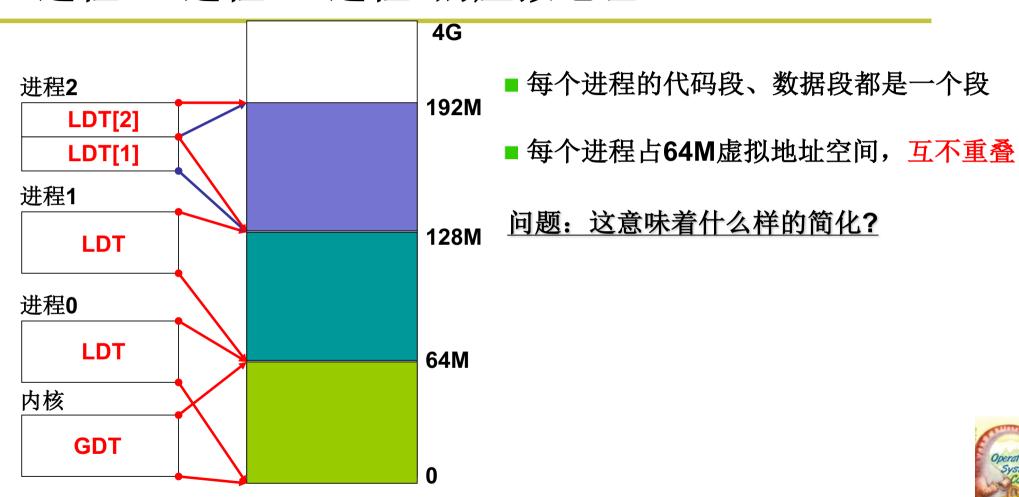
■ 现在开始分析当时那个神秘的copy_mem了



进程0、进程1、进程2的虚拟地址

虚拟内存

Operating Systems





- 10 -

接下来应该是什么了?分配内存、建页表

```
int copy_mem(int nr, task_struct *p)
{    unsigned long old_data_base;
    old_data_base=get_base(current->ldt[2]);
    copy_page_tables(old_data_base,new_data_base,data_limit);

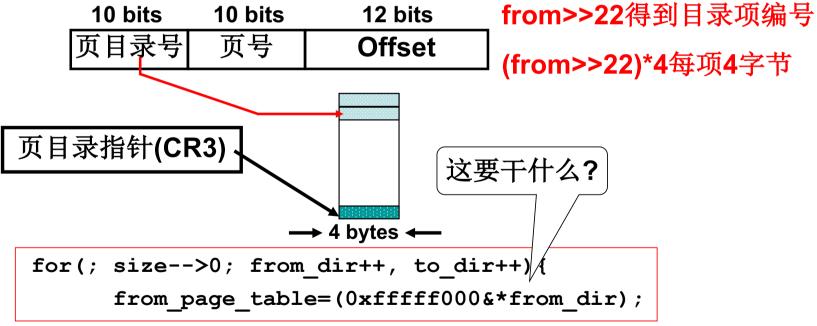
int copy_page_tables(unsigned long from,unsigned long to, long size)
{    from_dir = (unsigned long *)((from>>20)&0xffc);
    to_dir = (unsigned long *)((to>>20)&0xffc);
    size = (unsigned long)(size+0x3fffff)>>22;
    for(; size-->0; from_dir++, to_dir++){
        from_page_table=(0xfffff000&*from_dir);
        to_page_table=get_free_page();
```



这里的from_dir, to_dir是什么?

```
from_dir = (unsigned long *)((from>>20)&0xffc);
to_dir = (unsigned long *)((to>>20)&0xffc);
size = (unsigned long)(size+0x3fffff)>>22;
```

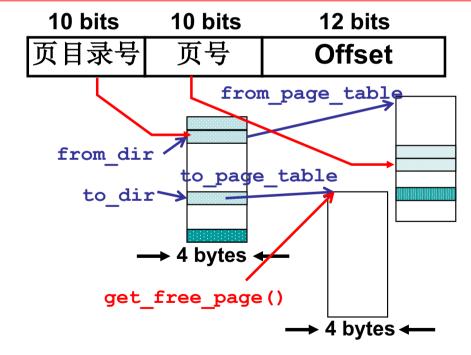
■ from是? 32位虚拟地址,这个地址的格式是否还记得?





from_page_table与to_page_table?

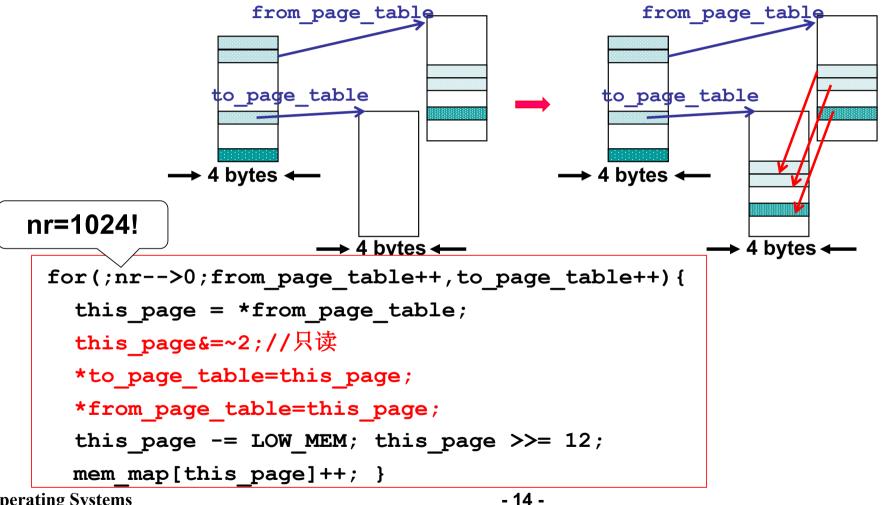
```
for(; size-->0; from_dir++, to_dir++) {
    to_page_table=get_free_page();
    *to_dir=((unsigned long)to_page_table)|7;
```



```
unsigned long
get_free_page(voi
{ register unsign
long _res asm("ax
_asm_("std; repne
scasb\n\t"
   "movl %%edx,%%eax
   "D"(mem_map+PAGIG
GES-1));
return _res; }
```

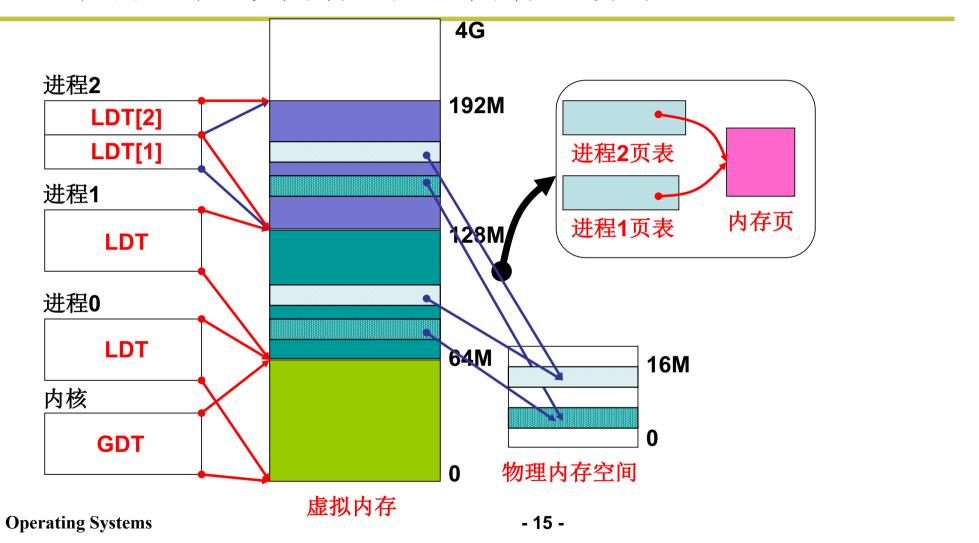


接下来干什么应该猜也猜的到...





程序、虚拟内存+物理内存的样子





*p=7? 父进程*p=7、子进程*p=8? 读写内存*p=7

